Community Structure in Networks

CS224W: Social and Information Network Analysis Jure Leskovec, Stanford University

http://cs224w.stanford.edu



How the Class Fits Together

Observations

Small diameter, Edge clustering

Patterns of signed edge creation

Viral Marketing, Blogosphere, Memetracking

Scale-Free

Densification power law, Shrinking diameters

Strength of weak ties, Core-periphery

Models

Erdös-Renyi model, Small-world model

Structural balance, Theory of status

Independent cascade model, Game theoretic model

Preferential attachment, Copying model

Microscopic model of evolving networks

Kronecker Graphs

Algorithms

Decentralized search

Models for predicting edge signs

Influence maximization, Outbreak detection, LIM

PageRank, Hubs and authorities

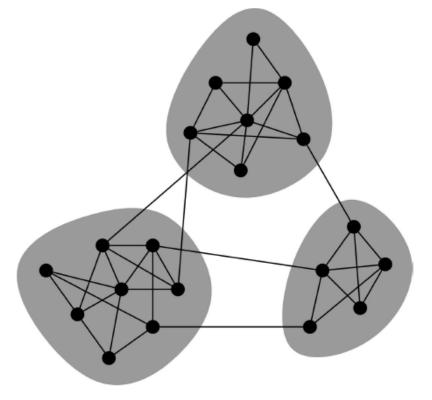
Link prediction,
Supervised random walks

Community detection: Girvan-Newman, Modularity

Networks & Communities

We often think of networks "looking"

like this:



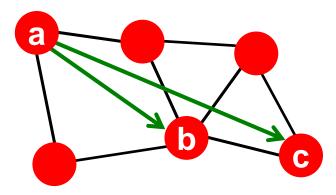
What lead to such conceptual picture?

Networks: Flow of Information

- How information flows through the network?
 - What structurally distinct roles do nodes play?
 - What roles do different links (short vs. long) play?
- How people find out about new jobs?
 - Mark Granovetter, part of his PhD in 1960s
 - People find the information through personal contacts
- But: Contacts were often acquaintances rather than close friends
 - This is surprising: One would expect your friends to help you out more than casual acquaintances
- Why is it that acquaintances are most helpful?

Granovetter's Answer

- Two perspectives on friendships:
 - Structural: Friendships span different parts of the network
 - Interpersonal: Friendship between two people is either strong or weak
- Structural role: Triadic Closure

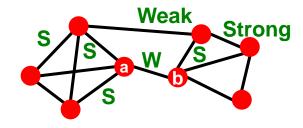


Which edge is more likely a-b or a-c?

If two people in a network have a friend in common there is an increased likelihood they will become friends themselves

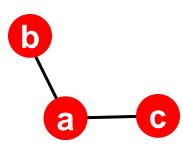
Granovetter's Explanation

- Granovetter makes a connection between social and structural role of an edge
- First point:
 - Structurally embedded edges are also socially strong
 - Edges spanning different parts of the network are socially weak
- Second point:
 - The long range edges allow you to gather information from different parts of the network and get a job
 - Structurally embedded edges are heavily redundant in terms of information access



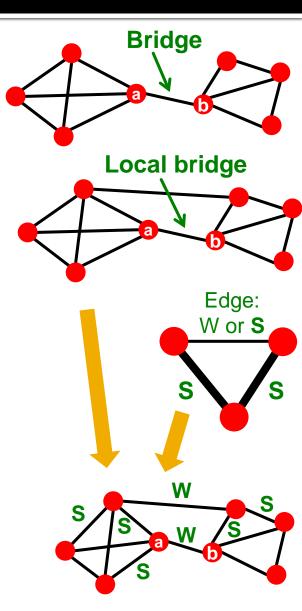
Triadic Closure

- Triadic closure == High clustering coefficient Reasons for triadic closure:
- If B and C have a friend A in common, then:
 - B is more likely to meet C
 - (since they both spend time with A)
 - B and C trust each other
 - (since they have a friend in common)
 - A has incentive to bring B and C together
 - (as it is hard for A to maintain two disjoint relationships)
- Empirical study by Bearman and Moody:
 - Teenage girls with low clustering coefficient are more likely to contemplate suicide



Granovetter's Explanation

- Define: Bridge edge
 - If removed, it disconnects the graph
- Define: Local bridge
 - Edge of Span > 2
 (Span of an edge is the distance of the edge endpoints if the edge is deleted. Local bridges with long span are like real bridges)
- Define: Two types of edges:
 - Strong (friend), Weak (acquaintance)
- Define: Strong triadic closure:
 - Two strong ties imply a third edge
- Fact: If strong triadic closure is satisfied then local bridges are weak ties!

