



# COMPUTER COMMUNICATION NETWORKS

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Department of Electronics and Communication Engineering

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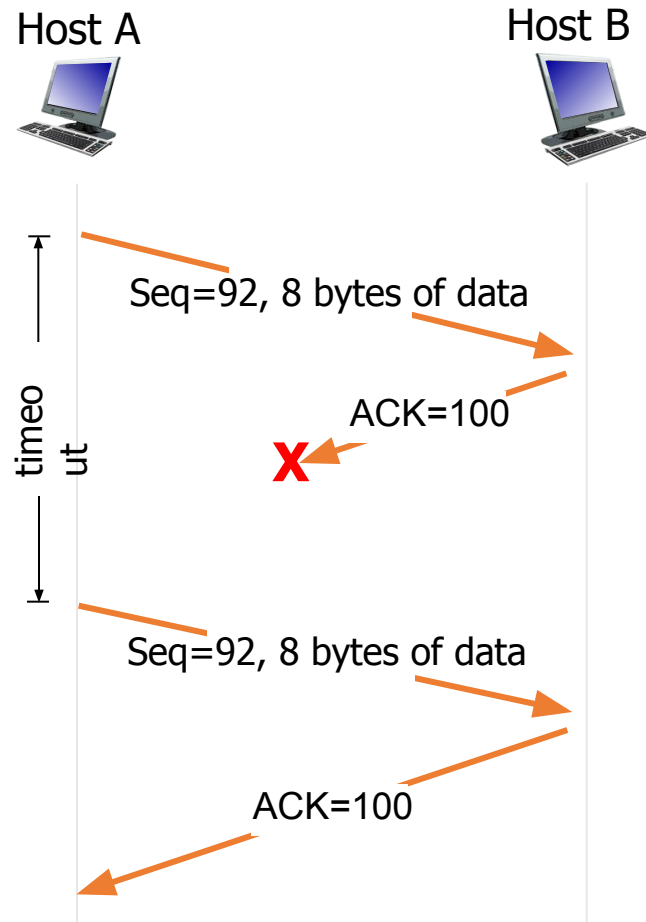
## TCP : Duplicate ACKs & Retransmissions Flow Control

**Dr. Arpita Thakre**

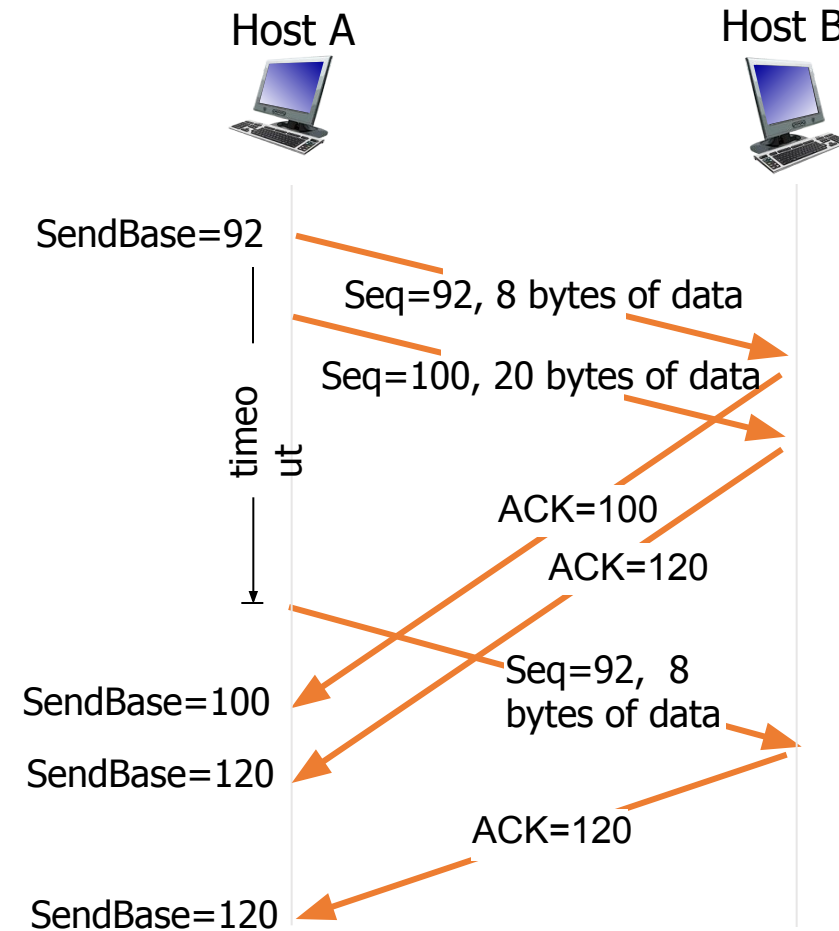
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## TCP Retransmission Scenarios



