732A54 Big Data Analytics 6hp

http://www.ida.liu.se/~732A54

Relational databases



Literature

Elmasri, Navathe, Fundamentals of Database Systems, 7th edition, Addison Wesley, 2016. Chapters 3-6 and 9; section 7.1.

Database methods

1. Representation and storage of data



Databases

One (of many) way(s) to store data in electronic form

 Used in every-day life: bank, reservation of hotel or journey, library search, shops



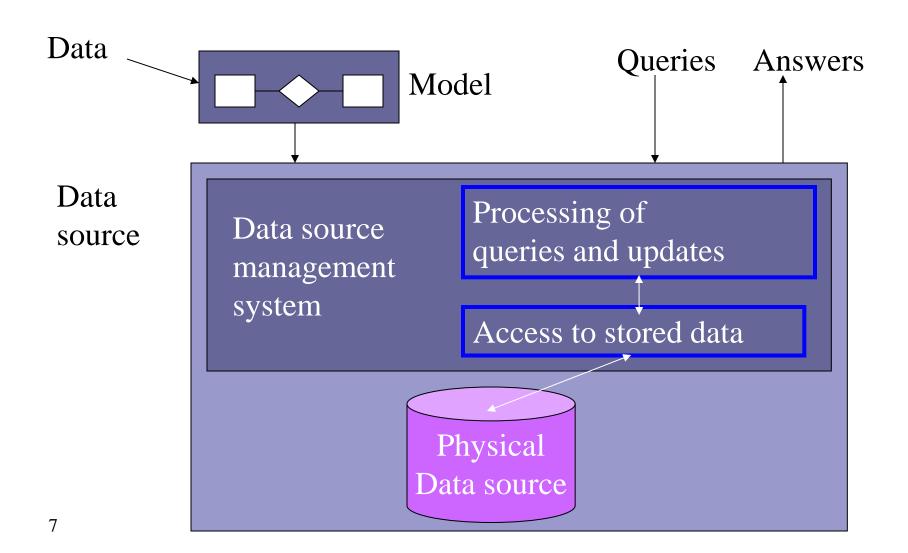
Databases

 Database management system (DBMS): a collection of programs that supports a user to create and maintain a database

database system = database + DBMS



Data Sources





Persons

- Data source administrator
- Data source designer
- 'end user'
- application programmer

- DBMS designer
- tool developer
- operator, maintenance



Issues - this course

- What information is stored?
- How is the information stored?
 (high level and low level)
- How is the information accessed?
 (user level and system level)

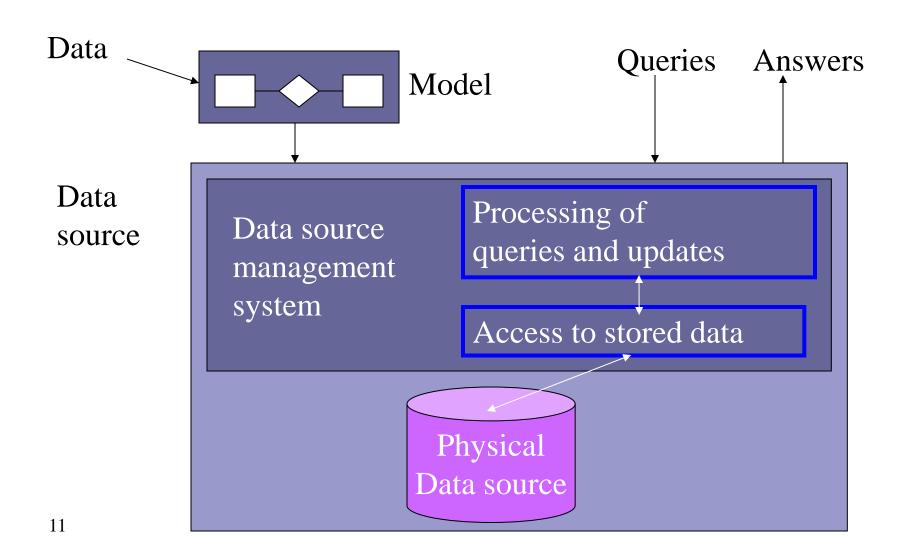


Other issues

- How to optimize performance of a data source?
- How to recover a data source after crash?
- How to access information from multiple data sources?
- How to allow multiple users to access a data source?



Data Sources



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Which information is stored?

- Model of reality
 - Entity-Relationship model (ER)
 - Unified Modeling Language (UML)



ER/EER diagram

- structured way to model data, independent of type of data source
- notions:
 - entities and entity types
 - attributes
 - key attributes
 - relationships and cardinality constraints
 - sub-types (EER)



