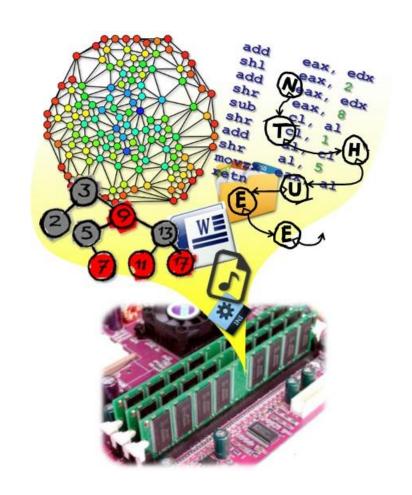
Data Structures

CH3 Stacks & Queues

Prof. Ren-Shuo Liu NTHU EE Spring 2019



Outline



- 3.1 Templates in C++
- 3.2 The stack ADT
- 3.3 The queue ADT
- 3.4 Subtyping and inheritance in C++
- 3.5 A mazing problem
- 3.6 Evaluation of expressions

Observations



- Many codes look the same for different types
 - Sorting functions that handle
 - 32-bit integers
 - 64-bit integers
 - float
 - ...
 - Sparse matrix classes that handle
 - 32-bit integers
 - 64-bit integers
 - float
 - ...

Non-Template Solutions



- Implement the same behavior over and over
 - Hard to maintain code
 - Hard to globally modify code
- Write general code for a common base type
 - Lose the benefits of compiler's type checking
 - Incurs overhead
- Use macros (#define)
 - Sacrifice readability
 - Sacrifice debuggability

Template



- Template can be instantiated to any data type
 - So called "parameterized types"
- C++ language supports
 - Template functions
 - Template classes

Template Function Example



```
void SelectionSort (int *a , const int n )
{
    for (int i = 0 ; i < n ; i++ )
    {
        int j = i;
        for (int k = i + 1 ; k < n ; k++ )
            if (a[k] < a[j]) j = k;
            swap (a[i], a[j]);
        }
}</pre>
```

```
template <class T>
void SelectionSort (T *a , const int n )
  for (int i = 0 ; i < n ; i++ )
     int j = i;
     for ( int k = i + 1; k < n; k++)
       if ( a[k] < a[j] ) j = k;
          swap (a[i], a[j]);
```

- template <class T> is identical to template <typename T>
- It is a convention to use "T", but one can use any other name

Bag Class (for integers)



```
class Bag
public:
   Bag ( int bagCapacity = 10 );  // constructor
   ~Bag( );
                                   // destructor
   int Size( ) const; // return number of elements in bag
   int Element( ) const; // return an element that is in the bag
   void pop();
                         // delete an integer in the bag
private:
   int *array;
                                             const member function
   int capacity; // capacity of array
   int top; // array position of top element
                                             Specifies that the function
};
                                             does not modify the object for
                                             which it is called.
                                             const Bag emptyBag;
                                             emptyBag.size(); //valid
                                             emptyBag.push(1); //error
```

Bag Class (for integers)

```
Bag::Bag (int bagCapacity)
:capacity ( bagCapacity )
{
   if ( capacity < 1 )
       throw "Capacity must be > 0";

   array = new int [ capacity ];
   top = -1;
}
```

Bag::~Bag ()

{ delete [] array; }

{ return top + 1; }

inline int Bag::Size() const

Initialization list



initialize member variables when they are created rather than afterwards

Bag Class (for integers)



```
inline int Bag::Element ( ) const
{
    if ( IsEmpty ( ) )
        throw "Bag is empty";
    return array [0]; // always return 0<sup>th</sup> element
void Bag::Push (const int x)
{
    if (capacity == top + 1) {
        ChangeSize1D (array, capacity, 2 * capacity);
        capacity *= 2;
    array[++top] = x;
```

