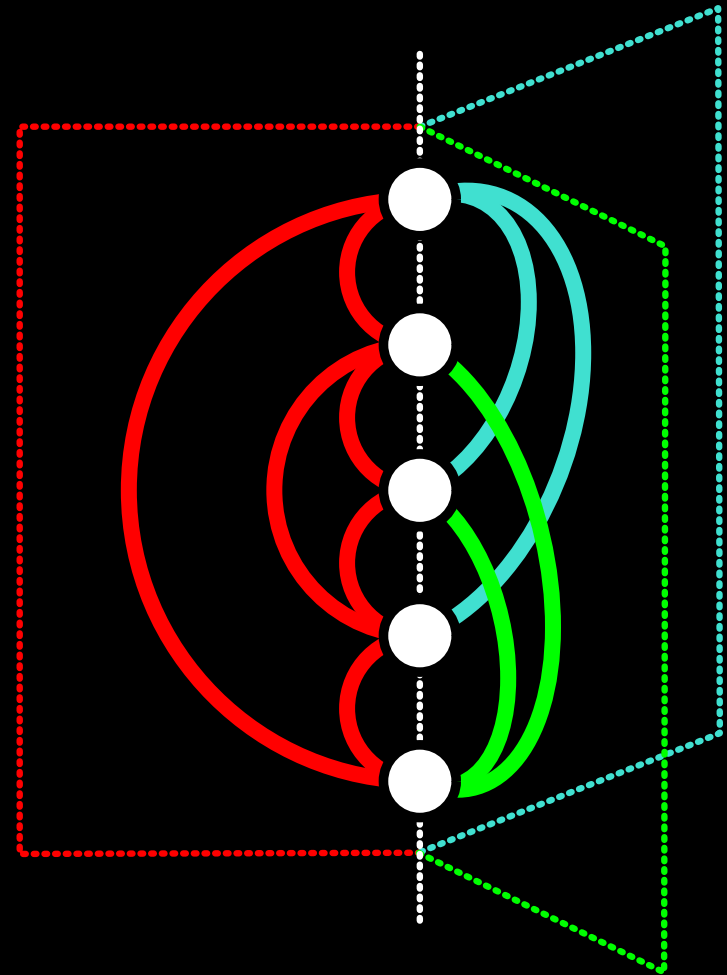
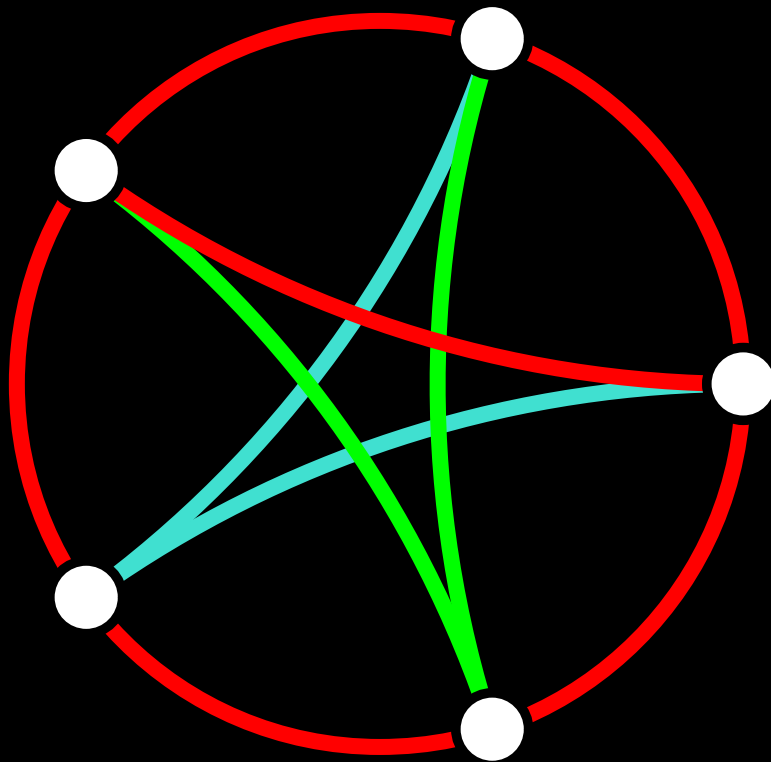


