

ECE326

PROGRAMMING LANGUAGES

Lecture 19 : Variance and Data Types

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Enumerated Type

- A data type consisting of named values
- Each named value behaves as a constant
- A variable of an enumerated type can be assigned any of the named value

```
enum Suit {  
    CLUB,  
    DIAMOND,  
    HEART,  
    SPADE,  
};
```

```
Suit suit = SPADE;  
  
// can be compared  
if (suit > HEART) {  
    win();  
}
```

Enumerated Type

- In principle, their in-memory representation should not be exposed by the programming language
- In practice, C++ treats them like integers (mostly)

```
int a = CLUB;  
Suit b = HEART;
```

```
cout << a << " " << b << endl;    // prints 0 2
```

```
b = (Suit)5;           // this is OK in C++
```

```
Suit * sp = &b;  
int * ip = sp;         // this is not OK without a cast
```

C Enum

- Named values can be assigned an integer expression
- The type's size can be altered via base specifier
 - If you prefer your enum to take up less memory

```
enum Suit : char {  
    CLUB = 1,  
  
    // 2, always previous value + 1 if not specified  
    DIAMOND,  
    HEART = 5,  
  
    // this is very bad but is allowed: DIAMOND == SPADE  
    SPADE = CLUB + 1,  
};
```

Enum Class

- A more type-safe version of C Enum
- Disallows coercion to integer and other enums
- Named values are also “scoped”

```
enum class Suit {  
    CLUB,  
    DIAMOND,  
    HEART,  
    SPADE = 5,  
};  
  
// requires name resolution  
Suit suit = Suit::SPADE;
```

```
// no longer allowed  
int a = Suit::CLUB;
```

```
enum Foo { A = 1 };  
enum Bar { B = 1 };
```

```
// allowed  
bool b = A == B;
```

```
// not allowed  
b = A == Suit::DIAMOND;
```

Union

- Stores different data types in same memory location
- Type-unsafe construct, but can save memory
- Can be used in C to implement private members

```
// size of union is always the largest member
// in this case it's the double (8 bytes)
union Mangle {
    int i;
    unsigned char c;
    double d;
};

union Private {
    const int readonly;
    int __writeable;
};

Mangle m = { .i = 0xffff };
cout << (unsigned)m.c << endl; // prints 255
```

Union Cast

- Used to bypass strict aliasing rules
- C/C++ does not type check other union members
- Used to see memory representation of other types

```
union Pun {  
    double d;  
    unsigned char c[sizeof(double)];  
};
```

```
Pun p = { .d = 1.0 };  
for (unsigned i = 0; i < sizeof(double); i++)  
    cout << (unsigned)p.c[i] << " ";
```

0 0 0 0 0 0 240 63

Tagged Union

- A union that has a tag field to indicate which union member is being used
- Also known as variant type, or algebraic type

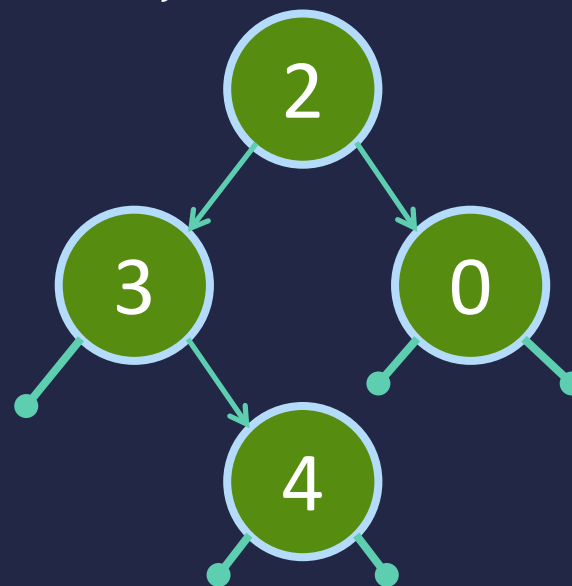
```
struct Tagged {  
    enum { INT, STR } tag;  
    union {                                // anonymous union (members  
        int * i;                          // can enter parent scope)  
        string * s;  
    };  
    Tagged(int i) : tag(INT), i(new int(i)) {}  
    Tagged(const char * s) : tag(STR), s(new string(s)) {}  
    ~Tagged() {  
        if (tag == INT) delete i; else delete s;  
    }  
};
```


Rust Enums

- Natively supports tagged union by mixing it with enum

```
let tree = Tree::Node(2,  
    Box::new(Tree::Node(0,  
        Box::new(Tree::Leaf),  
        Box::new(Tree::Leaf),  
    )),  
    Box::new(Tree::Node(3,  
        Box::new(Tree::Leaf),  
        Box::new(Tree::Node(4,  
            Box::new(Tree::Leaf),  
            Box::new(Tree::Leaf),  
        )),  
    )),  
);
```

