

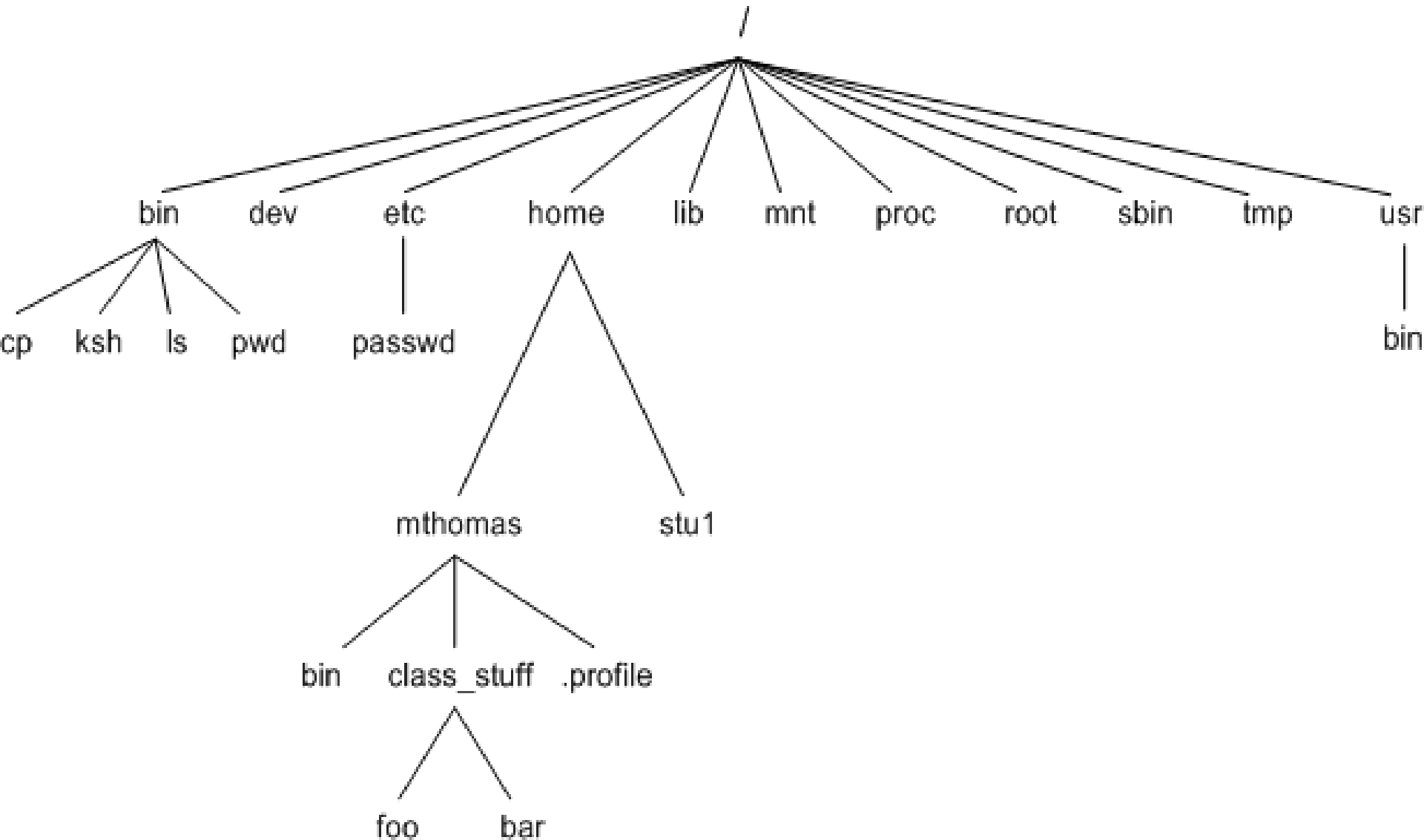
CS-3103 : Operating Systems : Sec-A (NB) : **File Systems.....**

OPERATING SYSTEM

A hand holding a black pen is pointing at the word 'SYSTEM' in the large red text. The background is a light gray with various computer-related terms like 'LAPTOP', 'SMARTPHONE', 'SERVER', 'DATA', 'COM', 'ACCESS', 'KST', 'SYN', 'TECHNO', 'ITY', 'TOP', 'FILE', 'SMARTPHONE', 'SERVER', 'DATA', 'COM', 'ACCESS', 'KST', 'SYN', 'TECHNO', 'ITY', 'TOP', 'FILE' in a lighter gray font.

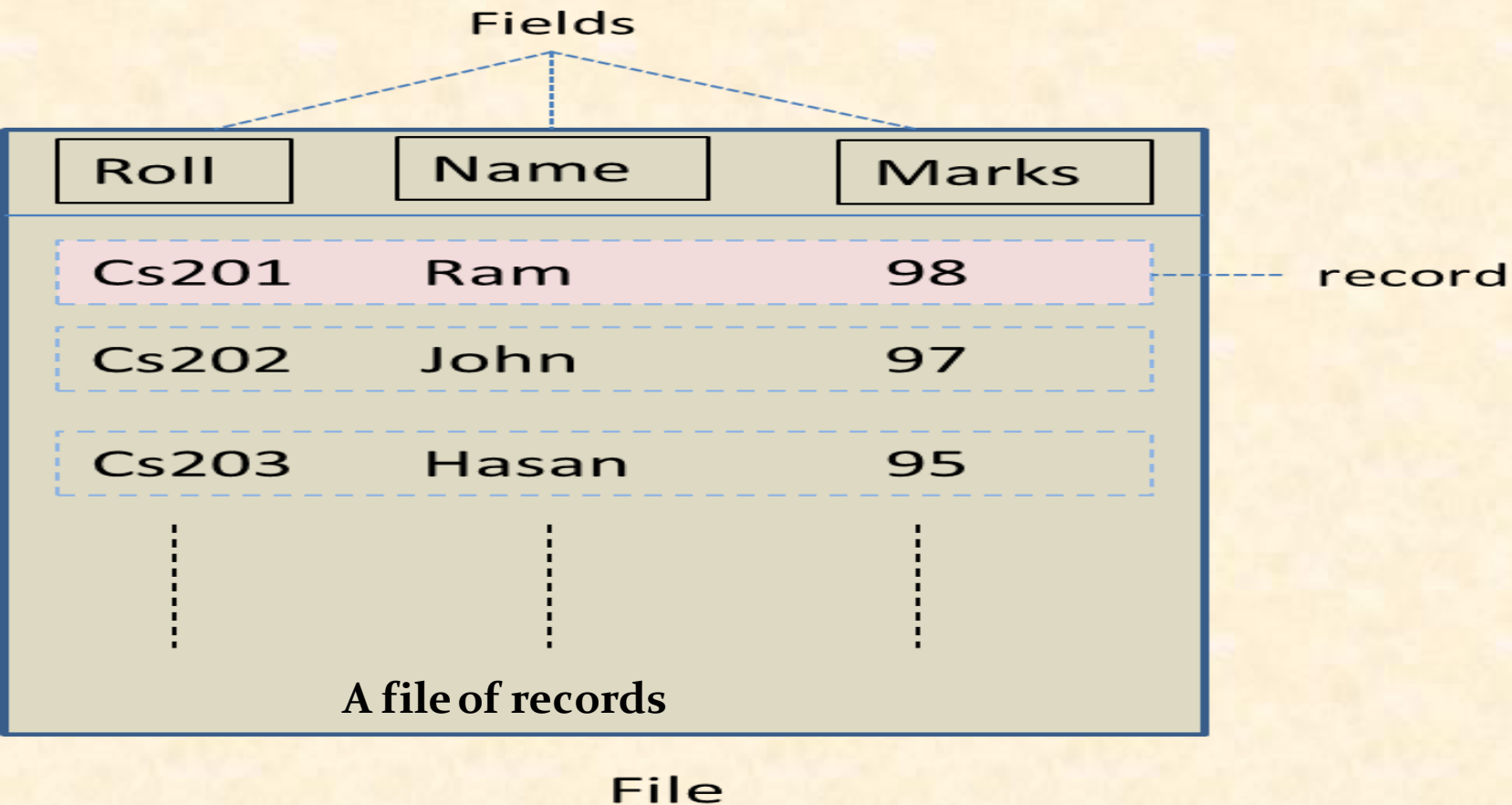
Computer Operating Systems: OS Families for Computers

- When information is stored on permanent magnetic media, like a tape or hard disk, it is called a **file**.
- Ex: The source code of a C program is an instance of a text file.



File Organisation

- A set of data items that collectively define an entity is called a **record**.
- The entity can be a person, event, place, etc. An individual data item within a record is called a **field**.
- The field that uniquely identifies a record within a file is called a **key field**.



Serial Files

- *Records are stored in the order of their arrival.*
- **Example:** In a booking window, the records of the customers are stored in a file in the order of arrival of the customers.
- A serial file is useful for situations where almost all records of the file need to be processed.

Roll	Name	Time
Cs205	Sham	10:00: 07
Cs213	Rohit	10:01:32
Cs208	Mary	10:30:21
⋮	⋮	⋮

Serial File

Sequential Files

- *The records in a sequential file are stored in a meaningful order of a key field.*
- The only way to reach a particular record is by searching the file from the beginning and skipping the records till the desired record is reached.

