Alternative Comparison in Underspecified Degree Operators

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Today I will

▶ Revisit the recurrent ambiguities between

Comparison, Additivity, and Continuation

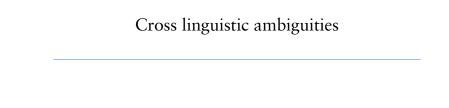
 Demonstrate an analysis based on a comparative meaning that compares structurally derived alternatives

Today I will

▶ Revisit the recurrent ambiguities between

Comparison, Additivity, and Continuation

Demonstrate an analysis based on a comparative meaning that compares structurally derived alternatives



The amount comparative in English, more, has an additive reading:

- (1) John bought three apples. ... Mary bought more (apples).
 - a. ⋄→ Mary bought more than three apples. (comparative reading)
 - b. \rightsquigarrow Mary bought apples, in addition to what John bought.

(additive reading)

In German, this additive meaning can be expressed by *noch*, which also has a different, continuative reading:

(2) Otto had **noch** einen Schnapps getrunken.

Otto had noch one Schnapps drunk

"Otto had another Schnapps."

(additive reading)

(3) Es regnet noch.

It raining noch

"It is still raining."

(continuative reading)

Romanian *mai* has a *three-way* ambiguitiy:

- Ion e mai intelligent decat Petre.
 - John is mai intelligent than Petre. "John is more intelligent than Petre."
- mai citi un roman. (5) Ion va John AUX mai read a novel
 - "John will read another nove."
- Ion mai merge la biblioteca. (6)
- John *mai* goes at library.
 - "John still goes to the library." (continuative reading)

(comparative reading)

(additive reading)

Empirical landscape (Thomas 2018):

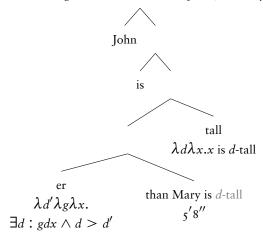
- Ambiguitities between comparison, additivity, and continuation are attested in a diverse set of languages.
- None of these languages allows for ambiguity between comparison and continuation to the exclusion of additivity.

	comparison	additivity	continuation
✓ (Vietnamese)	A	В	С
✓ (English, French)	A B		В
✓(German, Hungarian)	A	В	
*(unattested)	A	В	A
✓(Romanian)	A		

Brief review of the literature

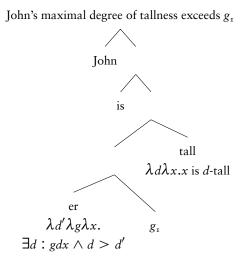
For the majority of the linguistic literature, comparative constructions express a relation between two degrees:

John's maximal degree of tallness exceeds 5'8" (i.e. Mary's height)



Implication on *incomplete comparatives*?

A fair hypothesis is the overt standard is replaced by a degree pro-form:



- □ additive *more* can be captured as a derived measure function of events:
 (cf. Greenberg 2010, Thomas 2010)
 - (7) $\begin{bmatrix} \text{more } \end{bmatrix}^{\text{T}} := \\ \lambda d\lambda Q\lambda P\lambda e. \exists x : [Qx \land P(x, e) \land \mu(h(e)) = d] \land \\ \partial (\exists e', P', d', y : [Qy \land P'(y, e') \land \mu(h(e')) = d']) \land \\ \exists e'' \exists P'' \exists z : [Qz \land P''(z, e'') \land z = x + y \land \mu(h(e'')) = d + d']$
 - 8) [John bought two more apples] \simeq
- continuative operators like *still* are typically associated with a scale determined by its containing context:
 - (9) [] noch/still [] := $\lambda S \lambda x' \lambda x \lambda P \cdot \partial (x' \prec_S x \wedge Px') \wedge Px$

¹In this talk I use the partiality operator ∂ (Beaver & Krahmer 2001) to indicate presupposition: $\partial(p) = 1$ iff p = 1, otherwise $\partial(p) = \#$.

The issue: wildly different lexical entries (but see Feldscher 2017², hard to see how to establish any logical connection between the three meanings and explain the recurrent ambiguities.

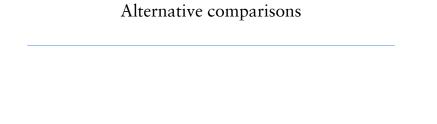
²Feldscher (2017) proposes a way to derive the additive reading from the comparative reading, but didn't discuss the continuative readings.

