



# COMPUTER COMMUNICATION NETWORKS

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# COMPUTER COMMUNICATION NETWORKS

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## Concept of Layer Based Communication

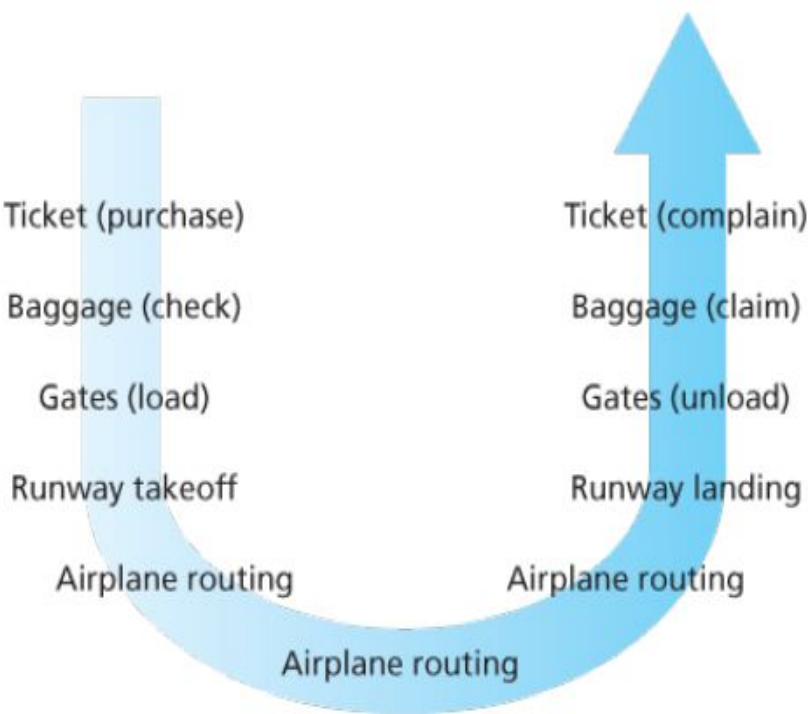
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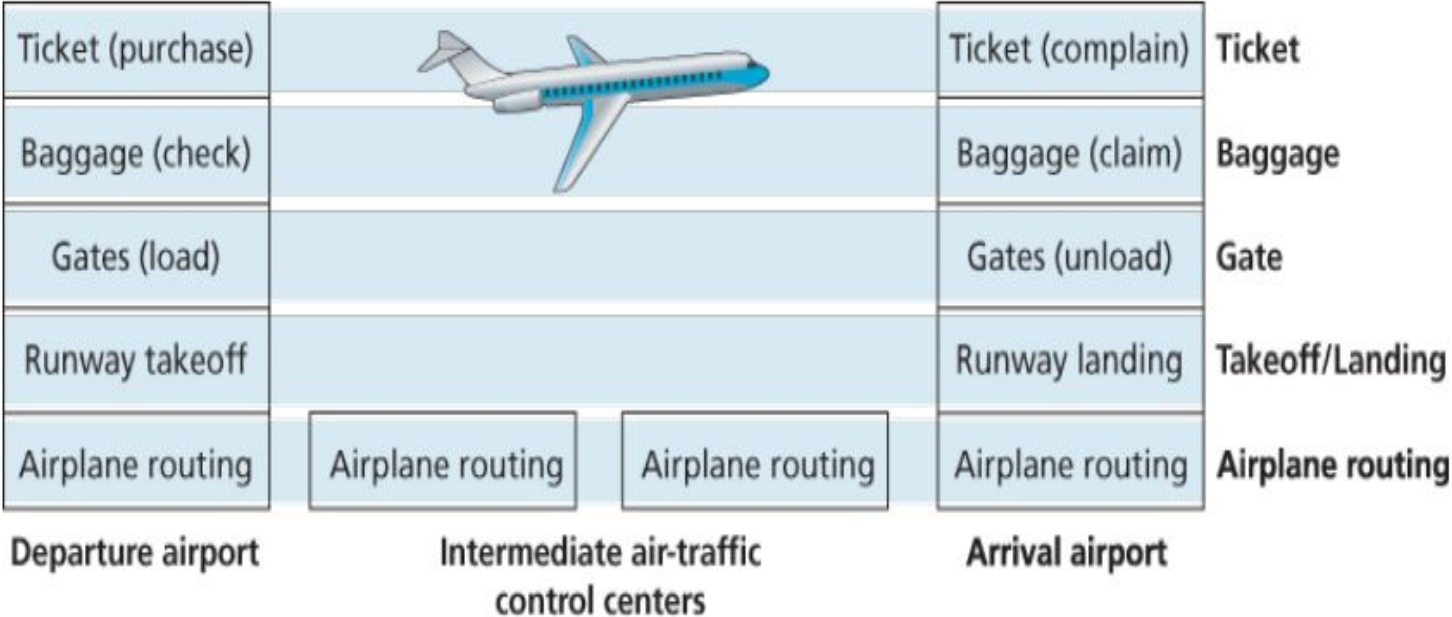
- Data exchange between two hosts over a communication network is a complex task
- The complex task is divided into smaller sub-tasks
  - Maintain simplicity for network devices
  - Put burden on the hosts
- The sub-tasks are completed sequentially
- The entire process can be visualized as layers arranged top to bottom, where
  - Each layer performs its own unique sub-task
  - On the sender side, each layer waits till the above layer finished
    - its sub-task
  - On the receiver side, each layer waits till the below layer finished its sub-task