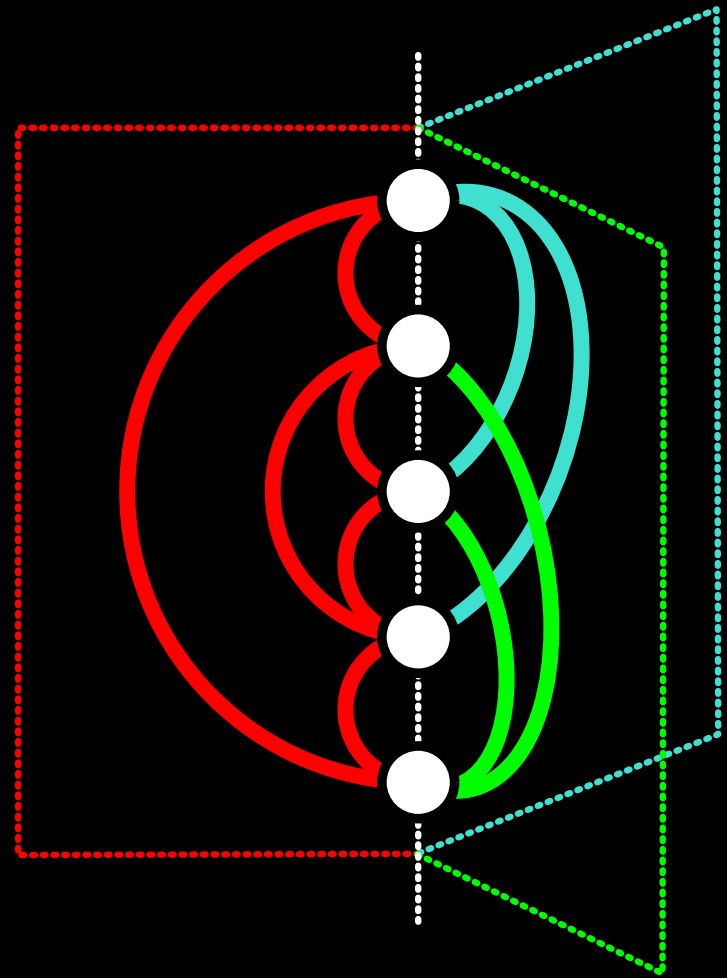
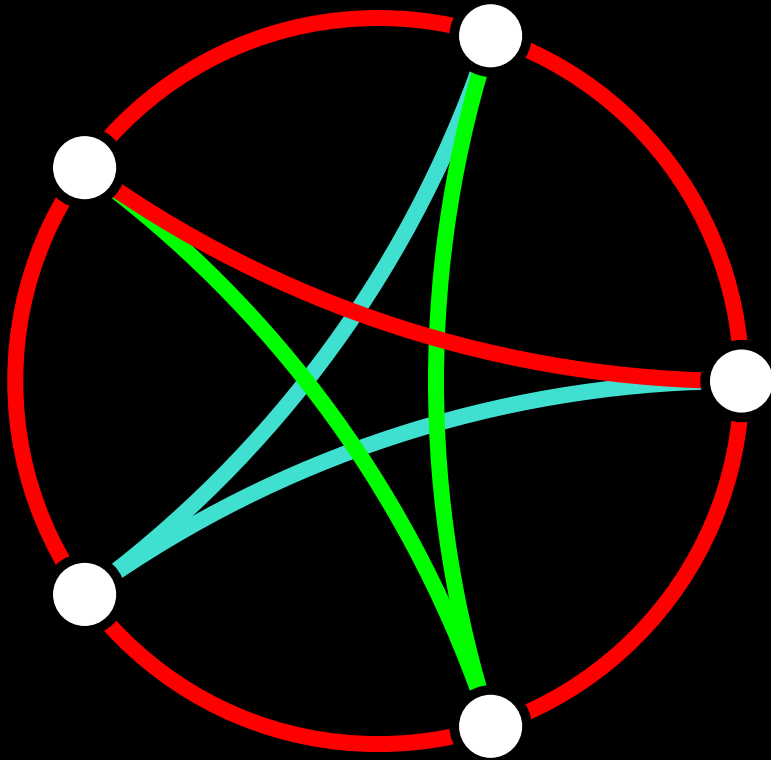
Fixed-Parameter Tractable Algorithms for Crossing Minimization in Book Embeddings