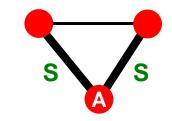


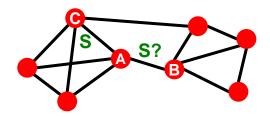
Local Bridges and Weak ties

• <u>Claim</u>: If node *A* satisfies **Strong Triadic Closure** and is involved in at least **two strong ties**, then any **local bridge** adjacent to *A* must be a **weak tie**.

Proof by contradiction:

- A satisfies StrongTriadic Closure
- Let A B be local bridge and a strong tie
- Then B − C must exist because of Strong
 Triadic Closure
- But then A B is not a bridge!





Tie strength in real data

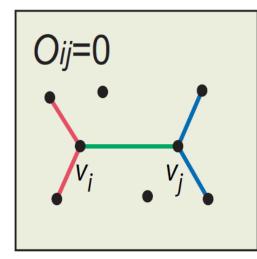
- For many years the Granovetter's theory was not tested
- But, today we have large who-talks-to-whom graphs:
 - Email, Messenger, Cell phones, Facebook
- Onnela et al. 2007:
 - Cell-phone network of 20% of country's population
 - Edge strength: # phone calls

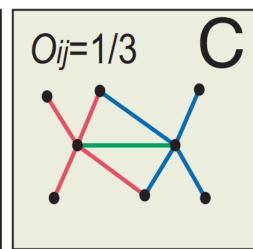
Neighborhood Overlap

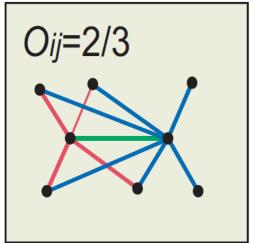
Edge overlap:

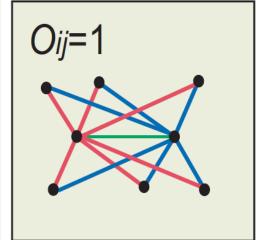
$$O_{ij} = \frac{N(i) \cap N(j)}{N(i) \cup N(j)}$$

- N(i) ... a set of neighbors of node i
- Overlap = 0when an edge is a local bridge









Phones: Edge Overlap vs. Strength

Cell-phone network

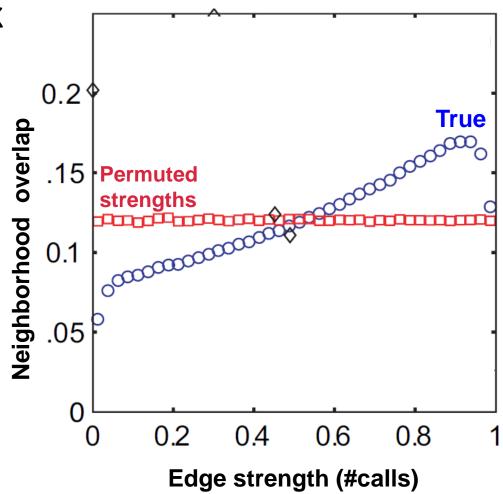
Observation:

Highly used links have high overlap!

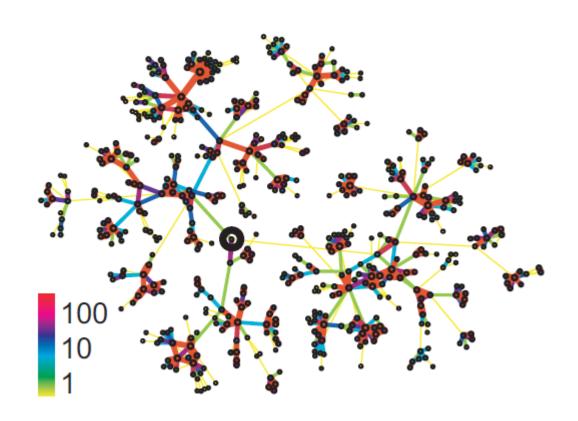
Legend:

True: The data

 Permuted strengths: Keep the network structure but randomly reassign edge strengths