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## Concept of Protocol Layers



Taking an airplane trip: actions



Horizontal layering of airline functionality

- Communication between two hosts requires the **same layers** to be implemented **in both hosts**.
- Changes to one layer should not result in changes to other layers
- The **peer layers** (i.e., sub-task in sender and its counterpart in the receiver) communicate with one other using formatted blocks of data that obey a set of rules or conventions known as a **protocol**.  
A protocol layer can be implemented in **software**, in **hardware**, or in a **combination of the two**.

**Protocol layering** has conceptual and structural **advantages**:

- **Layering** provides a **structured way** to discuss system components.
- **Modularity** makes it **easier to update** system components.

### Potential Drawbacks of Layering

- One layer may duplicate lower-layer functionality.
- Functionality at one layer may need information (for example, a timestamp value) that is present only in another layer; this violates the goal of separation of layers.

### Basics requirements of a Protocol:

- **Syntax:** Concerns the format of the data blocks
- **Semantics:** Includes control information for coordination and error handling
- **Timing:** Includes speed matching and sequencing

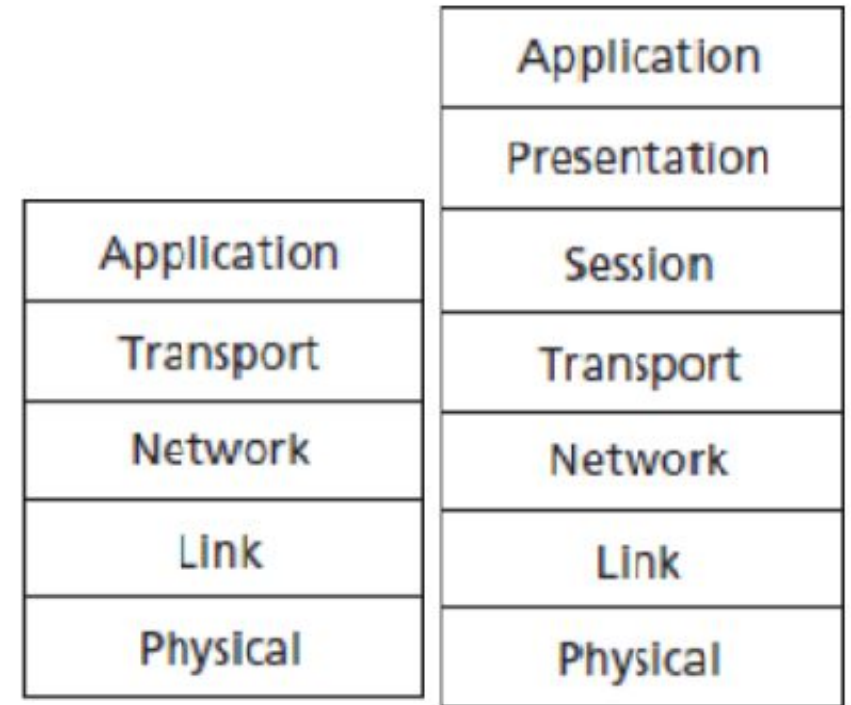
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## Concept of Protocol Layers

- Arranged vertically, the layers on the systems collectively constitute the **protocol architecture**
- Two types of protocol architecture were proposed
  - **TCP/IP**
  - **OSI**
- TCP/IP model or TCP/IP protocol suite
  - Resulted from protocol research under ARPANET
  - Consists of large collection of protocols for various layers called as **protocol stack**.
  - The Internet protocol stack consists of five layers: Application layer, Transport (host-to-host) layer, Network layer (IP layer), Link layer (network access layer), Physical layer



