Semester I

ME010101 Abstract Algebra

1.1 Introduction to Abstract Algebra

'Abstract' Consider the collection of of all students in a class. A person can't be repeated in this collection so you might think this collection is always a set. But in this collection, there may be two with name starting with an 'S'.

Thus, though the collection of students is a reality, set is an abstract idea which depends on the details we are interested in. And if a particular collection doesn't qualify to be a 'set', we can't use set-theoretic results on it. (which is, everything below)

Definitions 1.1.1. Set-theoretical foundation of Abstract Algebra

- A cartesian product, $A \times B = \{(a, b) : a \in A, b \in B\}$
- A relation on a set A is a subset of $A \times A$.
- A function f from A into B, $f:A \to B$ is a relation such that every element of A is related to some unique element of B (well defined).

Definitions 1.1.2. Functions:

- A function $\phi: A \to B$ is a **homomorphism** if for any two elements $x, y \in A$, $\phi(xy) = \phi(x)\phi(y)$
- A function $\phi: A \to B$ is an **isomorphism** if ϕ is a bijective, homomorphism. If two binary structures are isomorphic, then they have the same (algebraic) structure.
- A binary operation on a set A is a function $*: A \times A \rightarrow A$.

[&]quot;A binary operation on a set A gives an algebra on A."

Abstract Algebra It is the study of algebraic structures. We are interested in a few algebraic structures: 1. Group 2. Ring 3. Integral Domain 4. Field

Definitions 1.1.3 (Group). A set G together with a binary relation * satisfying the following three axioms -G1, G2, &G3 is a group.

- G1 Associativity For any three elements $a, b, c \in G$, a * (b * c) = (a * b) * c.
- G2 Identity element There exists a unique element $e \in G$ such that a * e = a = e * a for any element $a \in G$.
- G3 Inverse elements For any element $a \in G$ there exists a unique element a^{-1} such that $a * a^{-1} = e = a^{-1} * a$.

Definitions 1.1.4. Group terminologies:

- A group is abelian if the binary operation is commutative. a * b = b * a
- The order of a group G, * is the number of elements in G.

Definitions 1.1.5. Subgroup

- H is a **subgroup** of G if $H \subset G$, and $H, *_{|H}$ is a group.
- G is the improper subgroup of G and every other subgroup is proper.
- $\{e\}$ is the **trivial** subgroup of G and every other subgroup is **non-trivial**.
- The subgroup generated by $g \in G$ is the subgroup $\{g^n : n \in \mathbb{Z}\}$.
- ullet The **order** of an element g is order of the subgroup generated by g.
- An element $g \in G$ is a **generator** of G if g generates G.
- A group is **cyclic** if it has a generator.

*	e	\mathbf{a}	b	\mathbf{c}				
e	е	a	b	С	-	* _н е b	e	b
a	a	b	\mathbf{c}	e	\rightarrow	е	e	b
b	b	\mathbf{c}	e	\mathbf{a}		b	b	e
$^{\mathrm{c}}$	c	\mathbf{e}	\mathbf{a}	b			•	

Figure 1.1: Subgroup H obtained from group G

Remark (Results). Cyclic Groups:

- $\bullet \ \ Cyclic \ groups \ are \ abelian.$
- Subgroups of cyclic groups are cyclic.

1.1.1 Some Proof Techniques

Uniqueness Suppose there are two elements that qualify our conditions. We show (using the conditions) that they are the same, which is means it is unique.

For example, $3+a=\pi$. Suppose a=x,y. Then $3+x=3+y \implies x=y$, provided that the values of a comes from a set in which left cancelation law can be applied. Then a is unique.

Remember: We usually don't care to show what this unique element is. It may be also be the case that there is no such element.

Applications:

- Uniqueness of the identity element in a group
- Uniqueness of the inverse element of any element in a group

Existence Existence is usually proved by construction. Suppose we want to prove that there is a rational number between n and m, $(n \neq m)$. We come up with (n+m)/2. This number is a rational number and it comes between n and m.

• Existence of generator : Cyclic group test

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1.2 Direct Products and Finitely Generated Abelian Groups

Definitions 1.2.1 (Cartesian product of sets). Let S_1, S_2, \dots, S_n be a sets. Their cartesian product,

$$S_1 \times S_2 \times \dots \times S_n = \prod_{i=1}^n S_n = \{(a_1, a_2, \dots, a_n) : a_i \in S_i\}$$

For example, $\{A, B, C\} \times \{1, 2\} = \{(A, 1), (A, 2), (B, 1), (B, 2), (C, 1), (C, 2)\}.$

Question 1. How many elements in $\mathbb{Z}_3 \times \mathbb{Z}_{10} \times \mathbb{Z}_9$?

Theorem 1.2.2 (Direct product of Groups). Let G_1, G_2, \dots, G_n be groups. Then their cartesian product is a group with the binary operation *,

$$(a_1, a_2, \cdots, a_n) * (b_1, b_2, \cdots, b_n) = (a_1 * b_1, a_2 * b_2, \cdots, a_n * b_n)$$

where the binary operation in $a_i * b_i$ is the binary operation of the group G_i .

Proof. $\prod G_i$ is a group if it satisfies the group axioms.

G1 Associativity

$$(a_1, a_2, \dots, a_n)((b_1, b_2, \dots, b_n)(c_1, c_2, \dots, c_n))$$

$$=(a_1, a_2, \dots, a_n)(b_1c_1, b_2c_2, \dots, b_nc_n)$$

$$=(a_1(b_1c_1), a_2(b_2c_2), \dots, a_n(b_nc_n))$$

$$=((a_1b_1)c_1, (a_2b_2)c_2, \dots, (a_nb_n)c_n)$$

$$=(a_1b_1, a_2b_2, \dots, a_nb_n)(c_1, c_2, \dots, c_n)$$

$$=((a_1, a_2, \dots, a_n)(b_1, b_2, \dots, b_n))(c_1, c_2, \dots, c_n)$$

G2 Existence of a unique identity element in $\prod G_i$ Let e_i be the identity element in G_i . Then (e_1, e_2, \dots, e_n) is the identity element in $\prod G_i$.

$$(a_1, a_2, \dots, a_n)(e_1, e_2, \dots, e_n)$$

= $(a_1e_1, a_2e_2, \dots, a_ne_n)$
= (a_1, a_2, \dots, a_n)

G3 Existence of unique inverse element for each element in $\prod G_i$ Let (a_1, a_2, \dots, a_n) be in $\prod G_i$. Then it has the inverse element $(a_1^{-1}, a_2^{-1}, \dots, a_n^{-1})$ existing in $\prod G_i$.

$$(a_1, a_2, \cdots a_n)(a_1^{-1}, a_2^{-1}, \cdots, a_n^{-1})$$

$$= (a_1 a_1^{-1}, a_2 a_2^{-1}, \cdots, a_n a_n^{-1})$$

$$= (e_1, e_2, \cdots, e_n)$$

Remark. We usually write ab instead of a * b and relevent binary operations are used in different contexts. Student should be able to recongnise the difference from the context.

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\mathbb{Z}_n = \{0, 1, \cdots, (n-1)\} is a group with +_n. (addition modulo n)
For example, Consider (1,1) \in \mathbb{Z}_2 \times \mathbb{Z}_3. We have, (1,1) + (1,1) = (0,2)
since 1 +_2 1 = 0 and 1 +_3 1 = 2.
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Theorem 1.2.3. $\mathbb{Z}_m \times \mathbb{Z}_n$ is cyclic and is isomorphic to \mathbb{Z}_{mn} if and only if m and n are relatively prime.

Proof. Sufficient part:

Consider the cyclic subgroup¹ H generated by $(1,1) \in \mathbb{Z}_m \times \mathbb{Z}_n$. It is enough to prove that the order of this cyclic subgroup H is mn. The order of H is the smallest power of (1,1) that gives the identity (0,0). The first component gives 0 for multiples of m. And the second component gives 0 for multiples of n. Since m, n are relatively prime, mn is the smallest power of (1,1) that will give (0,0). Thus $\mathbb{Z}_m \times \mathbb{Z}_n = H$ and is cyclic.

Necessary part:

Suppose $\gcd(m,n)=d>1$. Then mn/d is the smallest integer divisible by both m and n. Consider $(r,s)\in\mathbb{Z}_m\times\mathbb{Z}_n$. r gives 0 in mn/d since it is a multiple of m. Similarly, s gives 0 in mn/d since it is a multiple of n. Thus, $\frac{mn}{d}(r,s)=(0,0)$. And the cyclic group generated by any element of $\mathbb{Z}_m\times\mathbb{Z}_n$ is a proper subgroup. Therefore $\mathbb{Z}_m\times\mathbb{Z}_n$ has no generators and it is not cyclic. \square

Corollary 1.2.3.1. $\prod \mathbb{Z}_{m_i}$ is cyclic and is isomorphic to $Z_{m_1m_2\cdots m_n}$ if and only if any two of the numbers m_i are relatively prime.

Question 2. Prove: For any non-negative integer n, there exists a cyclic group of order n, which is unique upto isomorphism.

Theorem 1.2.4. Let $(a_1, a_2, \dots, a_n) \in \prod G_i$. And a_i are of finite order r_i in G_i . Then the order of $\prod G_i$ is the least common multiple of r_i s.

Proof. Least common multiple of r_i s is the smallest positive integer d which is a multiple of all r_i s. For each i, the r_i th multiple of a_i gives 0 (identity). Thus, the order of the cyclic subgroup generated by (a_1, a_2, \dots, a_n) is the least common multiple of all the r_i s.