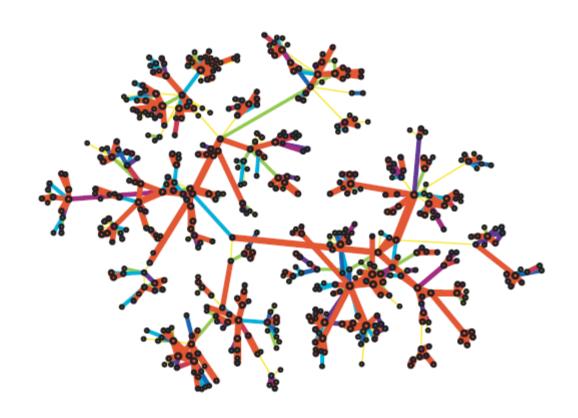


Real Network, Real Tie Strengths



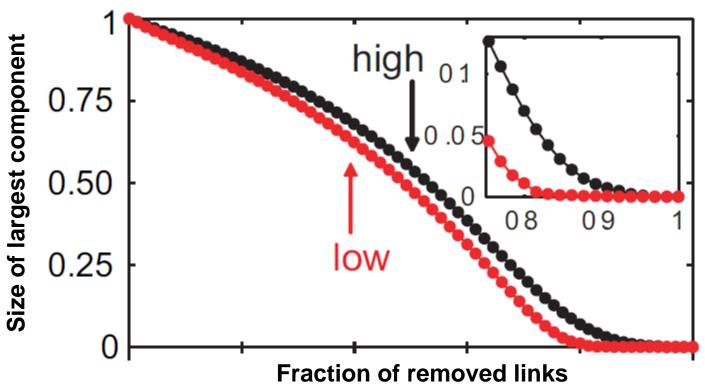
- Real edge strengths in mobile call graph
 - Strong ties are more embedded (have higher overlap)

Real Net, Permuted Tie Strengths



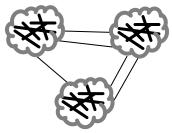
 Same network, same set of edge strengths but now strengths are randomly shuffled

Link Removal by Strength



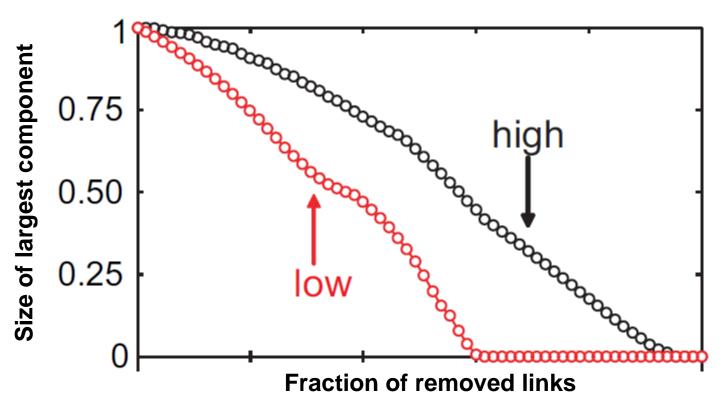
Low disconnects the network sooner

- Removing links by strength (#calls)
 - Low to high
 - High to low



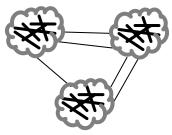
Conceptual picture of network structure

Link Removal by Overlap



Low disconnects the network sooner

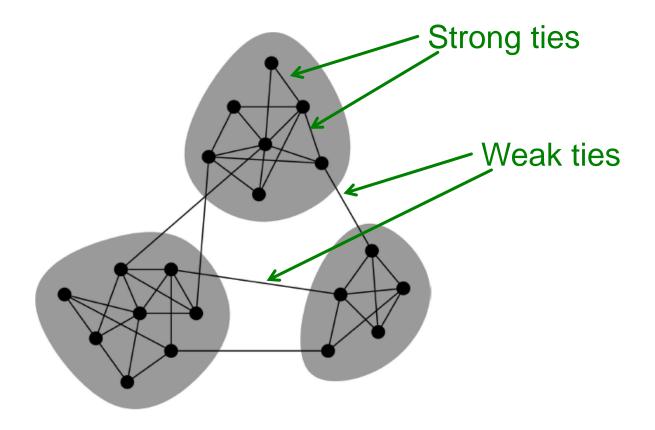
- Removing links based on overlap
 - Low to high
 - High to low