DEFINITION

ACCESSION

SOURCE ORGANISM

REFERENCE

AUTHORS

TITLE

REFERENCE

AUTHORS

TITLE

Homo sapiens adrenergic, beta-1-, receptor

NM_000684

human

1

Frielle, Collins, Daniel, Caron, Lefkowitz,

Kobilka

Cloning of the cDNA for the human

beta 1-adrenergic receptor

2

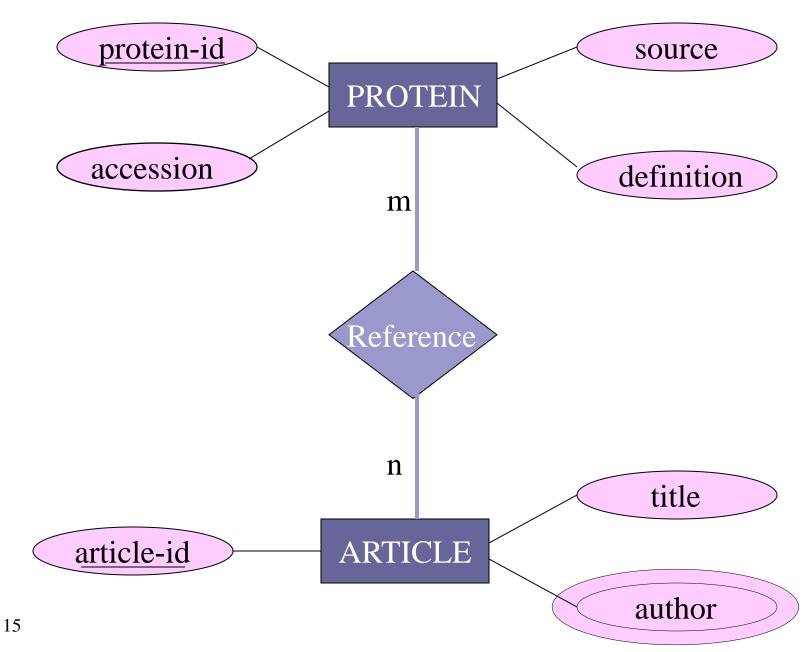
Frielle, Kobilka, Lefkowitz, Caron

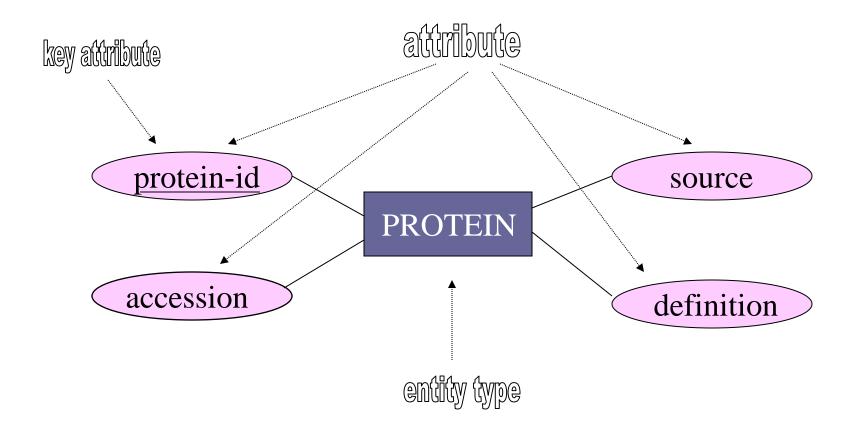
Human beta 1- and beta 2-adrenergic

receptors: structurally and functionally

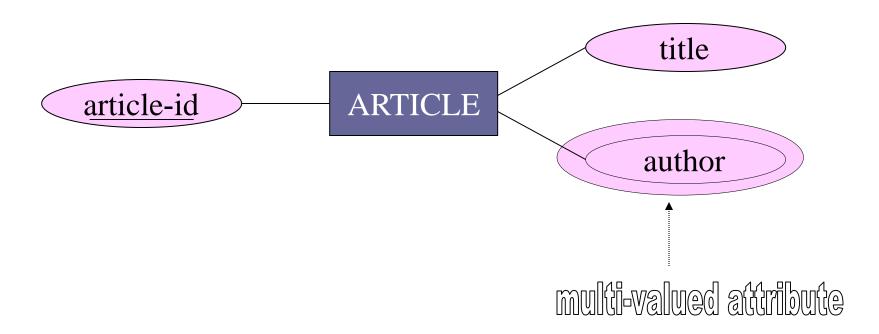
related receptors derived from distinct

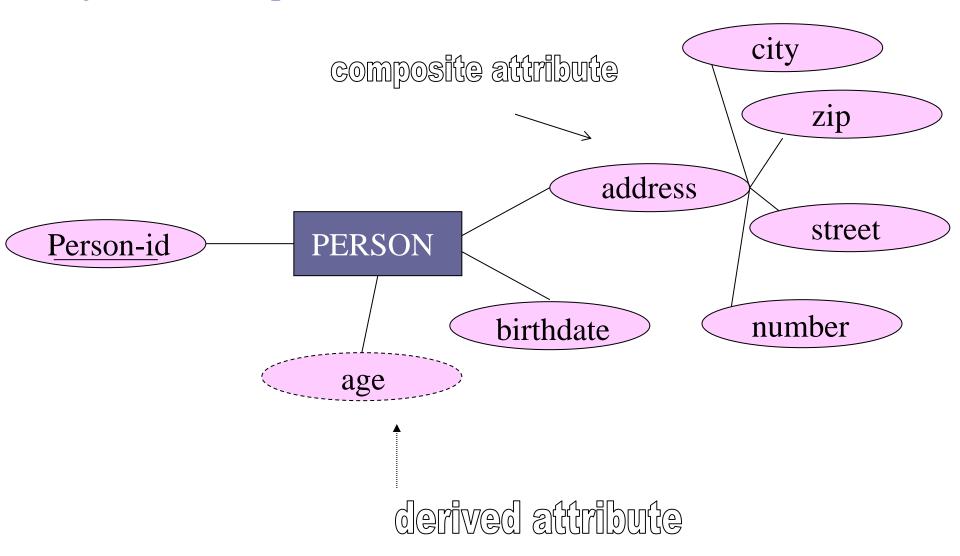
genes

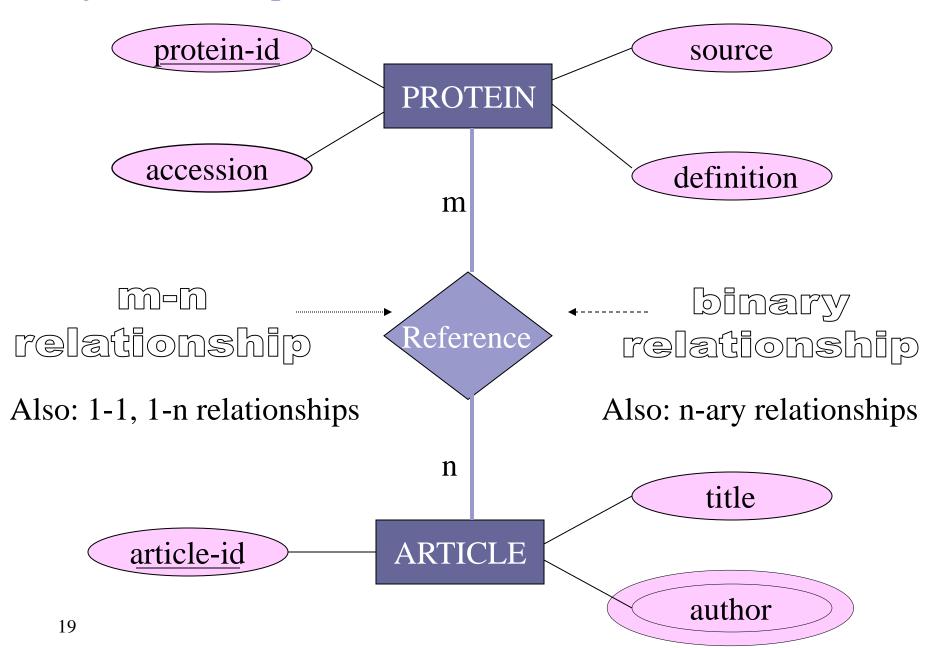


