CSCI 3104 Assignment 2

10:00 - 10:50 Wanshan

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1. (a)
                   \max = a[0]
                   \min = a[0]
                   for (i in a) {
                     if (a[i] > max) {
                       \max = a[i]
                     if (a[i] < min) {
                       \min = a[i]
                   return max, min
   (b)
                   function findExtremas(array a, int start, int end) {
                     // Base case
                     if (start - end < 2) {
                       return a[start], a[start] // max, min
                     int mid = (start+end)/2
                     int u = findExtremas(a, start, m)
                     int v = findExtremas(a, m+1, end)
                     if (u > v) {
                       return u, v //max, min
                     else {
                       return v, u //max, min
                     }
                   }
   (c)
                                  T(n) = 2T(\frac{n}{2}) + O(1)
```

(d) Since $f(n) = O(n^{\log_2 2 - 1})$ $T(n) = \Theta(n^1) = \boxed{\Theta(n)} \text{ by the Master method}$

2. (a) There is no majority element since the majority elements in the two subarrays are different. Therefore it is impossible for there to be another element with $\frac{n}{2} + 1$ frequency.

```
function majorEl(array a) {
(b)
                   // Base case(s)
                   if (a.length == 1) {
                     return a[0]
                   }
                   mid = a.length/2
                   11 as array = a[0] ... a[mid]
                   12 as array = a[mid+1] \dots a[n-1]
                   e1 = majorEl(l1)
                   e2 = majorEl(12)
                   if (e1 = e2) {
                     return e1
                   else {
                     if (getFrequency(a, e1) > mid + 1)
                        return e1
                     else if (getFrequency(a, e2) > mid + 1)
                        return e2
                     else
                        return NULL
                   }
                 }
(c)
                                T(n) = 2T(\frac{n}{2}) + O(n)
                                Since O(n) = O(n^{\log_2 2})
```

 $T(n) = \Theta(n \log n)$ by the Master method