# OSOROM Instruction Set Reference

The Moroso Project

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# Chapter 1

## CPU Architecture

#### 1.1 Instructions

Each instruction is a 32-bit word, stored in little-endian order. The top 2 bits of the instruction select a predicate register, and the third bit from the top optionally inverts it. The instruction is only executed if the resulting predicate is true.

### 1.2 Registers

There are 32 32-bit general-purpose registers, R0 through R31. R31 is used as the link register for BL instructions. The program counter is stored separately, and may only be modified through branch instructions and read through BL instructions.

There are three one-bit predicate registers, P0 through P2, that may be written to by compare instructions. A fourth predicate register, P3, is always true. Instructions predicated on P3 will always execute, and writes to P3 are ignored.

#### 1.3 ALU Instruction Formats

ALU instructions may accept immediate values for one of their input operands, which come in two forms: Short immediates are embedded in the instruction and consist of a 10-bit constant and a 4-bit rotate amount. To achieve any even rotation between 0 and 30, the rotate amount is first multiplied by 2, and the constant is rotated right by the result. ALU instructions that accept a single operand use the portion of the instruction normally used for the other operand register to extend the constant to 15 bits. Long immediate operands are 32 bit values, and occupy the memory word following the instruction.

If a register is used as the last source operand of an ALU instruction, it may optionally be shifted by a 5-bit immediate value, as specified with the SHF and SHIFTAMT fields of the instruction. The SHIFTAMT field is the amount to shift, between 0 and 31, and the SHF field specifies the shift type, as follows:

Encoding	Mnemonic	Shift Type
0 0	LSL	Logical Shift Left
0 1	LSR	Logical Shift Right
10	ASR	Arithmetic Shift Right
1 1	ROR	Rotate Right

ALU instructions taking a single operand may also specify a register shifted by an unsigned value contained in another register Rt. In this case, the same SHF types are available. If the value of the shift register is greater than 31, the following behavior is used: Logical shifts zero the result, arithmetic right shift extends the sign bit of the operand register across the entire result, and rotate rotates the value by Rt % 32.

#### 1.4 VLIW

At each time step, the CPU fetches a "packet" of 4 instructions located contiguously in memory and executes them in parallel. These packets must be aligned to their 16-byte size, and branch targets must also be so aligned.

There are three types of instructions: Control, Memory, and ALU instructions. Only one control instruction may be executed in any given packet; it must occupy the first (lowest-address) slot in its packet. Likewise, only the first two slots in a packet may execute memory instructions. ALU instructions may be located in any slot.

If a long immediate operand is specified for an ALU instruction, the following slot is interpreted as an operand and no operation is issued for that slot. The final slot in a packet may not specify a long immediate operand.

## 1.5 Virtual Memory

Virtual memory is accomplished through a hardware-filled TLB, with 4KB pages and a 2-level page-table structure. Page directories and page tables are one page in size, containing 1024 32-bit entries. Virtual addresses are 32 bits and are broken up into 3 parts: Page Directory Index, Page Table Index, and Page Offset, as follows:

	PD INDEX			PT	INDEX			OFFSET	
1		2	1			2	1		0
3		2	2			1	1		

The page directory base is stored in a coprocessor register, PTBR. Writes to this register will flush all entries not marked global from the TLB. Page table entries and page directory entries are each 32 bits and have the same format:

3	2	2				1	1						
1	9	8				2	1		4	3	2	1	0
R	SV		PHYSICAL	PAGE	BASE			UNUSED		G	K	W	Р

The four flags in each entry have the following meanings:

- G: Global page. Mappings with this bit set in either level will not be flushed from the TLB when the page table base register is modified.
- K: Kernel-only page. Mappings with this bit set in either level may only be read or written to while the processor is in kernel mode. Attempting to access such mappings from user mode will generate a page fault exception.
- W: Writeable page. Attempting to write to a page without this bit set in both levels will result in a page fault.

• P: Present. Attempting to access a page without this bit set in both levels will result in a page fault.

Since our FPGA has 512MB of physical memory, physical page bases are 17 bits, and offsets are 12 bits. The top bits of an entry are not used but should remain 0 in case we go to 32-bit physical addresses in the future. Bits 11 through 4 are free for programmer use.

## 1.6 Coprocessor Registers

The coprocessor registers are control registers accessible only in kernel mode via the MFC and MTC instructions. They provide information about exceptions to the kernel and are used to configure the processor.

#### 1.6.1 List of Coprocessor Registers

Encoding	Mnemonic	Purpose
00000	PFLAGS	Processor control flags. See below.
00001	PTB	Page Table Base Register. Specifies the start of the page directory for
		virtual address translation. Writes to this register flush the TLB.
00010	EHA	Exception Handler Address. Specifies a virtual address to which to trans-
		fer control when an exception or interrupt is encountered.
00011	EPC	Error PC. The virtual address of the last instruction packet to cause an
		exception, or the first instruction not executed due to an interrupt. LSBs
		contain system state when the exception occurred.
00100	EC0	Error Cause 0. The cause of an exception occurring on lane 0, or the type
		of a received interrupt. See Section 1.7 for a list of causes.
00101	EC1	Error Cause 1. The cause of an exception occurring on lane 1.
00110	EC2	Error Cause 2. The cause of an exception occurring on lane 2.
00111	EC3	Error Cause 3. The cause of an exception occurring on lane 3.
01000	EA0	Error Address 0. If lane 0 makes an invalid memory access, the virtual
		address it attempted to access will be stored in this register.
01001	EA1	Error Address 1. If lane 1 makes an invalid memory access, the virtual
		address it attempted to access will be stored in this register.
10000	SP0	Scratchpad register for system software usage.
10001	SP1	Scratchpad register for system software usage.
10010	SP2	Scratchpad register for system software usage.
10011	SP3	Scratchpad register for system software usage.