Template Bag



```
template<class T>
class Bag
public:
   Bag( int bagCapacity = 10 );  // constructor
   ~Bag( );
                                   // destructor
   int Size( ) const; // return number of elements in bag
   T Element() const; // return an element that is in the bag
   void push(const T&); // add an integer into the bag
   void pop();
private:
   T *array;
    int capacity;  // capacity of array
                        // array position of top element
   int top;
```

Template Bag



```
template<class T>
Bag<T>::Bag(int bagCapacity) : capacity (bagCapacity)
    if (capacity < 1)</pre>
        throw "Capacity must be > 0";
    array = new T [capacity];
    top = -1;
template <class T>
Bag<T>::~Bag( )
{delete [ ] array;}
template <class T>
inline int Bag<T>::Size( ) const
{ return top + 1; }
```

Template Bag



```
template <class T>
inline T Bag<T>::Element( ) const
    if ( IsEmpty() )
        throw "Bag is empty";
    return array [0];
template <class T>
void Bag<T>::Push(const T& x)
    if (capacity == top + 1) {
        ChangeSize1D (array, capacity, 2 * capacity);
        capacity *= 2;
    array [++top];
```

Use of the Template Bag



```
int main()
   Bag<int> myIntBag;
   myIntBag.push(1);
   myIntBag.push(9);
   cout << myIntBag.size << endl;</pre>
   cout << myIntBag.element();</pre>
   Bag<float> myFloatBag;
   for(int i=0; i<10; i++)</pre>
       myFloatBag.push(1.0/i);
   Bag<Bag<int> > myManyIntBag;
   myManyIntBag.push(myIntBag);
   return;
```

Outline

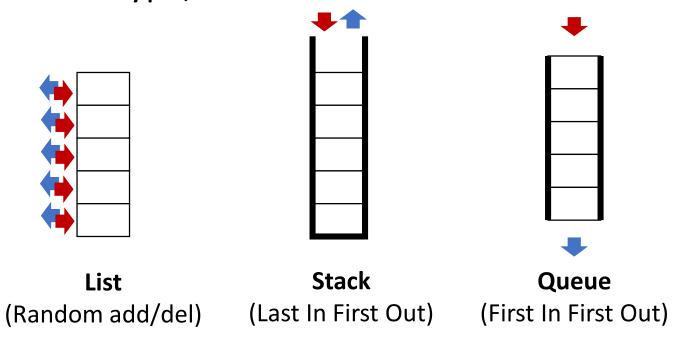


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Stacks and Queues



- Two frequently used data structures
- They are special cases of the more general data structure type, lists



Stack and Queue ADTs

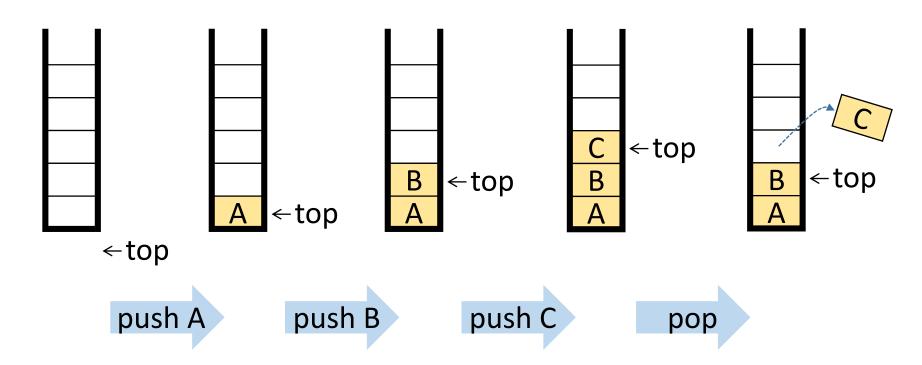


```
template < class T >
class Stack
public:
    Stack (int stackCapacity = 10);
    bool IsEmpty( ) const;
    void Push(const T& item);
    // add an item into the stack
    void Pop( );
    // delete an item
};
```

```
template < class T >
class Queue
public:
    Queue (int queueCapacity = 0);
    bool IsEmpty( ) const;
    void Push(const T& item);
    // add an item into the queue
    void Pop( );
    // delete an item
```