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Remark. Consider (3,6,12,16) \in \mathbb{Z}_4 \times \mathbb{Z}_{12} \times \mathbb{Z}_{20} \times \mathbb{Z}_{24}.
 Order of 3 \in \mathbb{Z}_4 is 4/\gcd(3,4) = 4 ie, < 3 >= \{3,2,1,0\}
 Order of 6 \in \mathbb{Z}_{12} is 12/\gcd(6,12) = 2 ie, < 6 >= \{6,0\}
 Order of 12 \in \mathbb{Z}_{20} is 20/\gcd(12,20) = 5 ie, < 12 >= \{12,4,16,8,0\}
 Order of 16 \in \mathbb{Z}_{24} is 24/\gcd(16,24) = 3 ie, < 16 >= \{16,8,0\}
 Order of (3,6,12,16) is lcm(4,2,5,3) = 2^2 3 5 = 60.
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Remark. Define $\overline{G_i} = \{(e_1, e_2, \dots, e_{i-1}, a_i, e_{i+1}, \dots, e_n) : a_i \in G_i\}$. Then $G_i \cong \overline{G_i}$. And $\prod G_i$ is the internal direct product of $\overline{G_i}$ s.

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For example, \mathbb{Z}_2 \times \mathbb{Z}_3 \cong (\mathbb{Z}_2 \times \{0\}) \otimes (\{0\} \times \mathbb{Z}_3)
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Question 3. Internal direct product form of $\mathbb{Z}_{12} \times \mathbb{Z}_{60} \times \mathbb{Z}_{24}$?

 $^{(1,1) \}in Z_m \times Z_n$, then the cyclic group generated by (1,1) has all its elements in $Z_m \times Z_n$. And therefore, it is a subgroup of $Z_m \times Z_n$

1.3 Fundamental Theorem

Definitions 1.3.1. A group G is **finitely generated** if G has a finite subset that generates G.

For example, $\{1, \sigma^2, \mu, \mu\sigma^2\}$ is a finitely generated subgroup of D_8 .

0	1	σ	σ^2	σ^3	μ	$\mu\sigma$	$\mu\sigma^2$	$\mu\sigma^3$
1	1	σ	σ^2	σ^3	μ	$\mu\sigma$	$\mu\sigma^2$	$\mu\sigma^3$
σ	σ	σ^2	σ^3	1	$\mu\sigma^3$	μ	$\mu\sigma$	$\mu\sigma^2$
σ^2	σ^2	σ^3	1	σ	$\mu\sigma^2$	$\mu\sigma^3$	μ	$\mu\sigma$
σ^3	σ^3	1	σ	σ^2	$\mu\sigma$	$\mu\sigma^2$	$\mu\sigma^3$	μ
μ	μ	$\mu\sigma$	$\mu\sigma^2$	$\mu\sigma^3$	1	σ	σ^2	σ^3
$\mu\sigma$	$\mu\sigma$	$\mu\sigma^2$	$\mu\sigma^3$	μ	σ^3	1	σ	σ^2
$\mu\sigma^2$	$\mu\sigma^2$	$\mu\sigma^3$	μ	$\mu\sigma$	σ^2	σ^3	1	σ
$\mu\sigma^3$	$\mu\sigma^3$	μ	$\mu\sigma$	$\mu\sigma^2$	σ	σ^2	σ^3	1

Figure 1.2: Dihedral Group D_8

Theorem 1.3.2 (fundamental theorem of finitely generated abelian groups). Every finitely generated abelian group G is isomorphic to a direct product of cyclic groups in the form

$$\mathbb{Z}_{p_1^{r_1}} \times \mathbb{Z}_{p_2^{r_2}} \times \cdots \times \mathbb{Z}_{p_n^{r_n}} \times \mathbb{Z} \times \mathbb{Z} \times \cdots \times \mathbb{Z}$$

where p_i are primes, not necessarily distict and r_i are positive integers. The direct product is unique, except for the possible rearrangement of the factors.

Proof. —proof is omitted—
$$\Box$$

For example, $G = \mathbb{Z}_{20} \times \mathbb{Z} \times \mathbb{Z}_{15} \times \mathbb{Z} \cong \mathbb{Z}_{2^2} \times \mathbb{Z}_3 \times \mathbb{Z}_5 \times \mathbb{Z} \times \mathbb{Z}$ In the above case, Betti number of G is 2 (number of \mathbb{Z} factors). For any finite abelian group, Betti number is 0.

Remark (finite abelian groups). Every finite group is finitely generated. And thus we can enumerate finite abelian group of any order.

Remark. There are precisely 6 different abelian groups of order $360 = 2^3 3^2 5$.

- 1. $\mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_3 \times \mathbb{Z}_3 \times \mathbb{Z}_5$
- 2. $\mathbb{Z}_2 \times \mathbb{Z}_{2^2} \times \mathbb{Z}_3 \times \mathbb{Z}_3 \times \mathbb{Z}_5$
- 3. $\mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_{3^2} \times \mathbb{Z}_5$
- 4. $\mathbb{Z}_{2^3} \times \mathbb{Z}_3 \times \mathbb{Z}_3 \times \mathbb{Z}_5$
- 5. $\mathbb{Z}_2 \times \mathbb{Z}_{2^2} \times \mathbb{Z}_{3^2} \times \mathbb{Z}_5$
- 6. $\mathbb{Z}_{2^3} \times \mathbb{Z}_{3^2} \times \mathbb{Z}_5$

Question 4. Group of order 360 with at least an element of order 8

Definitions 1.3.3. A group G is **decomposible** if it is isomorphic to a direct product of two proper, non-trivial subgroups. Otherwise G is **indecomposible**.

For example, $\mathbb{Z}_{2^3} \times \mathbb{Z}_3 \times \mathbb{Z}_3 \times \mathbb{Z}_5 \cong \mathbb{Z}_{24} \times \mathbb{Z}_{15}$ is decomposible.

Remark. $G = \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_3 \times \mathbb{Z}_3 \times \mathbb{Z}_5 \cong \mathbb{Z}_2 \times \mathbb{Z}_6 \times \mathbb{Z}_{30}$ is also decomposible. Since $G \cong (\mathbb{Z}_2 \times \mathbb{Z}_6) \times \mathbb{Z}_{30}$.

Theorem 1.3.4. The finite indecomposible abelian groups ar exactly the cyclic groups with order a power of a prime.

Proof. Necessary part:

Let G be a finite, indecomposible abelian group. By fundamental theorem of finitely generated abelian groups, G is isomorphic to a direct product of cyclic groups of prime power order.

$$G \cong \mathbb{Z}_{p_1^{r_1}} \times \mathbb{Z}_{p_2^{r_2}} \times \dots \times \mathbb{Z}_{p_n^{r_n}}$$

Thus for G to be indecomposible the direct product should be a cyclic group of prime power order. $G \cong \mathbb{Z}_{p_1^{r_1}}$

Sufficient part:

Let p be a prime and r a non-negative integer. Cyclic group of order p^r is isomorphic to \mathbb{Z}_{p^r} . Since every cyclic groups are abelian, \mathbb{Z}_{p^r} is an abelian group of finite order p^r . It is enough to prove that \mathbb{Z}_{p^r} is indecomposible.

A proper, non-trivial subgroup of \mathbb{Z}_{p^r} is of the form \mathbb{Z}_{p^i} where 0 < i < r. Suppose \mathbb{Z}_{p^r} is decomposible. Then $\exists i, j \in \mathbb{Z}^+$ such that $\mathbb{Z}_{p^r} \cong \mathbb{Z}_{p^i} \times \mathbb{Z}_{p^j}$ and i+j=r. Clearly, p^i and p^j are not relatively prime, thus $\mathbb{Z}_{p^r} \not\cong \mathbb{Z}_{p^i} \times \mathbb{Z}_{p^j}$. Therefore, cyclic groups of order prime power are indecomposible.

Theorem 1.3.5. If m divides the order of a finite abelian group G, then G has a subgroup of order m.

Proof. Let
$$G = \mathbb{Z}_{p_1^{r_1}} \times \mathbb{Z}_{p_2^{r_2}} \times \cdots \times \mathbb{Z}_{p_n^{r_n}}$$
. Then $|G| = p_1^{r_1} p_2^{r_2} \cdots p_n^{r_n}$. Suppose m divides $|G|$, then $m = p_1^{s_1} p_2^{s_2} \cdots p_n^{s_n}$ where $0 \le s_i \le r_i$. Define $H = \mathbb{Z}_{p_1^{s_1}} \times \mathbb{Z}_{p_2^{s_2}} \times \cdots \times \mathbb{Z}_{p_n^{s_n}}$. Then H is subgroup of order m .

Theorem 1.3.6. If m is square-free integer, then every abelian group of order m is cyclic.

Proof. Let m be a square-free integer and G be an abelian group of order m. By fundamental theorem of finitely generated abelian groups

$$G \cong \mathbb{Z}_{p_1^{r_1}} \times \mathbb{Z}_{p_2^{r_2}} \times \dots \times \mathbb{Z}_{p_n^{r_n}}$$

We have, m is square-free. Thus $r_i = 1$ and p_i are distinct. Therefore, $G \cong \mathbb{Z}_{p_1p_2\cdots p_n}$ is a cyclic group of order m.

1.4 Exercises §11

Remark. Prove the following results

- If group G is abelian, then its elements of finite order forms a subgroup.
 (hint: a, b ∈ G has finite order, then ab has finite order a ∈ G has finite order, then a⁻¹ has finite order)
- direct product of abelian groups is abelian.
- Subset of any abelian group G with Identity element and all elements of order 2 is a subgroup of G.
- Subset of any abelian group G with Identity element and all elements of order 3 is a subgroup of G.
- Subset of any abelian group G with Identity element and all elements of order 4 is not a subgroup of G.

Remark. Let G be an abelian group. The subset H of G with identity element and all elements of order n is subgroup of G if and only if n is a prime.

Proof. Suppose $a \in H$. Then every power of a has order n. Suppose n is not prime. Then d divides n and a^d has order n/d.

Definitions 1.4.1. The torsion group of an abelian group G is the subgroup of G containing only those elements of finite order. An abelian group is torsion free if identity element is the only element of finite order. G is Torsion free if Torsion group of G is trivial, $\{e\}$.

Definitions 1.4.2. The integers m_1, m_2, \dots, m_n are torsion coefficients of G such that $G \cong \mathbb{Z}_{m_1} \times \mathbb{Z}_{m_2} \times \dots \times \mathbb{Z}_{m_n}$ where m_i divides m_{i+1} .