Previous proposal: a re-analysis of the comparative couched in scale segment semantics (cf. Schwarzschild 2013)

A scale segment is an abstract entity, which provides a structured representation for degree-related meanings:

(10) A scale segment
$$\sigma$$
 is a quadruple $\langle u, v, >_{\sigma}, \mu_{\sigma} \rangle$
(11) $\llbracket \text{Mary is taller than John} \rrbracket := \exists \sigma. \text{START}(\sigma, \mu_{\sigma} \mathbf{j}) \land \nearrow \sigma \land \mu_{\sigma} = \text{HT} \land \text{END}(\sigma, \mu_{\sigma} \mathbf{m})$

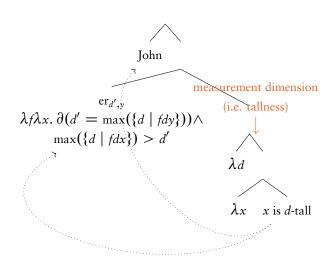
- We'll circle back to this proposal.
- ▶ It's still worthwhile to consider an apporach that does without scale segments.



Li (2021): Comparatives compare two things of the same type (i.e. two alternatives) on a locally derived measurement dimension (cf. Heim 1985, Bhatt & Takahashi 2007)

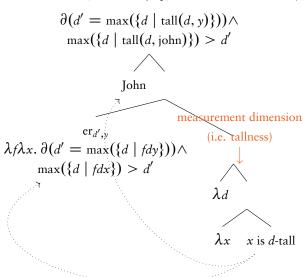
John $\operatorname{er}_{d',y}$ $\lambda f \lambda x. \, \partial (d' = \max(\{d \mid f d y\})) \wedge$ $\max(\{d \mid fdx\}) > d'$ λx x is d-tall

Li (2021): Comparatives compare two things of the same type (i.e. two alternatives) on a locally derived measurement dimension (cf. Heim 1985, Bhatt & Takahashi 2007)



Li (2021): Comparatives compare two things of the same type (i.e. two alternatives) on a locally derived measurement dimension

(cf. Heim 1985, Bhatt & Takahashi 2007)



er may be licensed by any scope-takers in the sentence, generating a comparison about the variables they bind:

- (12) Mary is 6 ft d' tall. ... Today I finally met a taller_{d',y} woman. \rightsquigarrow [a [er_{d',y} $\lambda x \lambda d$ [x[d-tall woman]]] determiner
- (13) John criticized^P five^{d'} books. ... He PRAISED more_{d',P}. \rightsquigarrow [PRAISED [er_{d',P} $\lambda P \lambda d$ [d-many books λz [He Pz]]] predicate
- (14) This boat is 20 ft^{d'} long. ... I thought it was longer_{d',w\infty}. \rightsquigarrow [I thought_{w\infty} [er_{d',w\infty} $\lambda w \lambda d$ [it was_w d-long]] intensional Op
- (15) John^y criticized^P five^{d'} books. ... Mary PRAISED more_{d',P,y}. \rightsquigarrow [Mary [PRAISED [er_{d',P,y} $\lambda P \lambda x \lambda d [d$ -many books $\lambda z [x P z]]]]

 multi licensors³$

$$\partial(d' = \max(\{d \mid fdy_0...y_n\})) \wedge \lambda f\lambda x_0...\lambda x_n.\max(\{d \mid fdx_0...x_n\}) > d'$$

³Technically, for this we need to adjust the meaning of *er* to a more general one:

Only restriction for possible comparisons: the standard degree must be the measurement of the standard alternative on the locally derived dimension.

Proposal for CAC ambiguities: we can compositionally derive the meaning of additive/continuative meaning from the comparative, because both meanings can be cashed out using alternative comparisons.

Deriving additivity by summing up the alternatives:

$$\partial(d' = \max\{d \mid y \text{ bought } d\text{-many apples}\}) \land$$

$$\max\{d \mid \text{ john bought } d\text{-many apples}\} > d'$$

$$\text{john}^u \qquad \lambda x. \partial(d' = \max\{d \mid y \text{ bought } d\text{-many apples}\}) \land$$

$$\max\{d \mid x \text{ bought } d\text{-many apples}\} > d'$$

$$\text{er}_{d',y}$$

$$\lambda d\lambda x \qquad x \text{ bought } d\text{-many apples}$$

(Mary bought three apples. ...) John bought more apples.
→ John bought more apples than Mary.

Deriving additivity by summing up the alternatives:

john"
$$\lambda x.\partial(d' = \max\{d \mid y \text{ bought } d\text{-many apples}\}) \land$$

$$\max\{d \mid x \oplus y \text{ bought } d\text{-many apples}\} > d'$$

$$\lambda x.\partial(d' = \max\{d \mid y \text{ bought } d\text{-many apples}\}) \land$$

$$\lambda f \lambda x. f(x \oplus y) \qquad \max\{d \mid x \text{ bought } d\text{-many apples}\} > d'$$

$$\operatorname{er}_{d',y} \qquad \dots$$

(Mary bought three apples. ...) John bought more apples. → John and Mary bought more apples than Mary alone. Deriving continuation as a presupposed additive comparison:

$$\operatorname{impf}(\operatorname{rain})(\operatorname{pres}) \wedge \partial(\operatorname{ADD}_{t'}(\operatorname{er}_{n',t'}(\lambda n \lambda t.\operatorname{impf}(\operatorname{rain})t \wedge n \leq_{\operatorname{impf}(\operatorname{rain})}t))(\operatorname{pres}))$$

$$\operatorname{ADD}_{t'}$$

$$\operatorname{er}_{n',t'} \qquad \operatorname{\lambda P \lambda f \lambda Q \lambda u. f u \wedge n} \leq_{f} u))(u) \qquad t \quad \operatorname{impf}(\operatorname{rain})$$

$$\begin{aligned} & \text{ADD}_{t'}(\text{er}_{n',t'}(\lambda n \lambda t.\text{impf}(\text{rain})t \land n \leq_{\text{impf}(\text{rain})} t))(\text{pres})^{4} \\ &= \partial (n' = \max \left\{ n | \text{impf}(\text{rain})t' \land n \leq_{\text{impf}(\text{rain})} t' \right\}) \land \\ &\max \left\{ n | \text{impf}(\text{rain})(\text{pres} \oplus t') \land n \leq_{\text{impf}(\text{rain})} (\text{pres} \oplus t') \right\} > n' \end{aligned}$$

 $^4n \leq_f u := fu \models_c fn$; for any two propositions $p, q, p \models_c q$ iff $\forall w \in c : pw \to qw$.

$$\longrightarrow \lambda t. \exists e : raine \land t \subseteq \tau(e)$$

ADD_{t'} (er_{n',t'} (
$$\lambda n \lambda t$$
, impf(rain) $t \wedge n \leq_{\text{impf(rain)}} t$)) (pres)⁴

$$= \partial (n' = \max \{ n | \text{impf(rain)} t' \wedge n \leq_{\text{impf(rain)}} t' \}) \wedge$$

$$\max \{ n | \text{impf(rain)} \text{ (pres } \oplus t') \wedge n \leq_{\text{impf(rain)}} \text{ (pres } \oplus t') \} > n'$$

$$= \partial (n' = \max \{ n | \exists e : \text{raine} \wedge t' \subseteq \tau(e) \wedge n \text{ is a subinterval of } t' \}) \wedge \max$$

$$\{ n | \exists e : \text{raine} \wedge \text{ (pres } \oplus t') \subseteq \tau(e) \wedge n \text{ is a subinterval of (pres } \oplus t') \} > n'$$

 $^{^4}n <_f u := fu \models_c fn$; for any two propositions $p, q, p \models_c q$ iff $\forall w \in c : pw \to qw$.

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$$\{ n | \exists e : \text{raine} \land \text{ (pres } \oplus t') \subseteq \tau(e) \land n \text{ is a subinterval of (pres } \oplus t') \} > n'$$

$$= \partial (\exists e : \text{raine} \land t' \subseteq \tau(e) \land n' = t') \land \exists e : \text{raine} \land \text{ (pres } \oplus t') \subseteq \tau(e) \land \max \{ n | n \text{ is a subinterval of (pres } \oplus t') \} > n'$$

 $^{^4}n \leq_f u := fu \models_c fn$; for any two propositions $p, q, p \models_c q$ iff $\forall w \in c : pw \to qw$.

```
ADD_{t'}(er_{n',t'}(\lambda n\lambda t.impf(rain)t \land n \leq_{impf(rain)} t)) (pres)^{4}
= \partial(n' = \max \{n | impf(rain)t' \land n \leq_{impf(rain)} t'\}) \land
```

$$\max \left\{ n \middle| \operatorname{impf}(\operatorname{rain}) \left(r \middle| \operatorname{rain} \right) \left($$

=
$$\partial(n' = \max\{n | \exists e : \text{raine} \land t' \subseteq \tau(e) \land n \text{ is a subinterval of } t'\}) \land \max\{n | \exists e : \text{raine} \land (\text{pres} \oplus t') \subseteq \tau(e) \land n \text{ is a subinterval of } (\text{pres} \oplus t')\} > n'\}$$

$$= \partial(\exists e : \text{raine} \land t' \subseteq \tau(e) \land n' = t') \land \exists e : \text{raine} \land (\text{pres} \oplus t') \subseteq \tau(e) \land (\text{pres} \oplus t') \land (\text{p$$

$$\tau(e) \wedge \max\{n \mid n \text{ is a subinterval of } (\text{pres} \oplus t')\} > n'$$

$$=\partial(\exists e : raine \wedge t' \subseteq \tau(e)) \wedge$$

$$\exists e : raine \land (pres \oplus t') \subseteq \tau(e) \land (pres \oplus t') > t'$$

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 $\rightarrow n$ is a sub-interval of

