



RISC V Architecture

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RISC V ARCHITECTURE

UNIT 2 – Instructions: The Language of Computer

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Instructions – Language of Computer

Signed and Unsigned Numbers



Let us consider a 32 bit (word) number representation

MSB

LSB

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	1	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

(32 bits wide)

Why didn't computers use decimal?

How many bit pattern it can detect ?

What number they represent in decimal ?

```
00000000 00000000 00000000 00000000two = 0ten
00000000 00000000 00000000 00000001two = 1ten
00000000 00000000 00000000 00000010two = 2ten
...
11111111 11111111 11111111 11111101two = 4,294,967,293ten
11111111 11111111 11111111 11111110two = 4,294,967,294ten
11111111 11111111 11111111 11111111two = 4,294,967,295ten
```

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Signed and Unsigned Numbers

Let us understand Unsigned Number Representation using 32 bit binary number

Binary Number								Hexa-Decimal
31							0	
0000	0000	0000	0000	0000	0000	0000	0000	(0x 0000 0000)
0000	0000	0000	0000	0000	0000	0000	0001	(0x 0000 0001)
0000	0000	0000	0000	0000	0000	0000	0010	(0x 0000 0002)
...								
...								
0111	1111	1111	1111	1111	1111	1111	1111	(0x 7FFF FFFF)
1000	0000	0000	0000	0000	0000	0000	0010	(0x 8000 0000)
1000	0000	0000	0000	0000	0000	0000	0001	(0x 8000 0001)
...								
....								
1111	1111	1111	1111	1111	1111	1111	1110	(0x FFFF FFFE)
1111	1111	1111	1111	1111	1111	1111	1111	(0x FFFF FFFF)

0_{ten}
1_{ten}
2_{ten}

4,294,967,293_{ten}
4,294,967,294_{ten}
4,294,967,295_{ten}

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Signed and Unsigned Numbers

How Signed Number Representation is different from Unsigned Number Representation?

31	30	0
Sign	Magnitude	

2's complement representation indicates negative numbers have a 1 in the MSB. Thus, hardware needs to test only MSB to decide whether the Number is positive or negative.

31	30									0	to decide whether the Number is positive or	
0	000	0000	0000	0000	0000	0000	0000	0000	0000	0000	(0x 0000 0000) 0	Positive Number Range
0	000	0000	0000	0000	0000	0000	0000	0000	0000	0001	(0x 0000 0001) 1	
0	000	0000	0000	0000	0000	0000	0000	0000	0000	0010	(0x 0000 0002) 2	
...												
...												
0	111	1111	1111	1111	1111	1111	1111	1111	1111	1111	(0x 7FFF FFFF) ..	Negative Number Range & They are in 2's complement form
1	000	0000	0000	0000	0000	0000	0000	0000	0000	0010	(0x 8000 0000) ..	
1	000	0000	0000	0000	0000	0000	0000	0000	0000	0001	(0x 8000 0001)	
...												
....												
1	111	1111	1111	1111	1111	1111	1111	1111	1111	1110	(0x FFFF FFFE)=-2	
1	111	1111	1111	1111	1111	1111	1111	1111	1111	1111	(0x FFFF FFFF) =-1	

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Signed and Unsigned Numbers



How Signed Number Representation is different from Unsigned Number Representation?

31	30	0
Sign	Magnitude	

2's complement representation indicates negative numbers have a 1 in the MSB. Thus, hardware needs to test only MSB to decide whether the Number is positive or negative.

2s-Complement Signed Integers

Given an n-bit number

$$x = -x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12^1 + x_02^0$$

Range: -2^{n-1} to $+2^{n-1} - 1$

Example

$$\begin{aligned} &1111\ 1111\ \dots\ 1111\ 1100_2 \\ &= -1 \times 2^{31} + 1 \times 2^{30} + \dots + 1 \times 2^2 + 0 \times 2^1 + 0 \times 2^0 \\ &= -2,147,483,648 + 2,147,483,644 = -4_{10} \end{aligned}$$

Using 64 bits: $-9,223,372,036,854,775,808$ to $9,223,372,036,854,775,807$

Instructions – Language of Computer

Signed and Unsigned Numbers



How Signed Number Representation is different from Unsigned Number Representation?