For example, $\mathbb{Z}_6 \times \mathbb{Z}_{12} \times \mathbb{Z}_{20}$ has torsion coefficients 2, 12, 60

Remark (Algorithm to find torsion coefficients of a group). Suppose G has a direct product form.

Step 1 Find power of each prime in the direct product form

Step 2 List power of each prime

Step 3 Append 1s on left to make all lists to equal length

Step 4 Product of ith number on each list gives m_i

For example, $G \cong \mathbb{Z}_6 \times \mathbb{Z}_{12} \times \mathbb{Z}_{20}$

Step 1 $\mathbb{Z}_6 \times \mathbb{Z}_{12} \times \mathbb{Z}_{20} \cong \mathbb{Z}_2 \times \mathbb{Z}_3 \times \mathbb{Z}_{2^2} \times \mathbb{Z}_3 \times \mathbb{Z}_{2^2} \times \mathbb{Z}_5$

Step 2 (2,4,4), (3,3), (5)

Step 3 (2,4,4), (1,3,3), (1,1,5)

Step 4 (2,12,60)

Remark. Let $G = H \times K$. $g \in G \implies g = (h, k) \in H \times K$. \Longrightarrow H is a subset of G, $H \times \{e\}$.

Thus
$$h \in H \subset G \implies h = (h, e)$$
. Similarly $k = (e, k)$
 $\implies hk = (h, e)(e, k) = (h, k) = (e, k)(h, e) = kh$

1.5 Cosets and Homomorphism

Definitions 1.5.1. Kernel of a function $\phi: G \to G'$ is the inverse image of the identity element in G'.

Definitions 1.5.2. Cosets and Normal Subgroup,

- A left coset gH is the subset $\{gh \in G : h \in H\}$ where $g \in G$ and H is a subgroup of G.
- A right coset Hg is the subset $\{hg \in G : h \in H\}$ where $g \in G$ and H is a subgroup of G.
- A subgroup H of group G is **normal** if gH = Hg, $\forall g \in G$.
- All subgroups of abelian groups are normal.

Remark. For example, $\{1, \sigma^2, \mu, \mu\sigma^2\}$ is a normal subgroup of D_8 . $\sigma H = \{\sigma, \sigma^3, \mu\sigma^3, \mu\sigma\} = H\sigma$. However $\sigma \mu \neq \mu\sigma$

Remark (Lagrange's Theorem). Let G be a finite group. If H is a subgroup of G, then order of H divides order of G.

- $aH \cap bH \neq \phi \implies aH = bH$.
- $\forall g \in G, g \in gH$.
- $\forall g \in G, |gH| = |H|.$

Remark (group homomorphism). Group homomorphism $\phi: G \to G'$ preserves identity, inverses and subgroups. And kernel of group homomorphism is a normal subgroup of G.

Remark. For example, $\phi: D_8 \to \mathbb{Z}_2$ defined by $\phi(\sigma) = 1$ and $\phi(\mu) = 0$ is a group homomorphism with $\ker(\phi) = \{1, \sigma^2, \mu, \mu\sigma^2\}$.

1.6 Factor Groups

Remark (factor group). Let H be a normal subgroup of a group G. Then the factor group of G over H, G/H is the group of cosets of H in G.

Remark. For example, $H = \{1, \sigma^2, \mu, \mu\sigma^2\}$ is normal subgroup of D_8 . And the factor group $D_8/H = \{1H, \sigma H\}$.

Remark. Notations:

• $\ker(\phi)$ is the kernel of a group homomorphism $\phi: G \to G'$.

$$\ker(\phi) = \{g \in G : \phi(g) = e'\}$$

• G/H is the factor group of G over H.

$$G/H = \{gH : g \in G, H \text{ is a normal subgroup of } G\}$$

• $\phi[G]$ is the range group of $\phi: G \to G'$.

$$\phi[G] = \{\phi(g) : g \in G, \ \phi \text{ is a group homomorphism}\}\$$

 \bullet (G:H) is the index of H in G ie, the number of left cosets of H in G

Theorem 1.6.1. Let $\phi: G \to G'$ be a group homomorphism with kernel H. Then cosets of H form a factor group, G/H where (aH)(bH) = (abH) Also, $\mu: G/H \to \phi[G]$ defined by $\mu(aH) = \phi(a)$ is an isomorphism.

Proof. Let $\phi: G \to G'$ be a group homomorphism with $\ker(\phi) = H$. We have, $\phi^{-1}(\phi(a)) = \{g \in G: \phi(g) = \phi(a)\}$.

Let $x \in aH \implies x = ah$ for some $h \in H$. $\phi(x) = \phi(ah) = \phi(a)\phi(h) = \phi(a)$ Thus, $x \in \phi^{-1}(\phi(a))$ and $aH \subset \phi^{-1}(\phi(a))$.

Let $x \in \phi^{-1}(\phi(a)) \Longrightarrow \phi(x) = \phi(a)$. $\phi(x) = \phi(a) \Longrightarrow \phi(a)^{-1}\phi(x) = e' \Longrightarrow \phi(a^{-1}x) = e' \Longrightarrow a^{-1}x \in \ker(\phi)$ There exists $h \in H$ such that $a^{-1}x = h \Longrightarrow x = ah$ for some $h \in H$ Thus, $x \in aH$ and $\phi^{-1}(\phi(a)) \subset aH$. Therefore, $\phi^{-1}(\phi(a)) = aH$.

Similarly, $\phi^{-1}(\phi(a)) = Ha$. Thus aH = Ha and H is a normal subgroup of G.

 $\mu: G/H \to \phi[G]$ is a one-one correspondence. ie, $aH \stackrel{\iota}{\longleftrightarrow} \phi(a)$ injective: $\mu(aH) = \mu(bH) \implies \phi(a) = \phi(b) \implies b \in aH \implies bH = aH$ surjective: If $\phi(a) \in \phi[G]$, then there exists aH such that $\mu(aH) = \phi(a)$

We have, $\mu(aH)=\phi(a),\ \mu(bH)=\phi(b),\ \mu((ab)H)=\phi(ab).$ $\mu((aH)(bH))=\mu((ab)H)=\phi(ab)=\phi(a)\phi(b)=\mu(aH)\mu(bH)$ Thus μ is a homomorphism.

Theorem 1.6.2. Let H be a subgroup of a group G. Then left coset multiplication is well-defined by (aH)(bH) = (ab)H if and only if H is a normal subgroup of G.

Proof. Necessary part:

Suppose (aH)(bH) = (ab)H is well-defined.

Let $a \in G$. It is enough to prove that aH = Ha.

Let $x \in aH$. Then $(xH)(a^{-1}H) = (xa^{-1})H$.

Also $(aH)(a^{-1}H) = eH = H$.

We have, coset multiplication is well-defined.

Thus $xa^{-1} = h \in H \implies x = ha \in Ha$.

Then, $aH \subset Ha$. Similarly, $Ha \subset aH$ and aH = Ha.

Therefore, H is a normal subgroup of G.

Sufficient part:

Suppose H is a normal subgroup of G, and let $x \in aH$ and $y \in bH$.

 $x \in aH \implies x = ah_1 \text{ for some } h_1 \in H$

 $y \in bH \implies y = bh_2 \text{ for some } h_2 \in H.$

Therefore $xy = (ah_1)(bh_2) = (ah_1)(h_2b) = a(h_1(h_2b)) = a((h_1h_2)b)$

Since H is a group, $h_1h_2 = h_3 \in H$

Thus, $xy = a(h_3b) = a(bh_3) = (ab)h_3 \in (ab)H$ for all $x \in aH$ and $y \in bH$ Thus (aH)(bH) = (ab)H.

Corollary 1.6.2.1. Let H be a normal subgroup of G. Then the cosets of H form a group G/H under the binary operation (aH)(bH) = (ab)H.

Proof. G1 Associativity (aH)[(bH)(cH)] = (aH)[(bc)H] = [a(bc)]H = [(ab)c]H = [(ab)H](cH) = [(aH)(bH)](cH)

G2 Existence of identity, eH (aH)(eH) = (ae)H = aH and (eH)(aH) = (ea)H = aH.

G3 Existence of inverse $(a^{-1}H)$ $(aH)(a^{-1}H)=(aa^{-1})H=eH$ and $(a^{-1}H)(aH)=(a^{-1}a)H=eH$.

Remark. $n\mathbb{Z}$ is a normal subgroup of \mathbb{Z} . And $\mathbb{Z}/n\mathbb{Z} \cong \mathbb{Z}_n$.

Remark. Let $c \in \mathbb{R}^*$, then the cyclic group generated by c is a normal subgroup of \mathbb{R} and $\mathbb{R}/\langle c \rangle \cong \mathbb{R}_c$.

Question 5. Let c = 0.31. Find the coset x + < 0.31 >containing 2.

1.6.1 Fundamental Homomorphism Theorem

Theorem 1.6.3. Let H be a normal subgroup of G. Then $\gamma: G \to G/H$ is defined by $\gamma(x) = xH$ is a homomorphism with kernel H.

Proof. ??pending revision $\gamma(x)\gamma(y) = (xH)(yH)$ Let $h_1, h_2 \in H$, $(xh_1)(yh_2) = xyh_3h_2 = xyh_4$ for some $h_3, h_4 \in H$. Therefore (xH)(yH) = (xy)H.

 $\gamma(xy) = (xy)H = \gamma(x)\gamma(y).$

 \Box

Theorem 1.6.4 (Fundamental Homomorphism). Let $\phi: G \to G'$ be a group homomorphism with kernel H. Then $\phi[G]$ is a group, and $\mu: G/H \to \phi[G]$ given by $\mu(gH) = \phi(g)$ is an isomorphism. If $\gamma: G \to G/H$ is the homomorphism given by $\gamma(g) = gH$, then $\phi(g) = \mu\gamma(g)$ for each $g \in G$.

1.6.2 Inner Automorphism

Theorem 1.6.5. Let H be a subgroup of G, then the following statements are equivalent:

- 1. $ghg^{-1} \in H, \forall g \in G, h \in H$
- 2. $gHg^{-1} = H, \forall g \in G$
- 3. $gH = Hg, \forall g \in G$

Proof.