ADD_{t'} (er_{n',t'} (\lambda n \lambda t.impf(rain) t \lambda n \leq_{impf(rain)} t)) (pres)⁴

$$= \partial(n' = \max \left\{ n | \text{impf}(\text{rain}) t' \wedge n \leq_{\text{impf}(\text{rain})} t' \right\}) \wedge$$

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$$= \partial(n' = \max \left\{ n | \exists e : \text{raine} \wedge t' \subseteq \tau(e) \wedge n \text{ is a subinterval of } t' \right\}) \wedge \max \left\{ n | \exists e : \text{raine} \wedge (\text{pres} \oplus t') \subseteq \tau(e) \wedge n \text{ is a subinterval of } (\text{pres} \oplus t') \right\} > n'$$

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$$= \partial(\exists e : \text{raine} \wedge t' \subseteq \tau(e)) \wedge \exists e : \text{raine} \wedge (\text{pres} \oplus t') \subseteq \tau(e) \wedge t' < \text{pres}$$

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$$= \partial (n' = \max \left\{ n | \exists e : raine \wedge t' \subseteq \tau(e) \wedge n \text{ is a subinterval of } t' \right\}) \wedge \max n' = t'$$

$$\{ n | \exists e : raine \wedge (pres \oplus t') \subseteq \tau(e) \wedge n \text{ is a subinterval of } (pres \oplus t') \right\} > n'$$

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 $= \partial(\exists e : raine \land t' \subseteq \tau(e)) \land \exists e : raine \land (pres \oplus t') \subseteq \tau(e) \land t' \prec pres$

 $=\partial(\exists e : raine \wedge t' \subset \tau(e)) \wedge$

 $\exists e : raine \land (pres \oplus t') \subseteq \tau(e) \land (pres \oplus t') > t'$

$$\operatorname{impf}(\operatorname{rain})(\operatorname{pres}) \wedge \partial(\operatorname{ADD}_{t'}(\operatorname{er}_{n',t'}(\lambda n \lambda t.\operatorname{impf}(\operatorname{rain})t \wedge n \leq_{\operatorname{impf}(\operatorname{rain})}t)$$



$$\operatorname{impf}(\operatorname{rain})(\operatorname{pres}) \wedge \partial(\operatorname{ADD}_{t'}(\operatorname{er}_{n',t'}(\lambda n \lambda t.\operatorname{impf}(\operatorname{rain})t \wedge n \leq_{\operatorname{impf}(\operatorname{rain})}t))(\operatorname{pres}))$$

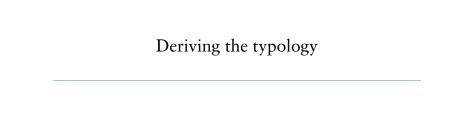
t))(pres))

$$t))(pres))$$
= $\exists e : raine \land pres \subseteq \tau(e) \land$

 $\partial(\partial(\exists e : raine \land t' \subseteq \tau(e)) \land \exists e : raine \land (pres \oplus t') \subseteq \tau(e) \land t' \prec pres)$

 $\operatorname{impf}(\operatorname{rain})(\operatorname{pres}) \wedge \partial(\operatorname{ADD}_{t'}(\operatorname{er}_{n',t'}(\lambda n \lambda t.\operatorname{impf}(\operatorname{rain})t \wedge n \leq_{\operatorname{impf}(\operatorname{rain})}t)$

- t))(pres)) $=\exists e : raine \land pres \subseteq \tau(e) \land$ $\partial(\partial(\exists e : raine \land t' \subseteq \tau(e)) \land \exists e : raine \land (pres \oplus t') \subseteq \tau(e) \land t' \prec pres)$
 - Assertion: it is raining now.
 - Presupposition: the raining has continued from an earlier time t'.
 - Implicature: the speaker can't assert that the rain will continue to a time later than now.



Distributed Morphology (Halle & Marantz 1993):

- ▶ The terminals of syntactic structures are **morphemes**: sets of features without phonological content.
- Subset Principle: a morpheme, i.e. a set of features, is spelt out by the lexical item that matches its greatest subset of features (Halle 2000).

CAC operators (e.g. more, noch, still)

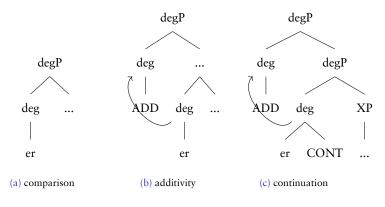
are phonological realizations of a deg head (cf. Thomas 2018).