Stack



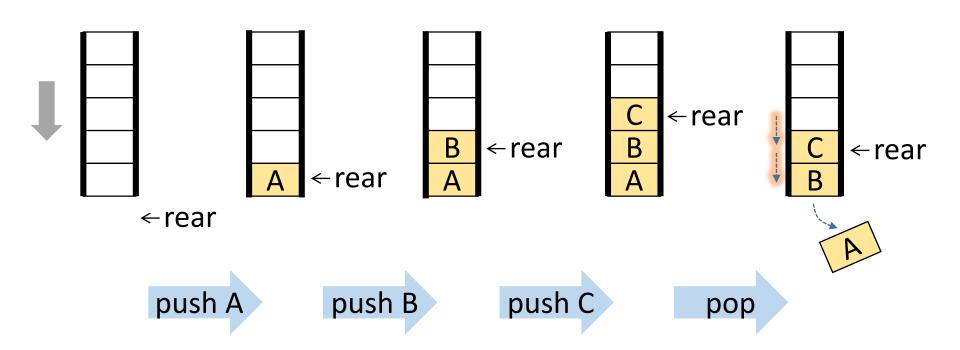


Time complexity

- push(): Θ(1)
- pop(): Θ(1)

Queue (Single Pointer)



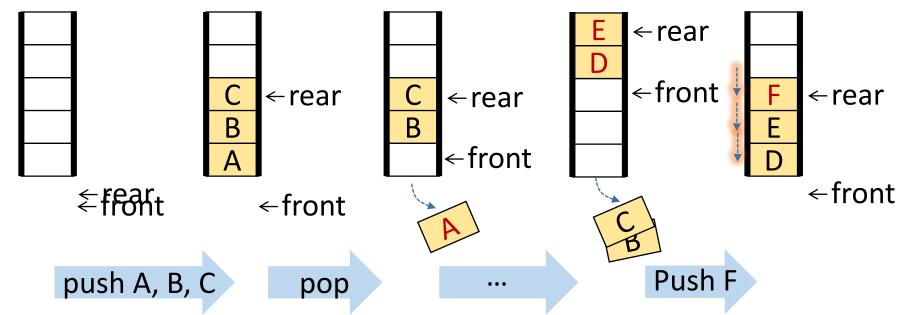


Time complexity

- push(): Θ(1)
- pop(): Θ(size)

Queue (Dual Pointers)





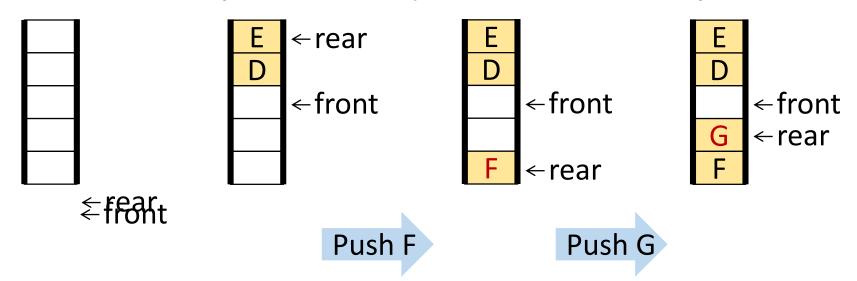
Time complexity

- push(): O(size)
 - When the rear pointer reaches the boundary and a push occurs, data need to be moved
- pop(): Θ(1)