Roll	Name	Marks
Cs201	Ram	98
Cs202	John	97
Cs203	Hasan	95
⋮	⋮	⋮

A sequential file

Disadvantages of Sequential Files

- The main disadvantage is that records cannot be deleted, inserted or modified in an existing sequential file.
- In order to overcome this limitation, the existing sequential file is given the name Old Master File (OMF).
- The records that need to be changed are read from the OMF into main memory.
- The records are then updated and stored, in the same sequence, in a separate file called the transaction file (TF).

Transaction Processing in Sequential Files

Old Master File

Roll	Name	Marks
Cs201	Ram	98
Cs202	John	97
Cs203	Hasan	95
.....

File

For the next run the new master file becomes old master file

New Master File

Transaction File

Roll	Name	Marks
Cs201	Ram	98
Cs202	John	97
Cs203	Hasan	95
.....

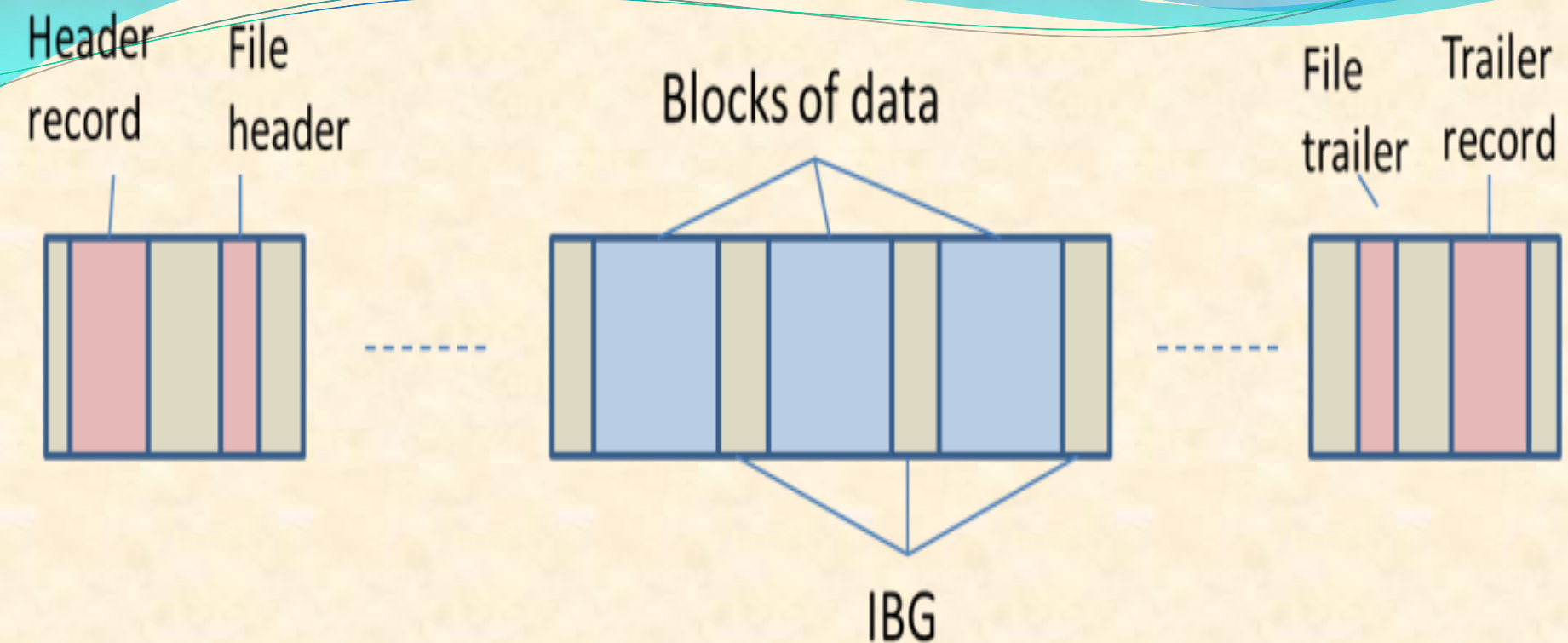
File

Update

Roll	Name	Marks
Cs201	Ram	98
Cs202	John	97
Cs203	Hasan	95
.....

File

A Sequential File on a Tape



- A file held on a magnetic tape has the following three components:
- ❑ A header record, containing identification information about the tape.
 - ❑ Blocks of data, where a block may contain one or more records depending upon the block size supported by the OS.
 - ❑ A trailer record marking the end of the tape.
 - ❑ Between two consecutive blocks of data, there is an interblock gap (IBG), that marks the end and beginning of blocks and allows the tape drive to detect the start and end of a block.

Advantages of Sequential Files

- Simple design
- One key field is enough to identify and search a record within the file
- Very efficient organisation
- Inexpensive way of storing data on tapes

Disadvantages of Sequential Files

- More time is taken to access the desired record.
- The entire file is read to reach the record.
- Processing and updating of records is performed in a separate file called the transaction file.
- Transactions have to be sorted before processing is done.
- The file is never up to date.
- Overall processing is slow.

Indexed Sequential Files

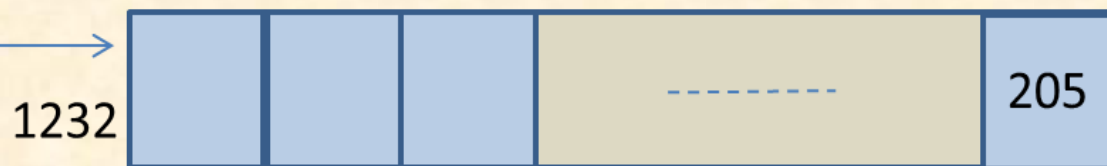
- The records are stored sequentially on separate blocks on the hard disk.
- Each block contains a specified number of records.
- A block can be reached through an entry into an index.
- The index is also stored on the disk.
- Each entry into the index contains only two items:
 - The key field of the last record present in a block
 - The starting disk address of the block

Indexed Sequential Files

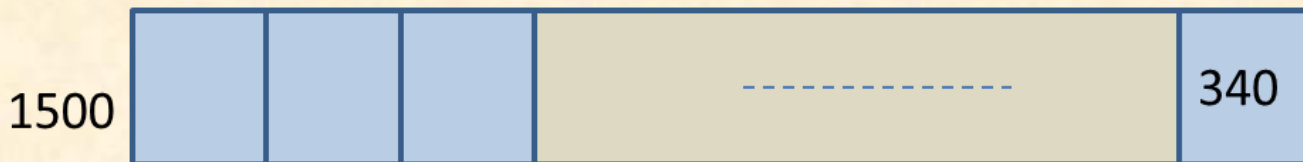
File index

Key field	Start block address
205	1232
340	1500
515	1870
840	2010
	3000

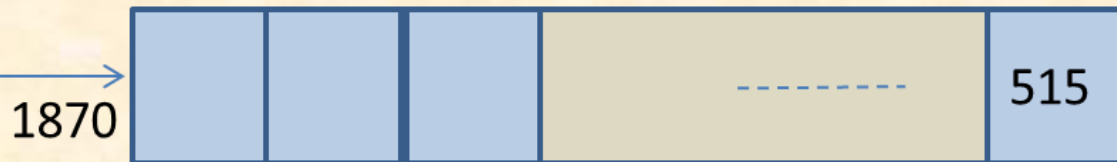
Block 1



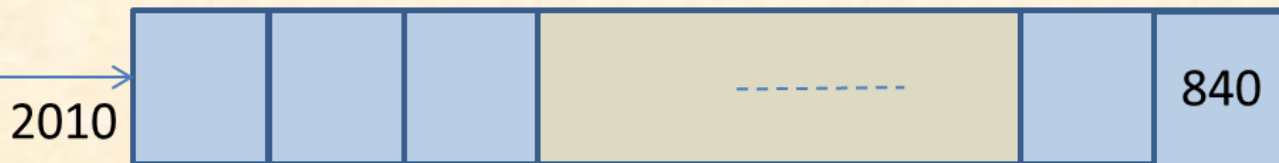
Block 2



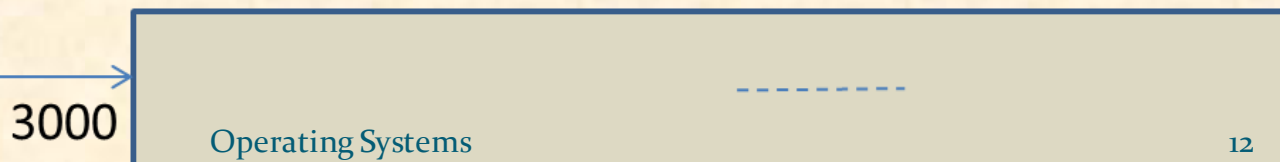
Block 3



Block 4



Overflow block



Advantages of Indexed Sequential Files

- Access to a record is faster than in sequential files
- Provides a good combination of direct and sequential processing of the same file
- Interactive processing is possible
- Very suitable for disk storage

Disadvantages of Indexed Sequential Files

- Indexed sequential files generate overflows because of insertion operations.
- During deletion, vacant storage locations are created in blocks.
- Extra storage space is required to store the 'index' of the file.
- Frequently, the file needs to be reorganised to remove the blank storage locations.

Direct File Organisation

- Allows records to be read directly from or written onto the file without proceeding from the beginning.
- For a given record, an address is generated by applying a mathematical function on the key field of the record.
- This mathematical function is called the **hash function**.
- The hash function generates a unique disk address for the given record and thereafter, the record is stored at that address.