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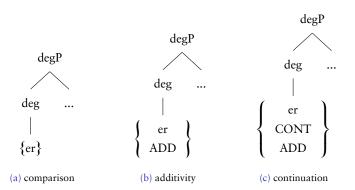


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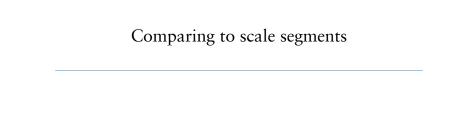


Generating the typological distribution:

Comp./Add. Cont.	English: $\{er\} \leftrightarrow er, \{er, CONT, ADD\} \leftrightarrow still$	
Comp. Add./Cont.	German: $\{er\} \leftrightarrow mehr$, $\{er, ADD\} \leftrightarrow noch$	
Comp. /Add. /Cont.	Romanian: {er} ↔ mai	
Comp. Add. Cont.	$\label{eq:Vietnamese: er} \mbox{Vietnamese: } \{\mbox{er}\} \leftrightarrow \mbox{hon, } \{\mbox{er, ADD}\} \leftrightarrow \mbox{n\~ua, } \{\mbox{er, CONT, ADD}\} \leftrightarrow \mbox{van}$	

Explaining the implicational universal:

- \bowtie are the phonological realization of both {er} and {er, CONT, ADD} $\rightarrow \alpha$ is the item matching the biggest subset of {er, ADD}
- ${\color{red} \triangleright} \ \ i.e. \ Comparison/Continuation \rightarrow Comparison/Additivity/Continuation \\$



In scale segment semantics:

▶ Instead of denoting a relation between degrees and individuals, adjectives denote a predicate of scale segments.

(16)
$$[tall] := \lambda \sigma. \mu_{\sigma} = \text{HEIGHT}$$

Components of the comparison (i.e. the target and the standard of the comparison, the differential) are treated as modifiers of the scale segment.

```
\exists \sigma. START(\sigma, \mu_{\sigma} j) \land \nearrow \sigma \land \mu_{\sigma} = HT \land
                     \Delta \sigma = 2 \text{in} \wedge \text{END}(\sigma, \mu_{\sigma} \text{m})
                                            \lambda \sigma.START(\sigma, \mu_{\sigma}i) \wedge \nearrow \sigma \wedge \mu_{\sigma} = HT
\lambda \sigma. \exists \sigma : f(\sigma)
                                                       \wedge \Delta \sigma = 2 \text{in} \wedge \text{END}(\sigma, \mu_{\sigma} \text{m})
                                                   Mary
                                                                                     \lambda_{\nu}\lambda_{\sigma}.START(\sigma, \mu_{\sigma}i) \wedge \nearrow \sigma \wedge
                                                                          \mu_{\sigma} = HT \wedge \Delta \sigma = 2in \wedge END(\sigma, \mu_{\sigma} y)
                                                      m
                                                                     END
                                                                                                          \lambda \sigma.START(\sigma, \mu_{\sigma}j) \wedge \nearrow \sigma \wedge \mu_{\sigma} = HT \wedge \Delta \sigma = 2in
                                                   \lambda_y \lambda_\sigma.END(\sigma, \mu_{\sigma y})
                                                                                                         \lambda \sigma. \Delta \sigma > 2in
                                                                                                                                                                         \lambda \sigma.\text{START}(\sigma, \mu_{\sigma} \mathbf{j}) \wedge \nearrow \sigma \wedge \mu_{\sigma} = \text{HT}
                                                                                             2 inches
                                                                                                                                      DIFF
                                                                                                                                                                                                                     \lambda \sigma. START(\sigma, \mu_{\sigma}) \wedge \nearrow \sigma
                                                                                                  2in
                                                                                                                          \lambda n \lambda \sigma . \Delta \sigma > n
                                                                                                                                                                         \lambda \sigma. \mu_{\sigma} = HT
                                                                                                                                                                                                                        RISE
                                                                                                                                                                                                                                                     \lambda \sigma.START(\sigma, \mu_{\sigma}j)
                                                                                                                                                                                                                  λσ. / σ
                                                                                                                                                                                                                                             START
                                                                                                                                                                                                                                                                                           John
                                                                                                                                                                                                                            \lambda x \lambda \sigma.START(\sigma, \mu_{\sigma} x)
```

(17) Mary is two inches taller than John

$$:= \exists \sigma. START(\sigma, \mu_{\sigma} \mathbf{j}) \land \nearrow \sigma \land \mu_{\sigma} = \mathsf{HT} \land$$

measurement and ends at Mary's measurement.