Presented by: Michael J Bannister

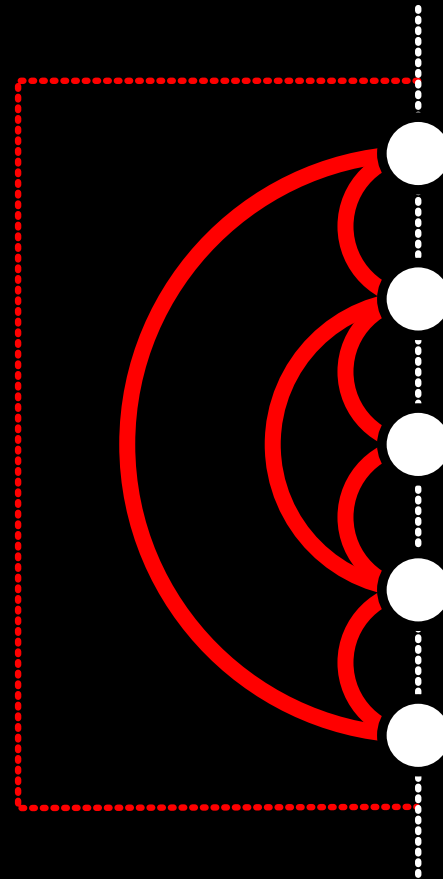
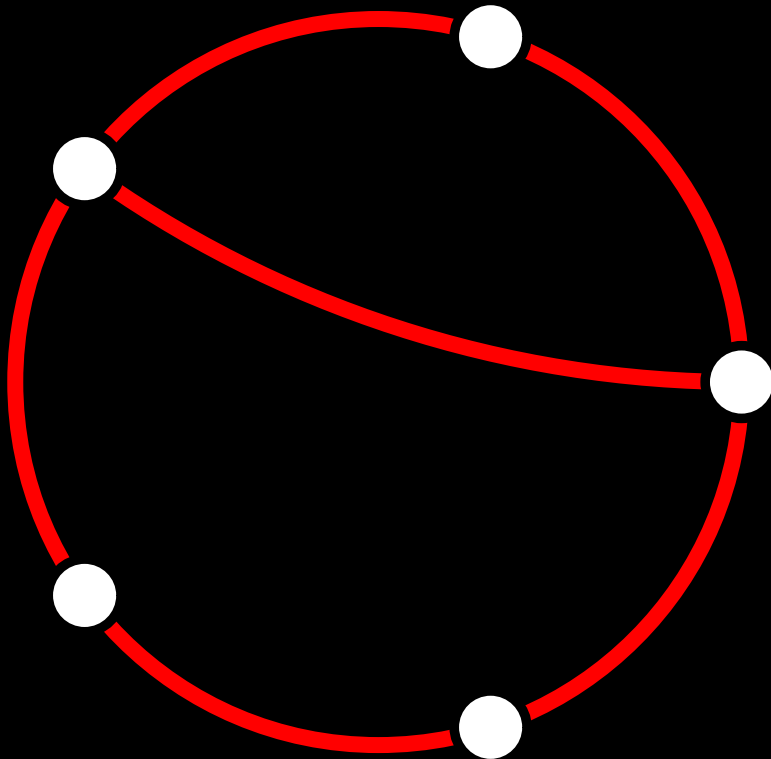
Joint work with: David Eppstein and Joseph A. Simons