Definitions 1.6.6. Automorphisms:

- An automorphism is an isomorphism $\phi: G \to G$
- The inner automorphism of G by $g \in G$ is the isomorphism $i_g : G \to G$ defined by $i_g(x) = gxg^{-1}$ for all $x \in G$
- The conjugate of x by g is the element $gxg^{-1} \in G$
- The conjugate subgroup of subgroup H, $i_a[H] = \{ghg^{-1} : h \in H\}$

Remark. For example, $i_{\sigma}: D_8 \to D_8$ defined by $i_{\sigma}(x) = \sigma x \sigma^{-1}$ is an inner automorphism. We have, $H = \{1, \mu\}$ is a subgroup of D_8 . The conjugate subgroup $i_{\sigma}[H] = \{1, \mu \sigma^2\}$.

Remark. Normal subgroups are invariant under any inner automorphism.

ME010102 Linear Algebra

- 2.1 Vector Spees
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7.1 Module I

Q: Why these axioms are called separation axioms? T_0, T_1, T_2 axioms separates points from points. $T_3, T_{3\frac{1}{2}}$ axioms separates points from closed subsets. T_4 axiom separates closed subsets from closed subsets.

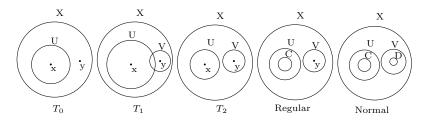


Figure 7.1: Separation Axioms

7.1.1 Compactness and Separation Axioms

[?, chapter 7 §2.1-10]

Proposition 7.1.1. Let X be a T_2 space, $x \in X$ and F is a compact subset of X not containing x. Then there exist open subsets U, V such that $x \in U$, $F \subset V$, $U \cap V = \phi$.

 T_2 spaces separates points from compact sets.

Proof. Let x be a points in X and F be a compact subset not containing x. X is $T_2 \implies \forall y \in F, \exists U_y, V_y \in \mathcal{T}, x \in U_y, y \in V_y, U - y \cap V_y = \phi.(\dagger^1)$

 $^{{}^{1}\}mathbf{Q}$: Do we need compactness for this result?

A: Given $x \in X$. For each point $y \in F$, we have two open sets namely U_y and V_y that separates these two points x and y in X. But, the intersection $\bigcap_{y \in F} \{U_y\}$ is not necessarily an

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 \mathcal{C} = \{V_y : y \in F\} \text{ is an open cover of compact subset } F. \\ \Longrightarrow \text{ there exists a finite subcover, } \mathcal{C}' = \{V_{y_i} : y_i \in F, i = 1, 2, \cdots, n\}. \\ \text{Define } U = \cap_{i=1}^n \{U_{y_i}\} \text{ and } V = \cup_{i=1}^n \{V_{y_i}\}. \\ \Longrightarrow x \in U, \ F \subset V, \ U \cap V = \phi.
```

Corollary 7.1.1.1. A compact subset in a T_2 space is closed.

Proof. Let $x \in X - F$.

```
By proposition, x \in U, F \subset V, U \cap V = \phi \implies x \in U \subset X - V \subset X - F
 X - F is a nbd of each of its points. Thus X - F is open. (\star^2)
```

Corollary 7.1.1.2. Every map from a compact space into a T_2 space is closed. And its range is a quotient space of the domain.

Proof. Suppose X is compact, Y is T_2 and $f: X \to Y$ is continuous.

Let C be a closed subset of X.

 \implies C is compact, since compact is weakly hereditary. (\star^3)

 $\implies f(C)$ is compact, since compactness is preserved by conti. functions. (\star^4) By corollary, f(C) is compact, Y is $T_2 \implies f(C)$ is a closed subset of Y.

Thus $f: X \to Y$ is closed. (\dagger^5)

 $f: X \to Y$ is closed $\implies f: X \to f(X)$ is a quotient map, since every closed, surjective map is a quotient map..

Corollary 7.1.1.3. A continuous bijection from a compact space onto a T_2 space is a homeomorphism.

Proof. Let G be an open subset of X.

Then X - G is closed.

By corollary, f is closed, since f: compact $\to T_2$.

f is closed $\implies f(X-G)$ is closed.

f(X - G) = f(X) - f(G) since f in injective.

 $\implies f(X-G) = Y - f(G)$ since f surjective.

 $\implies f(G)$ is open.

Thus f is an open map.

Every continuous bijective, open map is a homeomorphism.

Corollary 7.1.1.4. Every continuous, one-to-one function from a compact space into a T_2 space is an embedding.

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Proof. Let $f: X \to Y$ be continuous and injective.

 $f: X \to f(X)$ is surjective.

 $\implies f: X \to f(X)$ is continuous, surjective.

By corollary, $f: X \to f(X)$ is a homeomorphism.

Thus f is an embedding of X onto $f(X) \subset Y$.

open subset of X, since F is not necessarily a finite subset of X and an **arbitary** intersection of open subsets is not necessarily an open subset. Therefore, we have restrict this family into a finite family. We use compactness for this part.

²Neighbourhood characterisation of open subsets : Let $G \subset X$ be a nbd of each of its points. Then G is an open subset of X.

 $^{^3}$ Compactness is weakly hereditary : Suppose X is compact. Then every closed subset of X is compact.

⁴Compactness is preserved by continuous functions: Suppose G be a compact subset of a topological space X. And function $f: X \to Y$ be a continuous function. Then f(G) is a compact subset of the topological space Y.[?, 6.1.8]

⁵Closed function is a function which maps closed subsets into closed subsets.

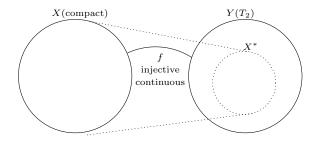


Figure 7.2: Embedding compact space into hausdorff space

Theorem 7.1.2. Every compact T_2 space is a T_3 space.

Proof. $T_2 \implies T_1$

It is enough to prove that compact, T_2 space is regular.

Let x be a point in X and C be a closed subset not containing x.

Then C is compact, compactness is weakly hereditary.

By proposition, T_2 space separates points from compact subsets.

Thus $T_2 + \text{compact} \implies T_3 \ (\dagger^6)$

Proposition 7.1.3. Let X be a regular space, C a closed subset of X and F a compact subset of X, such that $C \cap F = \phi$. Then there exist open subsets U, V such that $C \subset U$, $F \subset V$ and $U \cap V = \phi$.

Regular spaces separates closed subsets from compact subsets.

Proof. proof technique is same as T_2 separates points from compact subsets. Let C be a closed subset, F be a compact subset of X and $C \cap F = \phi$. X is regular $\Longrightarrow \forall y \in F, \ \exists U_y, V_y \in \mathcal{T}, \ C \subset U_y, \ y \in V_y, \ U_y \cap V_y = \phi$ Define $\mathcal{C} = \{V_y : y \in F\}$ is cover of compact subset F. $\Longrightarrow \exists \mathcal{C}'$ such that $\mathcal{C}' = \{V_{y_i} : i = 1, 2, \cdots, n\}$ is a finite subcover. Define $U = \cap_{i=1}^n \{U_{y_i}\}$ and $V = \cup_{i=1}^n \{V_{y_i}\}$. $\Longrightarrow C \subset U, \ F \subset V, \ U \cap V = \phi$.

Theorem 7.1.4. Every regular, Lindeloff space is normal.

Proof. Let C, D be two disjoint, closed subsets of a regular, lindeloff space X. X is regular $\Longrightarrow \forall x \in C, \ \exists U_x, V_x \ \text{such that} \ x \in U_x, \ D \subset V_x, \ U_x \cap V_x = \phi$ X is regular $\Longrightarrow \forall y \in D, \ \exists U_y, V_y \ \text{such that} \ C \subset U_y, \ y \in U_y, \ U_y \cap U_y = \phi$ Then $\{U_x : x \in C\}$ and $\{V_y : y \in D\}$ are open covers of C and D respectively. (\dagger^7) Let $\{U_n : n = 1, 2, \cdots\}$ and $\{V_n : n = 1, 2, \cdots\}$ be their countable subcovers. Define $G_n = U_n - \bigcup_{i=1}^n \overline{V_i} \ \text{and} \ H_n = V_n - \bigcup_{i=1}^n \overline{U_i}.$ Define $G = \bigcup_{i=1}^\infty G_n \ \text{and} \ H = \bigcup_{i=1}^\infty H_n.$ Claim: $C \subset G$ (similary $D \subset H$) $C \subset G$ (similary $C \subset G$) for some $C \subset G$ 0 some $C \subset G$ 1 is a cover of $C \subset G$ 2. Then $C \subset G$ 3 is a cover of $C \subset G$ 4 is a cover of $C \subset G$ 5. Then $C \subset G$ 6 is a cover of $C \subset G$ 6 claim: $C \subset G \subset G$ 6 is a cover of $C \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6. Then $C \subset G \subset G \subset G$ 6 is a cover of $C \subset G \subset G$ 6 is a cover of $C \subset G \subset G \subset G$ 6 is a cover of $C \subset G \subset G \subset$

 $^{^6}T_3 \not\Longrightarrow$ compact.

⁷Here, U_x , V_y are unrelated.

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Corollary 7.1.4.1. Every regular, second countable space is normal.

Proof. Every second countable space is lindeloff.
By theorem, every regular, lindeloff space is normal.

Corollary 7.1.4.2. Every compact, T_2 space is T_4 .

Proof. By proposition, every compact, T_2 space is regular. Every compact space is lindeloff since finite subcovers are countable. By theorem, every regular, lindeloff space is normal.

7.1.2 The Urysohn Characterisation of Normality

X is normal \iff there exist Urysohn functions \mathbf{Q} : Significance of Urysohn's Lemma?

- 1. There is no analog of Urysohn's Lemma for Regular spaces.
- 2. Constructs a nice real-valued function even on non-metrisable spaces.

Q: We know that, the completely regular space separates points from closed subsets using a real-valued function. Which space separates closed subsets from closed subsets using a real-valued function?

A: Normal space(by Urysohn's Lemma). Not completely normal(\star^8)

Proposition 7.1.5. Let A B be subsets of a space X and suppose there exists a continuous function $f: X \to [0,1]$, such that f(x) = 0, $\forall x \in A$ and f(x) = 1, $\forall x \in B$. Then there exists disjoint open subsets U, V such that $A \subset U$ and $B \subset V$.

Proof. Let
$$f$$
 be a continuous function, $f(x) = 0$, $\forall x \in A$ and $f(x) = 1$, $\forall x \in B$. Define $G = [0, \frac{1}{2})$, $H = (\frac{1}{2}, 1]$ and $U = f^{-1}(G)$, $V = f^{-1}(H)$. $\implies A \subset U$, $B \subset V$ and $U \cap V = \phi$ since $G \cap H = \phi$.

Corollary 7.1.5.1 (Urysohn's Lemma : Sufficient Condition). If X has the property that for any disjoint closed subsets A, B of X, there exists a continuous function $f: X \to [0,1]$ such that f(x) = 0, $\forall x \in A$ and f(x) = 1, $\forall x \in B$, then X is normal.

There exists a Urysohn function $f \implies X$ is normal

Proof. Let A, B be two closed subsets of X. By proposition, X normal.

Theorem 7.1.6 (Urysohn's Lemma). A topological space X is normal iff it has the property that for every mutually disjoint, closed subsets A, B of X, there exists a continuous function $f: X \to [0,1]$ such that f(x) = 0 for all $x \in A$ and f(x) = 1 for all $x \in B$

⁸Completely Normal Space separates any two subsets. [?, chapter 7 Exer. 1.11] That is these subsets need not be closed. Thus every completely normal space is normal, but not the other way. Thus T_1 + Completely Normal = $T_5 \supset T_4$.

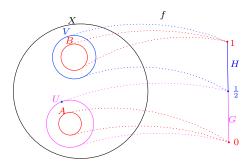


Figure 7.3: Urysohn's Lemma: Sufficient Condition

Proof. Urysohn's Lemma: necessary condition

X is normal \implies there exist Urysohn functions

Let A, B be two closed subsets in a normal space X.

Enumerate rationals in the unit interval.

 $\mathbb{Q} \cap [0,1] = \{0,1,\cdots\} = \{q_0,q_1,q_2,\cdots\}.$

Define $F_1 = F_{q_1} = X - B$. $A \subset X - B \implies \exists H \text{ such that } A \subset H \subset \overline{H} \subset X - B.(\star^9)$

Define $F_{q_0} = F_0 = H$.

 $\begin{array}{l} \overline{F_{q_0}} \subset F_{q_1} \Longrightarrow \text{ true for } n=1. \\ \text{Suppose}: \overline{F_{q_i}} \subset F_{q_j}, \ \forall j>i \text{ is true for } j=1,2,\cdots,n-1. \\ \text{Define } q_i = \sup\{q_k: q_k < q_n, k < n\} \text{ and } q_j = \inf\{q_k: q_k > q_n, k < n\} \ (\dagger^{10}) \\ \Longrightarrow q_i < q_n < q_j \text{ and } \overline{F_{q_i}} \subset F_{q_j}. \\ \overline{F_{q_i}} \subset F_{q_j} \Longrightarrow \exists H \text{ such that } \overline{F_{q_i}} \subset H \subset \overline{H} \subset F_{q_j}. \text{ Define } F_{q_n} = H. \\ \text{Therefore, } \overline{F_{q_j}} \subset F_{q_i}, \ \forall j>i. \end{array}$

Now, $\{F_t : t \in \mathbb{Q}\}$ has the properties required in lemma 2.

By lemma 2, there exists a function f. This is a Urysohn function. (\dagger^{11})

Lemma 7.1.7. Let $f: X \to [0,1]$ be continuous. For each $t \in \mathbb{R}$ let $F_t\{x \in X : t \in \mathbb{R} \}$ f(x) < t. Then the indexed family $\{F_t : t \in \mathbb{R}\}$ has the following properties

- 1. F_t is an open subset of X for each $t \in \mathbb{R}$
- 2. $F_t = \phi \text{ for } t < 0$
- 3. $F_t = X \text{ for } t > 1$
- 4. For any $s, t \in \mathbb{R}, s < t \implies \overline{F_s} \subset F_t$.

Moreover, for each $x \in X$, $f(x) = \inf\{t \in \mathbb{Q} : x \in F_t\}$.

Proof. Not needed.

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⁹Equivalent condition for normality: Let A be closed subset and G be an open subset containing A. Then there exists open subset H such that $A \subset H \subset \overline{H} \subset G$. [?, 7.1.16(3)]

¹⁰Let $\mathbb{Q} \cap [0,1] = \{0,1,0.3,0.7,0.8,\mathbf{0.5},0.6,\cdots\}$. Consider n=5. Then $q_n=q_5=0.5$ and in the respective induction step, $q_i = \sup\{0, 0.3\} = 0.3$ and $q_j = \inf\{1, 0.7, 0.8\} = 0.7$

¹¹You will have to replace "Urysohn function" with "There exists a continuous real-valued function, $f: X \to [0,1]$ such that $f(x) = 0, \ \forall x \in A, \ f(x) = 1, \ \forall x \in B$ ".

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Lemma 7.1.8. Let X be a topological space and suppose $\{F_t : t \in \mathbb{Q}\}$ is a family of sets in X such that

