

Lecture 18: Heaps & Heapsort

CS 62
Spring 2015
Kim Bruce & America Chambers

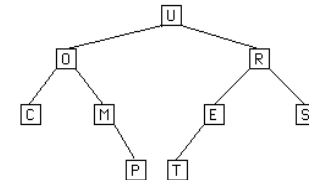
Assignments

- Lab: Compressed grid iterator
- Assignment: Darwin monster simulator
Due: 3/11 Right before spring break.
No mentor sessions on 3/12 or 3/13!

Array Representations of Trees

Array Representation

- `data[0..n-1]` can hold values in trees
- left subtree of node k in $2k+1$, right in $2k+2$,
- parent in $(k-1)/2$



Indices: 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14
data[]: U O R C M E S - - - P T - - -

Array Representation: Efficiency

- Tree of height h , takes $2^{h+1}-1$ slots, even if only has $O(h)$ elements
- Bad for long, skinny trees
- Good for full or complete trees.
- *Complete tree* is full except possibly bottom level and has all leaves at that level in leftmost positions.

PriorityQueue

```
public interface PriorityQueue<E extends Comparable<E>>
{
    /**
     * @pre !isEmpty()
     * @return The minimum value in the queue.
     */
    public E remove();
    public E getFirst();
    public void add(E value);
    public boolean isEmpty();
    public int size();
    public void clear();
}
```

Objects of type E can be compared

Implementations

Unsorted array: $O(1)$ add, $O(n)$ remove

Reverse sorted array: $O(n)$ add, $O(1)$ remove

Heap ordered array: $O(\log n)$ add, $O(\log n)$ remove

Min-Heap

A **min heap** is a complete full binary tree such that, the value in the root position is the smallest value and left and right subtrees of the root are both min heaps.

Typically a min heap is stored in an array using the array representation of trees.

See VectorHeap code

Heap Performance

- remove: $O(\log n)$
- getFirst: $O(1)$
- add: $(\log n)$
- isEmpty, size, clear: $O(1)$

Heapsort

1. Add all elements from vector into heap: $O(n \log n)$
2. Remove elements in order from heap: $O(n \log n)$
3. Total time: $O(n \log n)$

Can get about a factor of 2 speed up by building the heap “in place” in $O(n)$ time using heapify.