

Module-2
CSEN 3104
Lecture 17

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Interconnection Networks

Shuffle-Exchange and Omega Networks

- Based on two routing functions --- Shuffle (S) and Exchange (E)
- Let $A = a_{n-1} \dots a_1 a_0$ be the address of a Processing Element (PE)
- The **Shuffle** function is given by
$$S(a_{n-1} \dots a_1 a_0) = a_{n-2} \dots a_1 a_0 a_{n-1} \text{ where } 0 \leq A \leq (N-1) \text{ and } n = \log_2 N$$
- Corresponds to cyclic shifting of the bits in A to the **left** for 1 bit position
- **Show figure for perfect shuffle**
- This action corresponds to perfect shuffling a deck of N cards
- The inverse perfect shuffle does the opposite to restore ordering (**Show figure**)
- Corresponds to cyclic shifting of the bits in A to the **right** for 1 bit position
- The **Exchange** function is given by
$$E(a_{n-1} \dots a_1 a_0) = a_{n-1} \dots a_1 a_0'$$
- The Exchange function exchanges the data between two PEs with adjacent addresses
- It is to be noted that $E(A) = C_0(A)$, where C_0 was the cube routing function

Shuffle-Exchange and Omega Networks

- The Shuffle-Exchange function can be implemented as
 - Single stage network
 - Multistage network
- Single Stage recirculating shuffle-exchange network (Show figure)
- Dashed lines -> Shuffle Solid lines -> Exchange
- A number of parallel algorithms can be effectively implemented by using Shuffle-Exchange function. Examples:
 - Fast Fourier Transform (FFT)
 - Polynomial Evaluation
 - Sorting
 - Matrix Transposition etc...

Multistage Omega Networks

- To implement Shuffle-Exchange functions (Show figure)
- An $N \times N$ Omega network consists of $n (= \log_2 N)$ identical stages
- Perfect shuffle interconnection between two adjacent stages
- Each stage has $N/2$ numbers of 4-function (straight, exchange, upper broadcast and lower broadcast) switch boxes under independent box control
- The switch boxes can be repositioned without violating the perfect shuffle interconnection between stages (Show figure)
- The n -cube network has the same interconnection topology as the repositioned Omega
- However, they are different in the following two points:
 - Cube NW uses 2-function switch boxes, whereas Omega NW uses 4-function ones
 - The dataflow directions in the two NWs are opposite to each other i.e. the roles of the input-output lines are exchanged in the two networks

Routing Algorithm for Omega Network

- A source S (with address $s_{n-1} s_{n-2} \dots s_0$) has to be connected to a certain destination D (with address $d_{n-1} d_{n-2} \dots d_0$)
- Starting at input S , connect the input of the first switch [in the $(n-1)^{\text{th}}$ stage] that is connected to S to
 - the upper output of the switch when $d_{n-1} = 0$
 - otherwise, to the lower output
- In the same way, bit d_{n-2} determines the output of the switch located on the next stage
- This process continues until a path is established between S and D
- In general, the input of the switch on the i^{th} stage is connected to the upper output when $d_i = 0$; Otherwise, the switch is connected to the lower output
- Example: Source 2 (i.e., $S = 010$) and destination 6 (i.e., $D = 110$) (Show Figure)
- In addition to one-to-one connections, the omega network also supports broadcasting
- Show Figure to explain the paths between source 2 and destinations 4,5,6 and 7

Thank you