The Idea Behind Krylov Methods [1]

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Do they make sense?

Can they be practical?

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A few definitions before we begin...

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A matrix, A, is diagonalisable if there are matrices S and Λ such that $A = S\Lambda S^{-1}$, where Λ is diagonal.

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"a Krylov method solves $A\mathbf{x} = \mathbf{b}$ by repeatedly performing matrix-vector multiplications involving A" [1]



Krylov space of dimension k:

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Residual: $\mathbf{r}_k \equiv \mathbf{b} - A\mathbf{x}_k$

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General Krylov: number of iterations required depends upon eigenvalues

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The future of the talk...

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...BOARD...

Summarising Theorem:

Theorem ([1] p6)

If the minimal polynomial of the nonsingular matrix A has degree m, then the solution to $A\mathbf{x} = \mathbf{b}$ lies in the space $\mathcal{K}_m(A, \mathbf{b})$.

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How could this be better?

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Ammended Theorem:

Theorem ([1] p6 (adapted))

If matrix A is nonsingular and diagonalisable with d distinct eigenvalues, then the solution to $A\mathbf{x} = \mathbf{b}$ lies in the space $\mathcal{K}_d(A, \mathbf{b})$.

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Preconditioning

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- ► Good case: *PA* diagonalisable with small *d*
- ► Realistic expectation: PA has a few 'clusters' of eigenvalues

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$$\|\mathbf{r}_k\| \le \|S\| \|S^{-1}\| \|\mathbf{b}\| \rho^k.$$

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Numerical Example

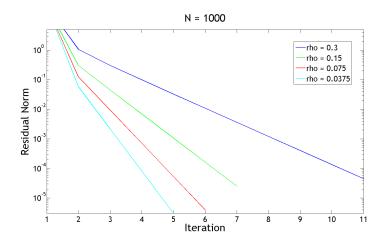
Take

$$A = \begin{pmatrix} 1 & -\rho & & & \\ -\rho & 1 & -\rho & & & \\ & -\rho & 1 & -\rho & & \\ & & \ddots & \ddots & \ddots \end{pmatrix}^{N \times N}$$

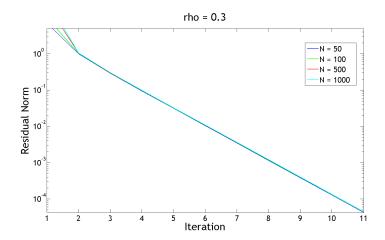
which has eigenvalues

$$\lambda_i = 2\cos\left(\frac{i\pi}{N+1}\right)$$

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S. L. Campbell, I. C. F. Ipsen, C. T. Kelley, and C. D. Meyer. Gmres and the minimal polynomial. *BIT*, 36:32–43, 1996.