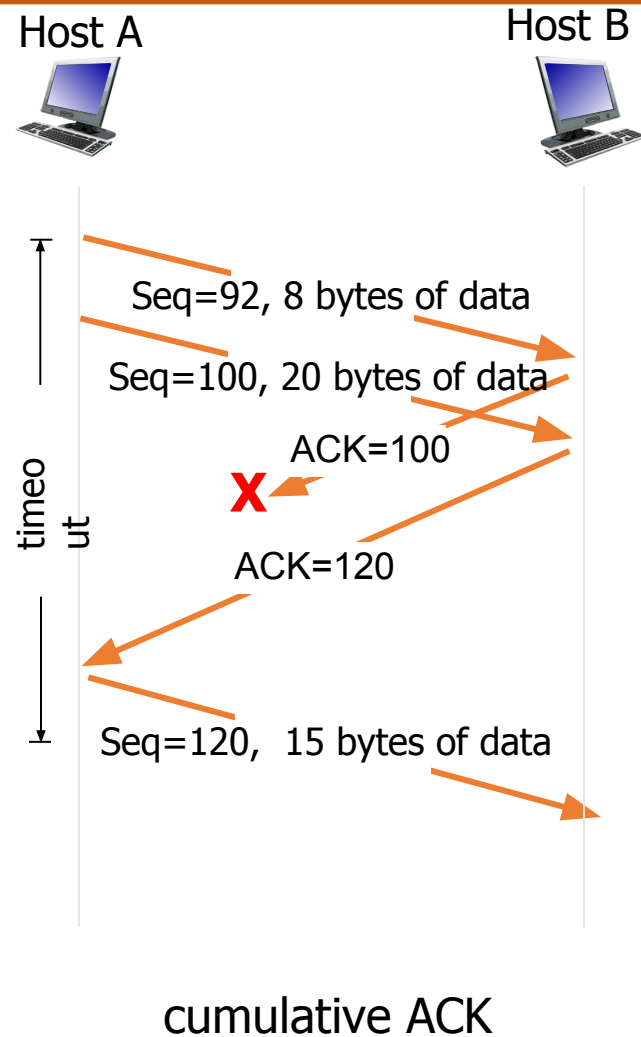
lost ACK scenario



premature timeout

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## TCP Retransmission Scenarios



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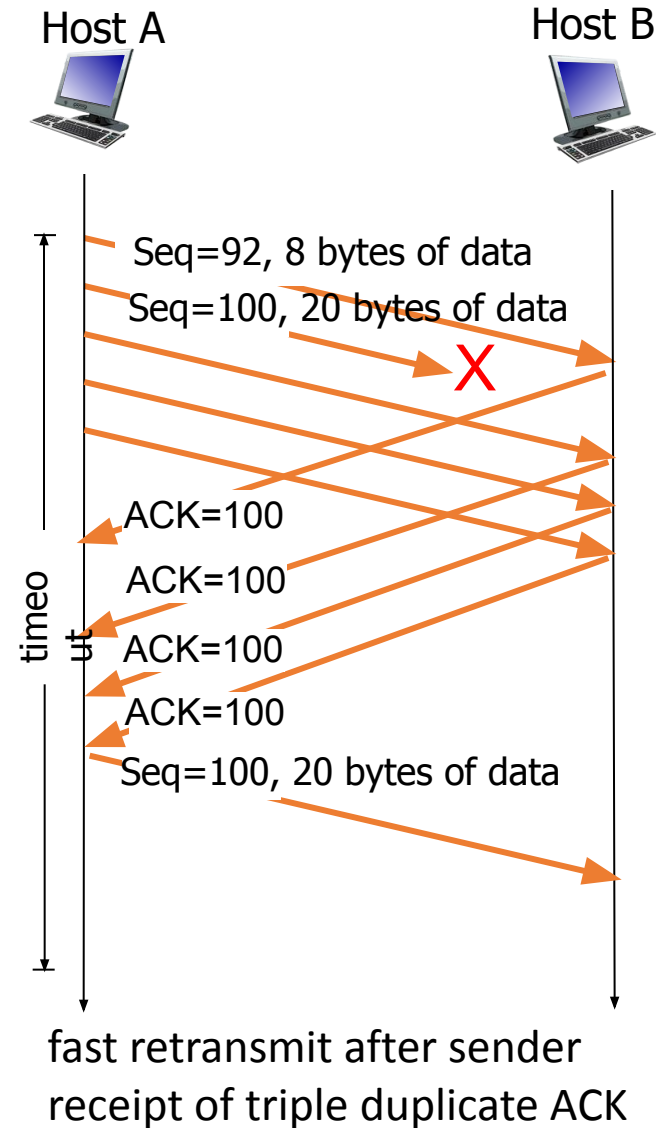
## TCP ACK Generation