Direct File Organisation

Roll	Name	Marks
Cs201	Ram	98
Cs202	John	97
Cs203	Hasan	95
⋮	⋮	⋮

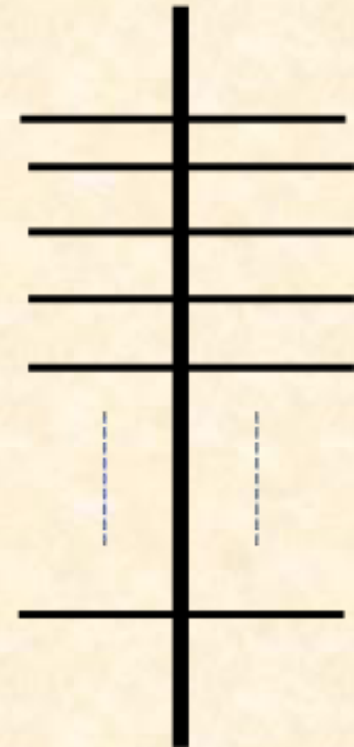
File

$K = \text{Cs203}$

$F(k)$

Hash Function

Disk address
of record



Advantages of Direct Files

- Any record can be directly accessed.
- File processing and updating are done in the same file without the creation of a New Master File (NMF).
- Speed of processing of large files is very quick.
- The files can be kept up-do-date because online updating of records is possible.
- Concurrent processing of many files is possible.



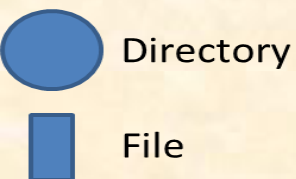
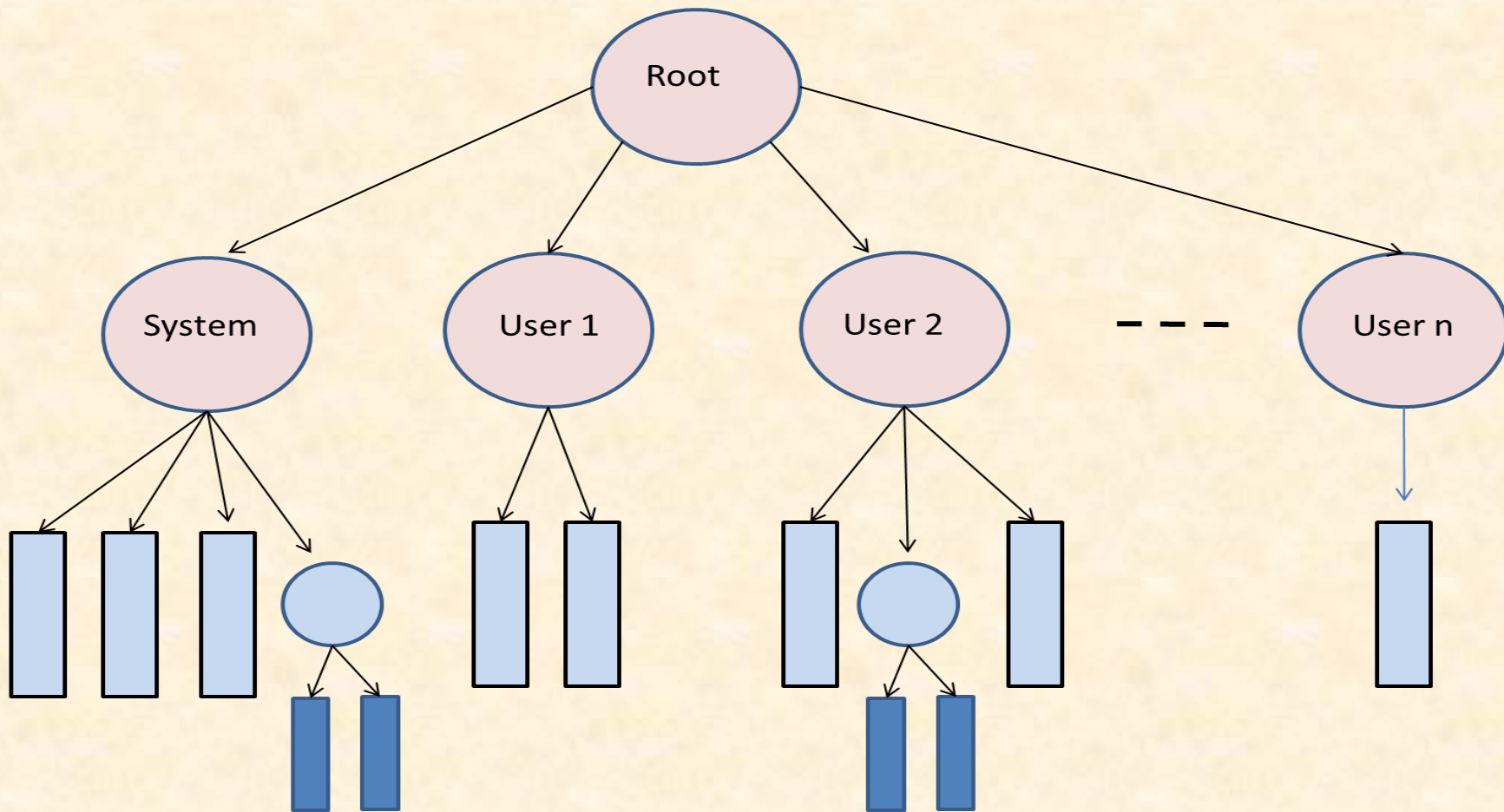
Disadvantages of Direct Files

- Direct or random files are prone to security and backup problems.
- A poor hash function may lead to excessive overflows.
- Poor access time and memory overheads.



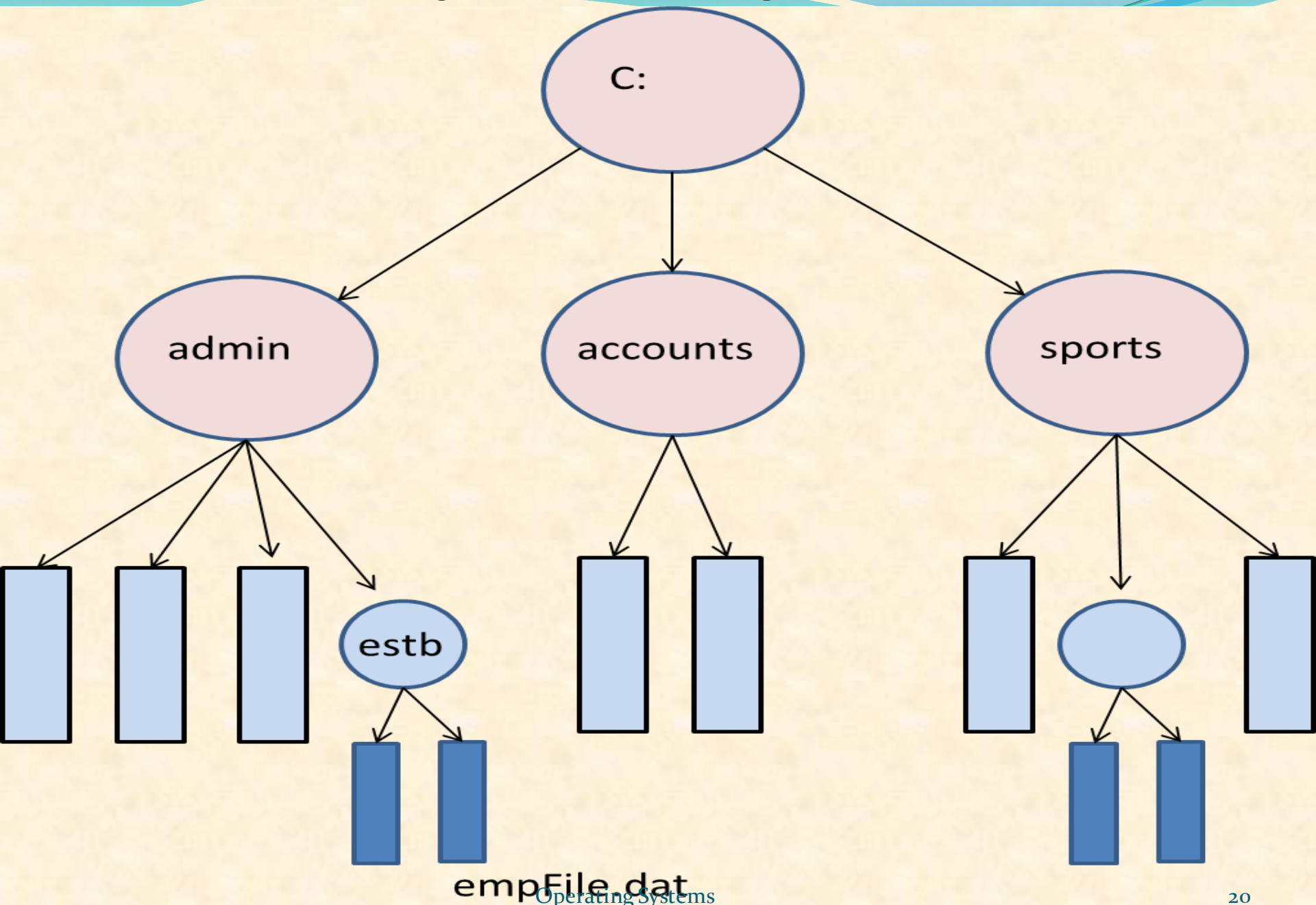
Files and Directories

- On a particular logical drive, the files are further divided into groups.
- Each group of files is stored in a separate area called the directory.