Conceptual picture of network structure

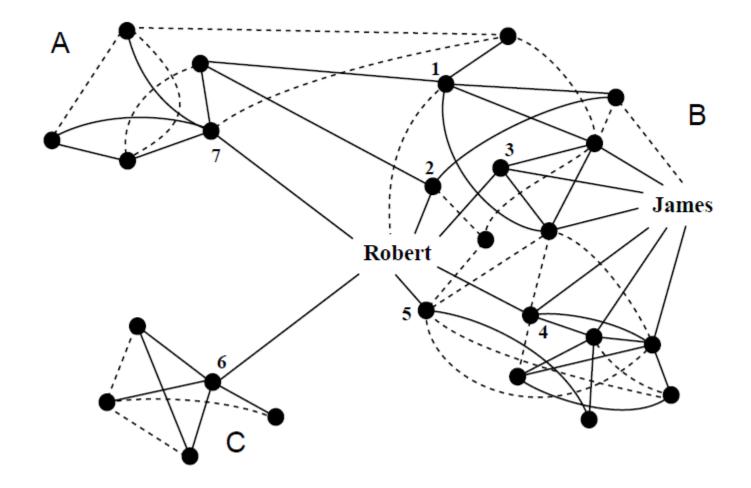
Conceptual Picture of Networks

 Granovetter's theory leads to the following conceptual picture of networks

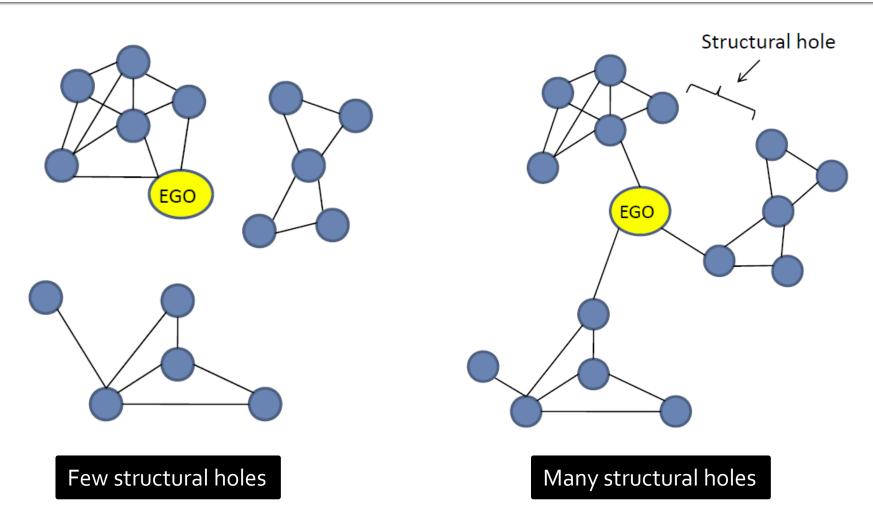


Small Detour: Structural Holes

Small Detour: Structural Holes



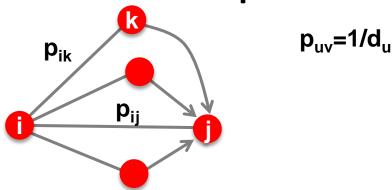
Structural Holes



Structural Holes provide ego with access to novel information, power, freedom

Structural Holes: Network Constraint

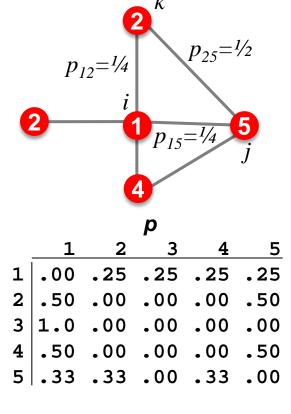
- The "network constraint" measure [Burt]:
 - To what extent are person's contacts redundant



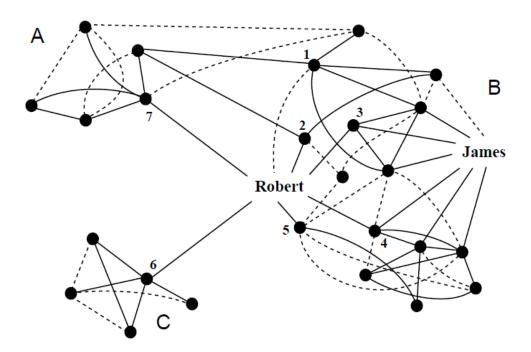
- Low: disconnected contacts
- High: contacts that are close or strongly tied

$$c_i = \sum_{j} c_{ij} = \sum_{j} \left[p_{ij} + \sum_{k} (p_{ik} p_{kj}) \right]^2$$

 $p_{uv} \dots$ prop. of u's "energy" invested in relationship with \boldsymbol{v}



Example: Robert vs. James

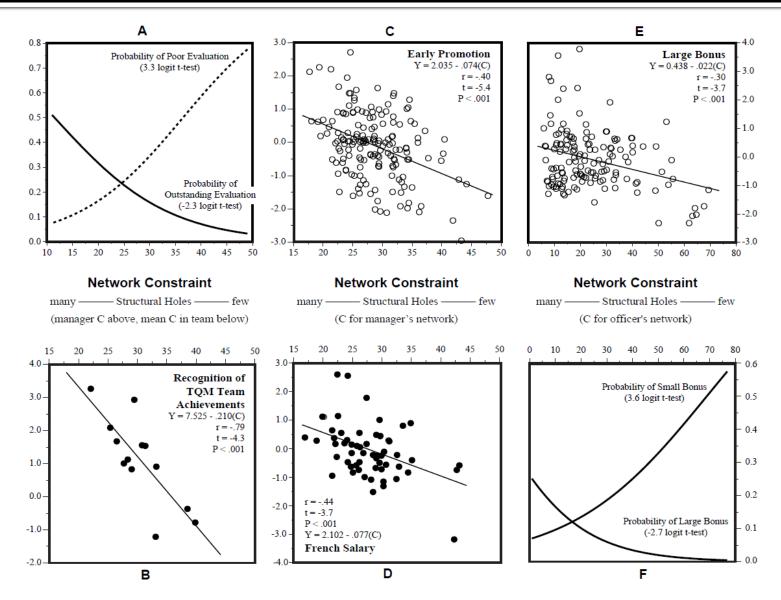


Network constraint:

- James: $c_I = 0.309$
- Robert: $c_R = 0.148$

- Constraint: To what extent are person's contacts redundant
 - Low: disconnected contacts
 - High: contacts that are close or strongly tied

Spanning the Holes Matters

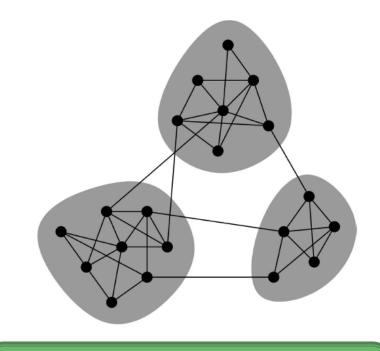


Network Communities

Network Communities

Granovetter's theory

 (and common sense)
 suggest that networks
 are composed of
 tightly connected
 sets of nodes



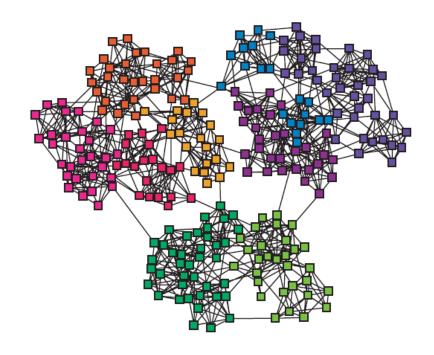
Communities, clusters, groups, modules

Network communities:

 Sets of nodes with lots of connections inside and few to outside (the rest of the network)

Finding Network Communities

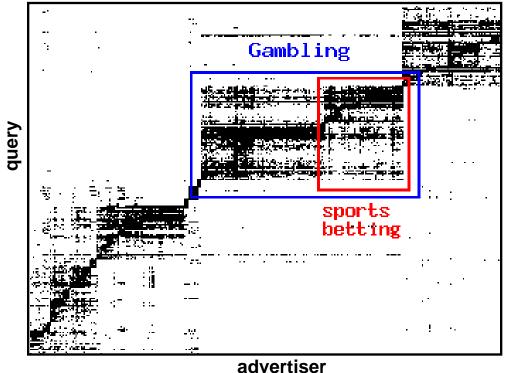
- How to automatically find such densely connected groups of nodes?
- Ideally such automatically detected clusters would then correspond to real groups
- For example:



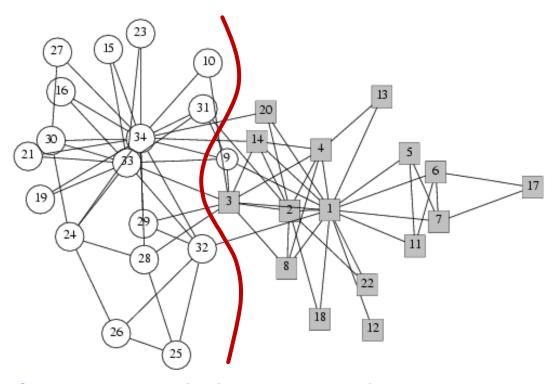
Communities, clusters, groups, modules

Micro-Markets in Sponsored Search

Find micro-markets by partitioning the "query x advertiser" graph:



Social Network Data

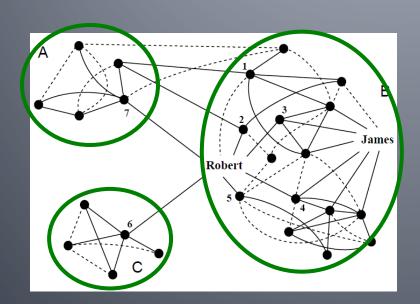


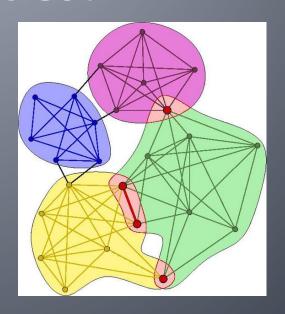
Zachary's Karate club network:

- Observe social ties and rivalries in a university karate club
- During his observation, conflicts led the group to split
- Split could be explained by a minimum cut in the network

Community Detection

How to find communities?

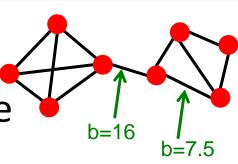




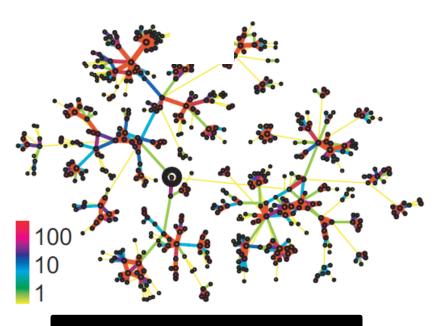
We will work with undirected (unweighted) networks