```
1. F_t is open in X for each t \in \mathbb{Q}
```

```
2. F_t = \phi for t \in \mathbb{Q}, t < 0
```

3.
$$F_t = X$$
 for $t \in \mathbb{O}$, $t > 1$

4.
$$\overline{F_s} \subset F_t$$
 for $s, t \in \mathbb{Q}, s < t$

For $x \in X$, let $f(x) = \inf\{t \in \mathbb{Q} : x \in F_t\}$. Then f is a continuous real-valued function on X and it takes values in the unit interval [0,1].

```
Proof. Function f is well-defined and Range(f) = [0, 1].
Define H = \{x \in X : f(x) < s\} = f^{-1}(-\infty, s) and
K = \{x \in X : f(x) > s\} = f^{-1}(s, \infty).
\mathcal{S} = \{(s, \infty), (-\infty, s) : s \in \mathbb{R}\} is a subbase for \mathbb{R} with usual topology.
Claim: H = \bigcup \{F_t : t \in \mathbb{Q}, t < s\}
 x \in H \implies f(x) < s \implies \inf G_x < s \implies \exists q \in \mathbb{Q} \ (q < s), \ x \in F_q
\implies x \in \bigcup \{F_t : t \in \mathbb{Q}, t < s\} \implies H \subset \bigcup \{F_t : t \in \mathbb{Q}, t < s\}.
x \in \bigcup \{F_t : t \in \mathbb{Q}, \ t < s\} \implies \exists t \in \mathbb{Q} \ (t < s), \ x \in F_t \implies \inf G_x < s\}
\implies f(x) < s \implies x \in H \implies \bigcup \{F_t : t \in \mathbb{Q}, \ t < s\} \subset H.
Claim: X - K = \bigcap \{ \overline{F_t} : t \in \mathbb{Q}, t > s \}.
 x \in X - K \implies x \notin K \implies f(x) \le s \implies \inf G_x \le s
Let t \in \mathbb{Q} (s < t) \Longrightarrow \exists q \in G_x \ (s < q < t) \Longrightarrow x \in \overline{F_q} \subset F_t \subset \overline{F_t}.