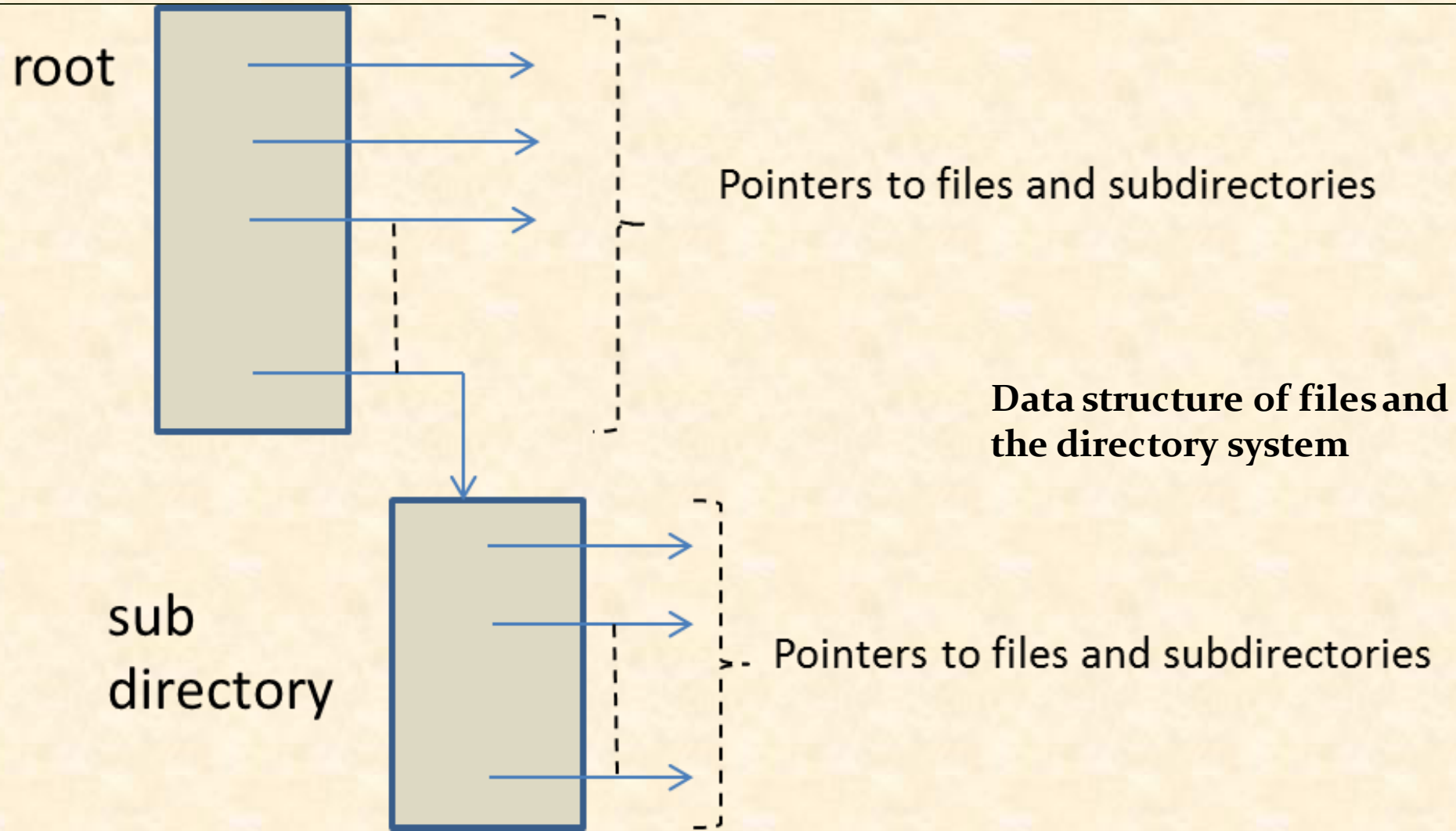
Files and directories

A Sample Directory Structure



File Meta-data

- The data structure of a directory is an array of records.
- Each record has many fields, including a pointer that points either to a file or to a subdirectory.



File System Functions

- Open
- Close
- Read
- Write
- Erase
- Append
 - Seek
- Rewind
- Rename
- Delete

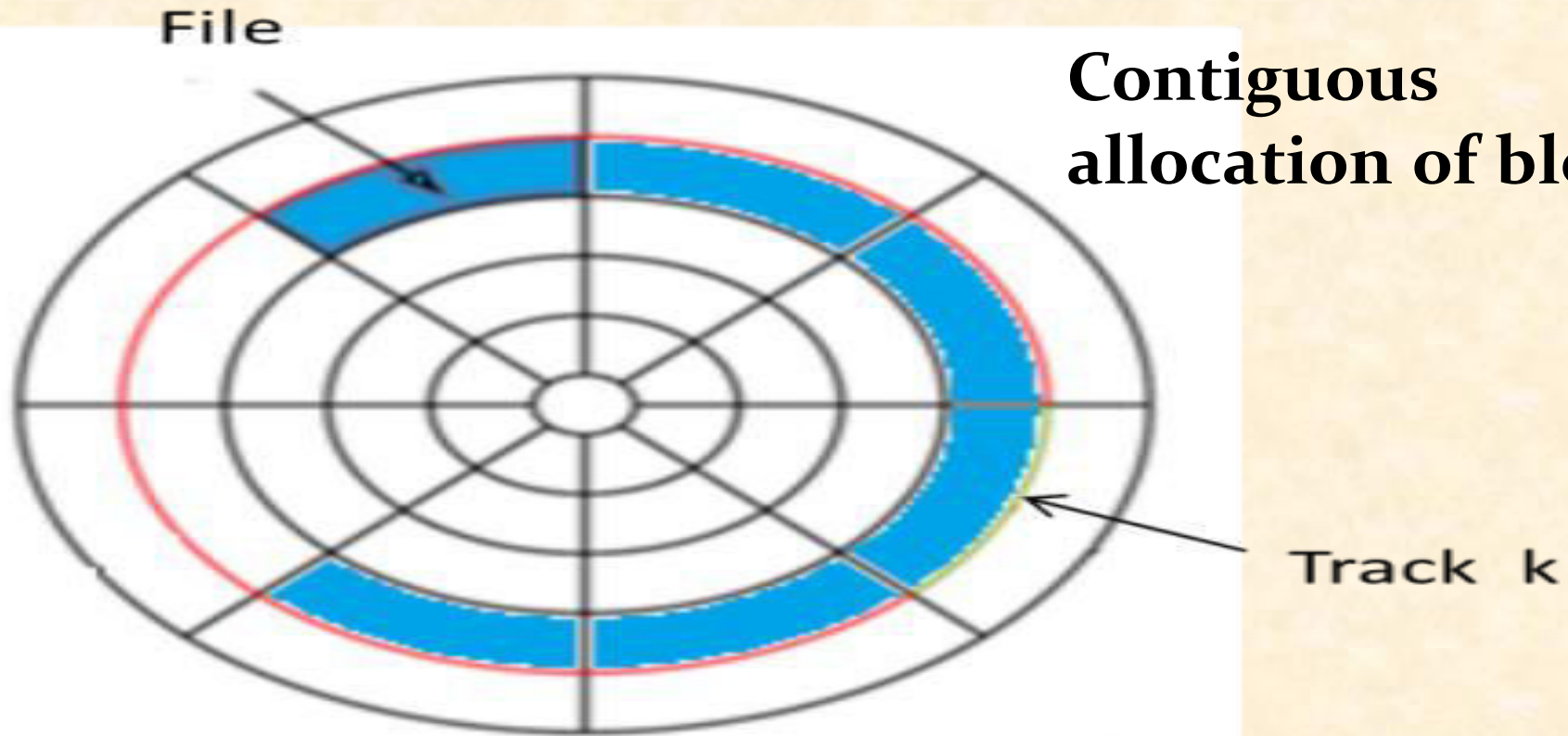
Allocation of Disk Space to Files

- The file system allocates disk space to files as per two broad policies:
 - Contiguous allocation
 - Non-contiguous allocation



Contiguous Allocation

- The space a file is allocated is in the form of contiguous blocks on the disk.
- If the file spans over some sectors, then contiguous sectors are located on the disk to store the file.
- If the file is larger than a track, then it spills over to the next platter on the same cylinder, and so on.



Advantages of Contiguous Allocation

- The size of the file remains fairly constant, contiguous allocation is efficient.
- The files can be processed at a fast rate.
- The file manager needs to know the address of the first block of the file. The address of the subsequent blocks can be easily computed by knowing the size of the block.