#### 1.6.2 The PFLAGS Register

At present, the PFLAGS register contains 2 flags: Bit 0 is 1 if interrupts are enabled, and bit 1 is 1 if paging is enabled. More control flags may be added here if needed.

#### 1.6.3 The EPC Register

The EPC register's MSBs contain the PC at which an exception occurred. (If an instruction fetch failed because of a paging exception, EPC[PC] contains the PC for which the fetch attempt failed.)



The M field contains the system mode at the time that the exception occurred: if EPC[M] is set, then the system was in kernel mode at the time, and otherwise, the system was in user mode. The EPC[IF] field contains the interrupt flag from PFLAGS at the time; if it is set, then interrupts were enabled, and if it is clear, than interrupts were disabled.

#### 1.6.4 Scratchpad registers

Scratchpad registers are provided for the convenience of system software, should it wish to save and restore user registers (or other data) without reserving user-architectural registers for that purpose. Their value is defined not to affect system state other than through reads and writes.

### 1.7 Exceptions and Interrupts

Exceptions and interrupts cause the processor to enter kernel mode and transfer control to the virtual address specified in the EHA register. If an interrupt occurs, the code corresponding to that interrupt will be written into the EC0 register. If one or more instructions in a packet cause exceptions, a code corresponding to the exception type will be written into the EC register corresponding to that instruction's location in the packet. If an instruction on lane 0 or 1 attempts to access an invalid memory address, the address will be written into the EA register for that lane. In all cases, the virtual address of the program counter where the exception occurred is written into the EPC register. For exceptions, this is the address of the instruction packet that caused the exception. For interrupts, this is the address of the first instruction packet not executed due to the interrupt. In all cases, interrupts are disabled. This is necessary in order to ensure the kernel can save the coprocessor registers before they are overwritten by another interrupt.

#### 1.7.1 Error Code Format



- EXNTYPE corresponds to a value in Table 1.1 and indicates the type of exception that occurred.
- If EXNTYPE indicates that an interrupt occurred, INT will correspond to a value in Table 1.2 and signify the type of interrupt.

Error Code	Type
00000	No Error
00001	Page Fault on Instruction Fetch
00010	Illegal Instruction
00011	Duplicate Destination
00100	Page Fault on Data Access
00101	Invalid Physical Address
00110	Divide by Zero (TODO are we doing other arithmetic exceptions??)
00111	Interrupt
01000	SYSCALL
01001	BREAK

Table 1.1: Exception Codes

Interrupt Code	Type
00	Timer
01	USB
10	Framebuffer
11	SD Card

Table 1.2: Interrupt Codes

# Chapter 2

# Instruction Set Summary

## 2.1 Instruction Set Encoding

	31 30 29																								
ALU 1-Op Short	Pred	0			]	mme	ediat	e Lo	w 10	)				Rot	ate		A	ALU	Opo	ode	Rd		Imm High 5		
ALU 2-Op Short	Pred	0			Immediate 10 Rotate ALU Opcode Rd Rs   Immediate 10 Rotate 0 1 1 Ctype Pd Rs															Rs					
Compare Short	Pred	0			Immediate 10 Rotate 0 1 1 1 Ctype Pd Rs   Shift amount Stype Rt ALU Opcode Rd Rs															Rs					
ALU Register	Pred	1	0	1		Shift amount Stype Rt ALU Opcode Rd Rs																			
Compare Register	Pred	1	0	1		Shift	am	ount		Styp	е			Rt			0	1	1	1	Ctype	Pd	Rs		
Load	Pred	1	0	0	1					О	ffset	12 LSU Opcode Rd Rs													
Store	Pred	1	0	0	1		Of	fset 1	Higł	ı 6				Rt			M	LS	UΟ	pcode	Offset L	ow 5	Rs		
Branch immediate	Pred	1	1	0	L											О	ffset	25							
Branch register	Pred	1	1	1	L									O	ffset	20							Rs		
Control	Pred	1	0	0	0	1	Ct	rl Ol	pcod	le Higl	h			Rt			Cti	rl Oj	pcode Low		Rd		Rs		
ALU 1-Op Regsh	Pred	1	0	0	0	0	0	0	1	Styp	е			Rt			I	ALU	Opo	ode	Rd		Rs		
ALU Long	Pred	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 ALU Opcode Rd						Rs	Imm 3	