Method 1: Strength of Weak Ties

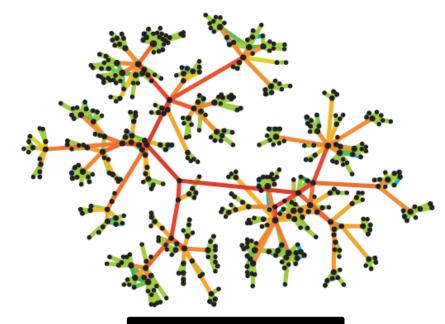
Edge betweenness: Number of shortest paths passing over the edge



Intuition:



Edge strengths (call volume) in real network



Edge betweenness in real network

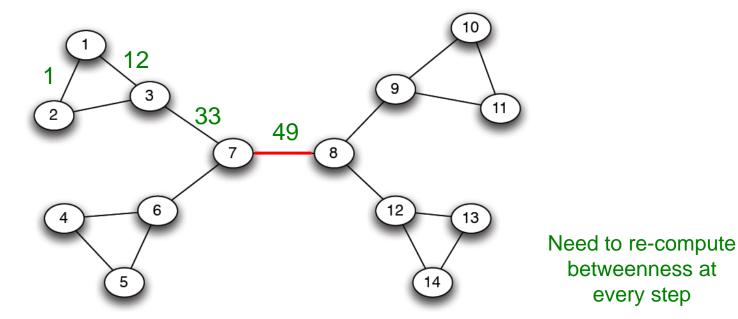
Method 1: Girvan-Newman

 Divisive hierarchical clustering based on the notion of edge betweenness:

Number of shortest paths passing through the edge

- Girvan-Newman Algorithm:
 - Undirected unweighted networks
 - Repeat until no edges are left:
 - Calculate betweenness of edges
 - Remove edges with highest betweenness
 - Connected components are communities
 - Gives a hierarchical decomposition of the network

Girvan-Newman: Example

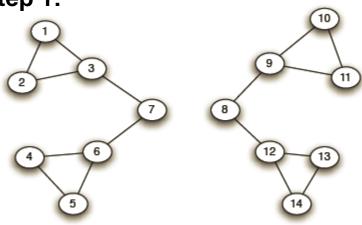


betweenness at

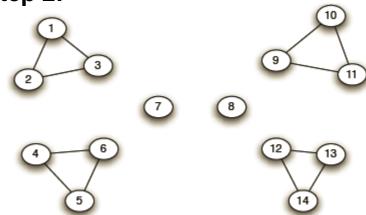
every step

Girvan-Newman: Example

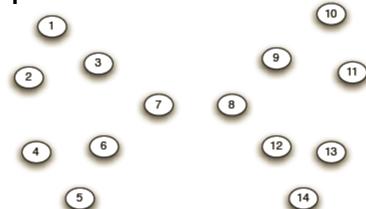
Step 1:



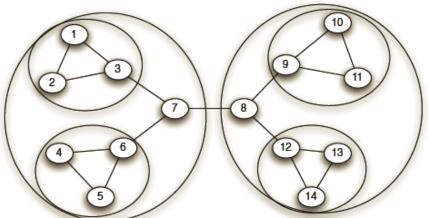
Step 2:



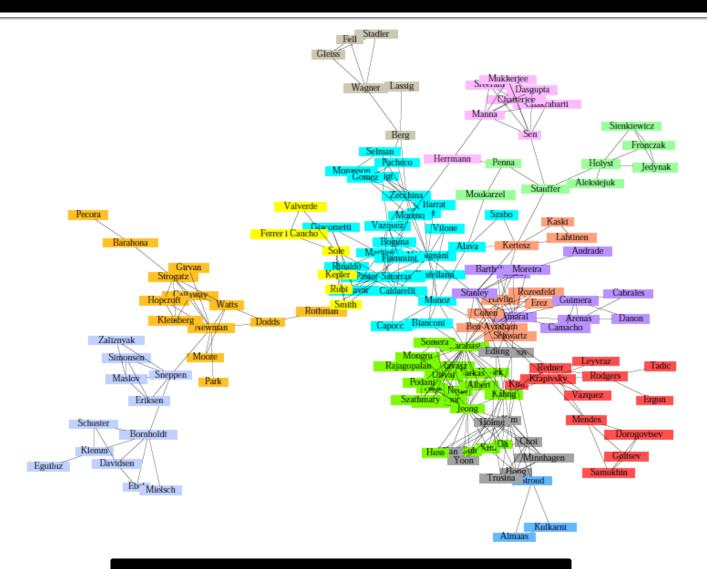
Step 3:



Hierarchical network decomposition:



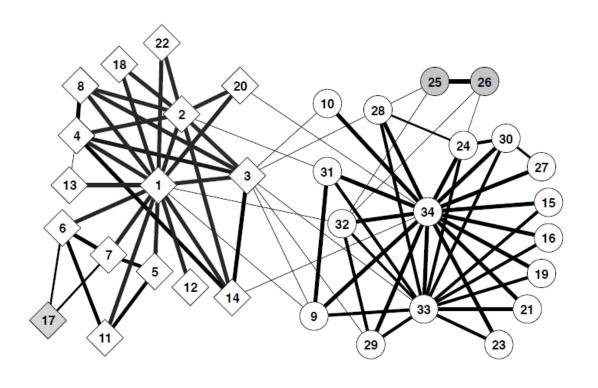
Girvan-Newman: Results

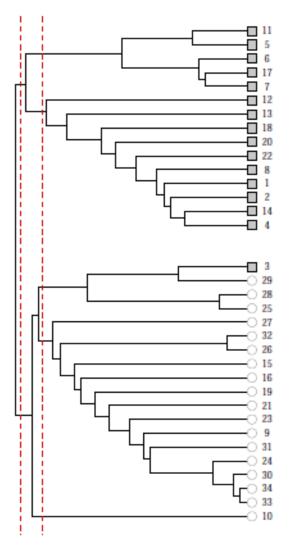


Communities in physics collaborations

Girvan-Newman: Results

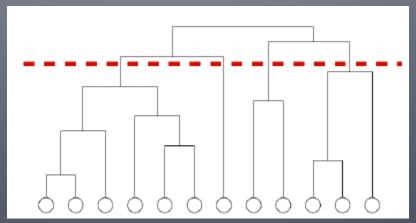
Zachary's Karate club:
 Hierarchical decomposition



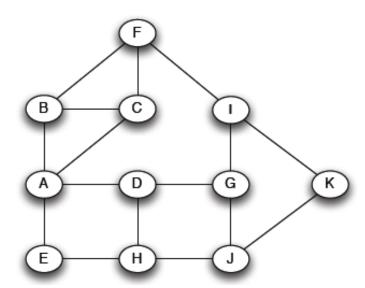


We need to resolve 2 questions

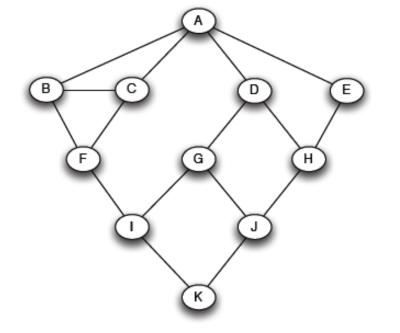
- 1. How to compute betweenness?
- 2. How to select the number of clusters?



 Want to compute betweenness of paths starting at node A

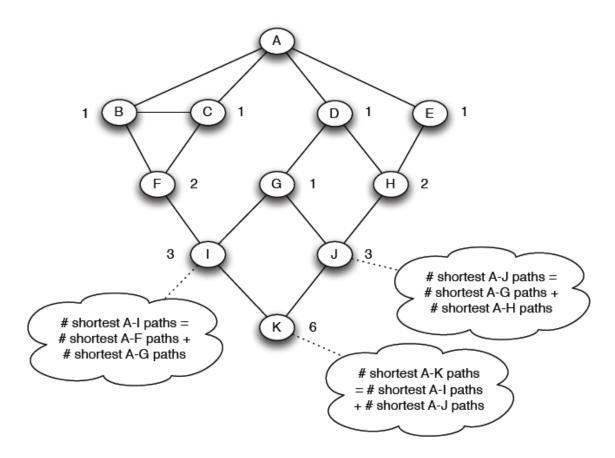


Breath first search starting from A:



11/12/2012

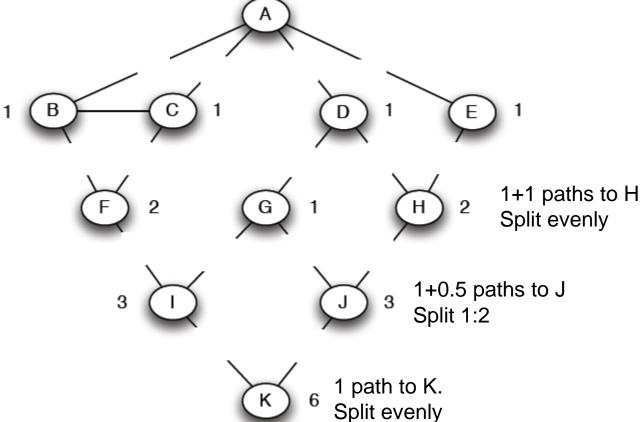
Count the number of shortest paths from A to all other nodes of the network:



Compute betweenness by working up the tree: If there are multiple paths count them fractionally

The algorithm:

- •Add edge **flows**:
 - -- node flow = 1+∑child edges
- -- split the flow up based on the parent value
- Repeat the BFS procedure for each starting node *U*