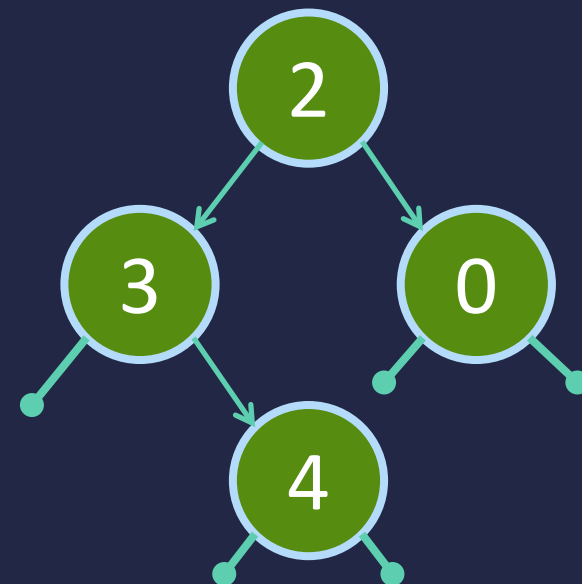
```
enum Tree {  
    Leaf,  
    Node(i64,  
        Box<Tree>,  
        Box<Tree>)  
}
```



Rust Enums

```
enum Tree {  
    Leaf,  
    Node(i64, Box<Tree>, Box<Tree>)  
}  
  
fn add_values(tree: Tree) -> i64 {  
    match tree {  
        Tree::Node(v, a, b) => {  
            v + add_values(*a) +  
              add_values(*b)  
        },  
        Tree::Leaf => 0  
    }  
}
```

```
assert_eq!(add_values(tree), 9);
```



You do not need to say **return** in Rust. Just leave an expression by the end of a function. In the case of multiple branches, all branches must return the same type.

Subtype Compatibility

and Contract Programming

Variance

- Compatibility of types and their subtypes
- Substitutability
 - If S is a subtype of T, then objects of type T may be replaced with objects of type S without altering program correctness
- Complication arises with dealing with complex types
- Refers to how subtyping relation changes
 - E.g. when used as a parameter in a virtual function
 - E.g. when used as a return type in a virtual function
 - E.g. when used as a parameterized type

Variance

- Suppose Apple is a subtype of Fruit
 - Is a list of apple a subtype of list of fruit?

```
template<typename T>                struct Fruit {
struct List<T> {                    int calories() const=0;
    T * get(unsigned i);           float weight() const=0;
    void add(T *);                 float sweetness() const=0;
    unsigned count() const;        };
};

int total_calories(const List<Fruit> & basket) {
    int total = 0;
    for (int i = 0; i < basket->count(); i++)
        total += basket->get(i)->calories();
    return total;
}
```

Covariance

- Allows use of more derived type than specified
 - Ordering of types is preserved
 - type can be substituted by subtype

```
int total_calories(const List<Fruit> & basket) {  
    int total = 0;  
    for (int i = 0; i < basket->count(); i++)  
        total += basket->get(i)->calories();  
    return total;  
}
```

```
List<Apple> apples = { new Apple("Gala", 0.3),  
                     new Apple("Fuji", 0.4), ... };
```

```
// covariant, but NOT allowed in C++  
cout << total_calories(apples) << endl;
```

Covariance

- For parameterized types, only safe if immutable

```
List<Apple> apples = { new Apple("Gala", 0.3),  
                      new Apple("Fuji", 0.4), ... };
```

```
List<Fruit> * fruits = &apples;
```

```
// THIS IS UNSAFE  
fruits->add(new Orange("Navel", 0.5));
```

- However, Java allows the above to occur
 - Will cause a runtime `java.lang.ArrayStoreException`
 - Unless `List<Fruit>` is actually implemented as a list of assorted fruits

Contravariance

- Allows use of more generic type than specified
 - Ordering of types is *reversed*

```
struct Canine {  
    virtual void eat(Meat *)=0;  
};  
  
struct Wolf : public Canine {  
    ...  
};  
  
struct Dog : public Canine {  
    virtual void eat(Food *) override;  
}
```

```
struct Food {};  
struct Meat : public Food {};  
struct Fruit : public Food {};  
struct Apple : public Fruit {};
```

`eat(Food *)` is a subtype of `eat(Meat *)`. In other words, `eat` is contravariant because the parameter of its subtype is more generic.

```
Dog fido;    Apple * apple = new Apple("Honeycrisp", 0.33);  
fido.eat(apple);    // contravariant, but NOT allowed in C++
```