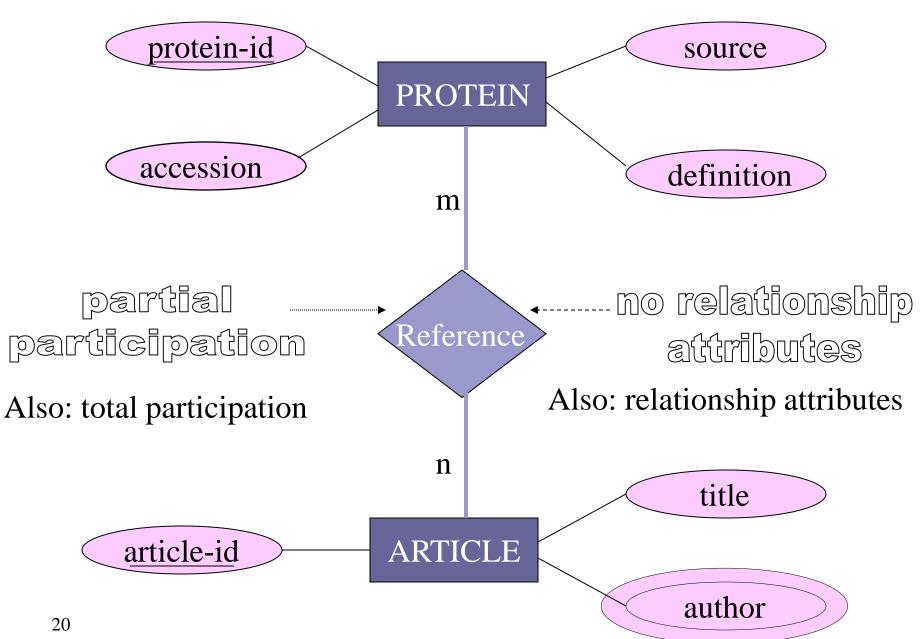


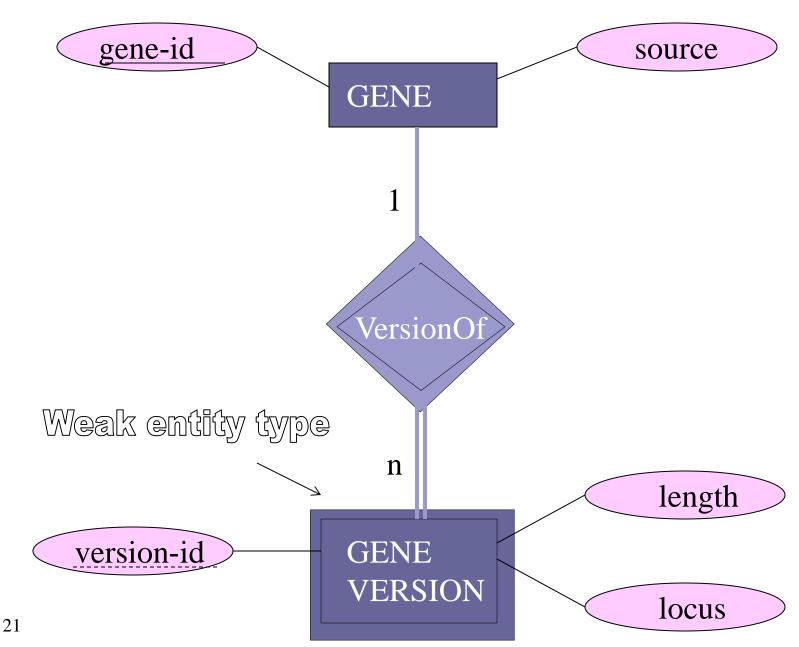
Weak entity type: type without key attribute





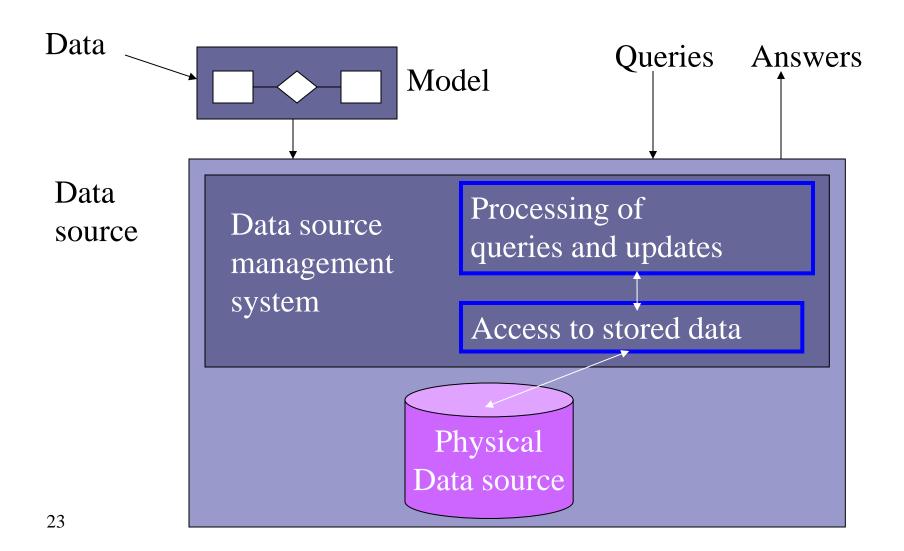






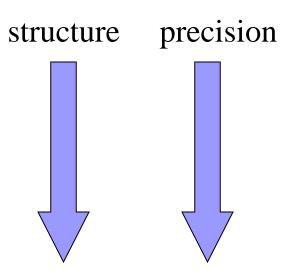
Enhanced Entity-relationship Birthdate P-id **PERSON** address name sub-type relationship department TEACHER

Data Sources



How is information stored? (high level) How is information accessed? (user level)

- Text (IR)
- Semi-structured data
- Data models (DB)
- Rules + Facts (KB)





Databases

- Relational databases:
 - model: tables + relational algebra
 - query language (SQL)