# 2.2 ALU Opcodes

Encoding	Mnemonic	Instruction	Action	Section(s)
0 0 0 0	ADD	Addition	Rd := Rs + Op2	3.1
$0\ 0\ 0\ 1$	AND	Bitwise And	Rd := Rs & Op2	3.2
$0\ 0\ 1\ 0$	NOR	Bitwise Logical NOR	Rd := ~(Rs   Op2)	3.28
$0\ 0\ 1\ 1$	OR	Bitwise Logical OR	Rd := Rs   Op2	3.29
$0\ 1\ 0\ 0$	SUB	Subtraction	Rd := Rs - Op2	3.34
$0\ 1\ 0\ 1$	RSB	Reverse Subtraction	Rd := Op2 - Rs	3.30
$0\ 1\ 1\ 0$	XOR	Bitwise Exclusive OR	Rd := Rs ^ Op2	3.39
$0\ 1\ 1\ 1$	COMPARE	Compare Instructions	Pd := Rs CMP Op2	3.6, 3.7, 3.8, 3.9, 3.10, 3.11, 3.12
$1\ 0\ 0\ 0$	MOV	Move	Rd := 0p2	3.23
$1\ 0\ 0\ 1$	MVN	Move and Invert	Rd := ~Op2	3.26
$1\ 0\ 1\ 0$	SXB	Sign-Extend Byte	Rd := SX Op2[7:0]	3.36
$1\ 0\ 1\ 1$	SXH	Sign-Extend Half-Word	Rd := SX Op2[15:0]	3.37
$1\ 1\ 0\ 0$	reserved	n/a	n/a	n/a
$1\ 1\ 0\ 1$	reserved	n/a	n/a	n/a
$1\ 1\ 1\ 0$	reserved	n/a	n/a	n/a
1111	reserved	n/a	n/a	n/a

# Chapter 3

# Instruction Listing

## 3.1 ADD - Addition

#### 3.1.1 Encoding

1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	
PRED	0				C	ONS	STA	NT				]	ROT	'AT	E	0	0	0	0			RD					RS			
PRED	1	0	1	1	SHI	FT	'AM'	Г	SI	ΗF			RT			0	0	0	0			RD					RS			
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			RD					RS			LONG

#### **3.1.2** Syntax

- ADD[PN] Rd, Rs, #Imm
- ADD[PN] Rd, Rs, Rt [SHF #Imm]

#### 3.1.3 Type

ADD is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.1.4 Behavior

Adds the two operands together into the destination register. If two register operands are specified, the second may optionally be shifted by an immediate value before the addition is performed.

TODO optional overflow exceptions???

#### 3.2 AND - Bitwise And

#### 3.2.1 Encoding

	1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	
ĺ	PRED	0				C	ONS	TA:	NT				I	ROT	'AT	E	0	0	0	1			RD					RS			
	PRED	1	0	1		SHI	FT	'AM'	Γ	SI	ΙF			RT			0	0	0	1			RD					RS			
ĺ	PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1			RD					RS			LONG

#### **3.2.2** Syntax

- AND[PN] Rd, Rs, #Imm
- AND[PN] Rd, Rs, Rt [SHF #Imm]

#### 3.2.3 Type

AND is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.2.4 Behavior

Computes the bitwise logical AND of the two operands, storing the result in the destination register. If two register operands are specified, the second may optionally be shifted by an immediate value before the computation is performed.

## 3.3 B - Branch

#### 3.3.1 Encoding

_		•			 	 	 	 •							_	 	 	 	 	 
PRED	1	1	0	0								OI	- H S	ET						
PRED	1	1	1	0					(	188	SE								RS	

#### 3.3.2 Syntax

- B[PN] <label>
- B[PN] (Rs + OFF)

#### 3.3.3 Type

Branch is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.3.4 Behavior

Transfers control to the address specified. The OFFSET field is shifted left by 4 places and sign-extended to 32 bits, then added to either the source register or, if none is specified, the address of the branch instruction. If the specified base register is not aligned to a 16-byte value, the target is rounded down to the nearest aligned value before control is transferred. Other instructions in the same packet as the branch instruction are executed before control is transferred.

## 3.4 BL - Branch with Link

## 3.4.1 Encoding

PRED	1	1	0	1	OFFSET	
PRED	1	1	1	1	OFFSET	RS

#### 3.4.2 Syntax

- BL[PN] <label>
- BL[PN] (Rs + OFF)

#### 3.4.3 Type

Branch with Link is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.4.4 Behavior

BL functions identically to the Branch instruction, except that the address of the BL instruction is placed in R31.

## 3.5 BREAK

## 3.5.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE	D	1	0	0	0	1	0	0	0	1	Х	X	Х	Х	X	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	X	Х	X	X

## 3.5.2 Syntax

• BREAK - Debug Breakpoint

## 3.5.3 Behavior

Equivalent to SYSCALL, except that the exception code stored in EC0 is different - See Section 1.7.