 $\Delta\sigma=2{
m in}\wedge{
m END}(\sigma,\mu_\sigma{
m m})$ \leadsto There is a rising scale segment of height that starts from John's

```
\exists \sigma. \nearrow \sigma \land \mu_{\sigma} = \text{COUNT} \land \text{START}(\sigma, \mu_{\sigma g_1}) \land
       END(\sigma, \mu_{\sigma}(\oplus(\{x \mid applesx \land john bought x\})))
                                         \lambda \sigma. \nearrow \sigma \wedge \mu_{\sigma} = \text{COUNT} \wedge \text{START}(\sigma, \mu_{\sigma}g_{i}) \wedge
    \lambda f.\exists \sigma : f(\sigma)
                                          END(\sigma, \mu_{\sigma}(\oplus(\{x \mid applesx \land john bought x\})))
                                                                                       \lambda x.john bought x
                                                                               apples
                           AMT
                                                                            \lambda x \lambda \sigma. \nearrow \sigma \wedge \mu_{\sigma} = COUNT \wedge
\lambda \Sigma \lambda P \lambda Q \lambda \sigma. \Sigma (\oplus (P \cap Q))(\sigma)
                                                                           START(\sigma, \mu_{\sigma g_1}) \wedge END(\sigma, \mu_{\sigma x})
                                                                 END
                                                                                                    \lambda \sigma. \nearrow \sigma \wedge \mu_{\sigma} = COUNT \wedge START(\sigma, \mu_{\sigma g_1})
                                                 \lambda x \lambda \sigma.END(\sigma, \mu_{\sigma} x)
                                                                                                              COUNT
                                                                                                                                                   \lambda \sigma. \nearrow \sigma \wedge START(\sigma, \mu_{\sigma g_1})
                                                                                                   \lambda \sigma, \mu_{\sigma} = COUNT
                                                                                                                                                         RISE
                                                                                                                                                                                \lambda \sigma.START(\sigma, \mu_{\sigma g_i})
                                                                                                                                                   λσ. 🖊 σ
                                                                                                                                                                             START
```

 $\lambda x \lambda \sigma.START(\sigma, \mu_{\sigma} x)$

(18) [John bought more apples] (comparative reading)

 $:= \exists \sigma. \nearrow \sigma \land \mu_{\sigma} = \text{COUNT} \land \text{START}(\sigma, \mu_{\sigma}g_1) \land \\ \text{END}(\sigma, \mu_{\sigma}(\bigoplus(\{x \mid \text{apples}x \land \text{john bought }x\})))$

 $\operatorname{END}(\sigma, \mu_{\sigma}(\bigoplus(\{x \mid \operatorname{apples} x \land \operatorname{john bought} x\}))$ \rightsquigarrow There is a rising scale segment of quantity that starts from the measurement of the apples John bought and ends at the measurement of some antecedent apples.

(19) ADD := $\lambda \Sigma \lambda \Sigma' \lambda x \lambda \sigma . \Sigma(\sigma)(g_1) \wedge \Sigma'(\sigma)(x \oplus g_1)$

```
\exists \sigma. \nearrow \sigma \land \mu_{\sigma} = \text{COUNT} \land \text{START}(\sigma, \mu_{\sigma}(g_i)) \land
     END(\sigma, \mu_{\sigma}(\bigoplus \{\{x \mid applesx \land john bought x\})) \bigoplus g_{\tau})
                                              \lambda \sigma. \nearrow \sigma \wedge \mu_{\sigma} = \text{COUNT} \wedge \text{START}(\sigma, \mu_{\sigma}(g_{\tau})) \wedge
      \lambda f. \exists \sigma : f(\sigma)
                                             END(\sigma, \mu_{\sigma}(\oplus(\{x \mid applesx \land john bought x\})) \oplus g_{\tau})
                                                                                                 \lambda x.john bought x
                                                                                         apples
                                                                                          \lambda x \lambda \sigma. \nearrow \sigma \wedge \mu_{\sigma} = COUNT \wedge
                             AMT
\lambda \Sigma \lambda P \lambda O \lambda \sigma. \Sigma (\oplus (P \cap O))(\sigma)
                                                                                START(\sigma, \mu_{\sigma}(g_x)) \wedge END(\sigma, \mu_{\sigma}(x \oplus g_x))
                                                                                                                         \lambda \Sigma' \lambda_x \lambda \sigma. \nearrow \sigma \wedge \mu_{\sigma} = COUNT \wedge
                                                                                    END
                                                                   \lambda x \lambda \sigma. END(\sigma, \mu_{\sigma} x)
                                                                                                                         START(\sigma, \mu_{\sigma}(g_{\tau})) \wedge \Sigma'(\sigma)(x \oplus g_{\tau})
                                                                                                                                                                                                  COUNT [RISE [START]]
                                                  \lambda \Sigma_{r \to lr} \lambda \Sigma'_{r \to lr} \lambda_x \lambda_\sigma . \Sigma(\sigma)(g_i) \wedge \Sigma'(\sigma)(x \oplus g_i)
                                                                                                                                                                      \lambda x \lambda \sigma. \nearrow \sigma \wedge \mu_{\sigma} = COUNT \wedge START(\sigma, \mu_{\sigma} x)
```

(20) [John bought more apples] (additive reading)

 $\exists \sigma. \nearrow \sigma \land \mu_{\sigma} = \text{COUNT} \land \text{START}(\sigma, \mu_{\sigma}(g_1)) \land \\ \text{END}(\sigma, \mu_{\sigma}(\bigoplus \{\{x \mid \text{apples}x \land \text{john bought } x\}\})) \bigoplus g_1)$

There is a rising scale segment of quantity that starts from the measurement of some antecedent apples and ends with the measurement of the antecedent apples and the apples John bought.