Object-oriented, extended-relational,
 NoSQL databases

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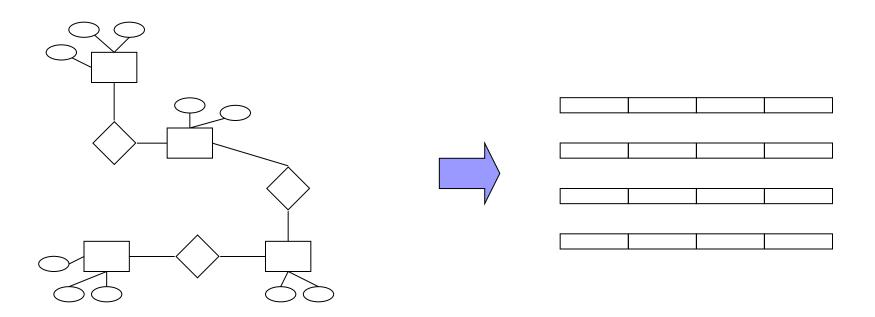
Relational Databases

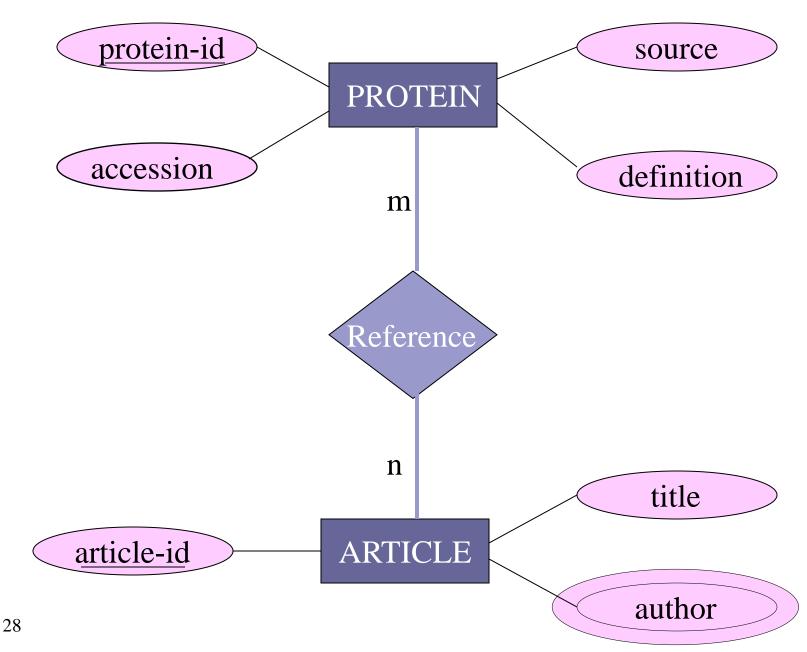
- Tables = relations (~ entity type)
 - row = tuple (~ entity)
 - column = attribute (~ attribute)

- primary keys
- foreign keys



ER/EER to database schema





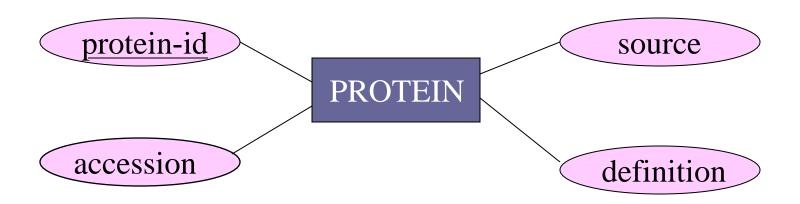


ER/EER to database schema

Step 1

For each (strong) entity E create a table R with the same simple attributes as the entity.





PROTEIN (PROTEIN-ID, ACCESSION, SOURCE, DEFINITION)





Author is multi-valued attribute.

ARTICLE-TITLE (ARTICLE-ID, TITLE)





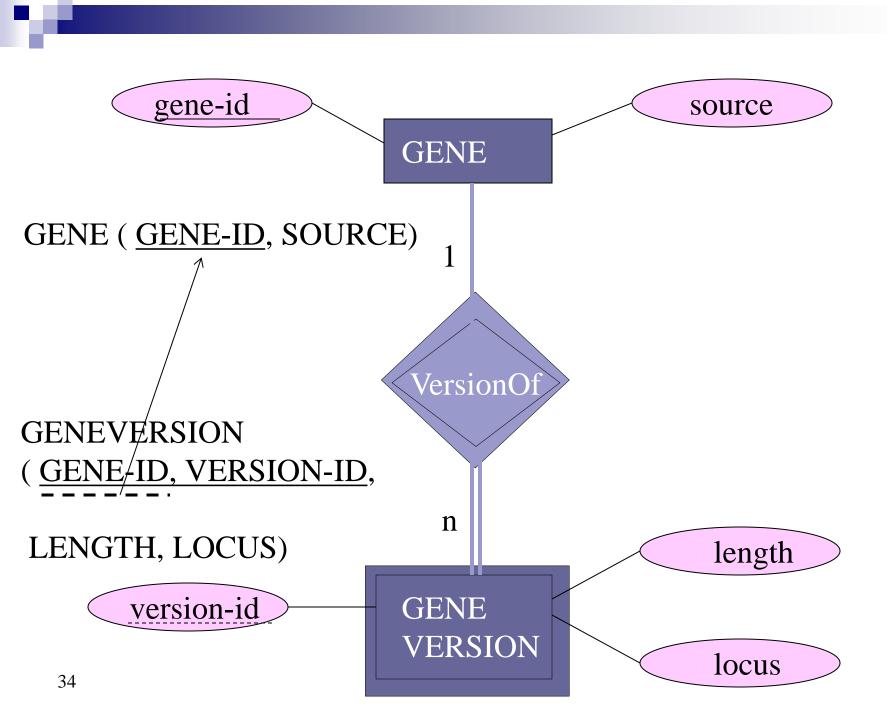
GENE (GENE-ID, SOURCE)



ER/EER to database schema

Step 2

For each weak entity W med owner entity E, create a table R with the same simple attributes as W and add the primary key attributes from the relation that corresponds to E.