## 3.6 CMPBC - Compare Bits Clear

#### 3.6.1 Encoding

1 0 9 8 7 6 5 4 3 2 1 0 9	7 6 5 4 3 2 1 0	9 8 7 6 5 4 3 2 1 0
---------------------------	-----------------	---------------------

PRED	0				CC	ONS	TA:	NT				I	ROT	ATI	Ξ	0	1	1	1	1	1	1	PD	RS	
PRED	1	0	1		SHI	FT	AM.	Γ	SI	HF			RT			0	1	1	1	1	1	1	PD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	PD	RS	LONG

#### **3.6.2** Syntax

- CMPBC[PN] Pd, Rs, #Imm
- CMPBC[PN] Pd, Rs, Rt [SHF #Imm]

#### 3.6.3 Type

CMPBC is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.6.4 Behavior

CMPBC checks each bit of Rs corresponding to a 1 bit in the second operand, writing 1 into the destination predicate if all such bits are 0, and writing 0 otherwise – that is, it computes the unary OR of the bits in ("Rs & OP2) and stores the result in the desired predicate.

## 3.7 CMPBS - Compare Bits Set

#### 3.7.1 Encoding

1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	
חשמת	$\overline{}$				~	חזוכי	ידי א	ידיזא				Т	חת	ית אי		$\overline{}$	1	1	1	1	1	$\overline{}$	וח	n			חמ			

PRED	0				CC	ONS	TA	NT				F	ROT	ATI	Ξ	0	1	1	1	1	1	0	PD	RS	
PRED	1	0	1		SHI	FT	AM'	Γ	SI	ΗF			RT			0	1	1	1	1	1	0	PD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	0	PD	RS	LONG

#### **3.7.2** Syntax

- CMPBS[PN] Pd, Rs, #Imm
- CMPBS[PN] Pd, Rs, Rt [SHF #Imm]

#### 3.7.3 Type

CMPBS is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.7.4 Behavior

CMPBS checks each bit of Rs corresponding to a 1 bit in the second operand, writing 1 into the destination predicate if any such bit is 1, and writing 0 otherwise – that is, it computes the unary OR of the bits in (Rs & OP2) and stores the result in the desired predicate.

## 3.8 CMPEQ - Compare Equal

#### 3.8.1 Encoding

	1 0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
ſ																															

PRED	0				CC	ONS	TA	NT				I	ROT	'AT	Ε	0	1	1	1	0	1	0	PD	RS	
PRED	1	0	1	1	SHI	FT	AM.	Γ	SI	ΙF			RT			0	1	1	1	0	1	0	PD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	0	1	0	PD	RS	LONG

#### 3.8.2 Syntax

- CMPEQ[PN] Pd, Rs, #Imm
- CMPEQ[PN] Pd, Rs, Rt [SHF #Imm]

#### 3.8.3 Type

CMPEQ is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.8.4 Behavior

CMPEQ tests whether its two operands are equal, writing 1 into the destination predicate if they are and 0 otherwise.

LONG

RS

## 3.9 CMPLES - Compare Less Than or Equal (Signed)

#### 3.9.1 Encoding

1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
PRED	0				CC	ONS	TAl	NT				I	ROT	'AT	Ε	0	1	1	1	1	0	1	P	D			RS		
PRED	1	0	1	5	SHI	FT	AM:	Γ	SI	HF			RT			0	1	1	1	1	0	1	P	D			RS		

#### 3.9.2 Syntax

PRED

- CMPLES[PN] Pd, Rs, #Imm
- CMPLES[PN] Pd, Rs, Rt [SHF #Imm]

#### 3.9.3 Type

CMPLES is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.9.4 Behavior

CMPLES interprets its operands as 32-bit two's-complement signed integers, and writes 1 into the specified predicate register if the value of Rs is less than or equal to the value of the second operand, and 0 otherwise.

## 3.10 CMPLEU - Compare Less Than or Equal (Unsigned)

#### 3.10.1 Encoding

1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

PRED	0				CC	ONS	TA:	NT				I	ROT	'AT	E	0	1	1	1	0	0	1	PD	RS	
PRED	1	0	1		SHI	FT	AM'	Γ	SI	HF			RT			0	1	1	1	0	0	1	PD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	0	0	1	PD	RS	LONG

#### 3.10.2 Syntax

- CMPLEU[PN] Pd, Rs, #Imm
- CMPLEU[PN] Pd, Rs, Rt [SHF #Imm]

#### 3.10.3 Type

CMPLEU is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.10.4 Behavior

CMPLEU interprets its operands as 32-bit unsigned integers, and writes 1 into the specified predicate register if the value of Rs is less than or equal to the value of the second operand, and 0 otherwise.

## 3.11 CMPLTS - Compare Less Than (Signed)

#### 3.11.1 Encoding

1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

PRED	0				CC	ONS	TA	NT				I	ROT	'AT	Ε	0	1	1	1	1	0	0	PD	RS	
PRED	1	0	1	1	SHI	FT	AM'	Γ	SI	HF			RT			0	1	1	1	1	0	0	PD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	0	0	PD	RS	LONG

#### 3.11.2 Syntax

- CMPLTS[PN] Pd, Rs, #Imm
- CMPLTS[PN] Pd, Rs, Rt [SHF #Imm]

#### 3.11.3 Type

CMPLTS is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.11.4 Behavior

CMPLTS interprets its operands as 32-bit two's-complement signed integers, and writes 1 into the specified predicate register if the value of Rs is less than to the value of the second operand, and 0 otherwise.