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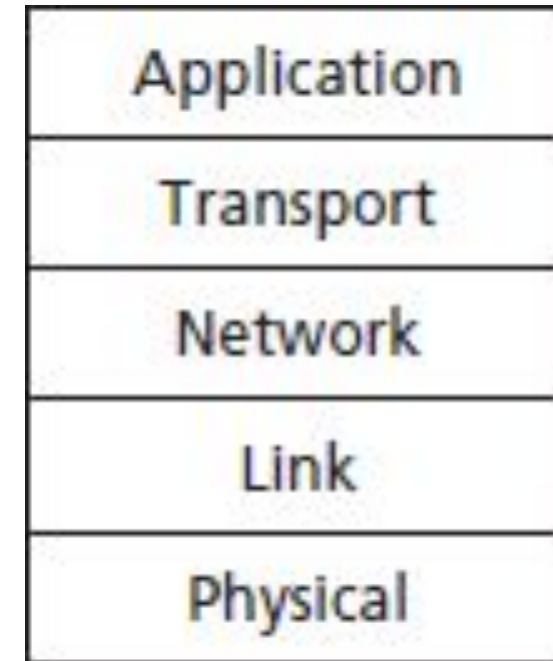


a. Five-layer  
Internet  
protocol stack

b. Seven-layer  
ISO OSI  
reference model

### Application layer :

- **Applications** (aka **processes**) running on hosts generate/receive data.
- The data here is referred to as **message**.
- An **active process** on one host initiates communication with an active process on another host.
- Messages are formatted according to the application layer protocol defined by the application developer.
- Messages can be big in size (e.g., audio; video).
- Applications can have QoS requirements.



TCP/IP model



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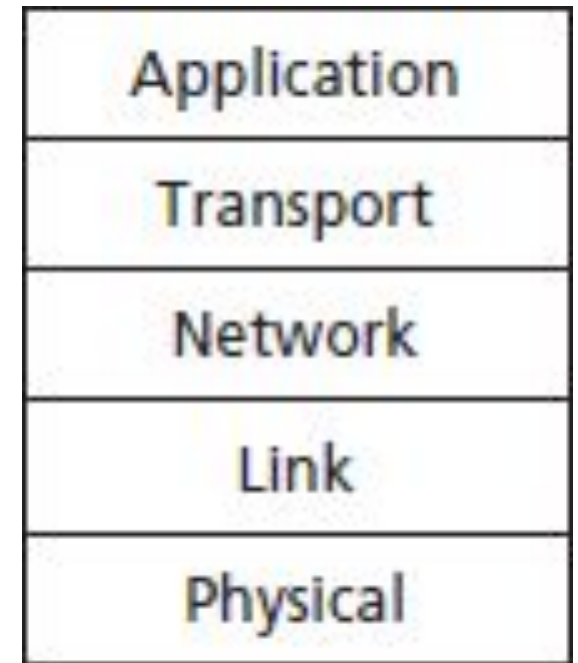
## Concept of Protocol Layer

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### Transport layer:

- Responsible for providing **QoS** for a pair of communicating processes
- Performs **multiplexing at the sender**
- Performs **demultiplexing at the receiver**
- **Maps** each message to a corresponding process inside a host
- This mapping is accomplished by unique identifiers called **ports or sockets**
  - Appends a new header to each message
  - Message plus header is called **segment**



TCP/IP model

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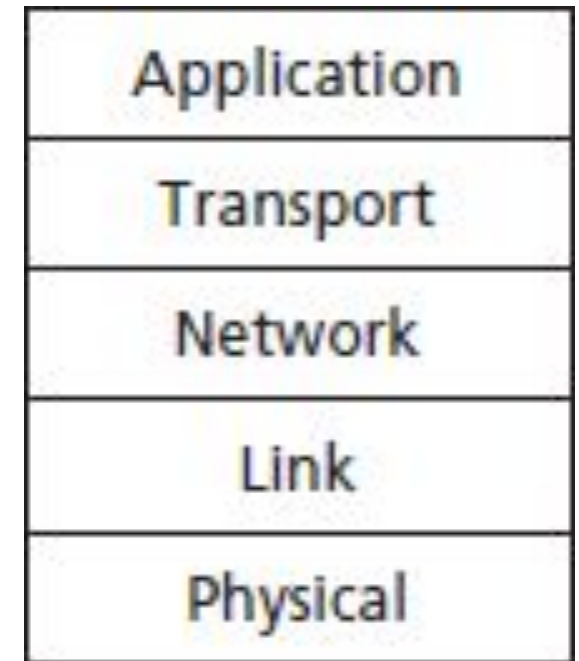
## Concept of Protocol Layers

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### Network layer:

