

Principles of Digital Signal Processing

Dr. B. Niranjana Krupa

Department of Electronics and Communication.

DSP



Discrete Fourier Transform

Dr. B. Niranjana Krupa

Department of Electronics and Communication.

DFT



Recovering x(n):

When x(n) has infinite duration, the equally spaced frequency samples do not represent x(n).

They (frequency samples) only correspond to a periodic sequence $x_p(n)$, the aliased version of x(n), as in

$$x_p(n) = \sum_{l=-\infty}^{\infty} x(n-lN)$$

DFT



When x(n) has finite duration of length $L \le N$, then

$$x_p(n) = \sum_{l=-\infty}^{\infty} x(n-lN)$$

$$x_p(n) = \begin{cases} x(n), & 0 \le n \le L - 1 \\ 0, & L \le n \le N - 1 \end{cases}$$

DFT



It is important to note that zero padding doesnt provide any additional information about the spectrum of the sequence. However padding zeros and computing N point DFT only results in better display.

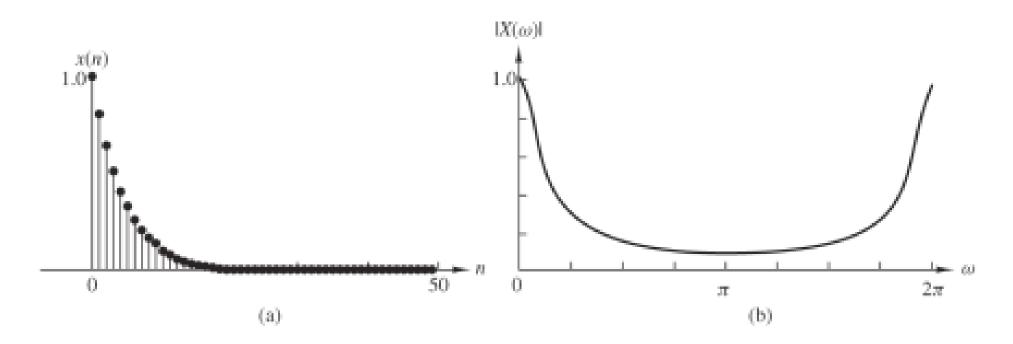


Figure 7.1.4 (a) Plot of sequence $x(n) = (0.8)^n u(n)$; (b) its Fourier transform

DFT



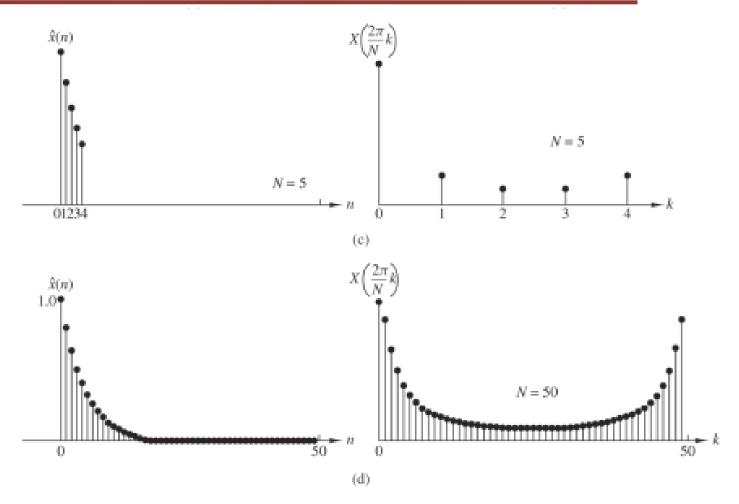


Figure 7.1.4 (a) Plot of sequence $x(n) = (0.8)^n u(n)$; (b) its Fourier transform (magnitude only); (c) effect of aliasing with N = 5; (d) reduced effect of aliasing with N = 50.

DFT



In summary, a finite-duration sequence x(n) of length L [i.e., x(n) = 0 for n < 0 and $n \ge L$] has a Fourier transform

$$X(\omega) = \sum_{n=0}^{L-1} x(n)e^{-j\omega n} \qquad 0 \le \omega \le 2\pi$$

where the upper and lower indices in the summation reflect the fact that x(n) = 0 outside the range $0 \le n \le L-1$. When we sample $X(\omega)$ at equally spaced frequencies $\omega_k = 2\pi k/N$, k = 0, 1, 2, ..., N-1, where $N \ge L$, the resultant samples are

$$X(k) \equiv X\left(\frac{2\pi k}{N}\right) = \sum_{n=0}^{L-1} x(n)e^{-j2\pi kn/N}$$

$$X(k) = \sum_{n=0}^{N-1} x(n)e^{-j2\pi kn/N}, \qquad k = 0, 1, 2, \dots, N-1$$

DFT



To summrise DFT and IDFT

DFT:
$$X(k) = \sum_{n=0}^{N-1} x(n)e^{-j2\pi kn/N}, \qquad k = 0, 1, 2, ..., N-1$$

IDFT:
$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) e^{j2\pi kn/N}, \qquad n = 0, 1, 2, \dots, N-1$$

DFT as linear transform



DFT and IDFT expressed as

$$X(k) = \sum_{n=0}^{N-1} x(n) W_N^{kn} \qquad k = 0, 1, \dots, N-1$$

$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) W_N^{-kn} \qquad n = 0, 1, \dots, N-1$$

Where,
$$W_N = e^{-j2\pi/N}$$

The Nth root of unity

DFT as linear transform



- Computational complexity of DFT
 - Each point: N complex multiplications and (N-1) complex additions
 - Hence, N-point DFT:
 - N*N complex multiplications and
 - N*(N−1) complex additions
- DFT and IDFT as linear transformations on x(n) and X(k)
 - Consider N-point vectors x_N of x(n) and X_N of X(k) and an NXN matrix W_N as:

DFT as linear transform



$$\mathbf{x}_{N} = \begin{bmatrix} x(0) \\ x(1) \\ \vdots \\ x(N-1) \end{bmatrix}, \qquad \mathbf{X}_{N} = \begin{bmatrix} X(0) \\ X(1) \\ \vdots \\ X(N-1) \end{bmatrix}$$

DFT as linear transform



$$\mathbf{W}_{N} = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & W_{N} & W_{N}^{2} & \cdots & W_{N}^{N-1} \\ & W_{N}^{2} & W_{N}^{4} & \cdots & W_{N}^{2(N-1)} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & W_{N}^{N-1} & W_{N}^{2(N-1)} & \cdots & W_{N}^{(N-1)(N-1)} \end{bmatrix}$$

DFT as linear transform



N-point DFT expressed in matrix form as:

$$\mathbf{X}_N = \mathbf{W}_N \mathbf{x}_N$$

 W_N is the matrix of linear transformation and is a symmetric matrix

If we assume that the inverse of W_N exists then,

$$(DFT) \rightarrow \mathbf{x}_N = \mathbf{W}_N^{-1} \mathbf{X}_N \qquad \mathbf{x}_N = \frac{1}{N} \mathbf{W}_N^* \mathbf{X}_N$$

DFT as linear transform

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Comparing

$$\mathbf{W}_N^{-1} = \frac{1}{N} \mathbf{W}_N^*$$

$$\mathbf{W}_N \mathbf{W}_N^* = N \mathbf{I}_N$$

 W_N is an orthogonal (unitary) matrix And its inverse exists, given as shown above



THANK YOU

Dr. B. Niranjana Krupa

Department of Electronics and Communication

bnkrupa@pes.edu

+91 80 6666 3333 Ext 777