(21) [it is still raining] :=

$$\exists e. \mathsf{rain}(e) \land t_e \subseteq au(e) \land$$

$$\partial(\exists \sigma \exists \epsilon \exists t' [\mathsf{rain}(\epsilon) \land t' \subseteq \tau(\epsilon) \land \neg \mathsf{INIT}(e, \epsilon) \land \mu_{\sigma} = \mathsf{STAGE}_{\epsilon} \land \nearrow \sigma \land \mu_{\sigma} = \mathsf{COUNT} \land \mathsf{START}(\sigma, \mu_{\sigma}(g_1))$$

$$SIAGE_{\epsilon} \land / \sigma \land \mu_{\sigma} = COUN1 \land SIAR1(\sigma, \mu_{\sigma}(g_1)) \land END(\sigma, \mu_{\sigma}(e \oplus g_1)])$$

- Assertion: now is within the duration of a raining event.
- Presupposition: there is a rising scale segment of event development that starts from the measurement of some antecedent event and ends with the sum of this antecedent event and the current event.

(22) CON :=
$$\lambda \Sigma \lambda R \lambda e \lambda t. R(e)(t) \wedge \partial (\exists \sigma \exists \epsilon \exists t' [R(\epsilon)(t') \wedge \neg INIT(e, \epsilon) \wedge \mu_{\sigma} = STAGE_{\epsilon} \wedge \Sigma(e)(\sigma)])$$

$$\exists e. \operatorname{rain}(e) \land t_e \subseteq \tau(e) \land \\ \partial (\exists \sigma \exists e \exists l' [\operatorname{rain}(e) \land t' \subseteq \tau(e) \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{STAGE}_{e} \land \\ \nearrow \sigma \land \mu \sigma = \operatorname{COUNT} \land \operatorname{START}(\sigma, \mu_{\sigma}(g_i)) \land \operatorname{END}(\sigma, \mu_{\sigma}(e \oplus g_i)]) \\ \Rightarrow \\ \partial (\exists \sigma \exists e \exists l' [\operatorname{rain}(e) \land t' \subseteq \tau(e) \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{STAGE}_{e} \land \\ \nearrow \sigma \land \mu \sigma = \operatorname{COUNT} \land \operatorname{START}(\sigma, \mu_{\sigma}(g_i)) \land \operatorname{END}(\sigma, \mu_{\sigma}(e \oplus g_i)]) \\ & \lambda R \lambda e \lambda t. R(e)(t) \land \\ \partial (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{STAGE}_{e} \land \\ \nearrow \sigma \land \mu \sigma = \operatorname{COUNT} \land \operatorname{START}(\sigma, \mu_{\sigma}(g_i)) \land \operatorname{END}(\sigma, \mu_{\sigma}(e \oplus g_i)]) \\ & \lambda \Sigma \lambda R_{r \to i \to j} \lambda e. \lambda t. R(e)(t) \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{STAGE}_{e} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{STAGE}_{e} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{COUNT} \land \\ \lambda (\exists \sigma \exists e \exists l' [R(e)(t') \land \neg \operatorname{INIT}(e, e) \land \mu_{\sigma} = \operatorname{INIT$$

Similarities with my proposal:

- ▶ The comparison meaning is captured as a comparison between two correlates.
- ▶ The logical connection between CAC meanings is derived by incrementally adding covert operators manipulating the correlates and the measurement dimension.

Difference: whether or not the measurement dimension is structurally derived.

Difference in prediction 1: infelicitous anaphoricity in amount comparatives.

Context: Mary bought three apples. [John bought more apples] :=

(23)
$$\exists \sigma. \nearrow \sigma \land \mu_{\sigma} = \text{COUNT} \land \text{START}(\sigma, \mu_{\sigma}(g_{\scriptscriptstyle 1})) \land \text{END}$$

 $(\sigma, \mu_{\sigma}(\bigoplus(\{x \mid \text{apples}x \land \text{john bought }x\})) \bigoplus g_{\scriptscriptstyle 1})$ (Thomas 2018) \rightsquigarrow comparing the first-mentioned three apples \bigoplus the apples John bought and the three apples.