## 3.12 CMPLTU - Compare Less Than (Unsigned)

#### 3.12.1 Encoding

1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	--

PRED	0				CC	ONS	TA	NT				I	ROT	'AT	Ε	0	1	1	1	0	0	0	PD	RS	
PRED	1	0	1	1	SHI	FT	AM'	Γ	SI	ΙF			RT			0	1	1	1	0	0	0	PD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	0	0	0	PD	RS	LONG

#### 3.12.2 Syntax

- CMPLTU[PN] Pd, Rs, #Imm
- CMPLTU[PN] Pd, Rs, Rt [SHF #Imm]

#### 3.12.3 Type

CMPLTU is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.12.4 Behavior

CMPLTU interprets its operands as 32-bit unsigned integers, and writes 1 into the specified predicate register if the value of Rs is less than to the value of the second operand, and 0 otherwise. P3 is pinned to 1. Writes into P3 from compare instructions are ignored.

## 3.13 DIV - Integer Division

## 3.13.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
_	-	-	-		-	-	_	-	_	_	-	-	-		-	-	_	-	_	_	-	-	-	•	-	-	_	-	_	_	0
P	RE:	D	1	0	0	0	1	1	0	0	1	Х			RT			Х	Х	Х	Х			RD					RS		

## 3.13.2 Syntax

• DIV[PN] Rd, Rs, Rt

#### 3.13.3 Type

DIV is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.13.4 Behavior

Performs the integer division Rs / Rt, writing the quotient into Rd and the remainder into the multiply/divide overflow register.

## 3.14 ERET - Return from Exception

#### 3.14.1 Encoding

3 3 2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
PRED	1	0	0	0	1	0	1	0	0	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	X	X	X

#### 3.14.2 Syntax

• ERET[PN]

#### 3.14.3 Type

ERET is a Control operation. It may only be placed as the first instruction in a packet. It may only be used when the processor is in kernel mode.

#### 3.14.4 Behavior

ERET returns the processor to the state specified by the EPC coprocessor register. Specifically, the program counter takes on the value from EPC[PC], the system mode flag takes on the value from EPC[M] (kernel mode if M is 1, otherwise user mode), and the interrupt enable flag takes on the value from EPC[IF] (interrupts enabled if IF is 1, otherwise disabled). Other instructions in the same packet as the ERET instruction are executed in kernel mode before control is transferred. This instruction also clears the memory link bit.

## **3.15 FENCE**

#### 3.15.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE	D	1	0	0	0	1	0	0	1	1	Х	Х	Х	X	X	X	X	X	X	Х	X	Х	X	X	Х	Х	Х	Х	X	X

## 3.15.2 Syntax

• FENCE[PN]

#### 3.15.3 Type

FENCE is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.15.4 Behavior

Our architecture doesn't reorder requests, so right now FENCE is a no-op. If we have to add behavior to it in the future, memory instructions in the same packet as the FENCE will probably go after it, but please don't count on that right now.

## 3.16 FLUSH - Flush L1 Caches

#### 3.16.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE:	D	1	0	0	0	1	0	1	0	1	С	Х	Х	Х	Х	Х	Х	Х	TY	PE	Х	Х	Х	X	Х			RS		

#### 3.16.2 Syntax

• FLUSH[PN] TYPE, Rs

#### 3.16.3 Type

FLUSH is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.16.4 Behavior

FLUSH invalidates the line in the specified L1 cache corresponding to the virtual address contained in Rs. For instruction or data cache flushes, attempting to flush a virtual address whose access would cause a page fault will also cause a page fault. If a dirty data cache line is flushed, the new value is immediately written out to the L2 cache, where it will be visible to DMA peripherals. Values for the TYPE field are as follows:

Value	Mnemonic	Type
0 0	DATA	L1 Data Cache
0 1	INST	L1 Instruction Cache
1 0	DTLB	Data TLB
1 1	ITLB	Instruction TLB

## 3.17 LB - Load Byte

## 3.17.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE.	D	1	0	0	1					(	)FF	SE	Γ					0	0	0			RD					RS		

#### 3.17.2 Syntax

• LB[PN] Rd, (Rs [+ OFFSET])

#### 3.17.3 Type

LB is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.17.4 Behavior

LB loads a byte from the specified address into Rd. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs. The high 24 bits of Rd are zeroed.

## 3.18 LH - Load Half-Word

#### 3.18.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE	D	1	0	0	1					(	)FF	SE'	Т					0	0	1			RD					RS		

#### 3.18.2 Syntax

• LH[PN] Rd, (Rs [+ OFFSET])

#### 3.18.3 Type

LH is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.18.4 Behavior

LH loads two bytes from the specified address into Rd. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs. Only 2-byte aligned addresses may be accessed - the low bit of the address is ignored. The high 16 bits of Rd are zeroed.

## 3.19 LL - Load Linked

#### 3.19.1 Encoding

3	3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1		0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
PRED 1 0 0 1 OFFSET														0	1	1			RD					RS								

#### 3.19.2 Syntax

• LL[PN] Rd, (Rs [+ OFFSET])

#### 3.19.3 Type

LL is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.19.4 Behavior

LL loads a word from the specified address into Rd. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs. Only 4-byte aligned addresses may be accessed - the low 2 bits of the address are ignored.

Additionally, LL sets the link bit. A future SC operation will only complete if no interrupt, exception, or ERET clears the link bit in the intervening period.