31	30	0
Sign	Magnitude	

2s-Complement Signed Integers

Some specific numbers

0: 0000 0000 ... 0000

-1: 1111 1111 ... 1111

Most-negative: 1000 0000 ... 0000

Most-positive: 0111 1111 ... 1111

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Signed and Unsigned Numbers



What is the decimal value of this 64-bit two's complement number?

11111111 11111111 11111111 11111111 11111111 11111111 11111111 11111000_{two}

- 1) -4_{ten}
- 2) -8_{ten}
- 3) -16_{ten}
- 4) $18,446,744,073,709,551,608_{\text{ten}}$

What is the decimal value if it is instead a 64-bit unsigned number?

Instructions – Language of Computer

Load and Store Instructions

Category	Instruction	Example	Meaning	Comments
Arithmetic	Add	add x5, x6, x7	$x5 = x6 + x7$	Three register operands; add
	Subtract	sub x5, x6, x7	$x5 = x6 - x7$	Three register operands; subtract
	Add immediate	addi x5, x6, 20	$x5 = x6 + 20$	Used to add constants
Data transfer	Load word	lw x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Word from memory to register
	Load word, unsigned	lwu x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Unsigned word from memory to register
	Store word	sw x5, 40(x6)	$\text{Memory}[x6 + 40] = x5$	Word from register to memory
	Load halfword	lh x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Halfword from memory to register
	Load halfword, unsigned	lhu x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Unsigned halfword from memory to register
	Store halfword	sh x5, 40(x6)	$\text{Memory}[x6 + 40] = x5$	Halfword from register to memory
	Load byte	lb x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Byte from memory to register
	Load byte, unsigned	lbu x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Byte unsigned from memory to register
	Store byte	sb x5, 40(x6)	$\text{Memory}[x6 + 40] = x5$	Byte from register to memory

Instructions – Language of Computer

Signed and Unsigned Numbers

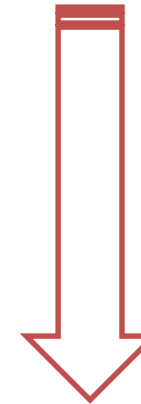
Sign extension

- Representing a number using more bits
Ex: Reading a variable of 16 bit (half word) size and loading into a 32 bit register
- The variable may be a Unsigned / signed number

Memory

15	12	11	8	7	4	3	0
0111	0101	1011	0001				

Loading 16 bit data from memory into a 32 bit register i.e., **Representing a number using more bits**



- ✓ Where is the 16 bit data from memory is loaded in destination register ???
- ✓ What will happen to remaining 16 bits of the destination register i.e. reg[31:16]???
- ✓ reg[31:16] will change the content and It differs based on whether variable is Signed and Unsigned ?