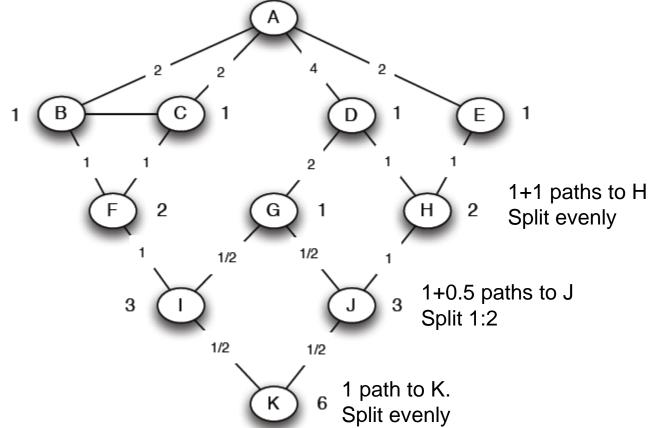


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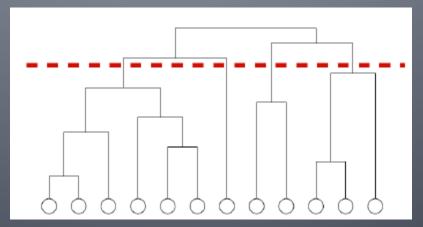
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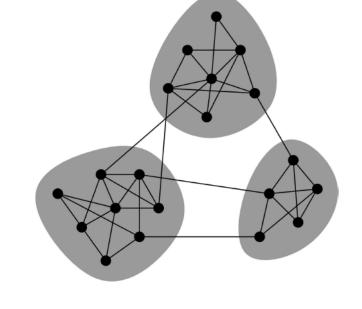
We need to resolve 2 questions

- 1. How to compute betweenness?
- 2. How to select the number of clusters?



Network Communities

- Communities: sets of tightly connected nodes
- Define: Modularity Q
 - A measure of how well a network is partitioned into communities
 - Given a partitioning of the network into groups $s \in S$:

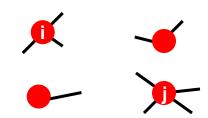


$$Q \propto \sum_{s \in S} [$$
 (# edges within group s) – (expected # edges within group s)]

Need a null model!

Null Model: Configuration Model

- Given real G on n nodes and m edges,
 construct rewired network G'
 - Same degree distribution but random connections



- Consider G' as a multigraph
- The expected number of edge between nodes i and j of degrees k_i and k_j equals to: $k_i \cdot \frac{k_j}{2m} = \frac{k_i k_j}{2m}$
 - The expected number of edges in (multigraph) G':

$$= \frac{1}{2} \sum_{i \in \mathbb{N}} \sum_{j \in \mathbb{N}} \frac{k_i k_j}{2m} = \frac{1}{2} \cdot \frac{1}{2m} \sum_{i \in \mathbb{N}} k_i \left(\sum_{j \in \mathbb{N}} k_j \right) =$$

$$= \frac{1}{4m} 2m \cdot 2m = m$$

Note: $\sum_{u \in N} k_u = 2m$

Modularity

Modularity of partitioning C of graph G:

• Q $\propto \sum_{s \in S} [$ (# edges within group s) – (expected # edges within group s)]

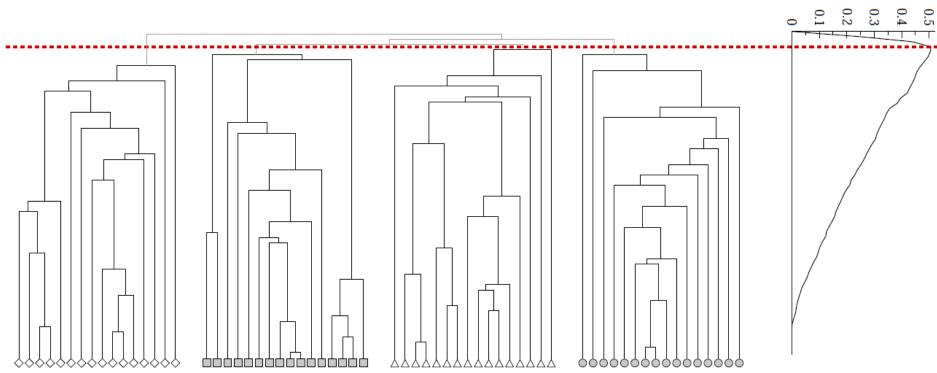
$$Q(G,S) = \underbrace{\frac{1}{2m} \sum_{S \in S} \sum_{i \in S} \sum_{j \in S} \left(A_{ij} - \frac{k_i k_j}{2m} \right)}_{\text{Normalizing cost.: -1$$

■ Modularity values take range [-1,1]

- It is positive if the number of edges within groups exceeds the expected number
- 0.3<Q<0.7 means significant community structure

Modularity: Number of clusters

 Modularity is useful for selecting the number of clusters:



Why not optimize modularity directly?

modularity