## 3.20 LW - Load Word

#### 3.20.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE.	D	1	0	0	1					(	)FF	'SE'	Т					0	1	0			RD					RS		

#### 3.20.2 Syntax

• LW[PN] Rd, (Rs [+ OFFSET])

#### 3.20.3 Type

LW is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.20.4 Behavior

LW loads a word from the specified address into Rd. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs. Only 4-byte aligned addresses may be accessed - the low 2 bits of the address are ignored.

## 3.21 MFC - Move From Coprocessor Register

#### 3.21.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE:	D	1	0	0	0	1	0	1	1	0	Х	X	Х	Х	Х	Х	Х	Х	Х	Х			RD				С	PR	3	

#### 3.21.2 Syntax

• MFC[PN] Rd, CPRs

#### 3.21.3 Type

MFC is a Control operation. It may only be placed as the first instruction in a packet. It may only be used when the processor is in kernel mode.

#### 3.21.4 Behavior

MFC moves the contents of the specified coprocessor register into the specified general-purpose register. See Section 1.6 for a description of the available coprocessor registers and their mnemonics.

## 3.22 MFHI - Move From Multiply/Division Overflow Register

#### 3.22.1 Encoding

3 3 2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
PRED	1	0	0	0	1	1	0	1	0	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х			RD			Х	Х	Х	X	X

#### **3.22.2** Syntax

• MFHI[PN] Rd

#### 3.22.3 Type

MFHI is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.22.4 Behavior

MFHI moves the contents of the multiply/division overflow register into the specified general-purpose register. This register stores the high bits of results of multiplication operations, or the remainders generated by division operations.

## **3.23** MOV - Move

#### 3.23.1 Encoding

1 0 9 8 7 6 5 4 3 2 1 0 9	7 6 5 4 3 2 1 0	9 8 7 6 5 4 3 2 1 0
---------------------------	-----------------	---------------------

PRED	0				CC	ONS	TL	OW				I	ROT	'AT	Ε	1	0	0	0	RD		CO	NST	ГНІ		
PRED	1	0	1		SHI	FT	AM'	Γ	SI	ΙF			RT			1	0	0	0	RD	Х	Х	Х	Х	Х	
PRED	1	0	0	0	0	0	0	1	SI	ΙF			RT			1	0	0	0	RD			RS			
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	RD			RS			LONG

#### 3.23.2 Syntax

- MOV[PN] Rd, #Imm
- MOV[PN] Rd, Rt [SHF #Imm]
- MOV[PN] Rd, Rt SHF Rs

#### 3.23.3 Type

MOV is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.23.4 Behavior

MOV moves the value of its source operand into the specified destination register. If the second operand is a register, it may be shifted by an immediate value or by the contents of a register before being stored.

## 3.24 MTC - Move To Coprocessor Register

#### 3.24.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE:	D	1	0	0	0	1	0	1	1	1	Х	X	Х	Х	X	Х	Х	Х	Х	Х		С	PR	D				RS		

#### 3.24.2 Syntax

• MTC[PN] CPRd, Rs

#### 3.24.3 Type

MTC is a Control operation. It may only be placed as the first instruction in a packet. It may only be used when the processor is in kernel mode.

#### 3.24.4 Behavior

MTC moves the contents of the selected general-purpose register into the specified coprocessor register. See Section 1.6 for a description of the available coprocessor registers and their mnemonics.

## 3.25 MTHI - Move To Multiply/Division Overflow Register

## 3.25.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE:	D	1	0	0	0	1	1	0	1	1	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х			RS		

## 3.25.2 Syntax

• MTHI[PN] Rs

#### 3.25.3 Type

MTHI is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.25.4 Behavior

 $\operatorname{MTHI}$  moves the value of the specified general-purpose register into the multiply/division overflow register.

# 3.26 MVN - Move and Invert

# 3.26.1 Encoding

. 0 9 8 7 6 5 4 3 2 1 0 9	7 6 5 4 3 2 1 0	9 8 7 6 5 4 3 2 1 0
---------------------------	-----------------	---------------------

PRED	0				CC	ONS	TL	DW				]	ROT	'AT	Ε	1	0	0	1	RD		CO	NS'	ΉΙ		
PRED	1	0	1		SHI	FT	AM'	Γ	SI	ΗF			RT			1	0	0	1	RD	Х	Х	Х	Х	Х	
PRED	1	0	0	0	0	0	0	1	SI	HF			RT			1	0	0	1	RD			RS			
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	1	RD			RS			LONG

## 3.26.2 Syntax

- MVN[PN] Rd, #Imm
- MVN[PN] Rd, Rt [SHF #Imm]
- MVN[PN] Rd, Rt SHF Rs

# 3.26.3 Type

MVN is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.26.4 Behavior

MVN inverts its source operand, storing the result in the specified destination register. If the second operand is a register, it may be shifted by an immediate value or by the contents of a register before being inverted and stored.

# 3.27 MULT - Integer Multiplication

# 3.27.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE	D	1	0	0	0	1	1	0	0	0	S			RT			Х	Х	Х	Х			RD					RS		

#### 3.27.2 Syntax

- MULT[PN] Rd, Rs, Rt
- MULTS[PN] Rd, Rs, Rt

# 3.27.3 Type

MULT is a Control operation. It may only be placed as the first instruction in a packet.