- Moves packets hop-by-hop in search of the destination
- E.g., router to router
- Search for the destination is done via unique identifiers called **IP addresses**
- Network level information is gathered to facilitate discovery of paths to destinations
- Appends a new header to each segment
- Segment plus header is called **datagram**



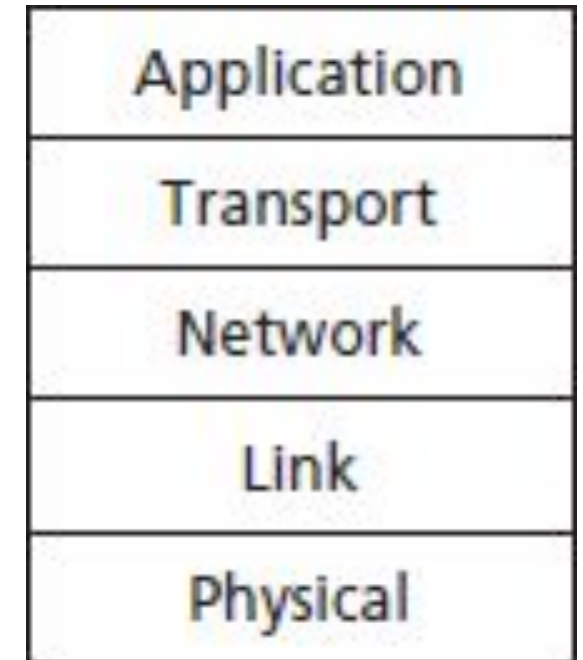
TCP/IP model

### Link layer:

- Ensures **reliable packet flow** across each link on the path between a pair of source and destination independently
- Using respective link layer protocol
- Reliable flow of packets across a link requires identifiers for that link

This is accomplished by **MAC address**

- Provides synchronization between sender and receiver of each link
- Appends a new header to the datagram
- Datagram plus header is called **frame**
- **Checks for errors** in frame



TCP/IP model

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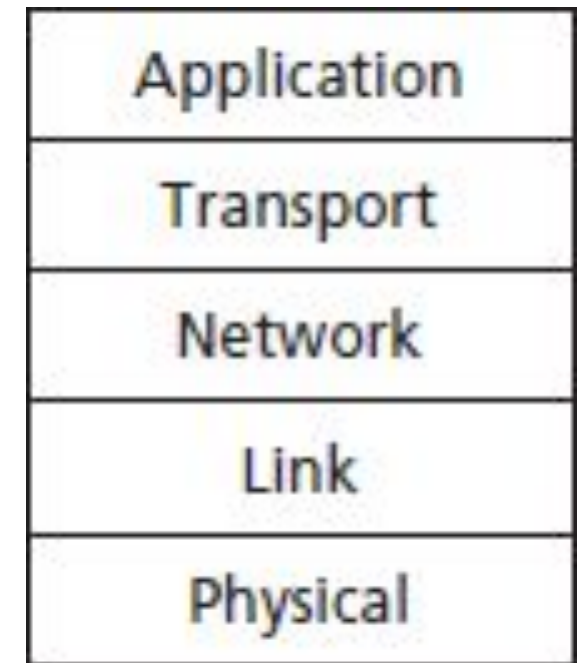
## Concept of Protocol Layers

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### Physical layer:

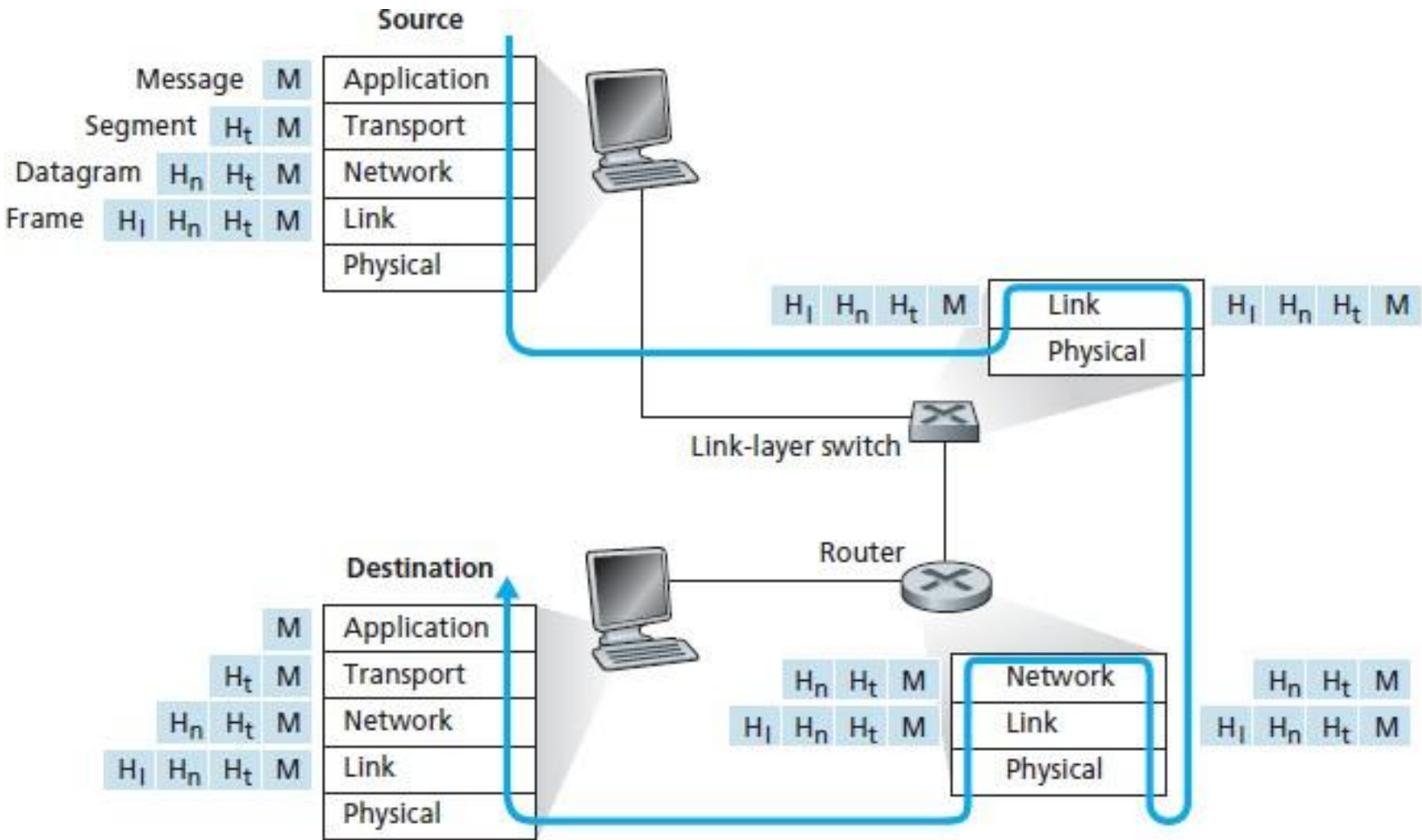
- Deals with the characteristics of the physical medium of a link
- Converts frames into signals (e.g., radio signal, optical signal)
- Performs signal modulation and demodulation
  - E.g.: Modem and Ethernet card, wireless adapter
- Manages transmit power



TCP/IP model

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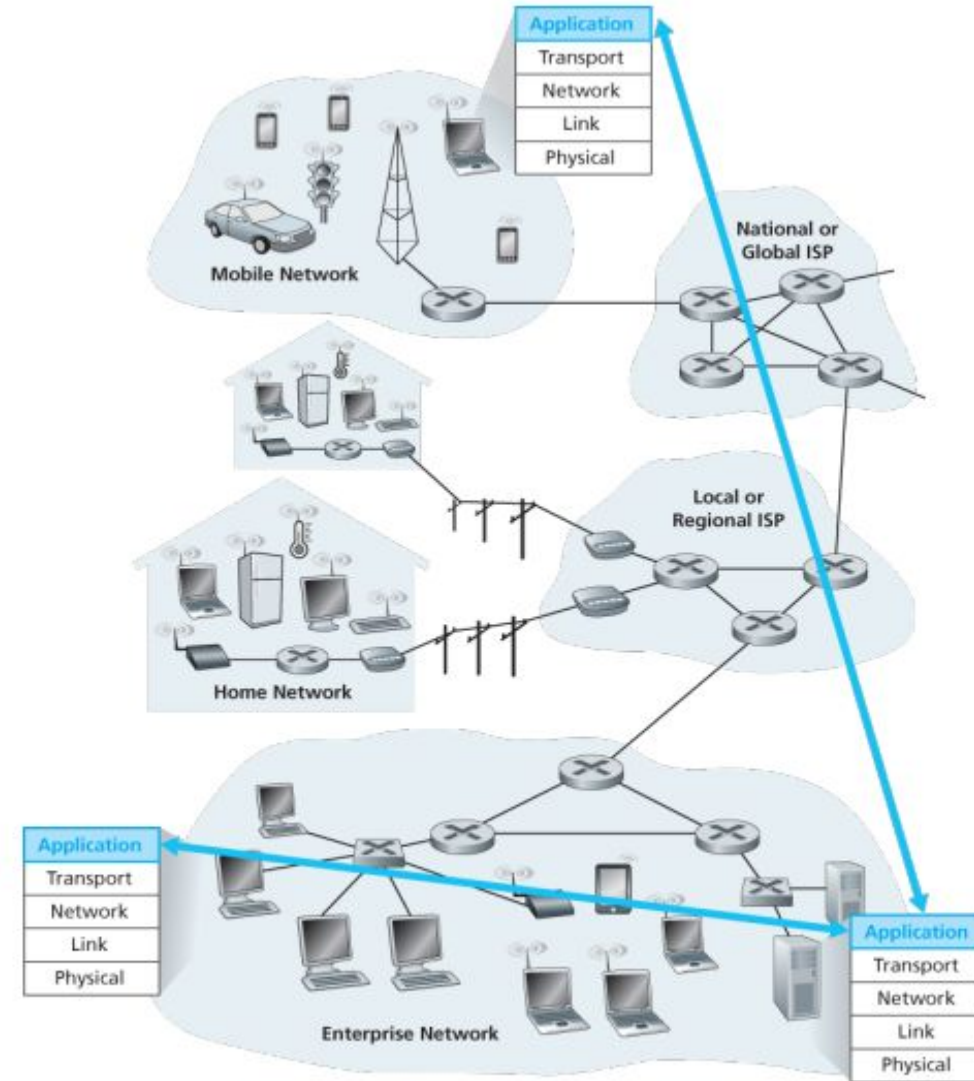
## Encapsulation and Decapsulation