Book Embedding



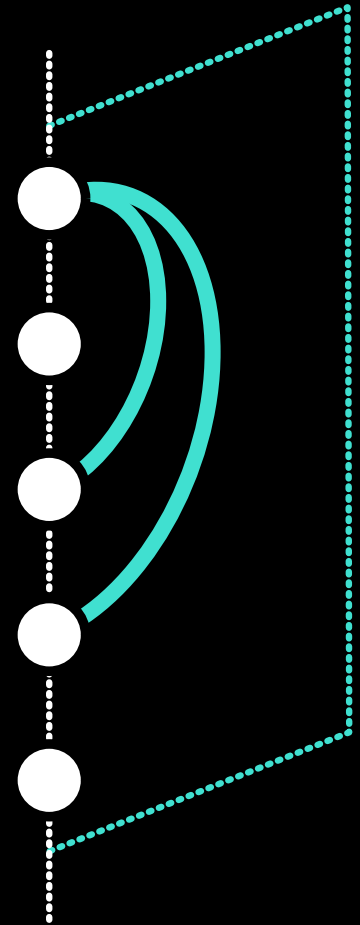
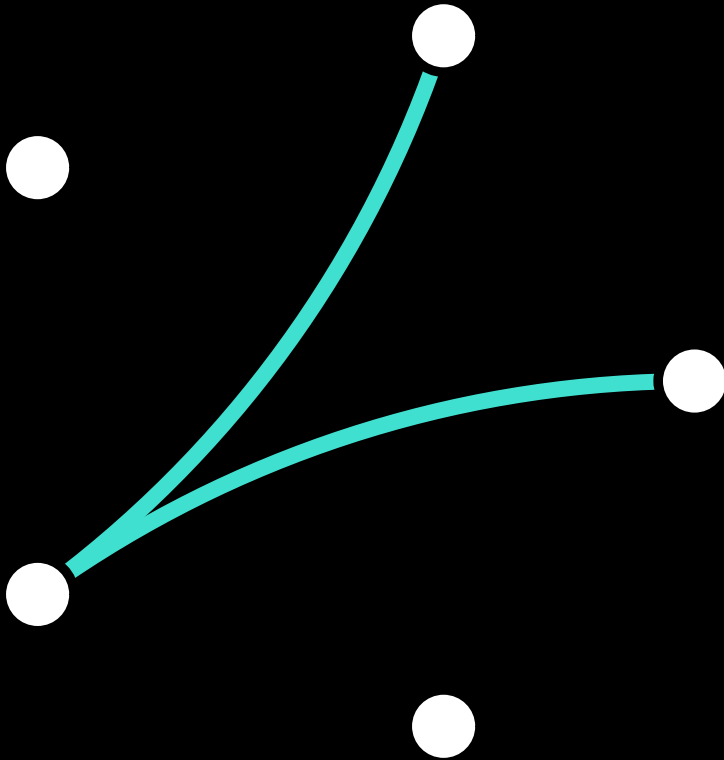
An embedding of K_5 in 3-pages