#### 3.27.4 Behavior

MULT performs an integer multiplication on its two register operands. If the S bit is set, the operands are interpreted as two's complement signed integers; otherwise they are interpreted as unsigned. The low 32 bits of the 64 bit result are placed in the specified destination register, and the high 32 bits are placed in the special-purpose multiply/division overflow register, where they are accessible via the MFHI instruction.

# 3.28 NOR - Bitwise Logical NOR

#### 3.28.1 Encoding

1 0 9	8	7	6	5 4	1 3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
DBED	Λ				מחוזכ	ו א ידיב	NТ				I	ΩП	ית א'	7	Λ	Λ	1	$\overline{\cap}$		1	תפ					ВG		

PRED	0				CC	ONS	TA:	NT				I	ROT	'AT	Ε	0	0	1	0	RD	RS	
PRED	1	0	1		SHI	FT	AM'	Γ	SI	HF			RT			0	0	1	0	RD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	RD	RS	LONG

# 3.28.2 Syntax

- NOR[PN] Rd, Rs, #Imm
- NOR[PN] Rd, Rs, Rt [SHF #Imm]

# 3.28.3 Type

NOR is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.28.4 Behavior

Computes the bitwise logical NOR of the two operands, storing the result in the destination register. If two register operands are specified, the second may optionally be shifted by an immediate value before the computation is performed.

# 3.29 OR - Bitwise Logical OR

# 3.29.1 Encoding

1 0 9	8	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	
חחחח	^			~	2270	1 m A	ATTT				Т.	О.	A 177	_	$\sim$	$\overline{}$	-	-			חח					Da			٦

PRED	0				CC	ONS	TAl	NT				F	ROT	'AT	Ε	0	0	1	1	RD	RS	
PRED	1	0	1	5	SHI	FT.	AM'	Γ	SI	ΗF			RT			0	0	1	1	RD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	RD	RS	LONG

# 3.29.2 Syntax

- OR[PN] Rd, Rs, #Imm
- OR[PN] Rd, Rs, Rt [SHF #Imm]

## 3.29.3 Type

OR is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.29.4 Behavior

Computes the bitwise logical OR of the two operands, storing the result in the destination register. If two register operands are specified, the second may optionally be shifted by an immediate value before the computation is performed.

# 3.30 RSB - Reverse Subtraction

## 3.30.1 Encoding

1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
DRED	0				רר	าพร	ТΔ1	NT				1	דחג	יד מי	F	n	1	Λ	1			ВD					RS		

PRED	0				CC	ONS	STA	NT				I	ROT	'ATI	Ε	0	1	0	1	RD	RS	
PRED	1	0	1		SHI	FT	'AM'	Γ	SI	HF			RT			0	1	0	1	RD	RS	
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	1	RD	RS	LONG

# 3.30.2 Syntax

- RSB[PN] Rd, Rs, #Imm
- RSB[PN] Rd, Rs, Rt [SHF #Imm]

# 3.30.3 Type

RSB is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.30.4 Behavior

Subtracts the value stored in RS from the value of the second operand, storing the result in the destination register. If two register operands are specified, the second may optionally be shifted by an immediate value before the computation is performed.

# 3.31 SB - Store Byte

# 3.31.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE	D	1	0	0	1		OFI	SE	TH:	IGH	Ī.			RT			М	1	0	0	0	FF	SET	LO	W			RS		

#### 3.31.2 Syntax

• SB[PN] Rt, (Rs [+ OFFSET])

# 3.31.3 Type

SB is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.31.4 Behavior

SB stores the low byte of Rt at the provided address. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs.

# 3.32 SC - Store Conditional

#### 3.32.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE:	D	1	0	0	1		OFF	SE	TH.	IGH	[			RT			М	1	1	1	0	FFS	SET	LO	W			RS		

#### 3.32.2 Syntax

• SC[PN] Rt, (Rs [+ OFFSET])

## 3.32.3 Type

SC is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.32.4 Behavior

SC stores the value in Rt at the provided address, if the processor's link bit is set. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs. Only 4-byte aligned addresses may be accessed - the low 2 bits of the address are ignored.

The link bit is set by executing the LL instruction, and is cleared by the SC or ERET instructions, or any time an exception or interrupt is dispatched. SC only completes if the link bit is set when it is executed - thus, if it completes, all instructions since the last LL instruction are known to have been executed atomically. The SC instruction writes a 1 into predicate register 0 if the store completes successfully, or a 0 if it fails.

# 3.33 SH - Store Half-Word

# 3.33.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
Р	RE	D	1	0	0	1		OFE	SE	TH:	IGH	I			RT			М	1	0	1	0	FF	SET	LO	W			RS		

#### 3.33.2 Syntax

• SH[PN] Rt, (Rs [+ OFFSET])

## 3.33.3 Type

SH is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.33.4 Behavior

SH stores the low two bytes of Rt at the provided address. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs. The address must be aligned to a multiple of 2; the low bit of the address is ignored.