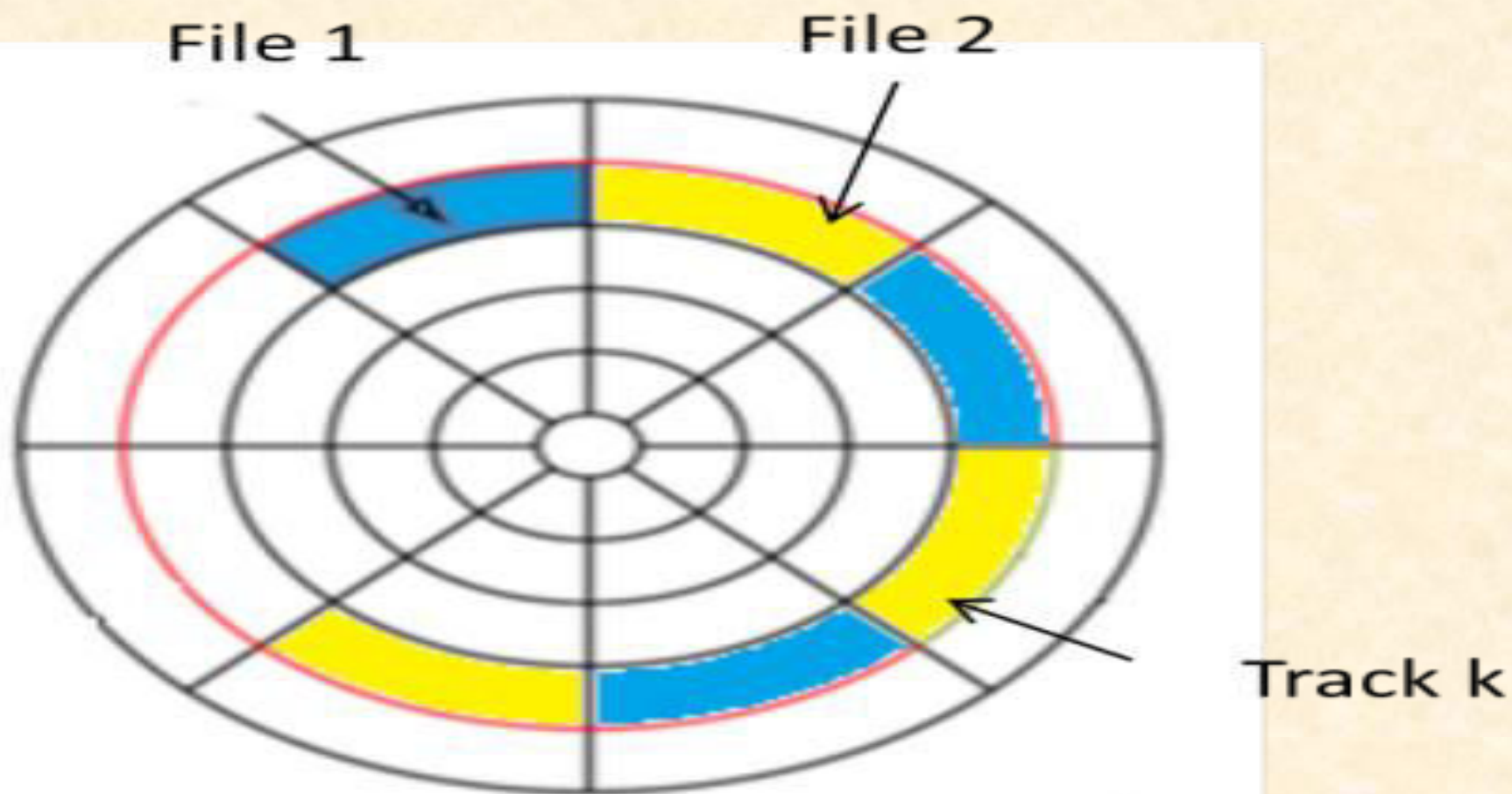
Disadvantages of Contiguous File Allocation

- File size varies with time, and it becomes difficult to find contiguous space that is large enough on the same track or cylinder to store the file.
- When a file becomes larger than the allocated space on the disk, it is moved to a new location. The changes are accordingly updated in the file directory entry, making the system even slower.
- Memory fragmentation
- Storage compaction



Interleaved Allocation

- Some disk controllers cannot read two physically adjacent sectors.
- By the time it detects the end of a sector, the R/W head crosses the beginning of the next sector.
- In such cases, the blocks of files are stored in interleaved fashion.



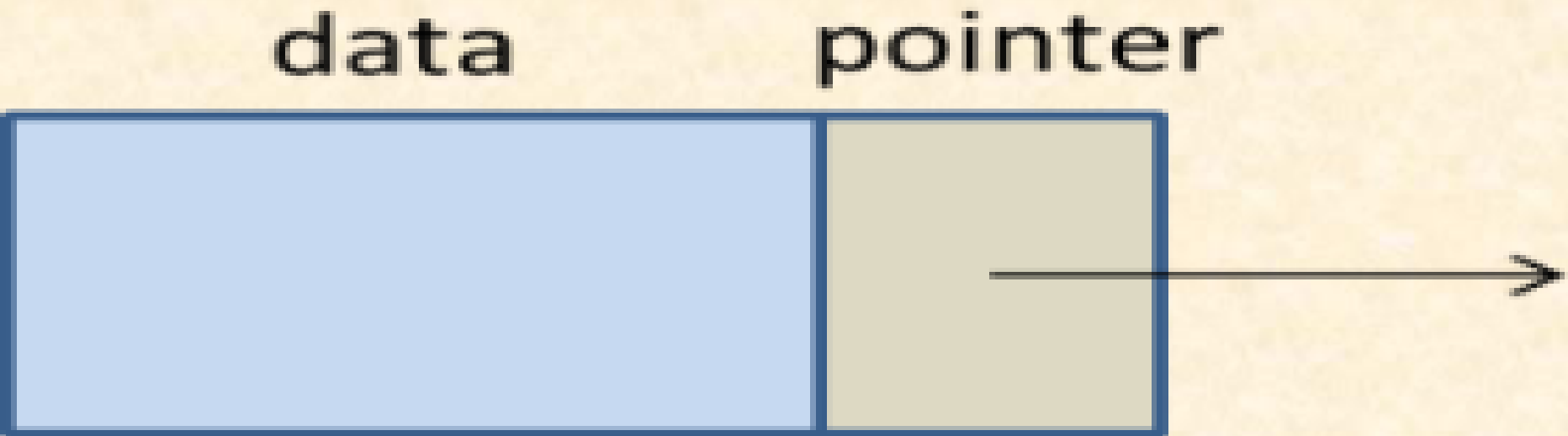
Non-contiguous Allocation

- Modern disks avoid contiguous allocation of space.
- The allocation of space to a file block is made as per the availability of space on the same platter or cylinder.
- Some non-contiguous allocation strategies are :
 - Linked allocation (block chaining)
 - Indexed allocation
 - Block-oriented file mapping



Linked Allocation (Block Chaining)

- The file on the disk is organised in the form of blocks.
- Each block contains two portions:
 - Data part
 - Pointer to the next block
- The data part contains the record(s) of the file and the pointer points to the next block of the file.