# COMPUTER COMMUNICATION NETWORKS

## Process Communication

- A **process** is a program that is running within an end system.
- When processes are running on the same end system, they can communicate with each other with **inter-process communication**, using rules that are governed by the end system's operating system.
- Processes on two different end systems communicate with each other by **exchanging messages** across the computer network. A **sending process creates and sends** messages into the network; a **receiving process receives** these messages and possibly responds by sending messages back.



Communication for a network application takes place between end systems at the application layer

- A **process** sends messages into, and receives messages from, the network through a software interface called a **socket**.
- A **socket** is the interface between the application layer and the transport layer within a host. It is also referred to as the **Application Programming Interface (API)** between the application and the network layer.
- The application developer has control of everything on the application-layer side of the socket but has little control of the transport-layer side of the socket.
- The only control that the application developer has on the transport-layer side is
  - (1) the choice of transport protocol and
  - (2) perhaps the ability to fix a few transport-layer parameters such as maximum buffer and maximum segment sizes



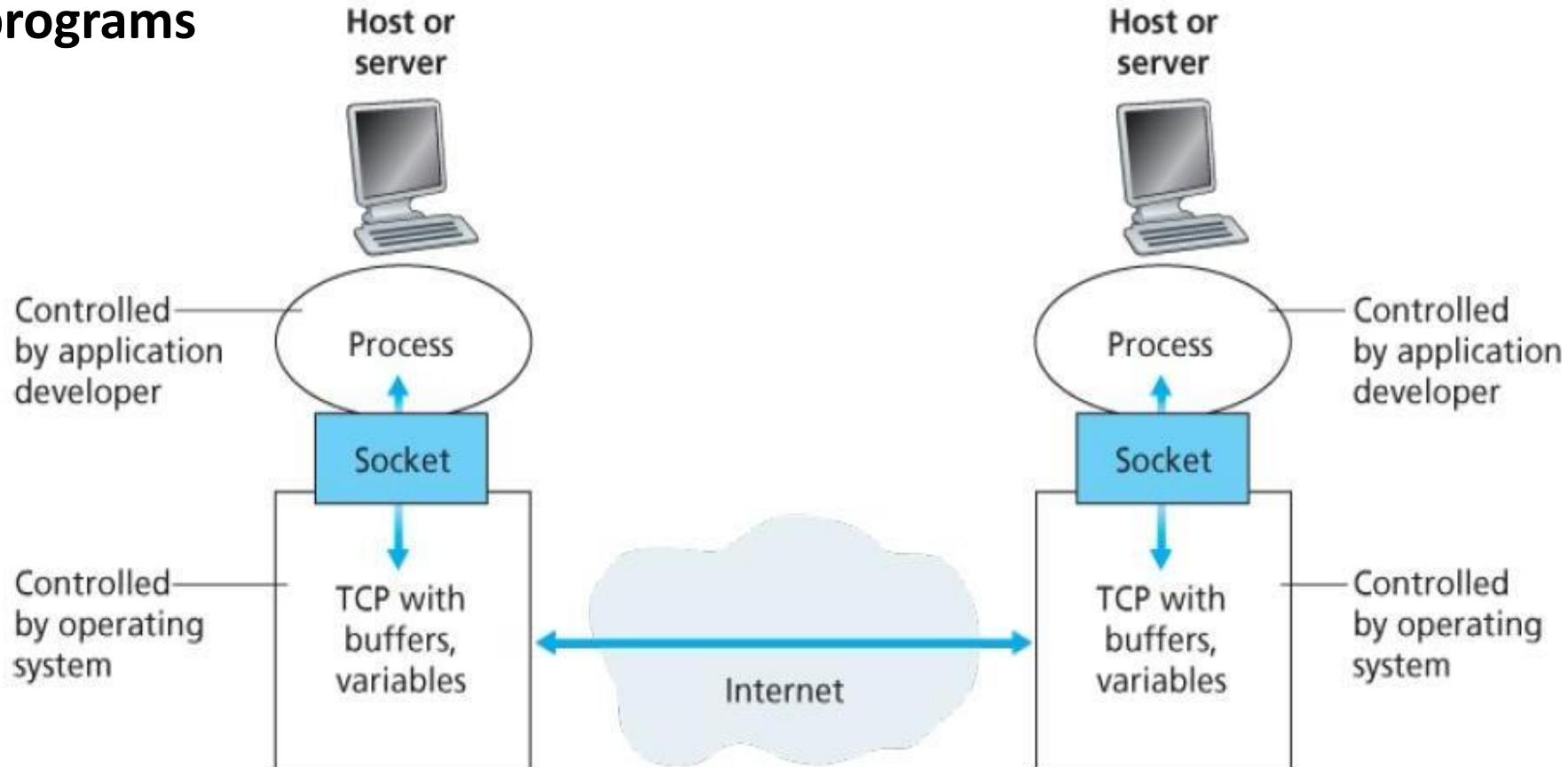
- Application developers provide **unique IDs** for **identifying processes** running in a host (e.g., HTTP server process 80; DNS server process 53)