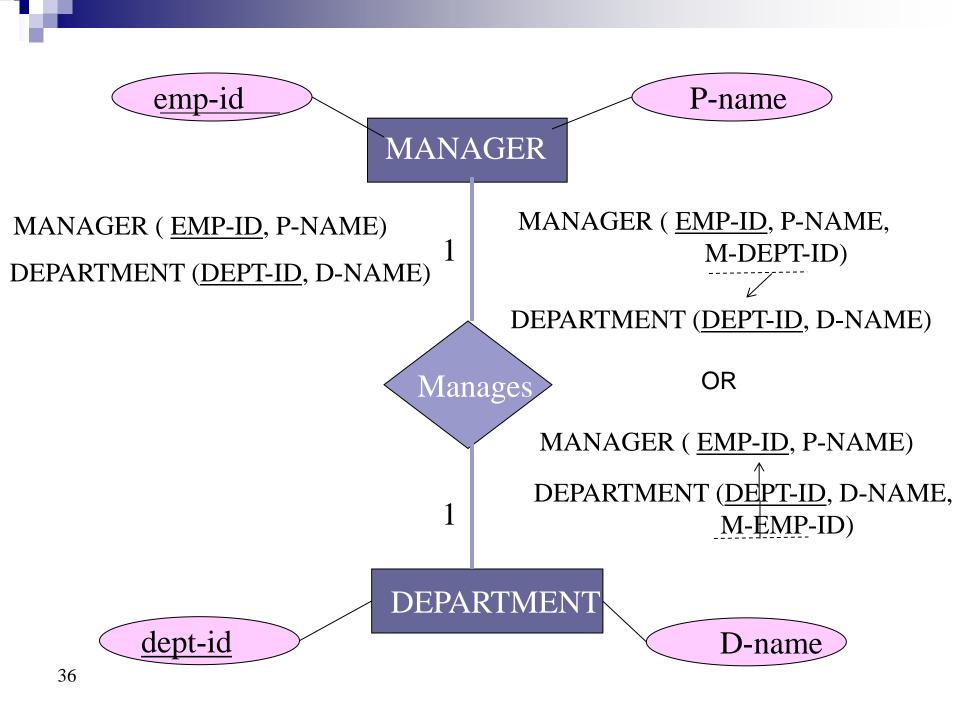




ER/EER to database schema

Step 3

For each binary 1:1 relationship between S and T, add the primary key of the table corresponding to one of S or T as a foreign key to the table corresponding to the other.

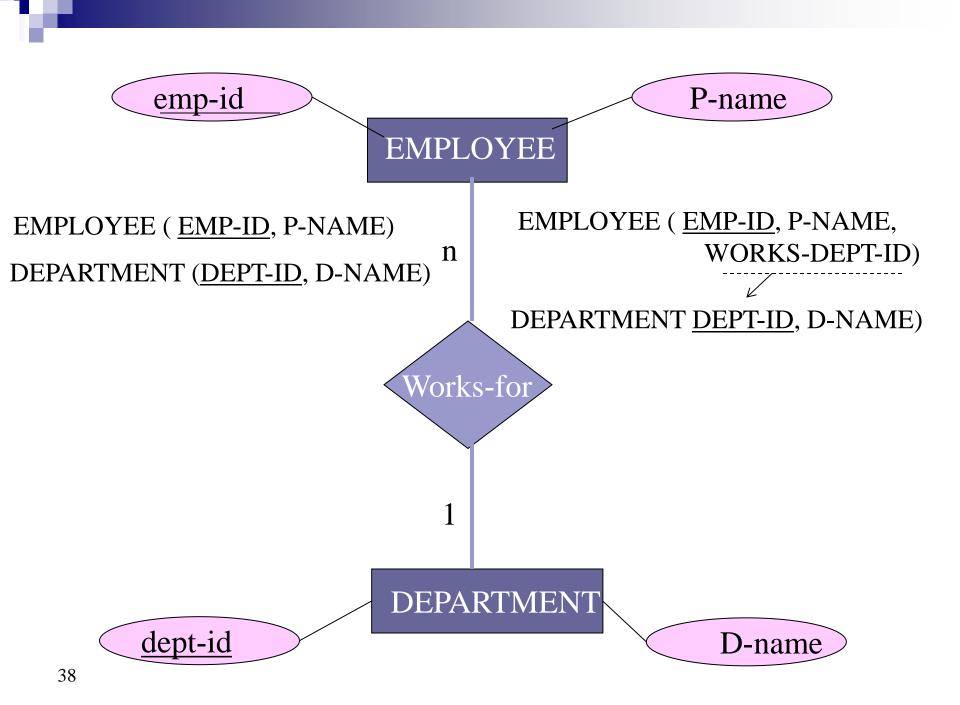




ER/EER to database schema

Step 4

For each binary 1: n relationship between S and T (every S is related to many T, every T is related to one S), add the primary key of the table corresponding to S as a foreign key to the table corresponding to T.

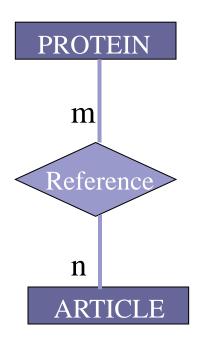




ER/EER to database schema

Step 5

For each binary m: n relationship between S and T create a table R with the primary keys of the tables corresponding to S and T as foreign keys. If the relationship has attributes, then add these to R.



REFERENCE (PROTEIN-ID, ARTICLE-ID)

PROTEIN (PROTEIN-ID) ARTICLE (ARTICLE-ID)

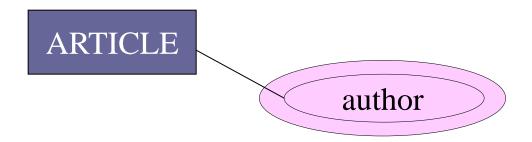


ER/EER to database schema

Step 6

For every multi-valued attribute A in R, create a new table that contains an attribute corresponding to A and the primary key of R as foreign key.





