1 Reduction and variable capture

1. Answer:

Non capture avoiding

$$(\lambda x.(\lambda y.x)) \ y$$
$$=_{\beta} \lambda y.y$$

Capture avoiding

$$(\lambda x.(\lambda y.x)) y$$

$$=_{\alpha} (\lambda x.(\lambda z.x)) y$$

$$=_{\beta} \lambda z.y$$

2. Answer:

Non capture avoiding

$$(\lambda x.(\lambda y.x)) (\lambda y.x)$$

= $_{\beta} \lambda y.(\lambda y.x)$

Capture avoiding

$$\begin{split} &(\lambda x.(\lambda y.x))\ (\lambda y.x) \\ &=_{\alpha}\ (\lambda x.(\lambda z.x))\ (\lambda y.x) \\ &=_{\beta}\ \lambda zy.x \end{split}$$

3. Answer:

Non capture avoiding

$$\begin{aligned} &(\lambda x.(\lambda y.x))\ (\lambda y.y) \\ &=_{\beta} \lambda y.(\lambda y.y) \end{aligned}$$

Capture avoiding

$$\begin{split} &(\lambda x.(\lambda y.x)) \ (\lambda y.y) \\ &=_{\alpha} \ (\lambda x.(\lambda z.x)) \ (\lambda y.y) \\ &=_{\beta} \ \lambda zy.y \end{split}$$

4. Answer:

Non capture avoiding

$$(\lambda xyz. \ \lambda fgh. \ f \ x \ (g \ y) \ (h \ z)) \ h \ (\lambda ab. \ a \ (g \ b)) \ f$$

$$=_{\beta} (\lambda yz. \ \lambda fgh. \ f \ h \ (g \ y) \ (h \ z)) \ (\lambda ab. \ a \ (g \ b)) \ f$$

$$=_{\beta} (\lambda z. \ \lambda fgh. \ f \ h \ (g \ (\lambda ab. \ a \ (g \ b))) \ (h \ z)) \ f$$

$$=_{\beta} \lambda fgh. \ f \ h \ (g \ (\lambda ab. \ a \ (g \ b))) \ (h \ f)$$

Capture avoiding

$$(\lambda xyz. \ \lambda fgh. \ f \ x \ (g \ y) \ (h \ z)) \ h \ (\lambda ab. \ a \ (g \ b)) \ f \\ =_{\alpha} (\lambda xyz. \ \lambda qgr. \ (q \ x) \ (g \ y) \ (r \ z)) \ h \ (\lambda ab. \ a \ (g \ b)) \ f \\ =_{\beta} (\lambda yz. \ \lambda qgr. \ (q \ h) \ (g \ y) \ (r \ z)) \ (\lambda ab. \ a \ (g \ b)) \ (r \ z)) \ f \\ =_{\beta} (\lambda z. \ \lambda qgr. \ (q \ h) \ (g \ (\lambda ab. \ a \ (g \ b))) \ (r \ z)) \ f \\ =_{\beta} \lambda qgr. \ (q \ h) \ (g \ (\lambda ab. \ a \ (g \ b))) \ (r \ f)$$

5. Answer:

$$(\lambda xy.z) z z$$

$$= (\lambda xy.z) (z z)$$

$$= \lambda y.z$$

6. Answer:

$$(\lambda x.(\lambda y.(x\ y))) =_{\eta} \lambda x.x$$

7. Answer:

$$S = \lambda xyz.((x z) (y z))$$

$$K = \lambda xy.x$$

$$I = \lambda x.x$$

S K S = (S K) S by expression left association. The following are done with call-by-name.

$$\begin{split} S \ K &= (\lambda xyz.((x \ z) \ (y \ z)))(\lambda xy.x) \\ &=_{\beta} (\lambda yz.(((\lambda xy.x) \ z) \ (y \ z))) \\ &=_{\beta} (\lambda yz.((\lambda y.z) \ (y \ z))) \\ &=_{\beta} \lambda yz.z \\ (S \ K) \ S &= (\lambda yz.z) \ (\lambda xyz.((x \ z) \ (y \ z))) \\ &= \beta(\lambda z.z) = \alpha(\lambda x.x) = I \end{split}$$

S K S hence shown reduces to I.

2 Variable Binding and Closure

- 1. Answer: The function is recursive and never terminates. This is because the x(2) call on line 2 refers to the definition of x on line 2 which makes it recursive.
- 2. Answer: x passed in as argument on line two binds to the previous definition on line 1 because at the time of reference, x on line 2 is not assigned yet since the function is not executed at that point.
- 3. Answer: You get the following error: Error reassigning x. This is because Haskel use static types. Variables are not allowed to change values.

3 Closure and access links

- 1. See scaned document
- 2. 18
- 3. This does not terminate because the function is infinitely recursive without a base case.

4 Memory management and high-order functions

- 1. See scaned document
- 2. The value is 10 because. Within function scope g, x is set to 7. This is passed in as argument to h which is f. Therefore the y argument in f resolves to 7. x is 5 by access link to the parent scope. The function f therefore returns 5+7-2 which is 10.

5 More substitution and variable capture

1. Answer:

$$((\lambda x.((\lambda x.x) \ 2) + x) \ x)[x := 3]$$

$$=_{\beta} (((\lambda x.x) \ 2) + x)[x := 3]$$

$$=_{\beta} (2 + x)[x := 3]$$

$$= 2 + 3$$

$$= 5$$

Explaination: In this case it would not matter because the only x that is not captured is one that is ok to replace. All other x are captured by lambda expressions. 3 can not possibly be captured by any lambda.

2. Answer:

$$\begin{split} &(\lambda y.(\lambda xyz.z)\ y\ z\ y)[z:=w]\\ &=(\lambda y.((((\lambda xyz.z)\ y)\ z)))[z:=w]\\ &=_{\beta}\ (\lambda y.(((\lambda yz.z)\ y)\ z))[z:=w]\\ &=_{\beta}\ (\lambda y.((\lambda z.z)\ y))[z:=w]\\ &=_{\beta}\ (\lambda y.y)[z:=w]\\ &=\lambda y.y \end{split}$$

Explaination: In this case it would not matter whether capture avoiding is taken into account. w is not caputured anywhere, therefore substitution is legal in this case without needing to avoid capturing w.

3. Answer:

$$\begin{split} &(\lambda p.(\lambda x.p(x\ x))(\lambda x.p))[x:=p]\\ &=_{\alpha}\ (\lambda z.(\lambda x.(z(x\ x)))(\lambda x.z))[x:=p]\\ &=_{\beta}\ (\lambda x.((\lambda x.z)(x\ x)))[x:=p]\\ &=_{\beta}\ (\lambda x.z)[x:=p]\\ &=_{\beta}\ (\lambda x.z) \end{split}$$

Explaination: Yes, capture avoiding is a must because p is captured by the lambda expression. This violates substitution rules therefore capture avoiding (alpha renaming) must be performed.