- Application developers build a **pair of processes** which usually run on different hosts.
  - Example: A browser program running on a laptop and a server program running on a web server
- When one process initiates an action it is transmitted in the form of a special message and the corresponding process which receives this message takes suitable action and returns another special message back to the initiator.
  - Example: A browser program sends a request for a webpage when a URL is entered by a user. The web server receiving the request may return the requested object over the internet to be displayed on the browser



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## Process communication

- These special messages are formatted as defined by the application developer, stored in buffers, transmitted/received using appropriate process identifiers. All these actions can be coded and referred to as **socket programs**



### Addressing Processes

- For a process running on one host to send packets to a process running on another host, the receiving process needs to have an address.
- To identify the receiving process, two pieces of information need to be specified:
  - the address of the host and
  - an identifier that specifies the receiving process in the destination host
- In the Internet, the host is identified by its IP address.
- The sending process identifies the receiving process (the receiving socket) running in the host using a destination port number.

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## Client - Server Model

### Server process:

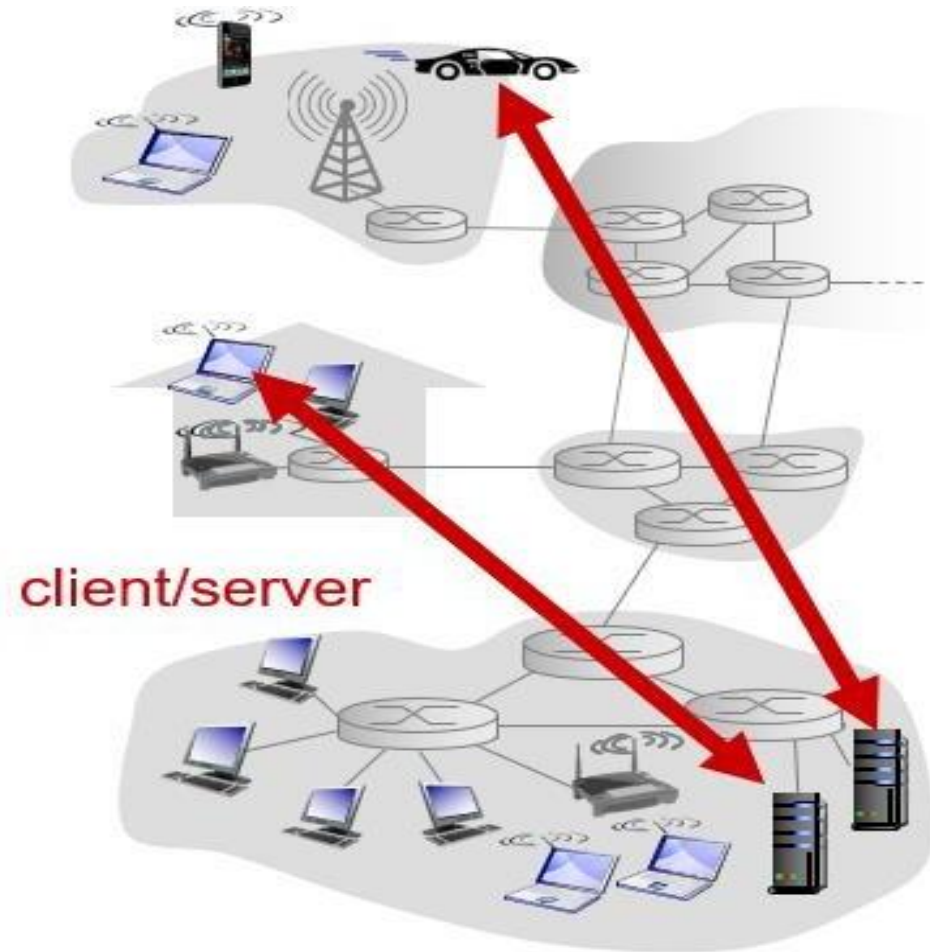
- ❖ Always running
- ❖ Uniquely identified by socket and IP address