ARTICLE-AUTHOR (<u>ARTICLE-ID</u>, <u>AUTHOR</u>)

ARTICLE (ARTICLE-ID)

Relational databases

PROTEIN				REFERENCE		
PROTEIN-ID	ACCESSION	DEFINITION	SOURCE	PROTEIN-ID	ARTICLE-ID	
1	NM_000684	Homo sapiens adrenergic, beta-1-, receptor	human	1 1	1 2	

ARTICLE-AUTHOR ARTICLE-TITLE

ARTICLE-ID	AUTHOR	ARTICLE-ID	TITLE
1 1 1 1 1 2 2 2 2 2	Frielle Collins Daniel Caron Lefkowitz Kobilka Frielle Kobilka Lefkowitz Caron	2	Cloning of the cDNA for the human beta 1-adrenergic receptor Human beta 1- and beta 2- adrenergic receptors: structurally and functionally related receptors derived from distinct genes

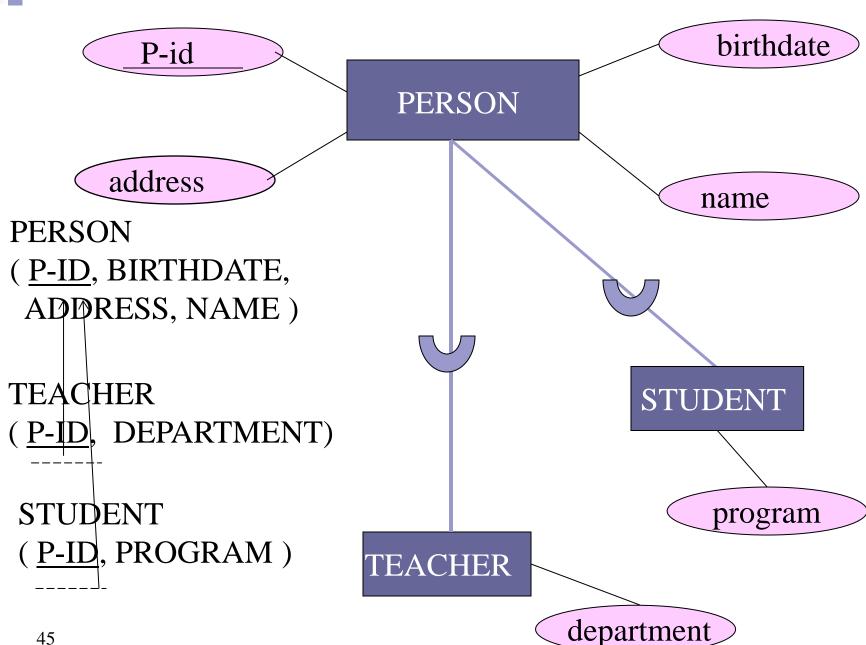


EER to database schema

Sub-type relationship

Assume type C has subtypes Si

option 1: Create a relation R for C with all attributes in C. Then create relation Ri for each Si with all attributes from Si and the primary key from R.





EER to database schema

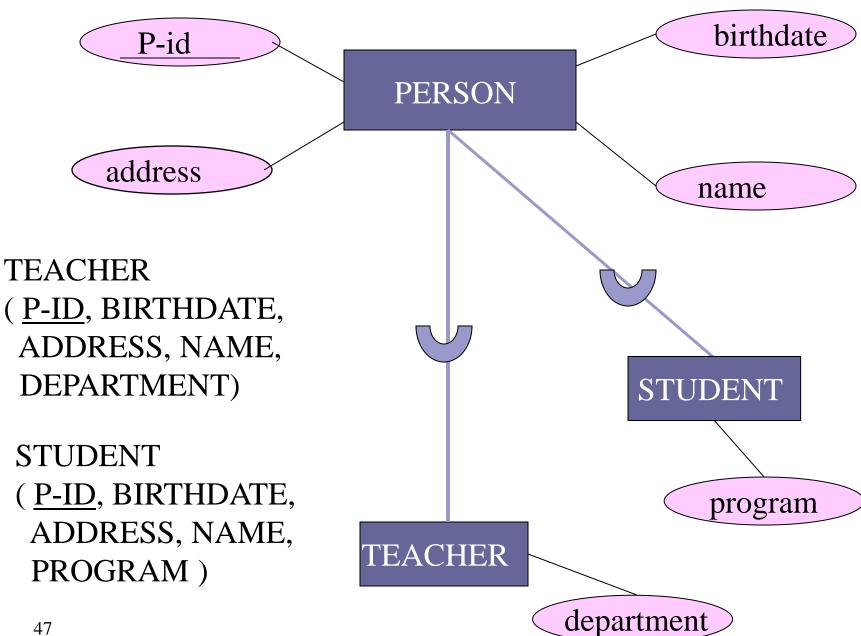
Sub-type relationship

Assume type C has subtypes Si

option 2: Create a relation Ri for each Si with as attributes all attributes from Si and from C.

Only works if every C belongs to a Si.







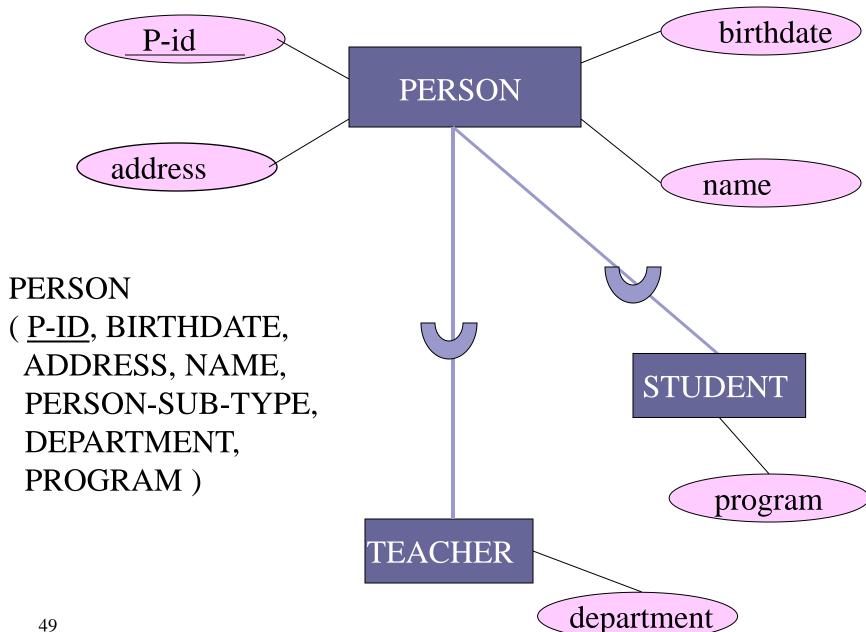
EER to database schema Sub-type relationship

Assume type C has subtypes Si

option 3: Create a relation R with all attributes from C and all attributes from the Si. Add an attribute to discriminate between the subtypes.

Only works for disjoint subtypes.







