

# Operating Systems: Homework 1

Jack Shi - A92122910

April 12, 2018

## 1. Missing mechanism:

**Without privileged instructions** Since manipulating protected control registers is a privileged instruction, without it, the user would be able to bypass the dual-mode operation and make their instructions kernel operations as they please. This can result in security issues especially if the user decide to manage I/O directly which is usually a system call through the kernel.

**Without memory protection** Programs would be able to access eachother's memory address space. This could allow malicious software to steal data or manipulate other programs by illegally access, change their memory data.

**Without timer interrupts** This would result in the inability to timeout process that take too long or stuck in an infinite loop.

## 2. System without privileged mode:

**Possible to be secure:** Programs could be forced to run in VMs such as the JVM where the enviornment is controlled and the program's access is restricted.

**Not possible:** In case that the program must run on bare-metal for high efficiency or performance reasons, the lack of a privileged mode could result in kernel corruption.

## 3. Which instructions should be privileged?

a) Set value of timer

**privileged:** Changing the timer could change the behavior of the interrupts which is a kernel level operation.

b) Read the clock

**privileged:** System time can be used for timing attacks. The system clock is also managed by the kernel.

c) Clear memory

**privileged:** Because of memory protection, processes should not be able to clear memory anywhere. However, within it's own allocated address space, it fine.

d) Turn off interrupts

**privileged:** If any process is allowed to turn off interrupts, then time-out timer interrupt would not be able to prevent malicious infinite loops.

e) Switch from user to monitor(kernel) mode

**privileged:** There is no point is separating modes if the user can just switch to kernel mode as they please.

## 4. Condition of failure:

**open** `pathname` refer to block device, and that block device is being used by the system. **open** fails with error `EBUSY`.

**read** `file descriptor` does not refer to a opened file. **read** fails with error `EBADF`.

**fork** **fork** fails with error `ENOMEM` when there is not sufficient memory for another process.

**exec** If the header of the file for **exec** is not recognized, **exec** will result in error `ENOEXEC`.

**unlink** If the `pathname` points to somewhere outside of the allotted address sapce, **unlink** result in error `EFAULT`.

5. Challenges when copying parameters from user to kernel modes:

**Challenge 1:** The program in user mode does not have the privilege to build a function stack in kernel mode and copy over the parameters.

**Challenge 2:** If the parameters contains pointerrs, dereferencing will require violating the user/kernel mode boundaries as well.

**Solution:** Temporary grant the caller kernel privileges.

6. Use hardware interval timer to keep track of the time of day:

**Solution:** Since time of day is usually descrtetized to seconds, the OS can set 1 second long interval timers that update the system time on 1Hz interrupt.

7. Divise subsitute for traps using interrupts and/or exceptions:

**Solution:** It is possible if the user get kernel privilege and implement it's own version of the interrupt handler.

8. Consider C program:

- a. 8 forks total are created.
- b. The statement is executed twice.