<i>event at receiver</i>	<i>TCP receiver action</i>
arrival of in-order segment with expected seq #. All data up to expected seq # already ACKed	delayed ACK. Wait up to 500ms for next segment. If no next segment, send ACK
arrival of in-order segment with expected seq #. One other segment has ACK pending	immediately send single cumulative ACK, ACKing both in-order segments
arrival of out-of-order segment higher-than-expect seq. # . Gap detected	immediately send <i>duplicate ACK</i> , indicating seq. # of next expected byte
arrival of segment that partially or completely fills gap	immediate send ACK, provided that segment starts at lower end of gap

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## TCP Fast Retransmit

### *TCP fast retransmit*



## TCP Flow control

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- TCP provides a flow-control service to its applications to eliminate the possibility of sender overflowing the receiver's buffer.
- Flow control is thus a speed-matching service – matching the rate at which sender is sending against the rate at which receiver application is reading.
- UDP does not provide flow control

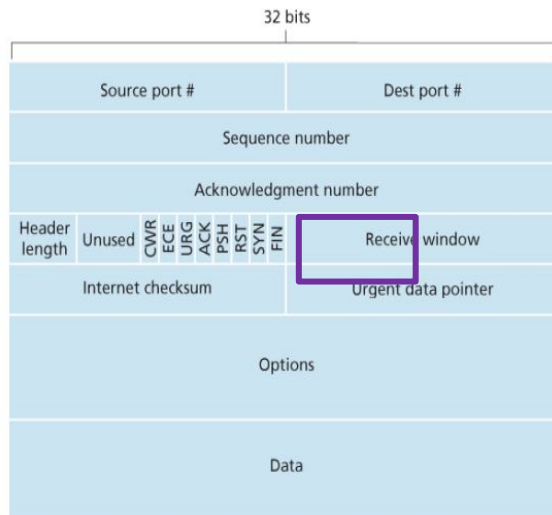
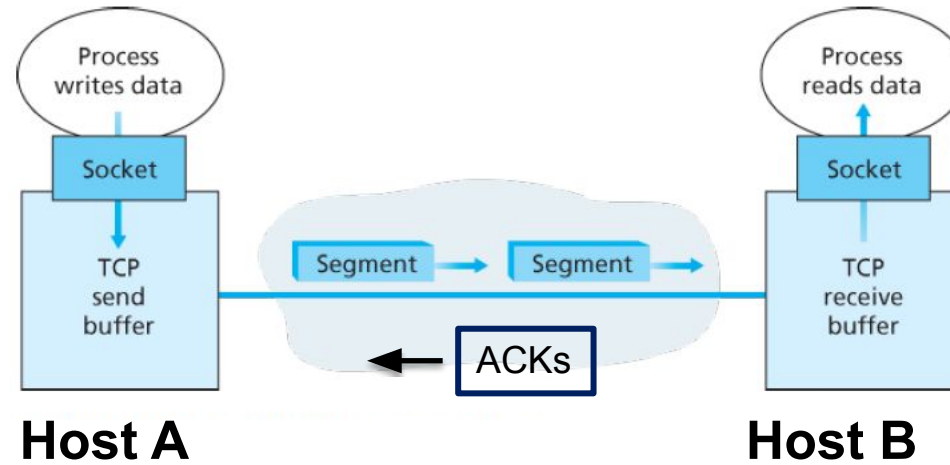