Circular Queue



Permit the queue to wrap around the end space



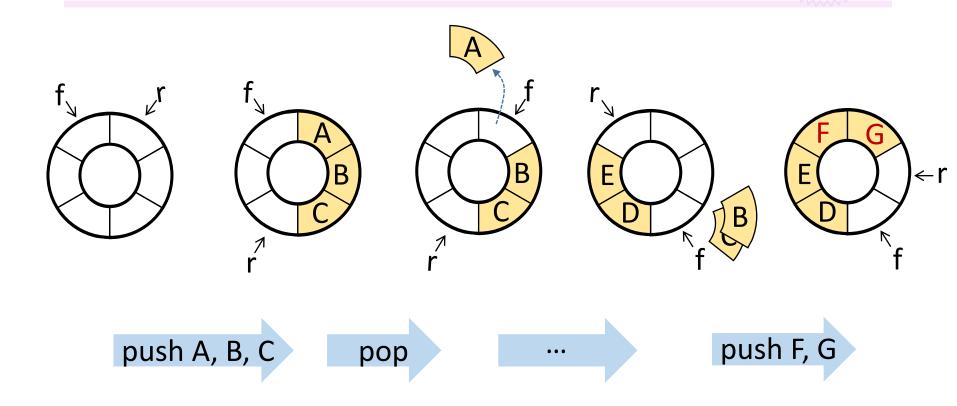
Time complexity

- push(): Θ(1)
- pop(): Θ(1)

Note that in this version of circular buffer, the position that the front pointer points to is a dead space. A slot is deliberately unused.

• Otherwise, we cannot determine whether the queue is empty or full.

Circular Queue (Circular Illustration)



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Relationships Between Things



- We abstract things on two key dimensions
 - IS-A relationship
 - HAS-A relationship
- Real world examples
 - iPhone is a smartphone. iPhone has a battery
 - NTHU is a university. NTHU has a Math department
- ADT examples
 - Rectangle is a Polygon. Rectangle has a height dimension
 - Stack is a Bag. Stack has a top pointer
 - Stack is a specialized bag that requires elements to be deleted in the LIFO order

Subtype / IS-A / Subclass

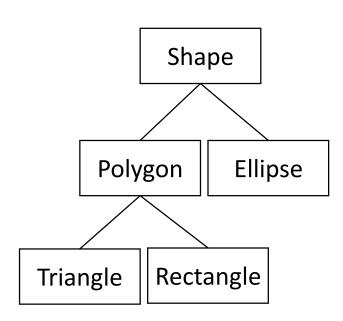


- Subtype
 - Equivalent concept to the IS-A relationship
 - Rectangle is a subtype of Polygon
 - Since C++ use classes to denote data types, subtypes are also widely referred to as subclasses
- Subtype is conceptual relationship between ADT specifications
 - "Stack IS A Bag" is true regardless of the implementation

Inheritance



- Use
 - Express IS-A relationships between classes
 - Derive a new class (derived class / sub type / sub class) from an existing class (base class)
- Objective
 - Eliminate redundant implementation
 - Members (data and functions) are by default inherited from a base class to a derived class
- Different inheritance styles
 - Public inheritance
 - Access levels (public/protected/private) of the members are also inherited
 - Protected inheritance
 - Private inheritance



Effects of Inheritance



- Stack inherits from Bag
 - Stack must redefine its constructors and destructors
 - Stack can redefine its unique data and functions (pop and top)
 - Stack inherits all the other data and functions of Bag

```
Class Bag
                                               class Stack: public Bag
public:
                                               public:
  Bag (int bagCapacity = 10);
                                                  Stack (int stackCapacity = 10);
  virtual Bag();
  virtual int Size( ) const;
                                                  ~Stack();
  virtual bool IsEmpty( ) const;
                                                  int Top( ) const;
  virtual int Element( ) const;
                                                  void Pop( );
  virtual void Push(const int);
  virtual void Pop( );
protected:
                                               protected:
  int *array;
  int top;
                                                                                       26
```