\forall t \in \mathbb{Q} \ (t > s), \ x \in \overline{F_t} \Longrightarrow x \in \cap \{\overline{F_t} : t \in \mathbb{Q}, \ t > s\}.
Thus, X - K \subset \cap \{\overline{F_t} : t \in \mathbb{Q}, \ t > s\}.
x \in \cap \{\overline{F_t} : t \in \mathbb{Q}, \ t > s\}
Suppose(1) x \in K \implies s < f(x) \implies \exists q, t \in \mathbb{Q} \ (s < q < t), \ t < f(x)
Suppose(2) x \in \overline{F_q} \implies x \in \overline{F_q} \subset F_t \implies \inf G_x < t \implies f(x) < t
is a contradiction (2). Thus x \notin \overline{F_q}.
x \notin \overline{F_q} \implies x \notin \cap \{\overline{F_t} : t \in \mathbb{Q}, \ t > s\} is a contradition (1). Thus x \notin K
Thus, \cap \{\overline{F_t} : t \in \mathbb{Q}, \ s > t\} \subset X - K.
Inverse images of subbase elements are open. Thus f is continuous.
```

Corollary 7.1.8.1. All T_4 spaces are completely regular and hence Tychonoff.

Proof. Let $x \in X$ and D be closed subset not containing x. We have X is a T_4 space. Therefore X is T_1 as well as normal.

Now the singleton set, $\{x\}$ is closed, since X is a T_1 space. And by **Urysohn's lemma** for disjoint, closed subsets $\{x\}$, D there exists a Urysohn function which is a continuous, real-valued function $f: X \to [0,1]$ such that f(x) = 0 and f(y) = 1 for all $y \in D$. Therefore the space X is completely regular and hence Tychonoff.

Remark (Urysohn function). The function whose existence is asserted by Urysohn's lemma is called a Urysohn function

7.1.3 Tietze Characterisation of Normality

Tietze Characterisation of normality

X is normal \iff there exist continuous extension of real-valued functions on closed subsets.

Proposition 7.1.9. Let A be a subset of a space X and let $f: A \to \mathbb{R}$ be continuous. Then any two extensions of f to X agree on \overline{A} . In other words, if at all an extension of f exists its values on \overline{A} are uniquely determined by values of f on A.

Proposition 7.1.10 (Tietze: necessary condition). Suppose a topological space X has the property that for every closed subset A of X, every continuous real valued function on A has a continuous extension to X. Then X is normal.

Definitions 7.1.11 (Pointwise Convergence). Let X be a topological space and (Y, d) a metric space. Then a sequence of functions $\{f_n\}$ from X to Y converges pointwise to f if for every $x \in X$ the sequence $\{f_n(x)\}$ converges to f(x) in Y.

In other words, given a very small value, $\epsilon > 0$, there exists some $\delta > 0$ such that for every $x \in X$ there exists $N_x \in \mathbb{N}$. This N_x may be different for different values of x and for every $n > N_x$, $d(f(x), f_n(x)) < \delta$.

Definitions 7.1.12 (Uniform Convergence). Let X be a topological space and (Y,d) a metric space. Then a sequence of functions $\{f_n\}$ from X to Y converges uniformly to f if given a small $\epsilon > 0$, there exists $\delta > 0$ such that there exists $N \in \mathbb{N}$. This N is independent of the value of x and for every n > N, $d(f(x), f_n(x)) < \delta$.

Proposition 7.1.13. Let X, (Y,d), $\{f_n\}$ and f be as above and suppose $\{f_n\}$ converges to f uniformly. If each f_n is continuous, then f is continuous.

Definitions 7.1.14 (Uniform Convergence of Series). Let X be a topological space and (Y,d) be a metric space. Then a series of function $\sum_{n=1}^{\infty} f_n$ converges uniformly to f if the sequence of partial sums converges uniformly to f.

In other words, let $g_m = \sum_{n=1}^m f_n$. Then $\sum_{n=1}^\infty f_n$ converges to f uniformly if the sequence of partial sums $\{g_m\}$ converges to f uniformly.

Proposition 7.1.15. Let $\sum_{n=1}^{\infty} M_n$ be a convergent series of non-negative real numbers. Suppose $\{f_n\}$ is a sequence of real valued functions on a space X such that for each $x \in X$ and $n \in \mathbb{N}$, $|f_n(x)| \leq M_n$. Then the series $\sum_{n=1}^{\infty} f_n$ converges uniformly to a real valued function on X.

Theorem 7.1.16. Let A be a closed subset of a normal space X and suppose function $f: A \to [-1, 1]$ is a continuous function. Then there exists a continuous function $F: X \to [-1, 1]$ such that F(x) = f(x) for all $x \in A$.

Theorem 7.1.17. Let A be a closed subset of a normal space X and suppose function $f: A \to (-1,1)$ is a continuous function. Then there exists a continuous function $F: X \to (-1,1)$ such that F(x) = f(x) for all $x \in A$.

Corollary 7.1.17.1 (Tietze: sufficient condition). Any continuous real-valued function on a closed subset of a normal space can be extended continuously to the whole space.

7.2 Products and Coproducts

7.2.1 Cartesian Products of Families of Sets

 ${\bf 7.2.2} \quad {\bf The\ Product\ Topology}$

7.2.3 Productive Properties

7.3 Embedding and Metrisation

7.3.1 Evaluation Functions into Products

7.3.2 Embedding Lemma and Tychonoff Embedding

7.3.3 The Urysohn Metrisation Theorem

7.4 Nets and Homotopy

7.4.1 Definition and Convergence of Nets

Definitions 7.4.1 (Directed Set). /?, 10.1.1/

A directed set D is a pair (D, \geq) where D is a nonempty set and ge is a binary relation on D such that

- 1. The relation 'follows' (\geq) is transitive. ie, $m \geq n, n \geq p \implies m \geq p$
- 2. The relation 'follows'(\geq) is reflexive. ie, For every $m \in D, m \geq m$
- 3. For any $m, n \in D$, there exists $p \in D$ such that $p \ge m$ and $p \ge n$.

sequence in a set X is a function f from the set of all integers into X.

Definitions 7.4.2 (Net). /?, 10.1.2/

A net in a set X is a function S from a directed set D into the set X.

Remark. The set $\mathbb N$ together with the relation 'less than or equal to' (\leq) is a directed set. Clearly, the relation 'less than or equal to' is reflexive and trasitive. And the third condition is true iff every finite subset E of D has an element $p \in E$ such that p follows each element of E. This is a weaker notion compared to the well ordering principle(\star^{12}) of the set of all integers. Thus $\mathbb N$ is a directed set and every sequence in X is also a net in X.

Remark (Significance of Net). A net on a set is a generalisation of 'a sequence on a set' obtained by simplifying the domain of the sequence into a directed set. The notion directed set is derived by assuming a few properties of \mathbb{N} .

The convergence of sequence is not strong enough to characterise topologies as the limit of convergent sequences are unique for both Hausdorff and Cocountable spaces. The notion of Net allows us to differentiate between Hausdorff spaces from Co-countable spaces in terms of convergence of nets. The limit of a convergent net on a topological space is unique iff it is a Hausdorff space. ie, We have removed a few restrictions, so that we will have some convergent nets (which are obviously not sequences) with multiple limit points for Co-countable spaces.

Remark. Examples of Directed Sets

- 1. Let X be a topological space and $x \in X$. Then the neighbourhood system \mathcal{N}_x is a directed set with the binary relation \subset (subset/inclusion).
 - (a) Let U, V, W be any three neighbourhoods of $x \in X$ such that $U \subset V$ and $V \subset W$. Then, clearly $U \subset W$.

 Therefore, $U \geq V, V \geq W \implies U \geq W$.
 - (b) Let U be any neighbourhood of $x \in X$, then $U \subset U$. Therefore, $U \geq U$.
 - (c) Let U, V be any two neighbourhoods of $x \in X$, then there exists their intersection $W = U \cap V$, which is a neighbourhood of x. Clearly $W \subset U$ and $W \subset V$.

 Therefore $\forall U, V \in \mathcal{N}_x, \exists W \in \mathcal{N}_x \text{ such that } W \geq U \text{ and } W \geq V$.

 $^{^{12}\}mbox{Well-ordering principle}$: Every subset of $\mathbb N$ has a least element in it.

- 2. Let \mathcal{P} be the set of all partitions on closed unit interval [0,1]. A partition $P \in \mathcal{P}$ is a refinement of $Q \in \mathcal{P}$ if every subinterval in P is contained in some subinterval of Q. Then \mathcal{P} with the binary relation refinement is a directed set.
 - (a) Suppose P, Q, R are three partitions of [0,1] such that P is a refinement of Q and Q is a refinement of R, then clearly P is a refinement of R since each subinterval of P is contained some subinterval of Q, which is contained in some subinterval of R.

 Therefore, $P \geq Q, Q \geq R \implies P \geq R$
 - (b) Suppose P is a partition of [0,1]. Then trivially, P is a refinement of itself since every subinterval of P is contained in the same subinterval of P.

 Therefore, $\forall P \in \mathcal{P}, P \geq P$
 - (c) Suppose P, Q be any two partition of [0,1]. Then $R=P\cup Q$ is a refinement of both the partitions. Therefore $\forall P,Q\in \mathcal{P},\ \exists R\in \mathcal{P}\ such\ that\ R\geq P\ and\ R\geq Q$

For example, let $P = \{0, 0.3, 0.7, 1\}$. Then the subintervals in P are [0, 0.3], [0.3, 0.7] and [0.7, 1]. Let $Q = \{0, 0.3, 0.5, 1\}$ and $R = \{0, 0.3, 0.5, 0.7, 1\}$. Then R is a refinement of P, but Q is not a refinement of P since there is a subinterval [0.5, 1] in Q which is not properly contained in any subinterval of P. However, R is a refinement of Q as well.

Remark. Examples of Nets

1. Let X be a topological space and $x \in X$. Let \mathcal{N}_x be the set of all neighbourhoods of x. Let $D = (\mathcal{N}_x, X)$ be the directed set given by $(N, y) \in (\mathcal{N}_x, X)$ if $N \in \mathcal{N}_x$ and $y \in N$ and $(N, y) \geq (M, z)$ if $N \subset M$. Then the function $S: (\mathcal{N}_x, X) \to X$ given by S(N, y) = y is a net on X.