31	28	15	12	11	8	7	4	3	0
1010	1111	1011	1000	0111	0101	1011	0001		

Data in Register in Hexa

0xAFB875B1

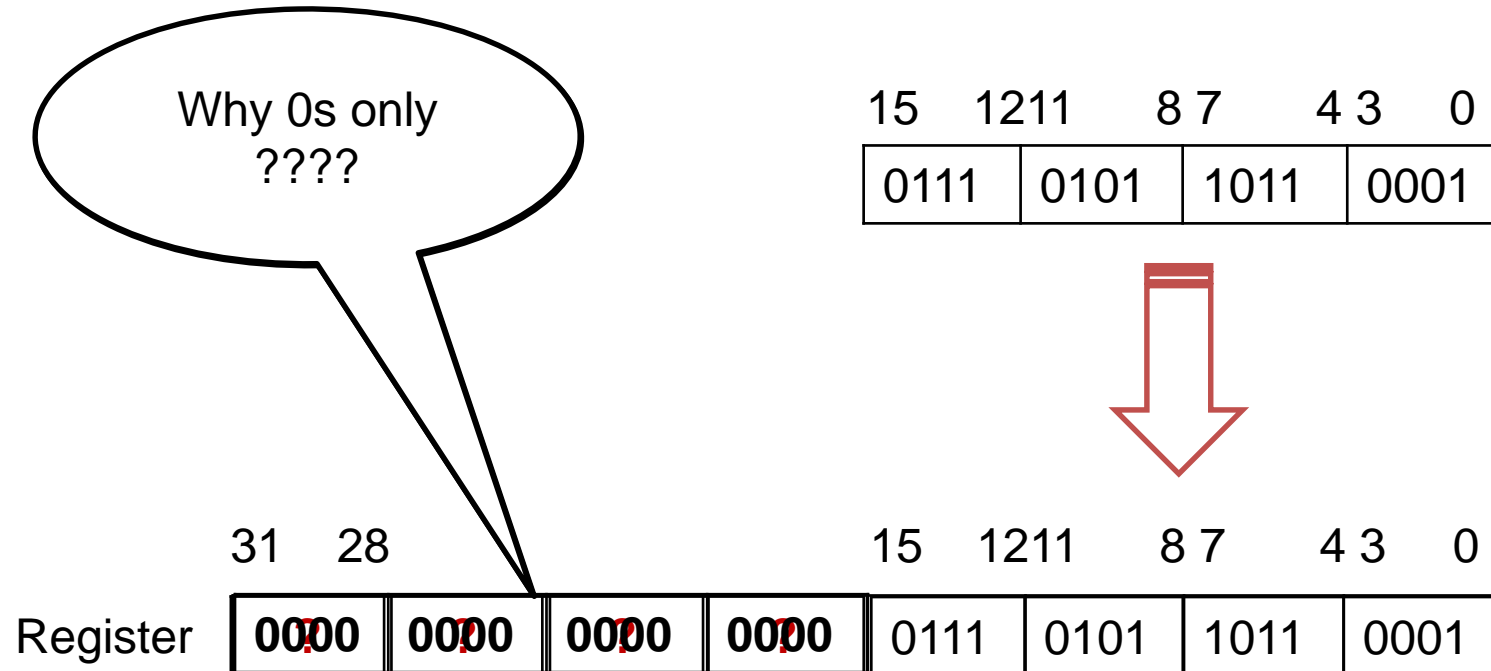
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Signed and Unsigned Numbers

Sign extension

How a Signed and Unsigned Load into a Register differs ?

Let us consider loading 32 bit register with a **16 bit unsigned value** ?



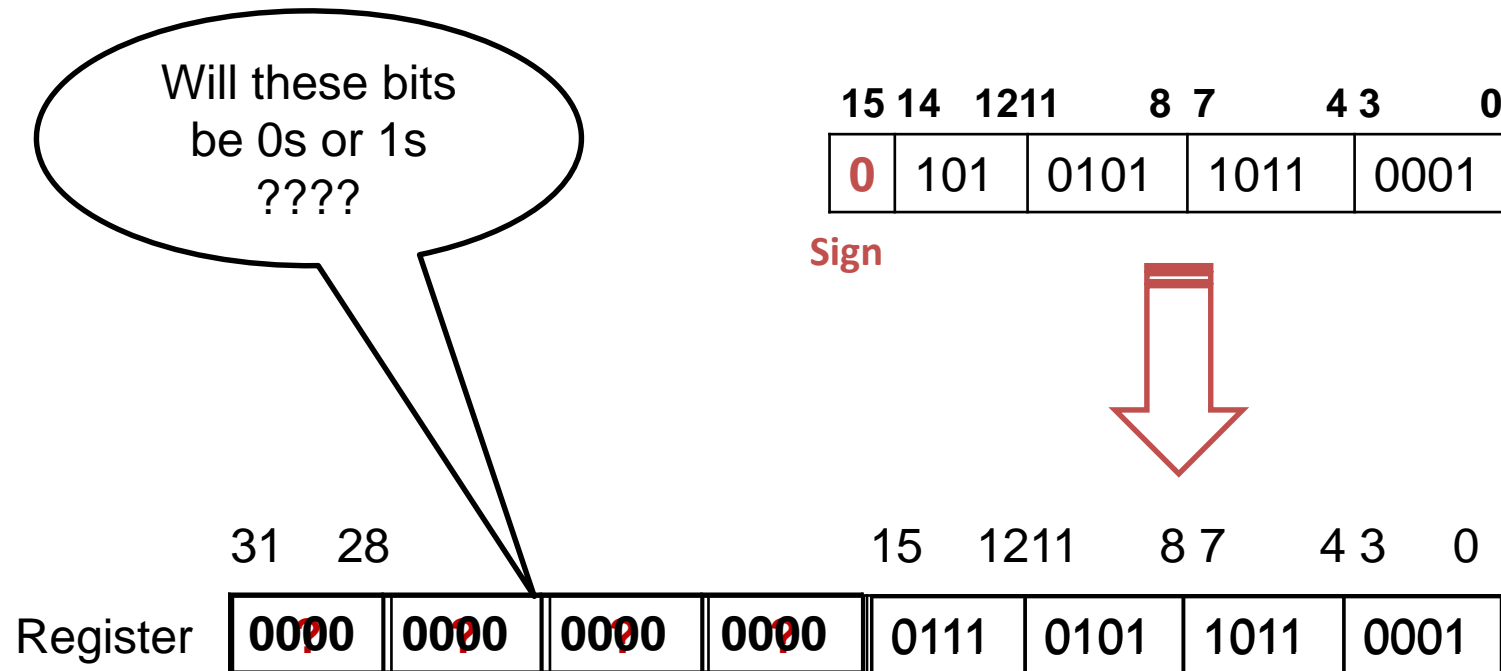
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Signed and Unsigned Numbers