Usage Example of Derived Classes

```
Bag b(4); // invoke Bag constructor
Stack s(7); // invoke Stack constructor, which also invokes Bag constructor
b.Push(2017); // use Bag::Push()
s.Push(330); // Stack does not contains a specialized Push(), so use Bag::Push
b.Pop(); // use Bag::Pop()
s.Pop(); // Stack contains a specialized Pop() overriding Bag::Pop(), so use Stack::Pop()
```

```
Class Bag
                                               class Stack: public Bag
public:
                                               public:
  Bag (int bagCapacity = 10);
                                                  Stack (int stackCapacity = 10);
  virtual Bag( );
  virtual int Size( ) const;
                                                  ~Stack();
  virtual bool IsEmpty( ) const;
                                                  int Top() const;
  virtual int Element( ) const;
                                                  void Pop( );
  virtual void Push(const int);
  virtual void Pop( );
protected:
                                               protected:
  int *array;
  int top;
```

Syntax of Implementing Derived Classes

```
Stack::Stack(int stackCapacity)
: Bag(stackCapacity)
// explicitly call to the Bag constructor that has arguments
    // here is code specifically for creating a stack, if any
int Stack::Stack( )
       // here is code specifically for destroying a stack, if any
//Bag destructor is automatically called when a stack is destroyed
int Stack::Top( ) const
{
        if (IsEmpty( )) throw "Stack is empty.";
        return array[top];
void Stack::Pop( )
{
        if (IsEmpty( )) throw "Stack is empty. Cannot delete.";
        top--;
```

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Evaluation of Expressions



- Arithmetic expressions
 - X = (A / B) C + D * E A * C
- Boolean expressions
 - X = (A == B) | | !(C>D)
- Expressions are made up of
 - Operands: A, B, C, D, E
 - Binary arithmatic operators: +, -, *, /, %
 - Unary arithmatic operators: -
 - Relational operators: <, <=, ==, !=, >=, >
 - Binary logical operators: &&, | |
 - Unary logical operators: !
 - Delimiters: (,)

Evaluation of Expressions



- Let's focus on an arithmetic expression
 - X = A / B C + D * E A * C
- Order of operations matters
 - Let A = 4, B = C = 2, D = E = 3
 - ((4/2)-2)+(3*3)-(4*2)=0+9-8=1
 - (4/(2-2+3))*(3-4)*2 = (4/3)*(-1)*2 = -2.666...
- How can computers uniquely define the order of an expression?

Priority of Operators



Priority is introduced to help defining the order

| | Priority | Operator | |
|--------|----------|--------------------|---|
| High | 1 | Unary minus (負號),! | |
| iligii | 2 | *, /, % | _ |
| | 3 | +, - | _ |
| | 4 | <, <=, >=, > | |
| | 5 | = =, != | |
| | 6 | && | |
| | 7 | | |

* and / have higher priority than + and -

- Tie break rule: left to right
- Example Two operators compete for one operand '/' and '*' win
 - A / B C + D * E A * C \rightarrow (A/B) C + (D*E) (A*C)
 - A/B*C/D \rightarrow ((A/B)*C)/D Tie-break rule

Infix and Postfix Notations



- Infix
 - Binary operators come in-between their operands
 - A*B/C
- Postfix
 - Binary operators appear after their operands
- Examples

• Infix:

A*B

• Postfix:

AB*

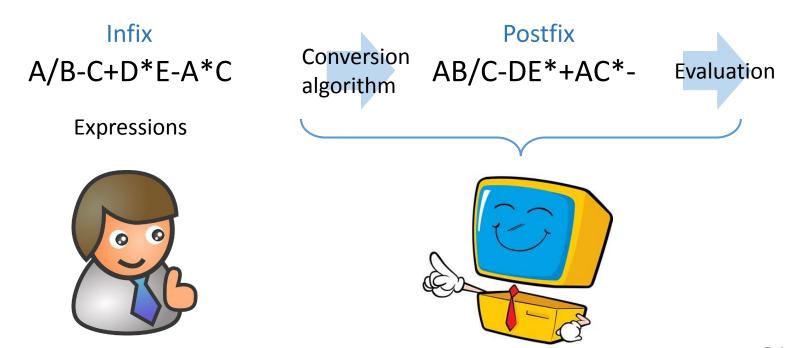
A/B-C+D*E-A*C

AB/C-DE*+AC*-

Two Essential Algorithms



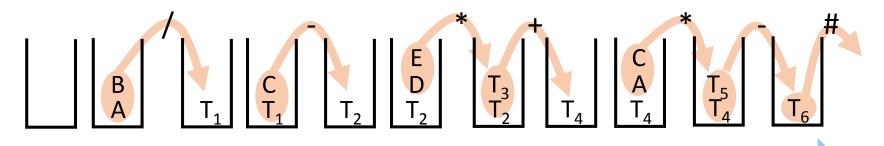
- Combining two algorithms enables computers to handle human-written expressions
 - Infix-to-Postfix conversion
 - Postfix evaluation (just mentioned)