For example, let $X = \{a, b, c, d\}$ and $\mathcal{T} = \{\{a\}, \{a, b\}, \{a, b, c\}, \{a, b, c, d\}\}$. Also let $S : (\mathcal{N}_b, X) \to X$ defined by S(N, y) = y. Suppose $C = \{a, b, c\}$. Then $C \in \mathcal{N}_b$. ie, C is a neighbourhood of b. Then S(C, c) = c.

2. Riemann Net - Let $D = (\mathcal{P}, \xi)$ where \mathcal{P} is the set of all partitions on [0, 1] and ξ is a finite sequence in [0, 1] such that consecutive terms belongs to consecutive subintervals of the partition. The set (\mathcal{P}, ξ) is directed set with \geq given by $(\mathcal{P}, \eta) \geq (\mathcal{Q}, \psi)$ iff \mathcal{P} is a refinement of \mathcal{Q} .

For example, let $P \in \mathcal{P}$ is given by $P = \{0, 0.3, 0.7, 1\}$ and $\eta = \{0.2, 0.6, 0.9\}$. Then $(P, \eta) \in (\mathcal{P}, \xi)$.

Let $f: \mathbb{R} \to \mathbb{R}$ be any function, then the function,

$$S: (\mathcal{P}, \xi) \to \mathbb{R}$$
 defined by $S(P, \eta) = \sum_{i=1}^{k} f(\eta_k)(a_k - a_{k-1})$

where $P = \{a_0, a_1, \dots, a_k\}$ is the Riemann Net with respect to the real function f.

For example, let $f(x) = x^2$ and P, η are same as above example, then $S(P, \eta) = 0.2^2(0.3 - 0) + 0.6^2(0.7 - 0.3) + 0.9(1 - 0.7) = 3.99$

Definitions 7.4.3 (Convergence of a Net). /?, 10.1.3/

A net $S: D \to X$ converges to a point $x \in X$ if for any $nbd\ U$ of x, there exists $m \in D$ such that $n \in D$, $n \ge m \implies S(n) \in U$. And x is a limit of the net S.

Remark. The choice of m depends on the choice of neighbourhood U.

 $S: D \to X, S \to x \iff (\forall U \in \mathcal{N}_x, \exists m_U \in D, \text{ such that } n \geq m_U \implies S(n) \in U)$

Theorem 7.4.4 (Net characterisation of Hausdorff space). [?, 10.1.4] A topological space is Hausdorff iff limits of all nets in it are unique.

Proof. Let X be a Hausdorff space. Suppose $S: D \to X$ is net on X such that S converges to two distinct points $x, y \in X$. Since X is a Hausdorff space and $x \neq y$, there exists open subsets U, V such that $x \in U$, $y \in V, U \cap V = \phi$.

The net S converges to $x \in X$, therefore $\exists m_x \in D$ such that $n \geq m_x \Longrightarrow S(n) \in U$ And, the net S converges to $y \in X$, therefore $\exists m_y \in D$ such that $n \geq m_y \Longrightarrow S(n) \in V$.

Since D is a directed set and $m_x, m_y \in D$, there exists $p \in D$ such that $p \ge m_x$ and $p \ge m_y$. Now, $n \ge p \implies n \ge m_x$, $n \ge m_y$, since \ge is transitive. (ie, $n \ge p$, $p \ge m_x \implies n \ge m_x$, and $n \ge p$, $p \ge m_y \implies n \ge m_y$).

We have $n \geq p \implies n \geq m_x$ and $n \geq m_x \implies S(n) \in U$. Therefore, $n \geq p \implies S(n) \in U$. Similarly, $n \geq p \implies n \geq m_y \implies S(n) \in V$. Therefore $S(n) \in U \cap V$ which is a contradiction, since $U \cap V = \phi$. Therefore, if a net S converges to two points x, y, then x = y. That is, if a net S in a Hausdorff space X is convergent then its limit is unique.

Conversely, suppose that X is a topological space and every convergent net in X has a unique limit. Suppose X is not a Haudorff space. Then there exists at least two distinct points $x, y \in X$ such that every neighbourhood of x intersects with every neighbourhood of y. Now consider the set $D = \mathcal{N}_x \times \mathcal{N}_y$ and relation \geq on D such that $(U_1, V_1) \geq (U_2, V_2)$ if $U_1 \subset U_2$ and $V_1 \subset V_2$.

By the axiom of choice, a function $S:D\to X$ such that $S(U,V)\in U\cap V$ is well defined, since every nbd of x intersects every nbd of y. Thus, S is a net in X. We claim that S converges to both x and y.

Let U be a nbd of x. Then $S(U',V') \in U' \cap V'$. We have $(U,X) \in D$ such that $(U',V') \geq (U,X) \implies U' \subset U$. Then, $S(U',V') \in U' \cap V' \subset U \cap X = U$. Thus, for any nbd U containing x, we have $(U,X) \in D$ such that $(U',V') \geq (U,X) \implies S(U',V') \in U$. Therefore, S converges to $x \in X$.

Similarly, Let V be a nbd of y. Then for any nbd V containing y, we have $(X,V)\in D$ such that $(U',V')\geq (X,V)\implies S(U',V')\in V$, since

 $S(U',V') \in U' \cap V' \subset X \cap V = V$. Therefore, S converges to $y \in X$ as well, where $x \neq y$. This is a contradition to the assumption that every convergent net in X has a unique limit. Therefore, for any two points $x,y \in X$, there should be some nbd of x that doesn't intersect some nbd of y. Therefore, X is a Hausdorff space.

Definitions 7.4.5 (Eventual Subset). [?, 10.1.5]

A subset E of a directed set D is an eventual subset of D if there exists $m \in D$ such that $n \ge m \implies n \in E$.

Remark. Let E be an eventual subset of D such that $n \ge m \implies n \in E$. Then $p \in E \not\Longrightarrow p \ge m$. ie, Subset E may contain elements that doesn't follow the above m.

Remark. /?, 10.1.6/

Let E be an eventual subset of D, then E is a directed set.

1.
$$m, n, p \in E, m \ge n, n \ge p \implies m, n, p \in D, m \ge n, n \ge p \implies m \ge p$$

$$2. m \in E \implies m \in D \implies m \ge m$$

3.
$$m, n \in E \implies m, n \in D \implies \exists p \in D \text{ such that } p \geq m \text{ and } \geq n$$
.

$$\exists m' \in D \text{ such that } n' \geq m' \implies n' \in E. \text{ (E is an eventual subset of D)}$$

$$p, m' \in D \implies \exists p' \in D \text{ such that } p' \geq p \text{ and } p' \geq m'.(D \text{ is a directed Set})$$

$$p' \ge m' \implies p' \in E$$
. (E is eventual subset of D with respect to m')

$$p' \ge p, \ p \ge m \implies p' \ge m \ and \ p' \ge p, \ p \ge n \implies p' \ge n.$$

Therefore $\forall m, n \in E, \exists p' \in E \text{ such that } p' \geq m \text{ and } p' \geq n.$

Definitions 7.4.6 (Net eventually in A). [?, 10.1.5]

Let $S: D \to X$ be a net in a topological space X. Then S is eventually in subset A of X if $S^{-1}(A)$ is an eventual subset of D.

Remark. Let $S:D\to X$ be a net in X. Then S converges to $x\in X$ if S is eventually in each $nbd\ U$ of x.

Definitions 7.4.7 (Cofinal subset). [?, 10.1.7]

A subset F of a directed D is a cofinal subset of D if for any $m \in D$, there exists $n \in F$ such that $n \ge m$.

Remark. Let X be a topological space and $x \in X$. Let \mathcal{N}_x be the set of all neighbourhood of x and \mathcal{L} be a local base of X at x. We have, (\mathcal{N}_x, \geq) is a directed set where $\forall U, V \in \mathcal{N}_x$, $U \geq V \iff U \subset V$, then \mathcal{L} is cofinal in \mathcal{N}_x .

Remark. /?, 10.1.8/

Let F be a cofinal subset of D, then F is a directed set.

1.
$$m, n, p \in F, m \ge n, n \ge p \implies m, n, p \in D, m \ge n, n \ge p \implies m \ge p$$

- $2. \ m \in F \implies m \in D \implies m \ge m$
- 3. $m, n \in F \implies m, n \in D \implies \exists p \in D \text{ such that } p \geq m \text{ and } \geq n$.

 $E \text{ is cofinal, } p \in D \implies \exists p' \in F \text{ such that } p' \geq p.$

 $p' \ge p, \ p \ge m \implies p' \ge m \ and \ p' \ge p, \ p \ge n \implies p' \ge n.$

Therefore $\forall m, n \in F$, $\exists p' \in F$ such that $p' \geq m$ and $p' \geq n$.

Definitions 7.4.8 (Net frequently in A). ??, 10.1.7

Let $S: D \to X$ be a net in a topological space X. Then S is frequently in subset B of X if $S^{-1}(B)$ is a cofinal subset of D.

Proposition 7.4.9. [?, 10.1.6]

Let $S: D \to X$ be a net in a topological space X. Let E be an eventual subset of D. Then, S converges to x iff $S_{/E}$ converges to x.[?, 10.1.6]

Proof. Let $S:D\to X$ be a net in X,E be an eventual subset of D, and $x\in X$. Then, $S_{/E}:E\to X$ is defined by $n\in E\implies S_{/E}(n)=S(n)$

Suppose S converges to x. Let U be a nbd of x, then S is eventually in U. ie, $S^{-1}(U)$ is an eventual subset of D. Then $\exists m \in D$ such that $n \geq m \implies n \in S^{-1}(U) \implies S(n) \in U$. Since set E is eventual subset of D, $\exists m' \in D$ such that $n \geq m' \implies n \in E$.

Since E is a directed set, $S_{/E}: E \to X$ is a net in X. And $m, m' \in D \Longrightarrow \exists p \in D$ such that $p \geq m$ and $p \geq m'$. We have, $p \geq m' \Longrightarrow p \in E$. And $n \geq' p \Longrightarrow n \geq p, \ p \geq m \Longrightarrow n \geq m \Longrightarrow S(n) \in U \Longrightarrow S_{/E}(n) \in U$. Therefore, $n \geq' p \Longrightarrow S_{/E}(n) \in U$. Since U is arbitrary, $S_{/E}$ converges to x.

Conversely, suppose that $S_{/E}$ converges to x. Let U be a nbd of x, then $S_{/E}$ is eventually in U. ie, $S_{/E}^{-1}(U)$ is an eventual subset of D. ie, $\exists m \in D$ such that $n \geq m \implies n \in S_{/E}^{-1}(U) \implies S_{/E}(n) \in U \implies S(n) \in U$. Therefore, $n \geq m \implies S(n) \in U$. Since, U is arbitrary, S converges to every nbd of x. ie, S converges to x.

Proposition 7.4.10. /?, 10.1.8/

Let $S: D \to X$ be a net in a topological space X. Let F be a cofinal subset of D. If S converges to x, then $S_{/_F}$ converges to x.

Proof. Let $S: D \to X$ be a net in X and S converges to $x \in X$. Also let F be a cofinal subset of D. Then $S_{/F}$ is also a net in X, since (F, \geq') is a directed set where $\forall m, n \in F, m \geq n \implies m \geq' n$.

Since S converges to x, for any nbd U of x, $\exists m \in D$, such that $n \geq m \Longrightarrow S(n) \in U$. Since F is cofinal, $\exists p \in F$ such that $p \geq m$. Thus $n \geq' p \Longrightarrow n \geq p$, $p \geq m \Longrightarrow n \geq m \Longrightarrow S(n) \in U \Longrightarrow S_{/F}(n) \in U$. Therefore, $\exists p \in F$ such that $n \geq' p \Longrightarrow S_{/F}(n) \in U$. Since U is arbitrary, $S_{/F}$ is eventually in every nbd of x. ie, $S_{/F}$ converges to x.

Remark. But converse of the above is not true. $S_{/F}$ converges to x does not imply that S converges to x, since cofinal subset F not necessarily contain every element following a particular m.