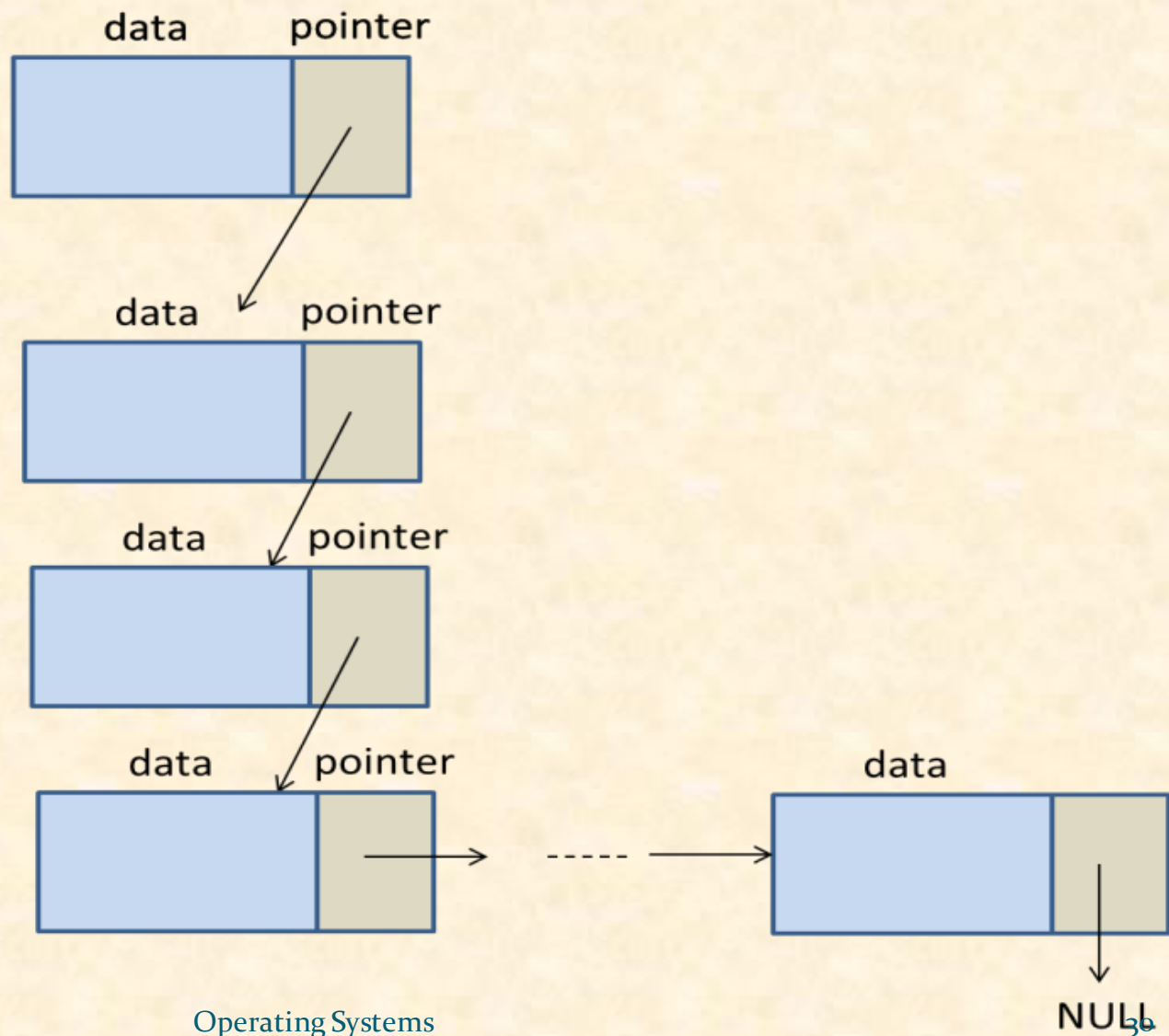


A linked file block

Block Chaining

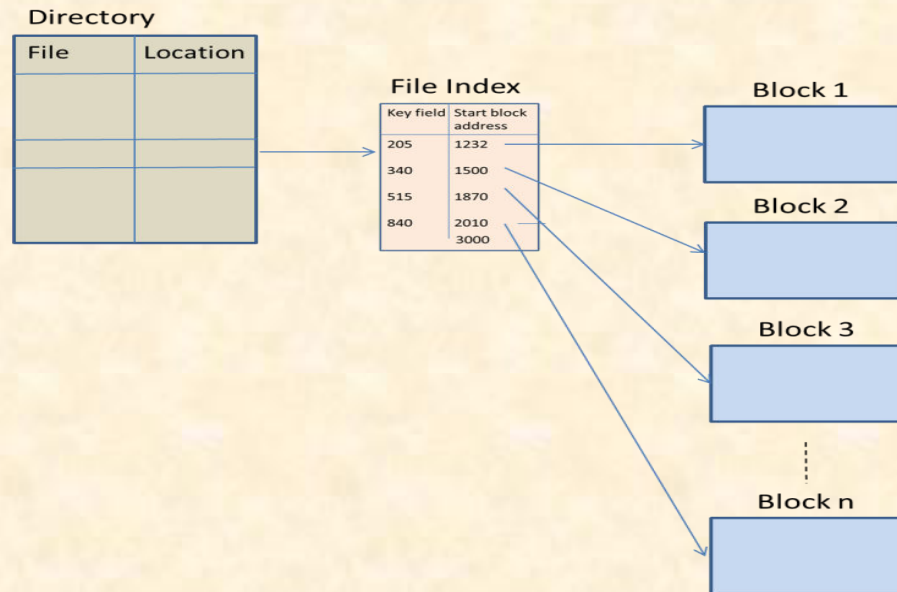
Directory

File	Location



Indexed Allocation

- An index for each file is maintained by the file manager.
- The index contains pointers to the various blocks of the file.



Indexed allocation



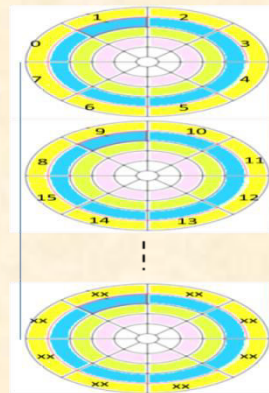
Drawback of Indexed Allocation

- The main drawback is that insertion of a record may lead to a cascade of changes in the index entries of the file.



Block-oriented File Mapping

- The disk is divided into blocks.
- Each block is provided a unique number.
- Instead of an index, this technique uses a **file map**.
- It is an array of integers.
- The subscripts of the array represent the block number and its contents can have three types of entries:
 - Free
 - Null
 - Integer