Postfix Evaluation



- Rules
 - Left to right scan
 - Push operands onto a stack
 - Evaluate operators using the required number of operands from the stack
 - Push the evaluating results onto the stack again
- AB/C-DE*+AC*-# (# denotes the end of an expression)



Advantages of Postfix Notation

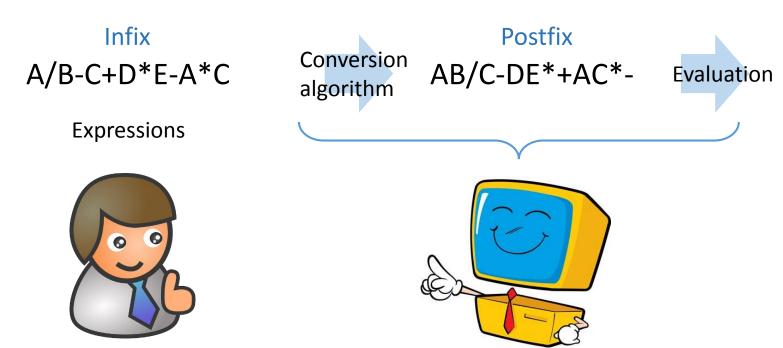
- Evaluation is simpler than infix notation
 - The need for parenthesis is gone
 - The need for operator priority is gone

```
void Eval(Expression e)
  Stack<Token> stack; // initialize a stack
  for (Token x = NextToken(e); x! = end of expression; x=NextToken(e))
    if (x is an operand) {
       stak.Push(x)
     } else {// x is an operator
       pop from the stack the correct number of operands for the operator;
       perform the operation x and store the result (if any) onto the stack;
```

Two Essential Algorithms



- Combining two algorithms enables computers to handle human-written expressions
 - Infix-to-Postfix conversion
 - Postfix evaluation (just mentioned)



Infix to Postfix Conversion

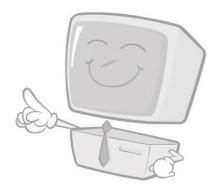


- Observations
 - Number of operands and operators do not change
 - Order of operands (A, B, C...) do not change









Infix to Postfix Conversion

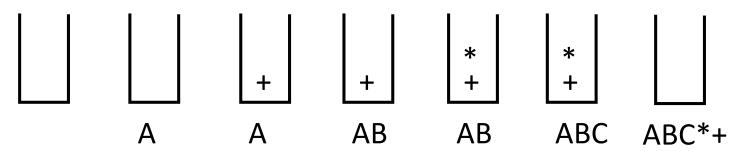


- Method 1
 - Fully parenthesize the expression (based on the operator priorities)
 - Move all operators so that they replace their corresponding right parentheses
 - Delete all parentheses

Infix to Postfix Conversion



- Stack-based algorithm
 - Create a stack
 - Scan the input infix expression left to right
 - Bypass each incoming operand to the output
 - For each incoming operator
 - First, continuously pop from the stack an operator (the top) if the top has equal or lower priority than the incoming operator
 - Then, push the incoming operator onto the stack
 - Pop all operators upon the end of an expression
- Example: A + B * C