### Client process:

- ❖ Activated on demand
- ❖ Many client process may communicate with one server process
- ❖ Each client process randomly chooses a socket for its ID

### Client-server model examples

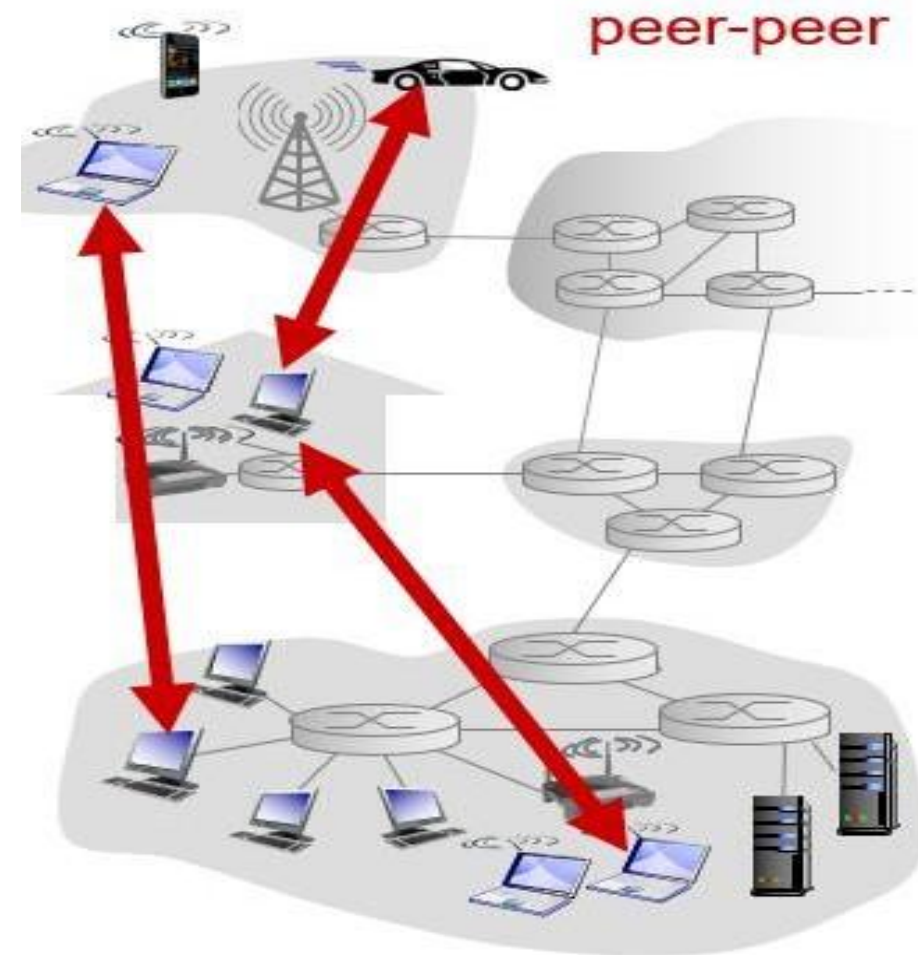
- ❖ Web, Email, File Transfer



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## Peer to Peer Model

- ❖ Peer processes are activated on demand
- ❖ Peers discover sockets and IP addresses of one another
- ❖ Involve complex process management to ensure good performance
- ❖ Applications based on peer-to-peer model exhibit self scalability and adaptability
- ❖ Best suited for file distribution
  - Examples: BitTorrent, Skype





# THANK YOU

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