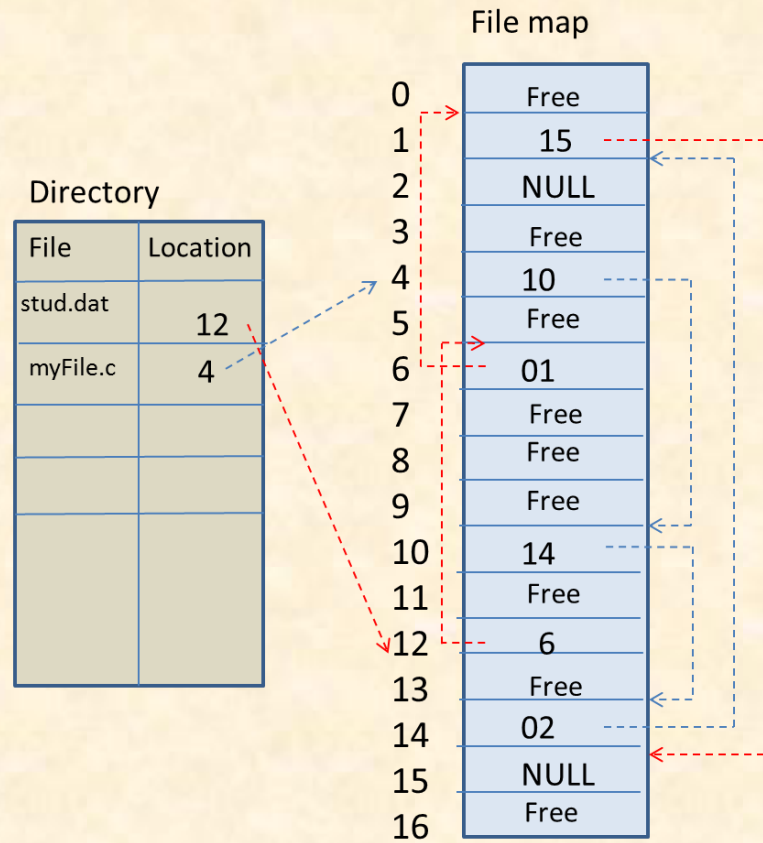


Cylinder

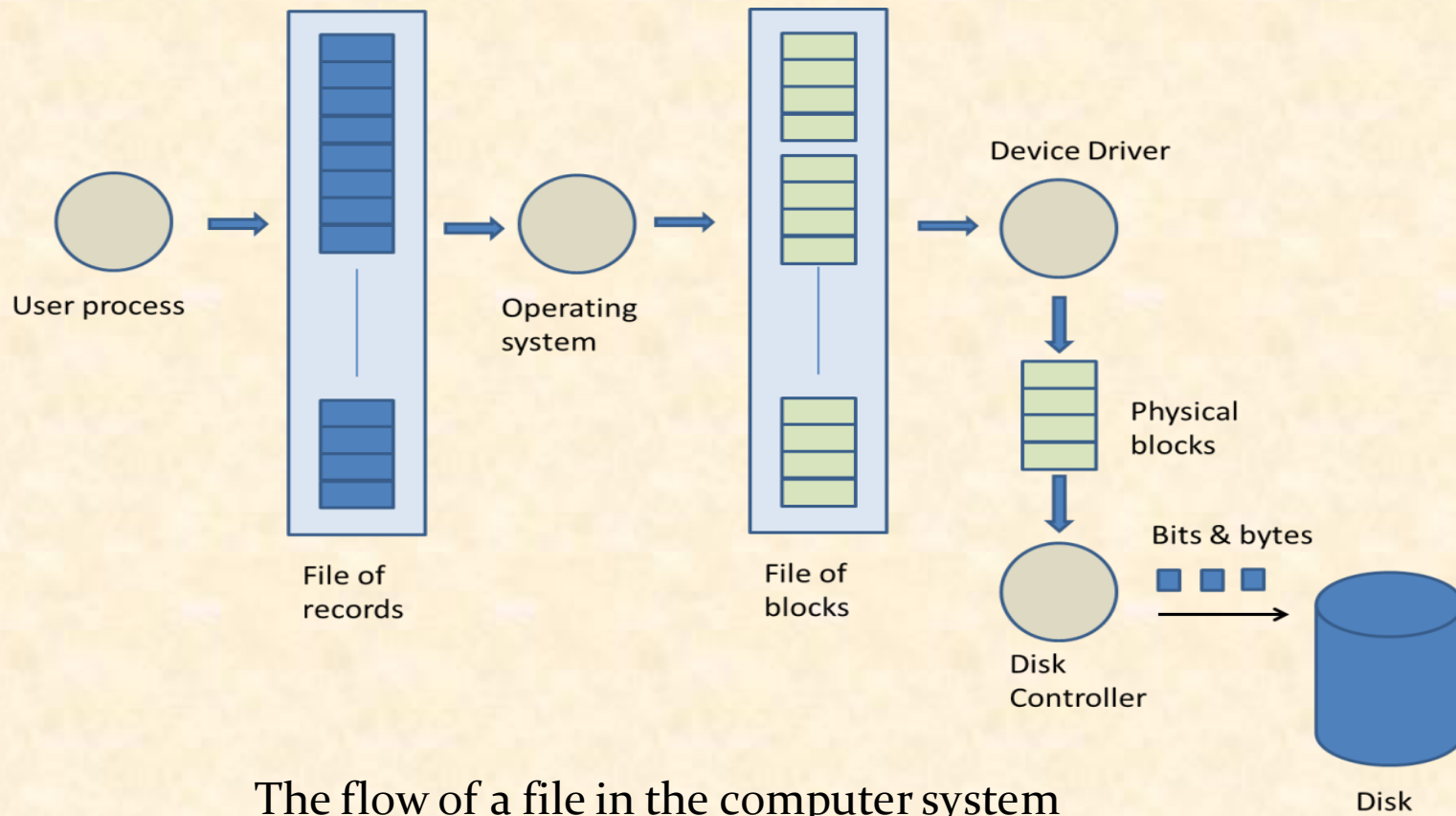
Numbered blocks on the disk



Block-oriented File Mapping

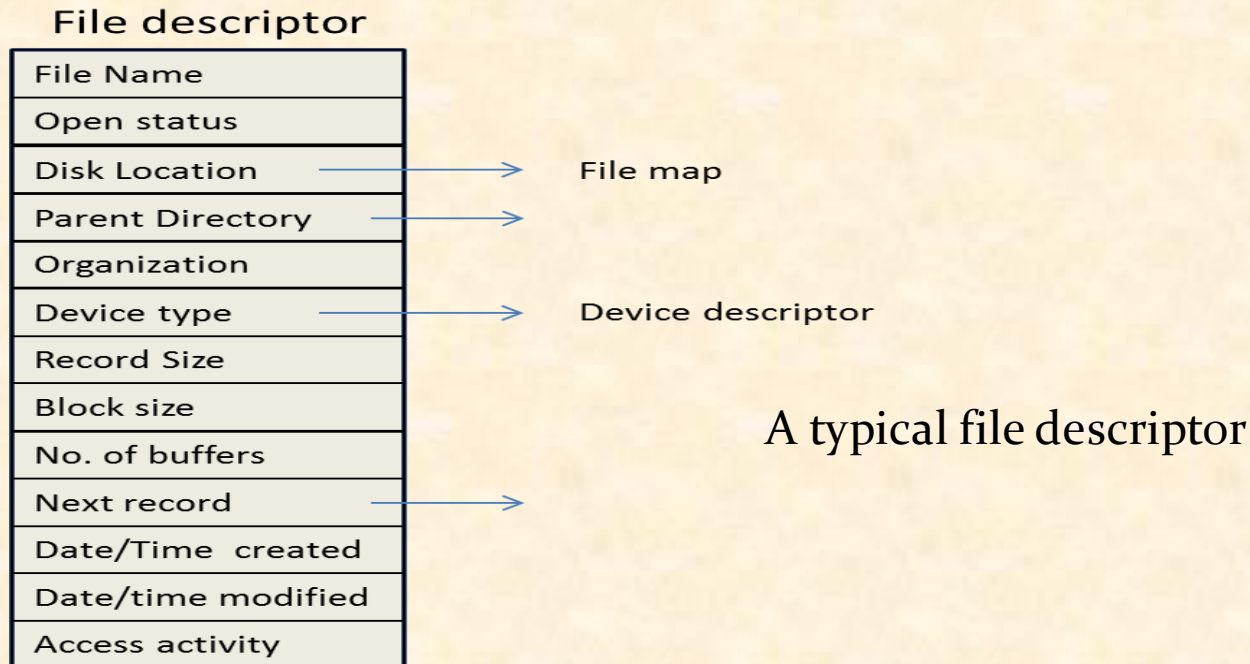


File System Architecture

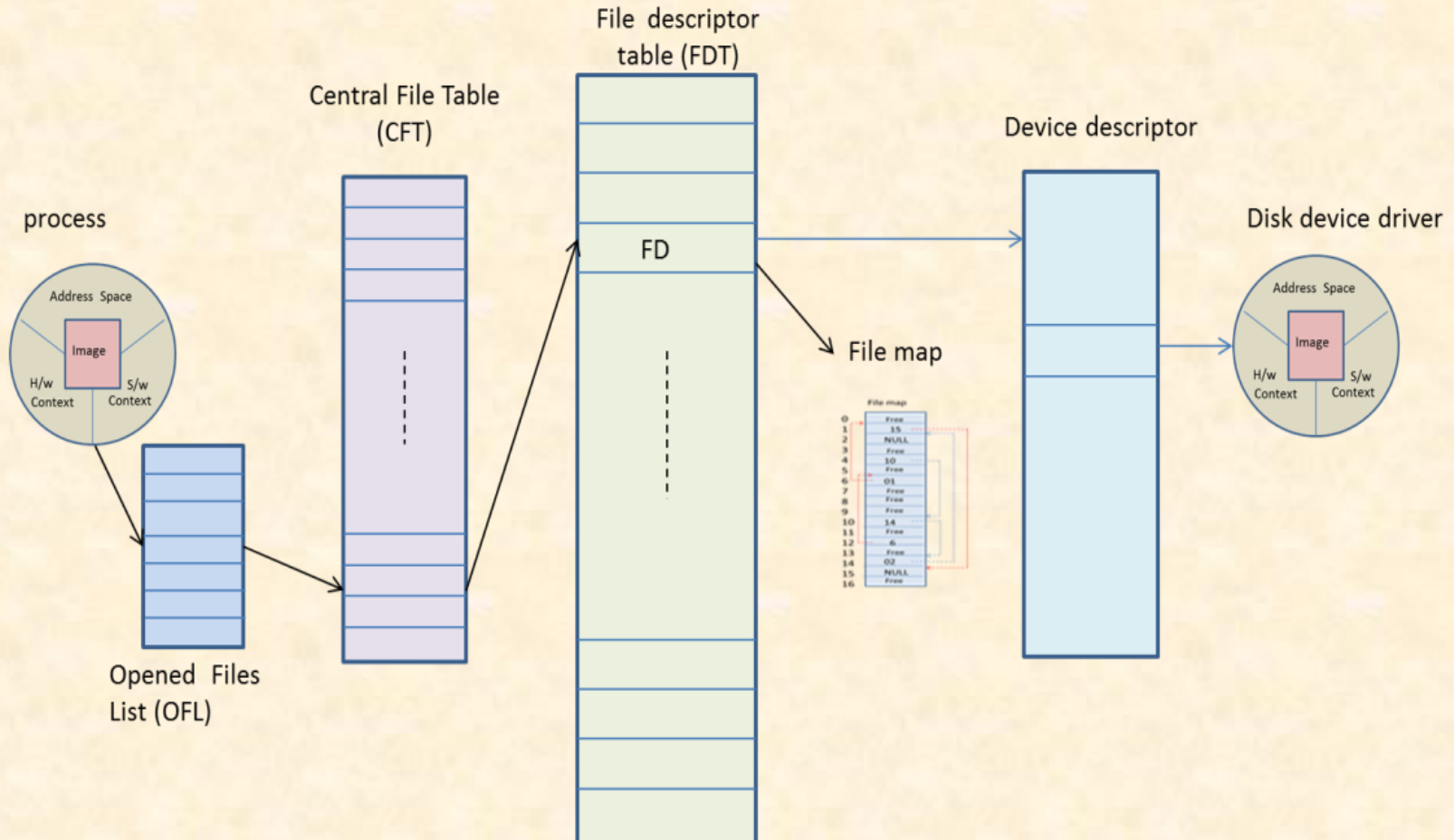


File System Data Structures

- For every file stored in the disk, the operating system maintains a file descriptor (FD).

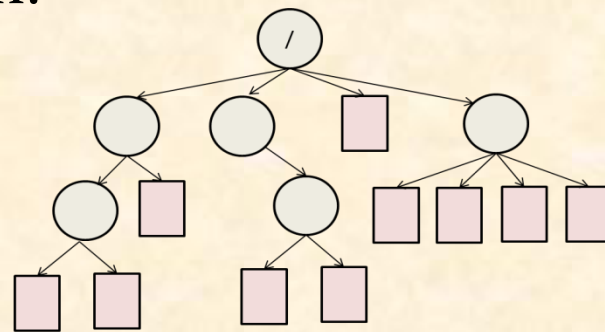


Data Structures of the File System

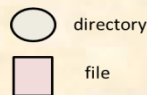


Unix File System

- Unix supports a hierarchical file system wherein it has a root node at the top of the inverted tree.
- The root directory contains a number of subdirectories and files.
- The subdirectories further contain subdirectories and files and so on.



