

# CS513 HW4

1. **Unit Ball Graph Dominating Set (3D):** Given a set of points in 3D, a unit ball graph is defined by connecting points of distance at most 1 away (a dominating set  $D$  is a subset of vertices such that every vertex in the graph is either in  $D$  or adjacent to a vertex in  $D$ ). Design a polynomial time algorithm to approximate the minimum dominating set by a constant factor.

## 2. Load Balancing (Greedy Makespan)

We are given  $m$  identical machines  $M_1, \dots, M_m$  and a set of  $n$  jobs. Each job  $j$  has a processing time  $t_j > 0$ . We need to assign each job to exactly one machine. Once assigned, a machine processes its jobs sequentially. Let  $L_i$  be the total load (sum of processing times) assigned to machine  $M_i$ . The objective is to minimize the **Makespan**, which is defined as the maximum load among all machines:

$$\text{Makespan} = \max_{i=1\dots m} L_i$$

(Intuitively, we want to finish all jobs as early as possible, so we need to minimize the completion time of the busiest machine.)

Now, consider the following greedy algorithm for this problem: Process the jobs in an arbitrary order  $1, \dots, n$ . For each job  $j$ :

1. Check the current load of every machine  $M_1, \dots, M_m$ .
2. Assign job  $j$  to the machine  $M_i$  that currently has the **minimum load**.
3. Update the load of  $M_i$ :  $L_i \leftarrow L_i + t_j$ .

### Questions:

- (a) Show an example with  $m = 2$  machines where this Greedy algorithm is not optimal.
- (b) What is the approximation ratio of this algorithm? Prove your answer.

## 3. Maximum 3D Matching

Given disjoint sets  $X, Y, Z$  (each of size  $n$ ) and a set  $T \subseteq X \times Y \times Z$  of triplets. A subset  $M \subseteq T$  is a *3D matching* if each element of  $X, Y, Z$  appears in at most one triplet in  $M$ . The goal is to find a 3D matching  $M$  of maximum size. Give a polynomial time algorithm to find a solution that is at least  $1/3$  of the optimal solution.

## 4. Facility Location (Supermarket Placement)

Given a set of  $n$  customers and a set of potential locations  $S$  for supermarkets, decide where to open the supermarkets to minimize a total cost function. The cost consists of two parts:

- **Opening Cost:** A cost  $f_i$  if we open a market at location  $s_i \in S$ .
- **Service Cost:** For each customer  $j$ , if they are served by a store at  $s_i$ , the cost is  $d_{ji}$ . Each customer connects to the closest open store.

Design an  $O(\log n)$  approximation algorithm for this problem.