Book Embedding



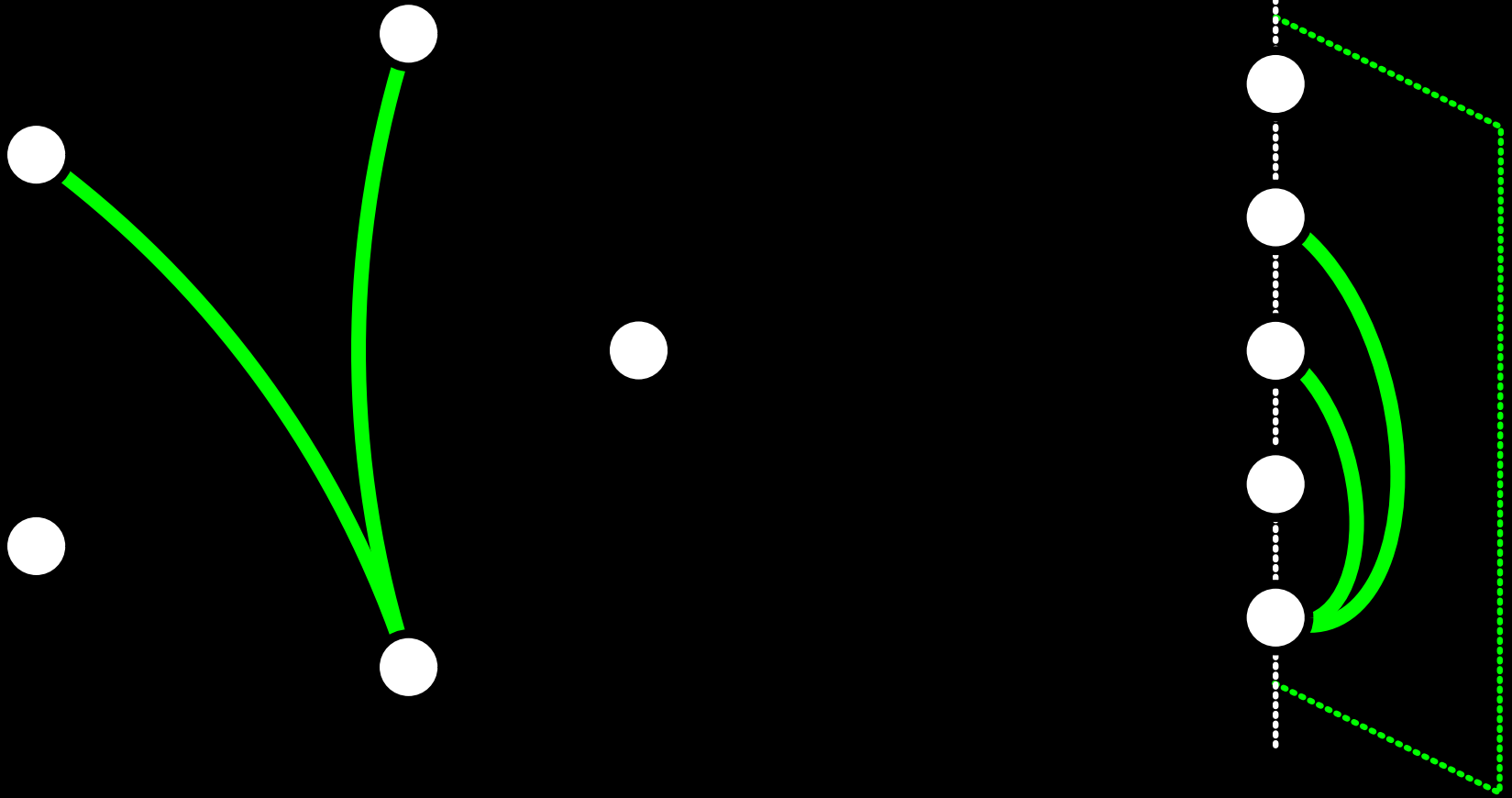
An embedding of K_5 in 3-pages

Book Embedding



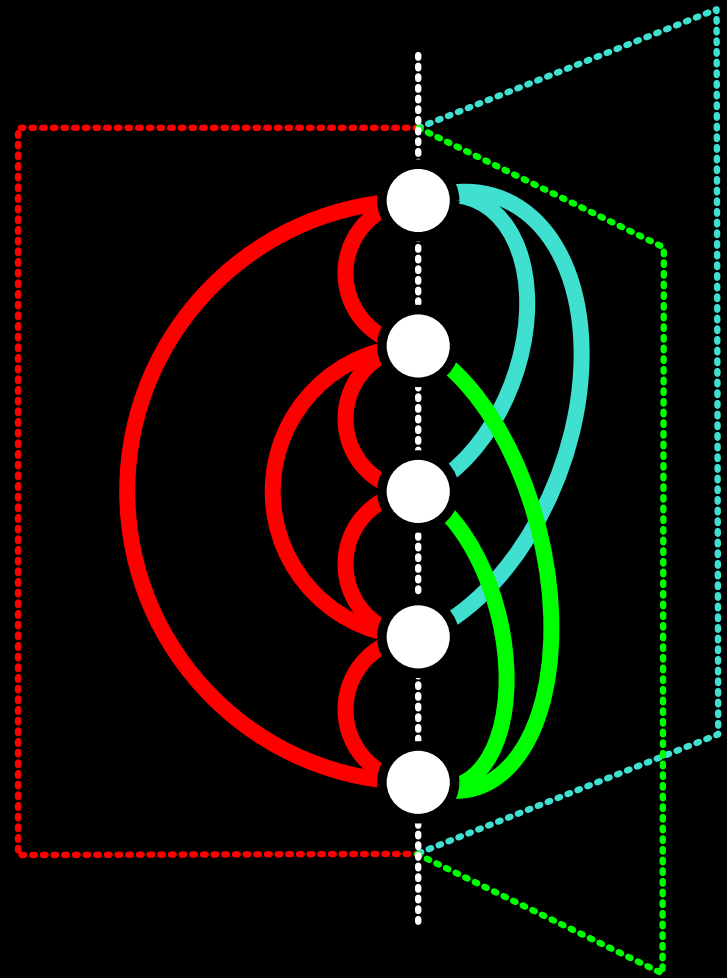
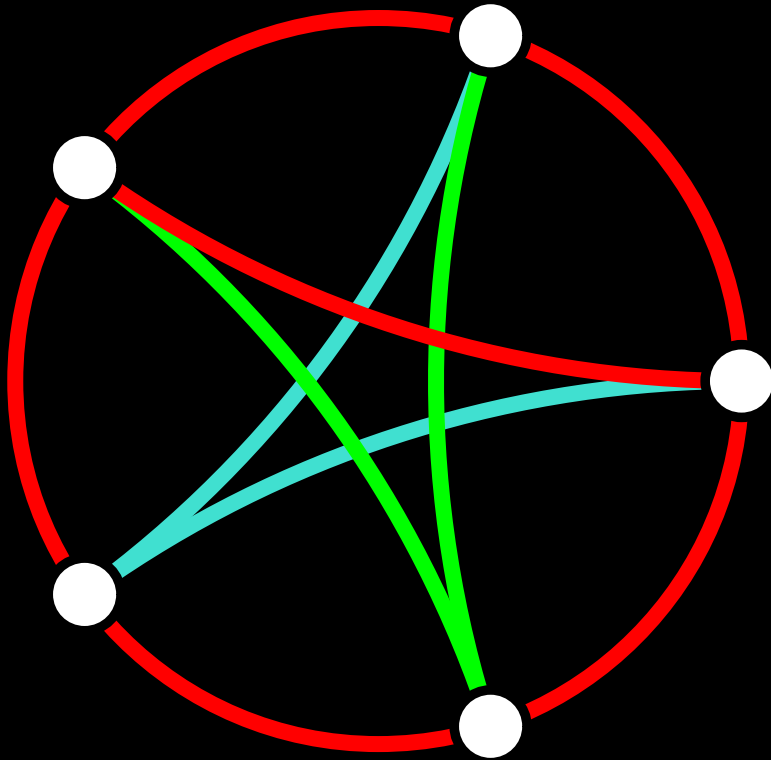
An embedding of K_5 in 3-pages

