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Crypto p3

Q1) $q = 71$ $d = 7$

a) $X_A = 5$ $Y_A = 7^5 \bmod 71 = 51$

b) $X_B = 12$ $Y_B = 7^{12} \bmod 71 = 4$

c) Shared Secret:

$S = (Y_B)^{X_A} \bmod 71 = 30$ $S = (Y_A)^{X_B} \bmod 71 = 30$

d) Diffie Helman is based off this property:

$$(x^a \bmod p)^b \bmod p = x^{ab} \bmod p$$

Therefore if you sent (x^x) you are invalidating the system. Also if it did work, you make it unsafe because you remove the discrete log.

Q2)

a) The idea here is you generate a bunch of valid and invalid packets and keep going till you have a valid packet that collides with an invalid one. Now you can get a signature for both, it works because when you try to collide with a hash it's super rare, but the chance that two hashes in the group will be the same is much more likely. The principle of the Birthday Problem.

b) For 1% chance you make:

e.g. $2^{hs/2} + 2^{hs/2}$ $hs = \text{hash size} = 64 \text{ bits}$

then m size: $m \cdot 2^{32} + m \cdot 2^{32}$

c) So they need to find $2^{32} + 2^{32}$ hashes:

it's $\frac{2^{32} + 2^{32}}{2^{20}} = 8192 \text{ Sec}$

d) when 128 it's
 $m \cdot 2^{64} + m \cdot 2^{64}$ Space
 and

$$\frac{2^{64} + 2^{64}}{2^{20}} = \boxed{3518437208832 \text{ sec}}$$

A lot longer in 2^{45}

Q3) generate key: "0d 010111"

$$B_i = a \cdot S_i \bmod p$$

$$B_1 = 1019 \cdot 5 \bmod 1999 = 1097$$

$$B_2 = 1019 \cdot 9 \bmod 1999 = 1175$$

$$B_3 = 1019 \cdot 21 \bmod 1999 = 1404$$

$$B_4 = 1019 \cdot 45 \bmod 1999 = 1877$$

$$B_5 = 1019 \cdot 103 \bmod 1999 = 1009$$

$$B_6 = 1019 \cdot 215 \bmod 1999 = 1194$$

$$B_7 = 1019 \cdot 450 \bmod 1999 = 779$$

$$B_8 = 1019 \cdot 946 \bmod 1999 = 456$$

encrypt:

$$c = 0 + 1(1175) + 0 + 1(1877) + 0 + 1(1194) + 1(779) + 1(456) = \boxed{5481}$$

decrypt: $a^{-1} = a^{m-2} \bmod p$ (rule)
 $a^{-1} = 1589$

$$1589 \cdot 5481 \bmod 1999 = \boxed{1665}$$

$$1665 - 946 = 719$$

$$719 - 450 = 269$$

$$269 - 215 = 54$$

$$54 - 45 = 9$$

$$9 - 9 = 0$$

Number is $\boxed{01010111}$

It worked ✓