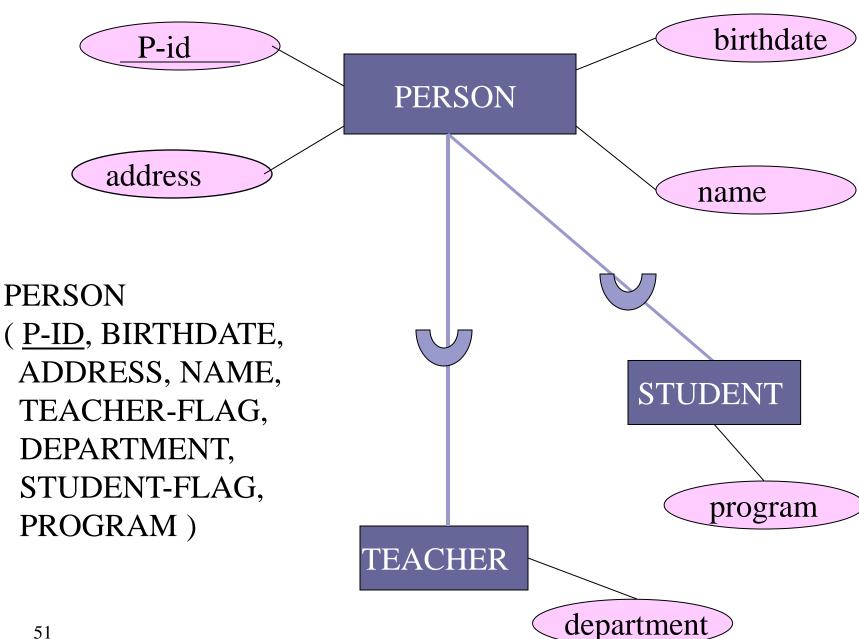
EER to database schema

Sub-type relationship

Assume type C has subtypes Si

option 4: Create a relation R with all attributes from C and all attributes from the Si. Add a flag attribute Fi to R for each Si.





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SQL

select source
from protein
where accession = 'NM_000684';

PROTEIN

PROTEIN-ID	ACCESSION	DEFINITION	SOURCE
1	NM_000684	Homo sapiens adrenergic, beta-1-, receptor	human



SQL

= reference.protein-id **and** reference.article-id

= article-title.article-id;

REFERENCE

ARTICLE-ID		
1		
2		

PROTEIN ARTICLE-TITLE

PROTEIN-ID	ACCESSION	DEFINITION	SOURCE	ARTICLE-ID	TITLE
1	1 NM 000684	Homo sapiens	human	1	Cloning of the
	1111_000001	adrenergic, beta-1-, receptor		2	Human beta 1

Database methods

2. Querying relational databases using SQL



SQL

- Developed by IBM Research as interface to System R. (197*, SEQUEL)
- QL, DDL and DML (queries, data definition, updates, views)
- used in many database systems
- table = relation, tuple = row, attribute = column



EMPLOYEE (FNAME, MINIT, LNAME, SSN, BDATE, ADDRESS, SEX, SALARY, SUPERSSN, DNO)

DEPT-LOCATIONS (<u>DNUMBER</u>, <u>DLOCATION</u>)

 DEPARTMENT (DNAME, <u>DNUMBER</u>, MGRSSN, MGRSTARTDATE)



■ WORKS-ON (ESSN, PNO, HOURS)

PROJECT (PNAME, <u>PNUMBER</u>, PLOCATION, DNUM)

DEPENDENT (<u>ESSN, DEPENDENT-</u> NAME, SEX, BDATE, RELATIONSHIP)

SQL syntax

SELECT attribute-list

FROM table-list

WHERE condition;

- attribute-list: attributes that are required
- table-list: tables that are needed to process the query
- condition: expression with logical operators (and, or, not) and equality and inequality operators; identifies the tuples that should be 58 retrieved



Q1: List SSN for all employees.

SELEC	TSSN
FROM	EMPLOYEE;

SSN
123456789
333445555
999887777 987654321
666884444
453453453
987987987
888665555

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 Q2: List birth date and address for all employees whose name is `John B. Smith'.

SELECT BDATE, ADDRESS
FROM EMPLOYEE
WHERE FNAME = 'John'
AND MINIT = 'B'
AND LNAME = 'Smith';

BDATE ADDRESS

1965-01-09 731 Fondren, Houston, TX



Q3: List all information about the employees of department 5.

SELECT FNAME, MINIT, LNAME, SSN, BDATE, ADDRESS, SEX, SALARY, SUPERSSN, DNO
FROM EMPLOYEE
WHERE DNO = 5;





Q4: List name and address for all employees that work at the research department.

SELECT FNAME, LNAME, ADDRESS FROM EMPLOYEE, DEPARTMENT WHERE DNAME = 'Research' AND DNUMBER = DNO;



SELECT * FROM EMPLOYEE, DEPARTMENT WHERE DNAME = 'Research' AND DNUMBER = DNO;

FNAME	DN	O DNAME	DNUME	BER	MGRSTARTDATE
John	. 5	Research	5	• • •	1988-05-22
Franklin	. 5	Research	5	• • •	1988-05-22
Ramesh	. 5	Research	5	• • •	1988-05-22
Joyce	5	Research	5	• • •	1988-05-22



Q5: List project number, department number and the name and address of the director of the department for all projects that are located in Stafford.

SELECT PNUMBER, DNUM, LNAME, ADDRESS

FROM PROJECT, DEPARTMENT, EMPLOYEE

WHERE PLOCATION = `Stafford'
AND DNUMBER = DNUM
AND SSN = MGRSSN;

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SELECT PROJECT. PNUMBER, PROJECT.DNUM, EMPLOYEE.LNAME, EMPLOYEE.ADDRESS

FROM PROJECT, DEPARTMENT, EMPLOYEE

WHERE PROJECT.PLOCATION = `Stafford'

AND DEPARTMENT.DNUMBER = PROJECT.DNUM

AND EMPLOYEE.SSN = DEPARTMENT.MGRSSN;



Q6: List first and last name for all employees together with first and last names of their bosses.

SELECT E.FNAME, E.LNAME, S.FNAME, S.LNAME FROM EMPLOYEE E, EMPLOYEE S WHERE E.SUPERSSN = S.SSN;



- SQL considers a table as a multi-set (bag), i.e. tuples can occur more than once in a table.
- Why?
- Removing duplicates is expensive.
- User may want information about duplicates.
- Aggregation operators.