Book Embedding



An embedding of K_5 in 3-pages

Book Embedding



An embedding of K_5 in 3-pages

Book Embedding

Book crossing number of a drawing D
= sum of the # crossings in each page

k -Page crossing number of a graph G
= min crossing # of a k -page drawing of G



1-Page crossing minimization

Known complexity results:

Planar: $O(n)$

Minimization: NP-hard

Exact minimization: $O(n!) = 2^{O(n \log n)}$

Approximation ratio: $O(\log^2 n)$

Heuristics work well in practice



2-Page crossing minimization

Known complexity results:

Planar: NP-complete

Minimization: NP-hard

Exact minimization: $O(n!2^m) = 2^{O(n \log n + m)}$

Approximation ratio: $O(\log^2 n)$

Heuristics do not work well in practice



Parameterized Complexity

NP-Hard \Rightarrow likely (at least) exponential

Is the problem easier for:

- graphs with few crossings?
- graphs that are “tree like”?
- social networks?

Fixed-parameter tractable if

- given a parameter p of the input
- \exists a computable function f
- \exists an exact algorithm running in $O(f(p)n^c)$ with $c = O(1)$



WE FEEL THE OFFICE SCAPEGOAT IS A
KEY COMPONENT OF TEAM-BUILDING,
AND YOU'RE A GREAT FIT FOR THE JOB.

Parameterized Complexity

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given a parameter p of the input
 \exists a computable function f
 \exists an exact algorithm running in
 $O(f(p)n^c)$ with $c = O(1)$



Blame the difficulty of the problem on p not n !

Parameters

k -Almost tree parameter

Distance from a tree by edge count

Max of $m - n + 1$ over all biconnected components

Treewidth parameter

Common FPT parameter

Measure of large scale tree structure

Natural parameter

Number of crossings

Not possible for 2-page

Methods

Kernelization

Step 1: Reduce the size of the input

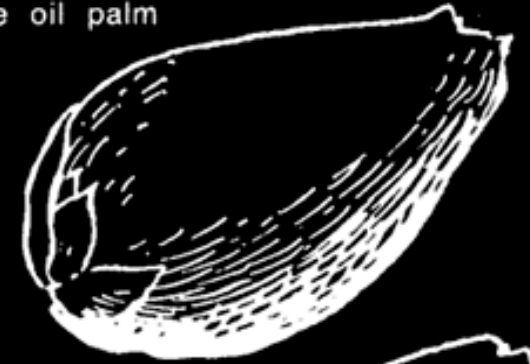
Step 2: Solve the smaller problem
with exact algorithms