(24)
$$\partial(d' = \max\{d \mid g_1 \text{ bought } d\text{-many apples}\}) \land \max\{d \mid \text{john} \oplus g_1 \text{ bought } d\text{-many apples}\} > d'$$
 (my proposal) \rightsquigarrow comparing Mary \oplus John and Mary in the apples they bought, presupposing the first mentioned quantity *three* is the number of apples Mary bought.

Only (24) makes the correct prediction in a context with added negation:

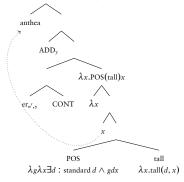
- (25) Mary didn't buy those three apples. ?? John bought more apples.
 - ▶ More in (25) doesn't have an additive (or comparative) reading.
 - (23) still generates the same felicitous meaning.

Difference in prediction 2: varieties of the continuative reading

- (26) Anthea is still tall.
 - a. \rightsquigarrow Anthea was tall at some earlier time. (temporal reading)
 - b. → Anthea is only marginally tall. (marginal reading)
 - Continuative operators like *still* across languages are systematically ambiguous between a variety of flavors.
 - ▶ The scale segment approach in Thomas (2018): unclear how to derive these different flavors of non-temporal continuation, as the measurement dimension (event development) is hard-wired into the meaning of CON.

▶ My proposal can derive the marginal reading of (26): change the scale of the presupposed comparison by changing the scope property of CONT.

 $POS(tall)(anthea) \land \\ \partial(ADD_{y}(cr_{n',y}(\lambda n\lambda x.POS(tall)x \land n \leq_{POS(tall)}x))(anthea))$



POS(tall)(anthea)
$$\wedge \partial$$
(ADD_y(er_{n',y}($\lambda n \lambda x$.POS(tall) $x \wedge n \leq_{POS(tall)} x$))(anthea))

 $\Rightarrow \exists d : \text{standard } d \land \text{tall}(d, \text{anthea}) \land \\ \partial(\max \{n \mid POS(\text{tall})(a \oplus y) \land n \leq_{POS(\text{tall})} (a \oplus y)\} > \\ \max \{n \mid POS(\text{tall})y \land n \leq_{POS(\text{tall})} y\})$

 $= \exists d$: standard $d \land tall(d, anthea) \land$

- a. Assertion: Anthea is tall.
 - b. Presupposition: An alternative individual *y* is tall and taller than Anthea.

 $\partial(\exists d : \text{standard } d \land \text{tall}(d, \text{anthea} \oplus y) \land y \text{ is taller than anthea})$

c. Implicature: People shorter than Anthea are not tall (i.e. Anthea is only marginally tall).

- > We can generate different readings of the same sentence by having different scope configurations, explaining the ambiguity of (28).
 - (28) can still explain Exercise two to Peter.
 - a. Focusing *Peter* \rightsquigarrow Paul is beyond my help.
 - b. Focusing $two \rightsquigarrow$ Exercise three is too hard.

[PETER [ADD_y [[er_{n',y} CONT]
$$\lambda x$$
 [I can explain ex. two to x]]]]

a. presupposed additive comparison:

a. presupposed additive comparison:
$$\max \{n \mid \text{I can explain ex. 2 to p} \oplus y \land n \leq_{\text{I can explain ex. 2 to }} (p \oplus y)\}$$

$$> \max \{n \mid \text{I can explain ex. 2 to } y \land n \leq_{\text{I can explain ex. 2 to }} y\}$$

- b. Assertion: I can explain ex. 2 to Peter.
 - Presupposition: I can also explain ex. 2 to an alternative individual y, and it is easier to do so than to Peter.
 - Implicature: for people who are ranked even lower on the scale (i.e. harder to teach than Peter), I may not be able to explain ex. 2 to them.

(30) [TWO [ADD_y [[er_{n',y} CONT]
$$\lambda x$$
 [I can explain ex.x to Peter]]]]

a. presupposed additive comparison:

max
$$\{n \mid I \text{ can explain ex. } 2 \oplus y \text{ to p } \land n \leq \lambda_{x,I \text{ can explain ex. } x \text{ to p }} (2 \oplus y)\}$$

> max $\{n \mid I \text{ can explain ex. } y \text{ to p } \land n \leq \lambda_{x,I \text{ can explain ex. } x \text{ to p }} y\}$

- - individual y, and it is easier to do so than to Peter.
 - ▶ Implicature: for people who are ranked even lower on the scale (i.e. harder to teach than Peter), I may not be able to explain ex. 2 to them.



Cross linguistically, we have degree operators ambiguous between comparison, additivity, and continuation.

A comparative meaning that compares correlates on a structurally alternatives, combined with the subset principle in Distributed Morphology, can explain the these ambiguities and their cross-linguistic distributions.

The proposal crucially differs from the previous analysis (Thomas 2018) in how the correlates and the measurement dimension is determined, and I have shown a structural approach makes better predictions.



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