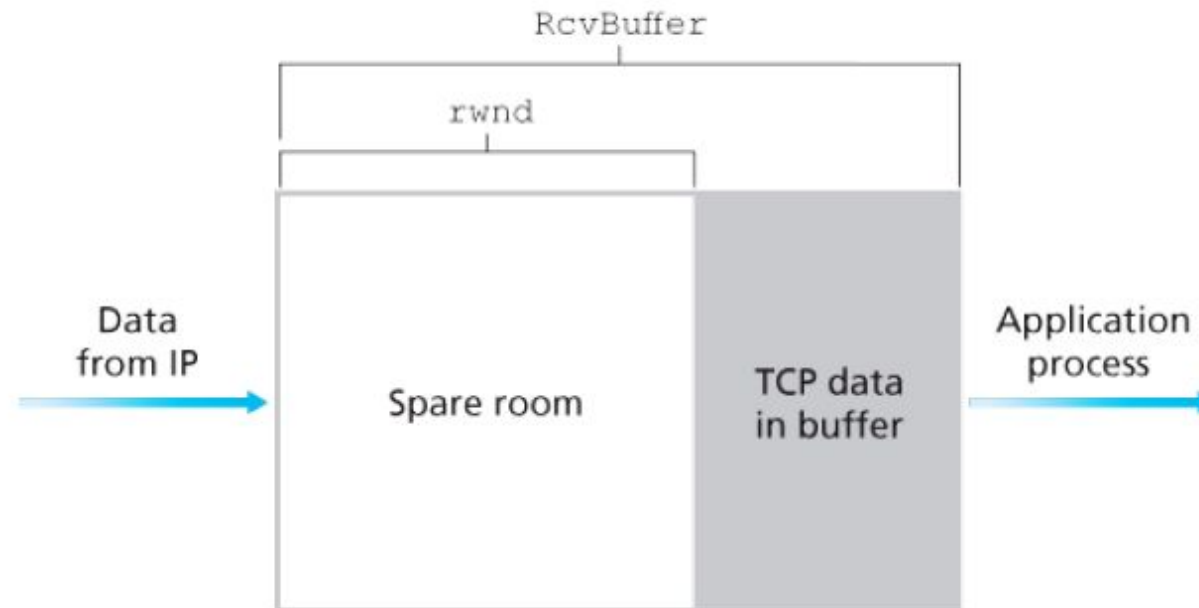
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## TCP Flow control

### Flow Control (contd.)

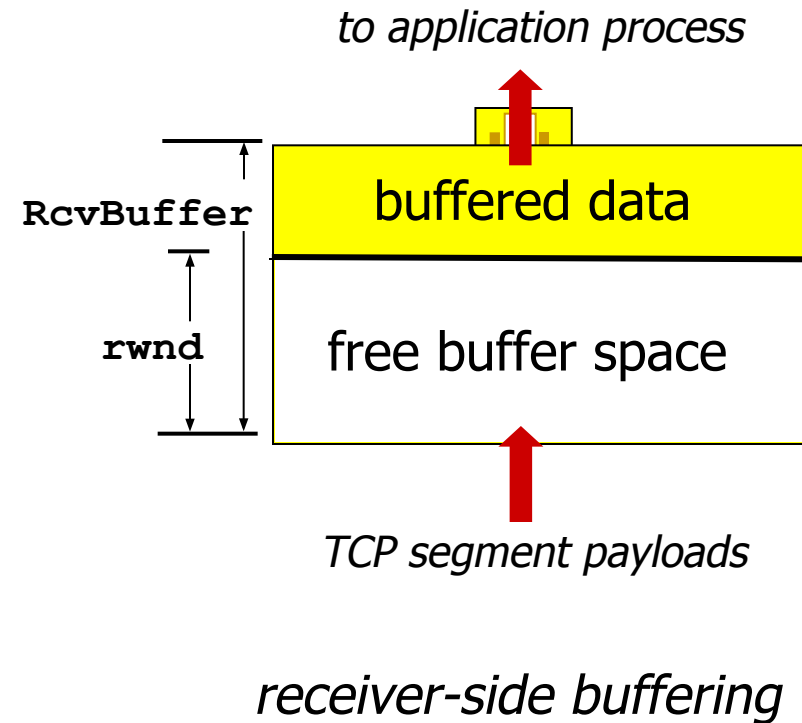


TCP segment



### *Flow Control (contd.)*

- receiver “advertises” free buffer space by including **rwnd** value in TCP header
  - **RcvBuffer** size set via socket options (typical default is 4096 bytes)
  - many operating systems autoadjust **RcvBuffer**
- sender limits amount of unACKed (“in-flight”) data to receiver’s **rwnd** value
- this is how sender guarantees receive buffer will not overflow
- **rwnd** is dynamic; it can be different at different times