Logical expressibility

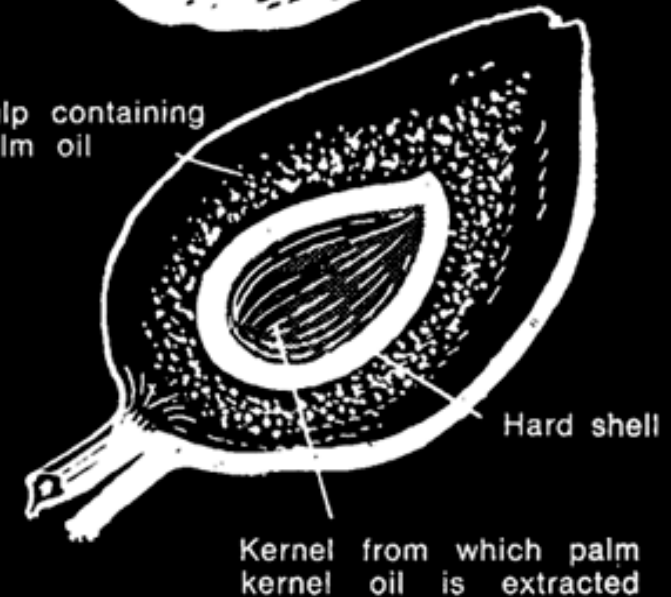
Step 1: Describe the problem in MSO ($\subseteq 2^{\text{nd}}$ order logic)

Step 2: Use powerful meta-theorems to produce algorithms

Fruit of the oil palm



Pulp containing
palm oil



Hard shell

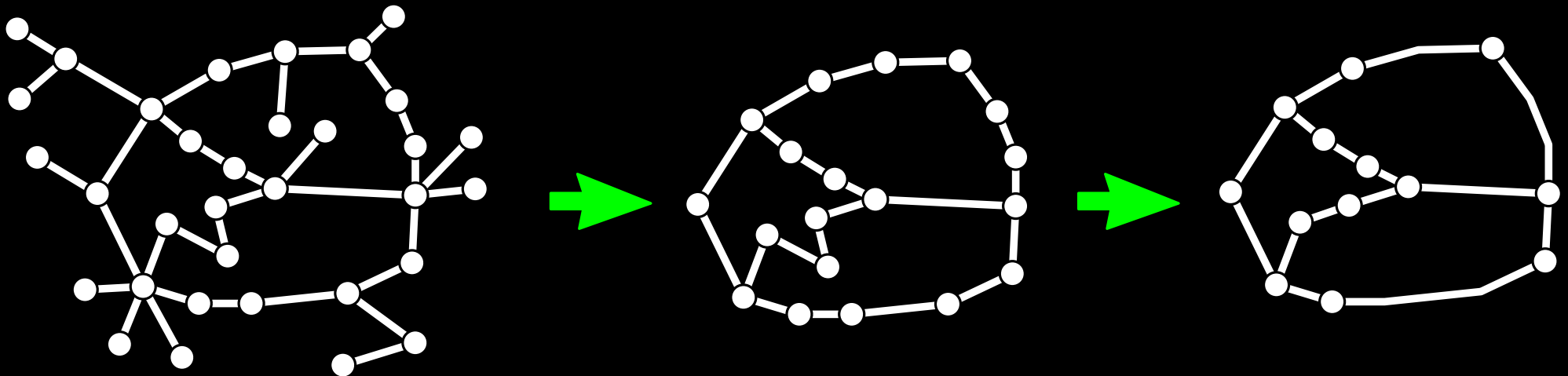
Kernel from which palm
kernel oil is extracted

k-Almost tree kernel

Step 1: Iteratively remove degree 1 vertices

Step 2: Reduce paths of degree 2 vertices to $\leq \ell(k)$

To be determined later ↗



Kernelization Results

Crossing minimization:

1-Page $O((5k)!n)$

2-Page $O(2^{6k^3} (6k^3)!n)$

Crossed edge minimization:

1-Page $O((5k)!n)$

2-Page $O(2^{6k^2} (6k^2)!n)$

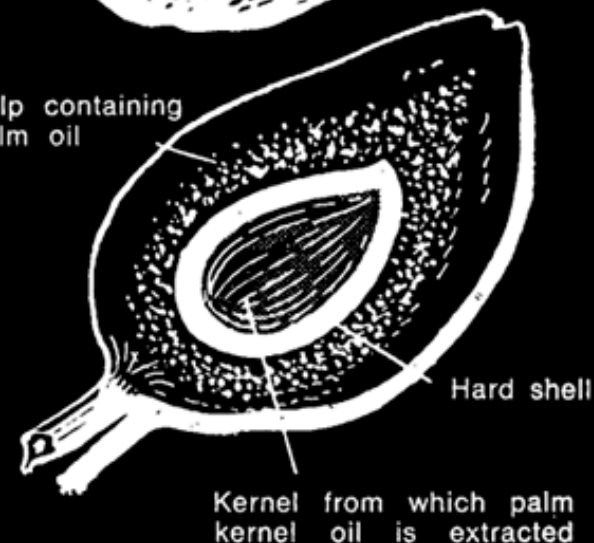
Additional Result:

1-Planar recognition is FPT

Fruit of the oil palm



Pulp containing palm oil



MSO₂ and Courcelle's theorem

Monadic second-order logic:

Vertex and edge variables: $v_0, v_1, \dots, e_0, e_1, \dots$

Vertex and edge set variables: $V_0, V_1, \dots, E_0, E_1, \dots$

Binary relations: $=, \in, \text{I}$

Propositional logic operations: $\neg, \wedge, \vee, \rightarrow$

Quantifiers: \forall, \exists

Examples of properties expressible in MSO₂

k -coloring

connectedness

hamiltonicity

minor containment

planarity

outerplanarity

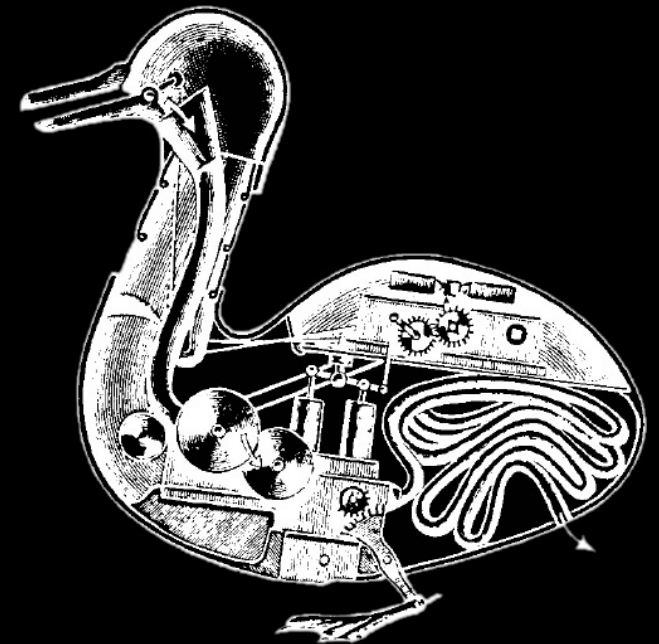
MSO_2 and Courcelle's theorem

Input: A graph G and an MSO_2 -formula ϕ

Parameter: $\text{treewidth}(G) + \text{length}(\phi)$

Output: Does G satisfy ϕ

Runtime: $f(k, \ell)n$



Application:

Crossing minimization is FPT in the $\#$ crossings

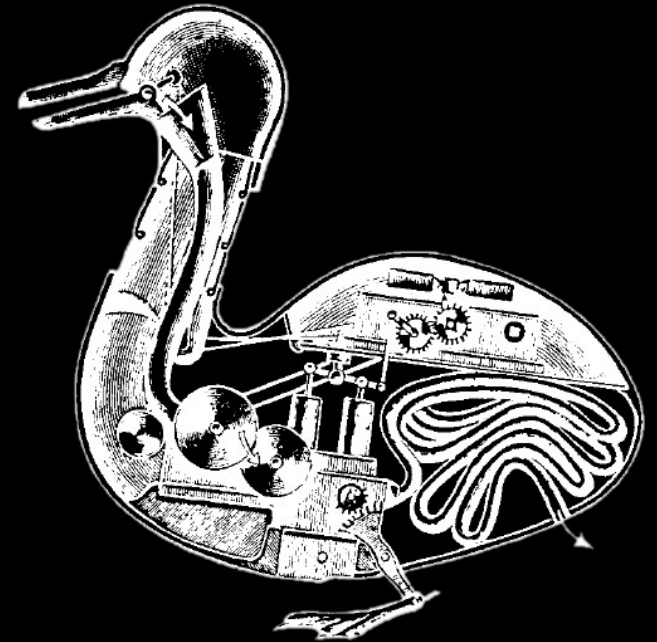
Grohe (2001), Kawarabayashi & Reed (2007)

Logical expressibility results

1-Page crossing minimization is FPT
parameter: crossing number

2-Page planarity is FPT
parameter: treewidth

2-Page crossing minimization is FPT
parameter: treewidth + crossing number



Open Problems

- Extend results to k -pages with $k > 2$
- Faster exact algorithms
- Close the approximation gap
- Practical FPT algorithms

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Thank You !