The hierarchical
file structure

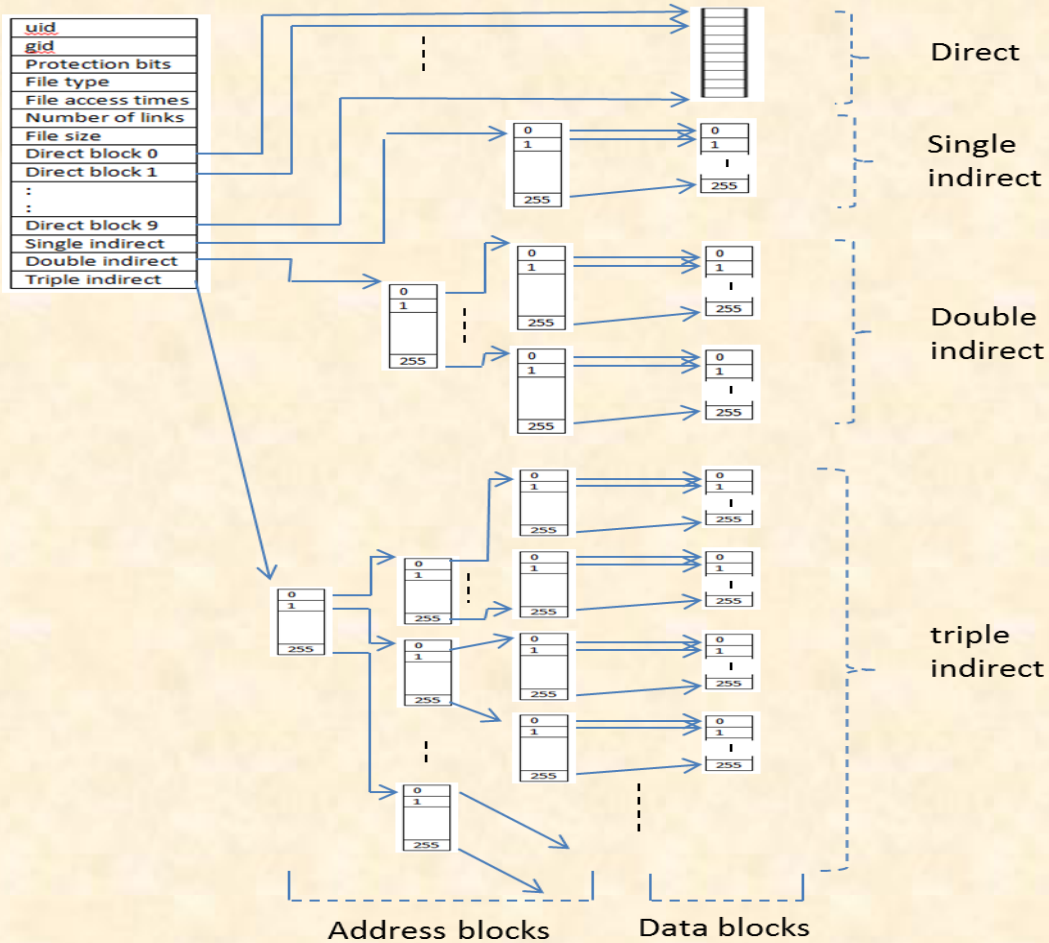


Unix File System

- Unix supports four types of files:
 - Ordinary
 - Directory
 - Special
 - FIFO
- Unix assigns a data structure to every file that contains information about the file. The data structure is called the 'inode'.

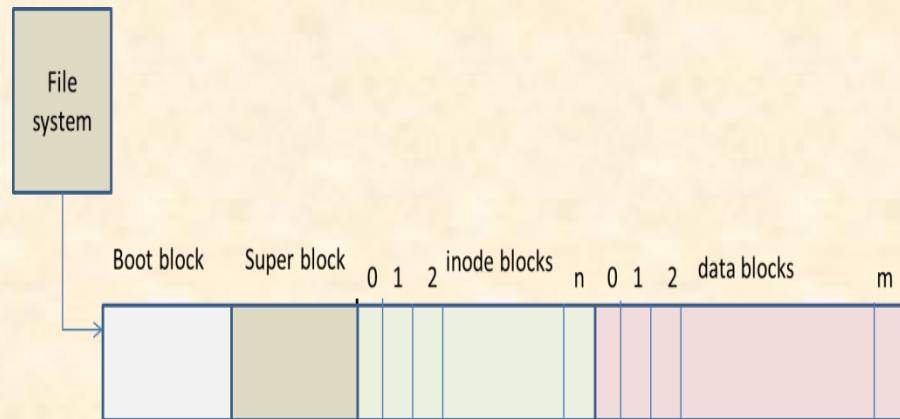


Inode Structure – Multiple Indirect Block Addressing



Unix File System Architecture

- Unix supports multiple file systems. The root file system is permanent and others can be removed.
- On the disk, a file system consists of four parts:
 - Boot block
 - Super block
 - Index blocks 0–n
 - Data blocks 0–m



Unix file system

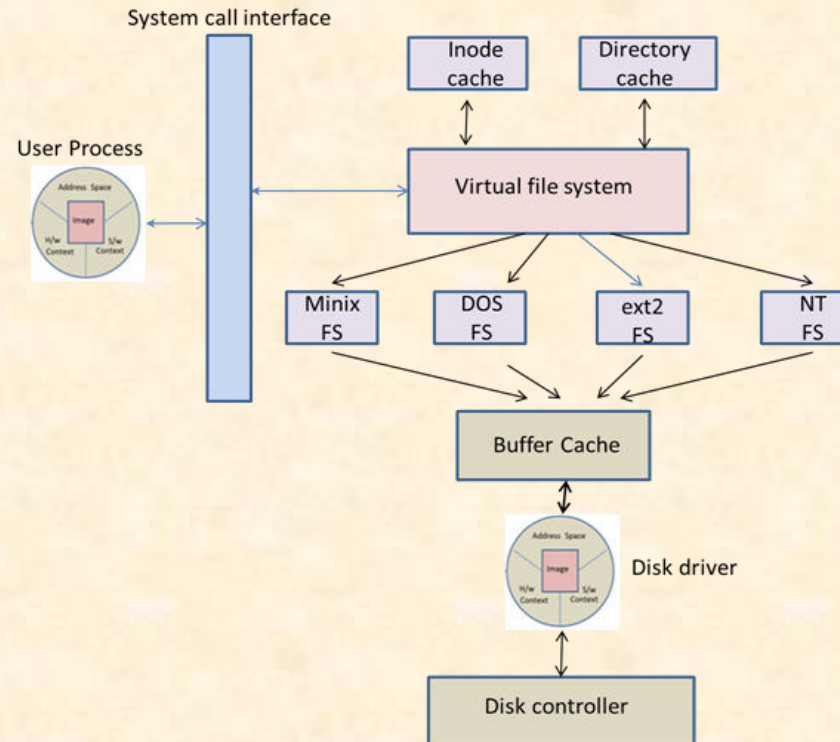


Linux File System

- The Linux file system was designed to support the Minix file system.
- The ability to concurrently support different file systems has been achieved with the help of the **virtual file system** (VFS).
- VFS offers a general interface capable of interacting with the different types of file systems mounted on a system at any given time.
- Every file is treated as an object and is assigned a symbolic name.



Virtual File System of Linux

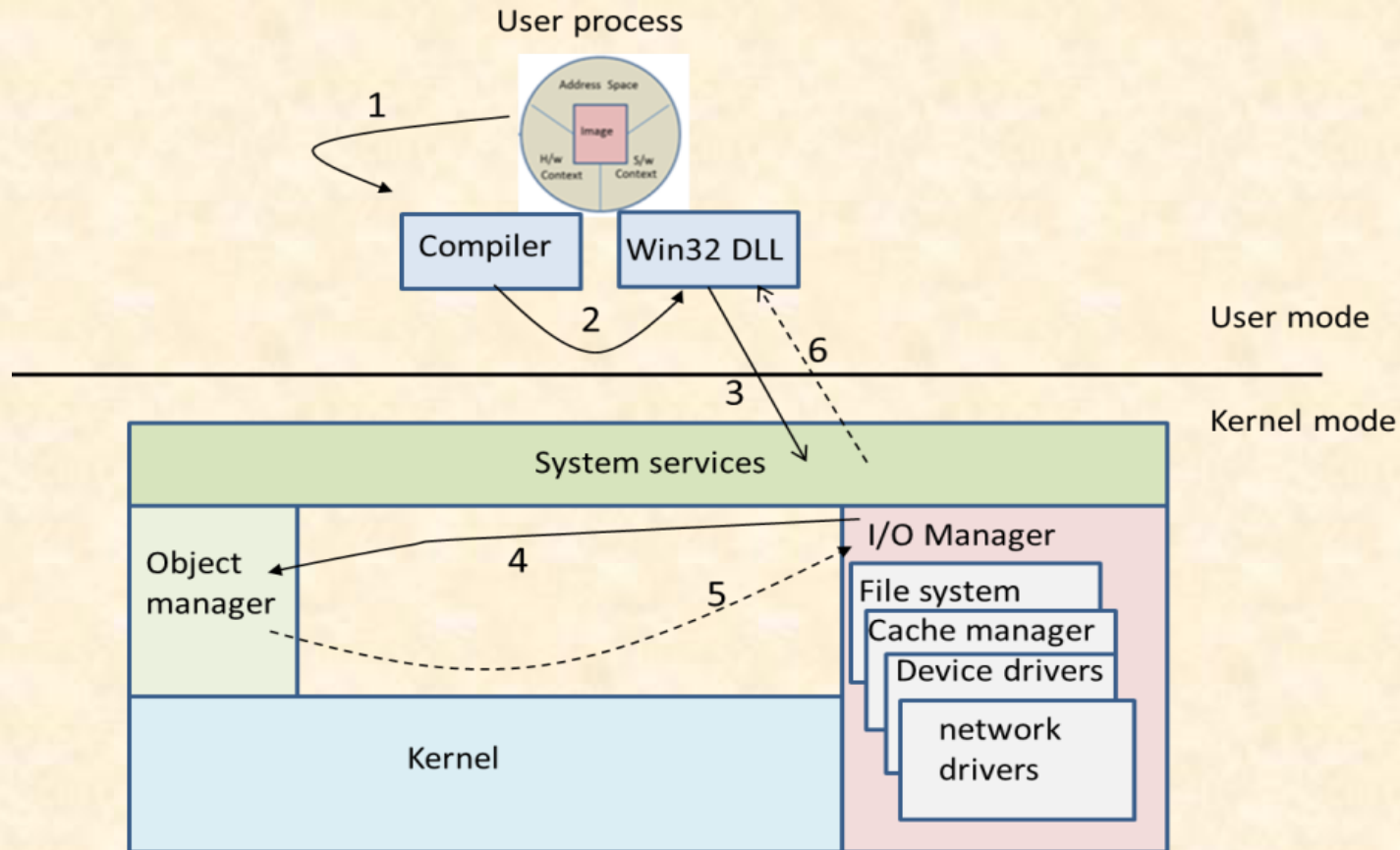


Windows File System

- Windows mainly support three types of file systems:
 - FAT (File allocation table)
 - NTFS (New technology file system)
 - HPFS (High performance file system)



Opening a File in Windows



Two-layered Driver System

User mode

Kernel mode