Parentheses Handling



- We want the stack algorithm to handle parentheses similarly to handling operators
- Specialized rules for left parenthesis
 - Incoming left parenthesis has the highest priority (i.e., always gets pushed onto the stack)
 - In-coming priority (ICP) = 0
 - Only gets popped from the stack upon a matched right parenthesis
 - Otherwise, behaves as one with the lowest priority
 - In-stack priority (ISP) = 8

| Priority | Operator | |
|----------|--------------------|--|
| 0 | In-coming (| |
| 1 | Unary minus (負號),! | |
| 2 | *,/,% | |
| 3 | +, - | |
| 4 | <, <=, >=, > | |
| 5 | = =, != | |
| 6 | && | |
| 7 | | |
| 8 | In-stack (| |

Example



• A*(B+C)/D

| Incoming token | Stack | Output | Note |
|----------------|-------|--------|------|
| Empty | Empty | Empty | |
| Α | | | |
| * | | | |
| (| | | |
| В | | | |
| + | | | |
| С | | | |
|) | | | |
| / | | | |
| D | | | |
| Done | | | |

Example



• A*(B+C)/D

| Incoming token | Stack | Output | Note |
|----------------|-------|---------|-------------------------------|
| Empty | Empty | Empty | |
| Α | Empty | Α | Bypass operands |
| * | * | | |
| (| *(| | ICP('(') higher than ISP('*') |
| В | *(| AB | Bypass operands |
| + | *(+ | | ICP('+') higher than ISP('(') |
| С | *(+ | ABC | Bypass operands |
|) | * | ABC+ | Pop until a left parenthesis |
| / | / | ABC+* | ICP('/') == ISP('*') |
| D | / | ABC+*D | Bypass operands |
| Done | Empty | ABC+*D/ | Pop all operators |

Recap Infix to Postfix Conversion



- A + B * C
- 1 left to right scan

incoming operands always bypasses the stack

ABC

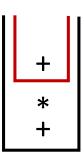
4 incoming operators always enters the stack

3 continuously pop top operator from the stack if it has equal or higher priority than the incoming ones

Recap Parenthesis Handling



- Incoming left parenthesis has the highest priority
 - It always enters the stack without popping any stacked operator
- In-stack left parenthesis has the lowest priority
 - It never gets popped from the stack until the right parenthesis appears
- Different perspective ¹
 - Left parenthesis creates an isolated, nested stack
 - Right parenthesis cleans up a nested stack



A+B*(C+D)

1. Contributed by Mr. 陳德暉 (101061132) on April 2, 2015

Infix to Postfix Algorithm



```
void Postfix(Expression e)
    Stack<Token>stack; // initialize the stack
    stack.Push('#');
    for (Token x = NextToken(e); x != '#'; x = NextToken(e))
        if (x is an operand) cout << x;</pre>
        else if (x == ')' ) { // pop until a left parenthesis
             for (;stack.Top( ) != '('; stack.Pop( ))
                 cout << stack.Top( );</pre>
             stack.Pop( ); // remove the left parenthesis
        } else { // x is a operator
             for (; isp(stack.Top( )) <= icp(x); stack.Pop( ))</pre>
                 cout << stack.Top( );</pre>
                                               // higher or equal priority
             stack.Push(x);
    // end of expression; empty the stack
    for ( ; !stack.IsEmpty( ); cout << stack.Top( ), stack.Pop( ));</pre>
    cout << endl;</pre>
```

Limitations of the Current Algorithm

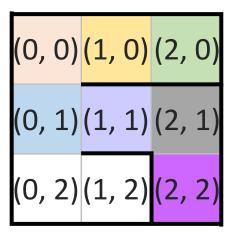
- Characters to tokens conversion (parser)
 - Energy = Mass * LightSpeed * LightSpeed
 - Area = 3.14*radius1*radius2
- Grammar
 - X = A B + -A
 computers need rules to differentiate the two minus
 symbols; Otherwise, the aforementioned postfix
 algorithm cannot work correctly.
- More techniques are available in a compiler course