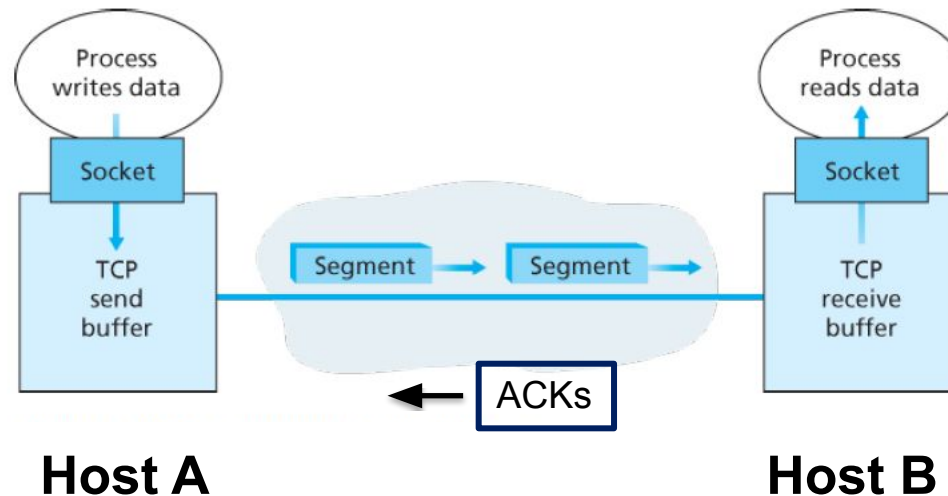


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## TCP Flow control

### Special case

- Host B advertises **rwnd** = 0 to host A & Host B has nothing to send to host A -> host B will not be able advertise **rwnd** value to host A even if **rwnd** changes after some time.
- Consequence: host A cannot send further packets to host B, host A is blocked !
- Solution: TCP protocol mandates that host A send segment of 1 data byte to host B even when **rwnd** = 0, and host B acknowledge this segment
- **rwnd** at host B will increase as time progresses and ACK from host B will now contain updated **rwnd** > 0



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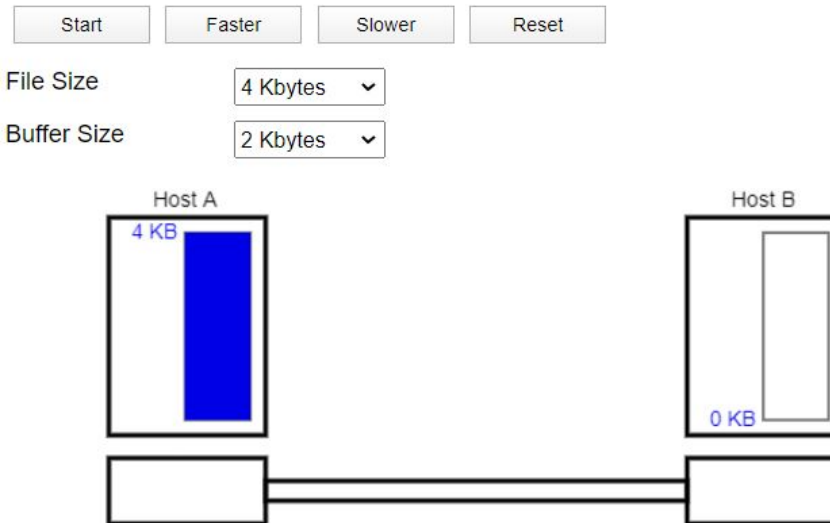
## TCP Flow control

**Animation Link:**

[https://media.pearsoncmg.com/aw/ecs\\_kurose\\_compnetwork\\_7/cw/content/interactiveanimations/flow-control/index.html](https://media.pearsoncmg.com/aw/ecs_kurose_compnetwork_7/cw/content/interactiveanimations/flow-control/index.html)

### Flow Control

This interactive animation shows the interaction between the sending application, the TCP send buffer, the TCP receive buffer, and the receiving application. The receiving application reads chunks of bytes at random times. When the receive buffer becomes full, the TCP receiver advertises a receive window of 0. As described in the text, the sender then continues to send segments with one byte of data.



#### Notes:

1. Host B consumes data in 2Kbyte chunks at random times.
2. When Host A receives an acknowledgment with WIN=0, Host A sends a packet with one byte of data. It is assumed for simplicity, that this one byte is not consumed by the receiver.



# THANK YOU

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