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Q7: List all salaries.SELECT SALARYFROM EMPLOYEE;

SELECT ALL SALARY **FROM** EMPLOYEE;



Q8: List all salaries without duplicates.
 SELECT DISTINCT SALARY
 FROM EMPLOYEE;



Q9: List all project numbers for projects in which an employee with name Smith works or where the leader of the department to which the project belongs is called Smith.



(SELECT DISTINCT PNUMBER FROM PROJECT, DEPARTMENT, **EMPLOYEE** WHERE DNUM = DNUMBER AND MGRSSN = SSN AND LNAME = 'Smith') UNION (SELECT DISTINCT PNUMBER FROM PROJECT, WORKS-ON, EMPLOYEE WHERE PNO = PNUMBER AND ESSN = SSN AND LNAME = 'Smith');



Q10: List all employees that live in Houston.

SELECT FNAME, LNAME
FROM EMPLOYEE
WHERE ADDRESS LIKE '%Houston%';

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Q11: List names for all employees that are born in the 1950's.

SELECT FNAME, LNAME FROM EMPLOYEE WHERE BDATE LIKE '__ 5%'; Q12: List names and salaries for all employees that work with ProductX in case they would receive a raise of 10%.

SELECT FNAME, LNAME, 1.1 * SALARY
FROM EMPLOYEE, WORKS-ON,
PROJECT
WHERE SSN = ESSN
AND PNO = PNUMBER
AND PNAME = `PRODUCTX';

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Q13: List all employees in department 5 with a salary between 30,000\$ and 40,000\$.

SELECT *
FROM EMPLOYEE
WHERE DNO = 5 AND
(SALARY BETWEEN 30000 AND 40000);

- Q14: List all employees and the projects they work with sorted with respect to department and within the department sorted alphabetically with respect to last name and then first name.
- **SELECT** DNAME, LNAME, FNAME, PNAME
- **FROM** DEPARTMENT, EMPLOYEE, PROJECT, WORKS-ON
- WHERE PNO = PNUMBER AND SSN = ESSN AND DNO = DNUMBER
- **ORDER BY** DNAME, LNAME, FNAME;



- **SELECT** DNAME, LNAME, FNAME, PNAME
- **FROM** DEPARTMENT, EMPLOYEE, PROJECT, WORKS-ON
- WHERE PNO = PNUMBER AND SSN = ESSN AND DNO = DNUMBER
- ORDER BY DNAME DESC, LNAME ASC, FNAME ASC;



Q15: List SSN for all employees that work with the same project at the same times as the person with SSN '123456789' (John Smith).

SELECT ESSN FROM WORKS-ON WHERE (PNO, HOURS) IN (SELECT PNO, HOURS FROM WORKS-ON WHERE ESSN = '123456789');



SELECT E.ESSN
FROM WORKS-ON E, WORKS-ON JS
WHERE JS.ESSN = '123456789'
AND E.PNO = JS.PNO
AND E.HOURS = JS.HOURS;



Q16: List all employees whose salary is higher than the salaries of the employees who work at department 5.

SELECT LNAME, FNAME
FROM EMPLOYEE
WHERE SALARY > ALL
(SELECT SALARY
FROM EMPLOYEE
WHERE DNO = 5);



Q17: List all employees whose salary is higher than the salary of some employee who works at department 5.

SELECT LNAME, FNAME
FROM EMPLOYEE
WHERE SALARY > SOME
(SELECT SALARY
FROM EMPLOYEE
WHERE DNO = 5);



Q18: List all employees that do not have a relative at the company.

SELECT LNAME, FNAME
FROM EMPLOYEE
WHERE NOT EXISTS
(SELECT *
FROM DEPENDENT
WHERE SSN = ESSN);



 Q19: List all department managers that have at least one relative at the company.

SELECT LNAME, FNAME FROM EMPLOYEE WHERE EXISTS (SELECT * FROM DEPARTMENT WHERE SSN = MGRSSN) **AND EXISTS** (SELECT * FROM DEPENDENT WHERE SSN = ESSN);

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Q20: List SSN for all employees that work with project 1, 2 or 3.

SELECT DISTINCT ESSN FROM WORKS-ON WHERE PNO IN (1, 2, 3); м

Q21: List all employees that do not have a boss.

SELECT FNAME, LNAME FROM EMPLOYEE WHERE SUPERSSN IS NULL; Q22: List the sum, the highest, lowest and average of the salaries of the employees of the research department.

SELECT SUM(SALARY),
MAX(SALARY), MIN(SALARY),
AVG(SALARY)

FROM EMPLOYEE, DEPARTMENT
WHERE DNAME = 'Research'
AND DNO = DNUMBER;

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Q23: List the number of employees.

SELECT COUNT(*) **FROM** EMPLOYEE;



Q24: List for each department the department number, the number of employees and the average salary.

SELECT DNO, COUNT(*), AVG(SALARY)

FROM EMPLOYEE GROUP BY DNO;

DNO COUNT(*) AVG_SALARY		
5	4	33250
4	3	31000
1	1	55000

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Q25: List for each project the project number, project name and the number of employees that work with the project.

SELECT PNUMBER, PNAME, COUNT(*)
FROM PROJECT, WORKS-ON
WHERE PNUMBER = PNO
GROUP BY PNUMBER, PNAME;

Q26: List for each project with at least 2 employees the project number, project name and number of employees that work with the project.

SELECT PNUMBER, PNAME, COUNT(*)
FROM PROJECT, WORKS-ON
WHERE PNUMBER = PNO
GROUP BY PNUMBER, PNAME
HAVING COUNT(*) > 1;



Creating new tables

CREATE TABLE DEPTS-INFO (DEPT-NUMBER integer primary key, DEPT-NAME varchar(15), NO-OF-EMPLOYEES integer, TOTAL-SAL integer, foreign key (DEPT-NUMBER) references DEPARTMENT(DNUMBER));



INSERT INTO DEPTS-INFO (DEPT-NUMBER, DEPT-NAME, NO-OF-EMPLOYEES, TOTAL-SAL) **SELECT** DNUMBER, DNAME, COUNT (*), SUM(SALARY) FROM DEPARTMENT, EMPLOYEE WHERE DNUMBER = DNO **GROUP BY** DNUMBER, DNAME;