Contravariant Parameter

```
struct Canine {  
    virtual void eat(Meat *)=0;  
};  
  
struct Dog : public Canine {  
    virtual void eat(Food *) override;  
}  
  
struct Food {};  
struct Meat : public Food {};  
struct Fruit : public Food {};  
struct Apple : public Fruit {};  
  
Dog fido;    Apple * apple = new Apple("Honeycrisp", 0.33);  
fido.eat(apple);    // this is type safe, a dog can eat any food  
  
Canine * canine = &fido;    // now fido is a canine  
  
// still type safe, fido can only eat meat as a canine,  
// but meat is also food so a dog can eat it  
canine->eat(new Meat("Beef", "10oz"));
```

Covariant Parameter

```
struct Canine {  
    virtual void eat(Food *)=0;  
};  
  
struct Wolf : public Canine {  
    virtual void eat(Meat *) override;  
}  
  
Wolf kevin;  
// this is type safe, a wolf eats meat only  
kevin.eat(new Meat("Venison", "8oz"));  
  
Canine * canine = &kevin;    // now kevin is a canine  
  
// NOT TYPE SAFE, wolf cannot eat anything but meat!  
canine->eat(new Apple("Red Delicious", 0.5));
```

Invariance

- Nonvariant
- Only the specified type is accepted
- Most conservative, least flexible
- C++ allows covariant return types

```
struct AnimalShelter {  
    virtual Animal * adopt();  
};
```

```
struct CatShelter : public AnimalShelter {  
    // OK - return type is a subtype of Animal  
    virtual Cat * adopt() override;  
}
```

Liskov's Substitution Principle

- Correctness of function subtyping is guaranteed if:
 - Method parameters are contravariant
 - Method return type is covariant
- Correctness of behavioural subtyping is guaranteed if:
 - Precondition
 - Condition that must be true before execution of some code
 - Cannot be strengthened in a subtype
 - Postcondition
 - Condition that must be true after execution of some code
 - Cannot be weakened in a subtype

Square Rectangle Problem

- Base class can incorrectly mutate its derive class
- Violation of Liskov's substitution principle
 - Square has a stronger precondition than Rectangle

```
class Rectangle {  
    int width, height;  
public:  
    Rectangle(int w, int h) : width(w), height(h) {}  
    void set_width(int w) { width = w; }  
    void set_height(int h) { height = h; }  
};  
  
struct Square : public Rectangle {  
    Square(int s) : Rectangle(s, s) {}  
};  
  
Square sq(5);          sq.set_width(6);          // oops
```

subclass cannot
properly adhere
to the interface
of the interface

Example

- Weakened post-condition

```
class Shipment {  
    double dimensions[3];  
public:  
    Shipment(double w, double h, double d);  
    virtual double cost() const;  
};  
  
class DiscountShipment : public Shipment {  
    double discount;  
public:  
    virtual double cost() const override;  
};  
  
Shipment * package = new DiscountShipment(...);  
  
package->cost();    // -$10.00 !?!?
```

Contract Programming

- Language support for specifying precondition, postcondition, errors, and invariants
- Allows business logic to be written more “aggressively”
 - Without having to type check or verify assumptions
- Opposite of “defensive programming”
- Fills the gap that type checking cannot accomplish
- Some can be done statically
- Most are done at runtime

Assertion

- Verifies that an expression is true at runtime

```
assert(i > 0);           // crashes program if expression not true
assert(j < 0, "j must be negative"); // supported by D language
```

- Static version available

```
// syntax in D programming language
static assert(Message.sizeof <= 1024,
    "size of message exceeds maximum packet size");
```

```
// in C++
static_assert(sizeof(Message) <= 1024);
```


Contracts

- Defining precondition and postcondition

```
int fun(ref int a, int b)
in {
    assert(a > 0);
    assert(b >= 0, "b cannot be negative!");
}
out (r) {    // r binds to the return value of fun
    assert(r > 0, "return must be positive");
    assert(a != 0);
}
do {
    a = (b - 3)*(b + 2) + 1;
    return b*b;
}
```

Invariants

- Characteristics of a class that must always be true
 - Except in methods, when temporary violation can occur

```
class Date {
    int day, hour;

    this(int d, int h) {
        day = d; hour = h;
    }

    void add_hours(int h) {
        hour += h;
        if (hour >= 24) {                // does not trip the invariant!
            d += 1; hour -= 24;
        }
    }

    invariant {
        assert(1 <= day && day <= 31);
        assert(0 <= hour && hour < 24, "hour out of bounds");
    }
}
```