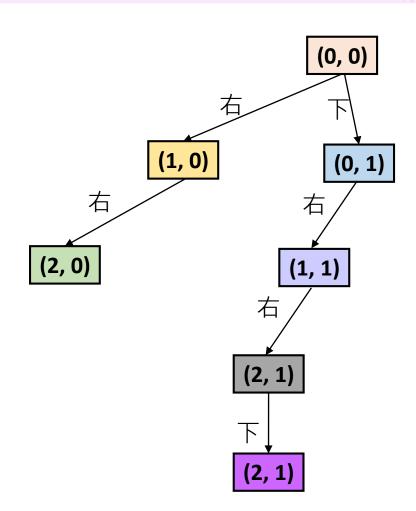
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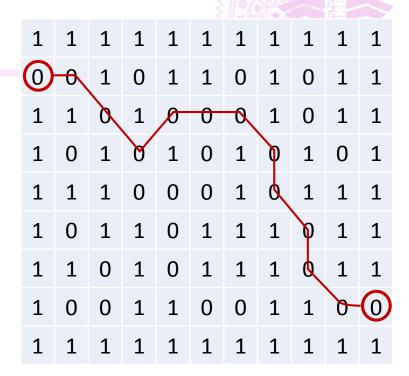
How a Computer Traverses a Maze

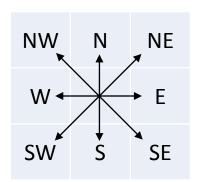




A Mazing Problem

- 2D array maze representation
 - '1' implies a blocked direction
 - '0' means otherwise
 - Borders are surrounded by '1'
- Allowable moves
 - Non-blocked squares of the eight neighboring squares
- How can a program get through the maze?



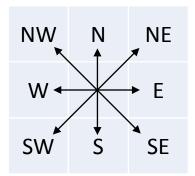


Algorithm (Pseudo Code)



```
struct Offsets
    int di, dj;
enum directions {N, NE, E,
SE, S, SW, W, NW};
Offsets move[8];
struct Items
    int x, y, dir;
```

| q | move[q].di | move[q].dj |
|----|------------|------------|
| N | -1 | 0 |
| NE | -1 | 1 |
| Е | 0 | 1 |
| SE | 1 | 1 |
| S | 1 | 0 |
| SW | 1 | -1 |
| W | 0 | -1 |
| NW | -1 | -1 |



Algorithm ()



```
initialize a stack // remember the point to retract
add the starting point, (0, 1, E), to the stack
while (the stack is not empty) { // there are still unexplored points
  (i, j, dir) = the top of the stack;
  remove the top of the stack;
  while(there are more move from (i, j)){
    (g, h) = nextPoint((i, j), dir);
    if ((g == m) && (h == p)) return success;
    if ((!maze [g][h]) && (!mark [g][h])) {
        dir = Next(dir);
        add (i, j, dir) to the stack; // prepare for a dead end
        (i, j, dir) = (g, h, N); // move to (g, h), start from dir N
        mark[i][j] = 1;

    Each position can be visited

                                               at most once.

    At most eight valid moves

                                               from each position
cout << "No path in maze." << endl;</pre>
                                            → O(size of the array) time
```

Algorithm (Pseudo Code)



```
void driver()
   if (findPath(0, 1))
       cout << "Success" << endl;</pre>
   else
       cout << "No path in maze." << endl;</pre>
   return;
bool findPath(int i, int j) // find a path starting from (i, j)
    for (all eight directions) { // explore all directions
        (g, h) = (i, j) + direction;
        if ((g == m) && (h == p)) return true;
        if ((!maze [g][h]) && (!mark [g][h])) {
            findPath(g, h); // keep finding a path...
    return false;
```

Stack Provided by C++ Library



```
#include <iostream>
#include <stack>
using namespace std;
int main()
    stack<int> s;
    for(int i=0; i < 5; i++){
        s.push(i);
    while(!s.empty())
        cout << s.size() << " ";</pre>
        cout << s.top() << endl;</pre>
        s.pop();
```

output

```
5 4
4 3
3 2
2 1
1 0
```

Reference of STL's Stack

http://en.cppreference.com/w/cpp/container/stack