Sign extension

How a Signed and Unsigned Load into a Register differs ?

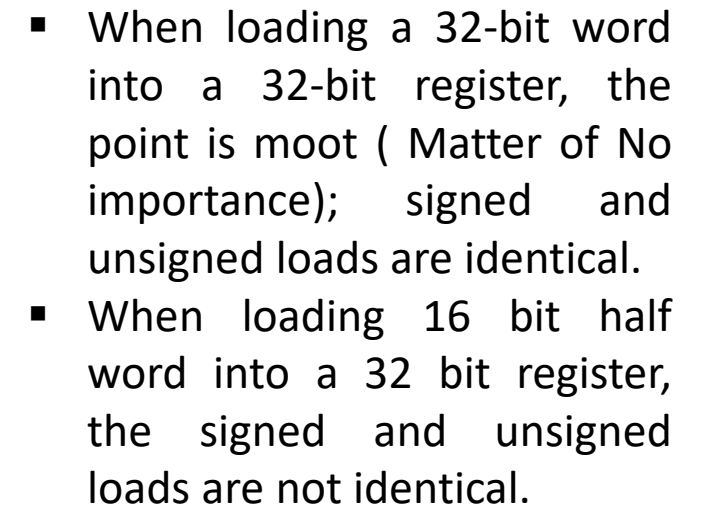
Let us consider loading 32 bit register with a **16 bit Signed value** ?



Signed and Unsigned Numbers



Let us consider loading 32 bit register with a **16 bit Signed value** ?



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Signed and Unsigned Numbers



Sign extension

- RISC-V does offer two flavors of byte loads:
 1. **load byte unsigned (lbu)** treats the byte as an unsigned number and thus zero-extends to fill the leftmost bits of the register,
 2. **load byte (lb)** works with signed integers.

Since C programs almost always use bytes to represent characters rather than consider bytes as very short signed integers, lbu is used practically exclusively for byte loads.

- The binary representation can be used as data and Address
Does Signed Number Representation makes any sense when it is used to address and data ????

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Signed and Unsigned Numbers



Working with two's complement numbers.

Negate a two's complement binary number ?

Ex: Negation of (-1) is (+1)

$X = (-1) = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111$ (0xFFFF FFFF)
 $(+1) = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001$ (0x0000 0001)

How can you find it quickly ????

Solution :

- a) $X' + 1 = -X$
- b) This shortcut is based on the observation that the **sum of a number and its inverted representation must be $(111 \dots 111)_2$** , which represents -1 . [$X + X' = -1$]

$X = (-1) = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111$ (0xFFFF FFFF)
 $X' = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000$ (0x0000 0000)

 $(-1) = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111$

Since $x + x' = -1$, therefore $x + x' + 1 = 0$ or $x' + 1 = -x$.



THANK YOU

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