# 3.34 SUB - Subtraction

## 3.34.1 Encoding

	1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	
	PRED	0				CO	ONS	TA:	NT				I	ROT	'AT	E	0	1	0	0			RD					RS			
	PRED	1	0	1		SHI	FT	'AM'	Γ	SI	HF			RT	ı		0	1	0	0			RD					RS			
ĺ	PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0			RD					RS			LONG

## 3.34.2 Syntax

- SUB[PN] Rd, Rs, #Imm
- SUB[PN] Rd, Rs, Rt [SHF #Imm]

## 3.34.3 Type

SUB is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.34.4 Behavior

Subtracts the value of the second operand from the value in Rs, storing the result in the destination register. If two register operands are specified, the second may optionally be shifted by an immediate value before the computation is performed.

# 3.35 SW - Store Word

# 3.35.1 Encoding

3 3 2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1 0 9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
PRED	1	0	0	1		OFF	SE	TH:	IGH	I			RT			М	1	1	0	0	FF	SE1	CLO	W			RS		

# 3.35.2 Syntax

• SW[PN] Rt, (Rs [+ OFFSET])

# 3.35.3 Type

SW is a Memory operation. It may only be placed in the first or second slot in an instruction packet.

#### 3.35.4 Behavior

SW stores the contents of Rt at the provided address. The address is computed by sign-extending the OFFSET field to 32 bits and adding it to the value of Rs. The address must be aligned to a multiple of 4; the low two bits of the address are ignored.

# 3.36 SXB - Sign-Extend Byte

# 3.36.1 Encoding

PRED	0				CO	ONS	TL	OW				]	ROT	CAT	E	1	0	1	0	RD		CO	NS.	THI		
PRED	1	0	1	5	SHI	FT	AM'	Γ	SI	ΗF			RT	•		1	0	1	0	RD	Х	Х	Х	X	Х	
PRED	1	0	0	0 0 0 0 1 S						HF			RT	•		1	0	1	0	RD			RS			
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	1	0	RD			RS			LONG

## 3.36.2 Syntax

- SXB[PN] Rd, #Imm
- SXB[PN] Rd, Rt [SHF #Imm]
- SXB[PN] Rd, Rt SHF Rs

# 3.36.3 Type

SXB is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.36.4 Behavior

SXB moves the low byte of its source operand into the specified destination register, sign-extending it across the rest of the register. If the second operand is a register, it may be shifted by an immediate value or by the contents of a register before being stored.

# 3.37 SXH - Sign-Extend Half-Word

## 3.37.1 Encoding

 $3 \ \ 3 \ \ 2 \ \ 2 \ \ 2 \ \ 2 \ \ 2 \ \ 2 \ \ 2 \ \ 1 \ \ 1 \ \ 1 \ \ 1 \ \ 1 \ \ 1 \ \ 1 \ \ 1 \ \ 1 \ \ 1$ 

1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

PRED	0				CO	ONS	TL	OW				I	ROT	'AT	Ξ	1	0	1	1	RD		CO	NST	ГНІ		
PRED	1	0	1		SHI	FT	AM'	Γ	SI	HF			RT			1	0	1	1	RD	Х	X	Х	Х	Х	
PRED	1	0	0	0	0	0	0	1	SI	HF			RT			1	0	1	1	RD			RS			1
PRED	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	1	1	RD			RS			LONG

## 3.37.2 Syntax

- SXH[PN] Rd, #Imm
- SXH[PN] Rd, Rt [SHF #Imm]
- SXH[PN] Rd, Rt SHF Rs

# 3.37.3 Type

SXH is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.37.4 Behavior

SXH moves the low two bytes of its source operand into the specified destination register, sign-extending it across the rest of the register. If the second operand is a register, it may be shifted by an immediate value or by the contents of a register before being stored.

# 3.38 SYSCALL - System Call Exception

# 3.38.1 Encoding

3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1										
1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0
P	RE	D	1	0	0	0	1	0	0	1	0	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	X	Х

# 3.38.2 Syntax

• SYSCALL

#### 3.38.3 Behavior

SYSCALL causes an exception to occur on the instruction packet of which it is a part. Its corresponding error code is placed in the EC0 coprocessor register - see Section 1.7 for details. Note that this causes other instructions in the packet to be canceled.

# 3.39 XOR - Bitwise Exclusive OR

#### 3.39.1 Encoding

	1 0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	9	8	7	6	5	4	3	2	1	0	
	PRE	D	0				C	ONS	TA	NT				]	ROT	'AT	Ε	0	1	1	0			RD					RS			
İ	PRE	D	1	0	1	:	SHI	FT	AM'	Г	SI	HF			RT	1		0	1	1	0			RD					RS			
ĺ	PRE	D	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0			RD					RS			LONG

#### 3.39.2 Syntax

- XOR[PN] Rd, Rs, #Imm
- XOR[PN] Rd, Rs, Rt [SHF #Imm]

## 3.39.3 Type

XOR is an ALU operation. It may be placed in any slot of a packet. If the long immediate form is used it must not be located in the last slot in a packet, and no instruction may be specified in the following slot.

#### 3.39.4 Behavior

Computes the bitwise logical XOR of the two operands, storing the result in the destination register. If two register operands are specified, the second may optionally be shifted by an immediate value before the computation is performed.