Definitions 7.4.11 (Cluster point). /?, 10.1.9/

Let $S: D \to X$ be a net in a topological space X. Then $x \in X$ is a cluster point of S, if S is frequently in each nbd U of x in X.

Proposition 7.4.12. /?, 10.1.10/

Let $S: D \to X$ be a net in a topological space X. Then $x \in X$ is a cluster points of X, if $S_{/F}$ converges to x for some cofinal subset F of D.

Proof. Let $S: D \to X$ be a net in X and (F, \geq') be a cofinal subset of (D, \geq) . Then $S_{/F}$ is also a net in X. Suppose $S_{/F}$ converges to $x \in X$. Let U be a nbd of x, then $\exists m \in F$ such that $n \geq' m \implies S_{/F}(n) \in U$.

Let $m' \in D$. Then $\exists p' \in F$ such that $p' \geq m'$, since F is a cofinal subset of D. We have, $m, p' \in F$, then $\exists p \in F$ such that $p \geq' m$ and $p \geq' p'$. Since $F \subset D$, we have $p, m \in F \implies p, m \in D$ and $p \geq' m \implies p \geq m$.

Also $p \geq' m \implies S_{/F}(p) \in U \implies S(p) \in U$. Therefore, $\forall m' \in D$, $\exists p \in D$ such that $p \geq m'$ and $S(p) \in U$. Since U, m' are arbitrary, S is frequently in every nbd of x. ie, x is a cluster point of S.

Definitions 7.4.13 (Subnet). /?, 10.1.11/

Let $S: D \to X$ be a net in a topological space X. Then a net $T: E \to X$ in X, is a subnet of S if there exists a function $N: E \to D$ such that $S \circ N = T$ and $\forall m \in D, \ \exists p \in E \ such \ that \ n \geq' p \implies N(n) \geq m$.

Remark. A net $T: E \to X$ is a subnet of $S: D \to X$ if $\exists N: E \to D$ such that $S \circ N = T$ and S is frequently in T(E). Let $T: E \to X$ be a subnet of $S: D \to X$ and A be a subset of X. If T eventually in A, then S is frequently in A.

Proposition 7.4.14. /?, 10.1.12/

Let $S: D \to X$ be a net in a topological space X. Then $x \in X$ is a cluster point of S iff there exists a subnet of S which converges to x.

Synopsis. Let (D, \geq) , (E, \geq') be two directed sets. And $T: E \to X$ be a subnet of $S: D \to X$.

If T converges to x, then T is eventually in each nbd U of x. And since T is a subnet of S, there exists $N: E \to D$ such that N(E) is a cofinal subset of D. Therefore, S is frequently in each nbd U of x. Thus, x is a cluster point of S.

If x is a cluster point of X, then S is frequently in every nbd of x. Let $N: E \to D$ be N(n,U) = n. Construct a directed subset E of $D \times \mathcal{N}_x$ such that $(n,U) \in E \iff S(n) \in U$. Now T is eventually in every nbd U of x, as those points with images outside U are removed by construction. Therefore, it is sufficient to show that $T: E \to X$ is a subnet of the net $S: D \to X$. Clearly, E is a directed set and $N: E \to D$ defined by N(n,U) = n satisfies both $S \circ N = T$ and $\forall n \in D$, $\exists p \in E$ such that $m \geq p \implies N(m) \geq n$.

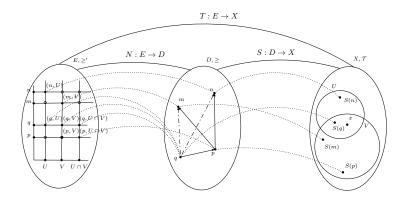


Figure 7.4: $\forall (n, U), (m, V) \in E, \exists (q, W) \in E \text{ such that } (q, W) \geq' (n, U), (q, W) \geq' (m, V)$

Proof. Let $S: D \to X$ be a net in X. Suppose there exists a subnet $T: E \to X$ that converges to $x \in X$. By the definition of subnet, we have $\exists N: E \to D$ such that $S \circ N = T$ and S is frequently in T(E).

We have, T convergers to x, thus for any neighbourhood U of x, there exists $m' \in E$ such that $n' \geq' m' \implies T(n') \in U$.

Also we have, T is a subnet of S. Then $\exists N : E \to D$ such that $\forall m \in D, \exists p' \in E$ such that $n' \geq p' \implies N(n') \geq m$.

Now, for any $m \in D$, we have $m', p' \in E$. Since E is a directed set, there exists $n' \in E$ such that $n' \geq m'$ and $n' \geq p'$.

Then, $n' > m' \implies T(n') \in U$ and $n' > p' \implies N(n') > m$.

Thus for any $m \in D$, there exists $N(n') = n \in D$ such that $S(n) = S(N(n')) = T(n') \in U$.

Thus S is frequently in any neighbourhood U of x. Therefore, x is a cluster point of S.

Conversely, suppose that x is a cluster point of S. We have to construct a directed set (E, \geq') and a function $N: E \to D$ such that T is a subnet of S and T converges to x. Let \mathcal{N}_x be the family of all neighbourhood of x in X.

Define $N: E \to D$ by N(n,U) = n. Again for any $(m,V) \in E$, there exists $m \in D$ such that $(n,U) \geq' (m,U)$ implies there exists $n \in D$ such that N(n,U) = n and $n \geq m$.

It remains to prove that, T converges to x. Let $U \in \mathcal{N}_x$ be a nbd of x in X. We have, x is a cluster points of S. Therefore, $\forall m \in D$, $\exists p \in D$ such that $S(p) \in U$. By the construction of E, we have $(p,U) \in E$. Suppose $(n,V) \geq' (p,U)$, then $n \geq p$, $V \subset U$, and $S(n) \in V$. Clearly $S(n) \in U$. Therefore, $\forall (n,V) \geq' (p,U)$, $T(n,V) \in U$. That is, T is eventually in every nbd of x. ie, subnet T is convergent to x. Therefore, for each cluster point of the net S, there exists some subset converging to it.

Remark. • Importance of Construction of E

If x is a cluster point of a net S in X, then S is frequently in some cofinal subset of D. Thus, if we consider any cofinal subset D' of D which is a direct set with \geq restricted to D'. Then $N:D'\to D$ defined by N(n)=n gives a subnet $T:D'\to X$ of the net S. However, this subnet need not converge to x. The strongest statement, we can make on T is that 'x is a cluster point of T', since $N:D\times \mathcal{N}(x)\to D$, N(n)=n is completely independent of U. This problem is overcame by constructing E which is dependent on each nbd U of x.

• Existence of $q \in D$ such that $q \geq' p$ and $S(q) \in U$. We have, p follows both n&m and $U \cap V$ is a subset of both U&V. However, since S is only frequently in U, p not necessarily be in U. But there is always someone following p which has its image in U. This q follows both n&m, since \geq' is transitive.

7.5 Variations of Compactness

7.5.1 Variations of Compactness

In this chapter, we have two other notions of compactness - countable compactness and sequential compactness.(\star^{13})

Compact A topological space is compact iff every open cover of it has a finite subcover. ([?, 6.1.1]) [Heine-Borel]

Countably Compact A topological space is countably compact iff every countable, open cover of it has a finite subcover. [?, 11.1.1]

Sequentially Compact A topological space is sequentially compact iff every sequence in it has a convergent subsequence. [?, 11.1.8] [Bolzano-Weierstrass]

Countable compactness is a weaker notion compared to compactness. (\star^{14}) However, sequentially compact and compact are not necessarily comparable. (\star^{15})

 $^{^{13} \}text{For } \mathbb{R},$ Compactness & Sequentially compactness are equivalent to the completeness axiom.

¹⁴Every compact space is countably compact.

 $^{^{15}\}mathcal{T}_1, \mathcal{T}_2$ are non-comparable, if $\mathcal{T}_1 \not\subset \mathcal{T}_2$ and $\mathcal{T}_2 \not\subset \mathcal{T}_1$.[?, 4.2.1]

We have seen earlier that compactness has the following properties 1. compactness is weakly hereditary. [?, 6.1.10] 2. compactness is preserved under continuous functions. [?, 6.1.8] 3. every continuous real functions on compact space is bounded and attains its extrema. [?, 6.1.6] 4. every continuous real function on a compact, metric space is uniformly continuous by Lebesgue covering lemma. [?, 6.1.7

Countably compact spaces, Sequentially compact spaces have all the four properites listed above.

7.5.2Countable compactness

Weakly hereditary property

A subspace $(A, \mathcal{T}_{/A})$ being countably compact doesn't imply that (X, \mathcal{T}) is countably compact. However, if (X, \mathcal{T}) is a countably compact space and A is a closed subset of X, then $(A,\mathcal{T}_{/A})$ is also a countably compact space. In other words, countably compactness is weakly hereditary.

Theorem 7.5.1. Countable compactness is weakly hereditary. [?, 11.1.3]

Synopsis. Let A be a closed subset of countably compact space, X. If A has a countable open cover U, then we can obtain a respective countable, open cover for X by attaching X-A to the extensions of members of \mathcal{U} to X. This cover has a finite subcover. Then restricting them to A, we get a finite subcover of U.

Proof. Suppose X is a countably compact space. And A is a closed subset of X. We need to show that A is countably compact. Without loss of generality, (\star^{16}) assume that A is a proper subset of X. Then X-A is a non-empty, open subset of X.

Let \mathcal{U} be a countable open cover of A. Then $\mathcal{U} = \{U_1, U_2, \cdots\}$ where each element $U_k \in \mathcal{U}$ is an open subset of A. Since A is a subspace of X, every open subset U_k in A is of the form $G \cap A$ for some open subset G in X. Therefore, there exists open subsets $V(U_k)$ for each U_k such that $A \cap V(U_k) = U_k.(\star^{17})$

Define $\mathcal{V} = \{X - A, V(U_1), V(U_2), \cdots\}$. Clearly, \mathcal{V} is a countable open cover (\star^{18}) of X. We have X is countably compact, thus V has a finite subcover, say \mathcal{V}' . Without loss of generality assume (\star^{19}) that $X - A \in \mathcal{V}'$. Suppose $X-A \notin \mathcal{V}'$, then we can define another finite subcover $\mathcal{V}' \cup \{X-A\}$. Thus $\mathcal{V}' = \{ X - A, \ V(U_{n_1}), \ V(U_{n_2}), \cdots, \ V(U_{n_k}) \}.$

Then the corresponding subcover $\mathcal{U}' = \{U_{n_1}, U_{n_2}, \cdots, U_{n_k}\}$ is a finite subcover of \mathcal{U} . Since countable open cover \mathcal{U} and closed subset A are arbitrary, every closed subset of X with relative topology is countably compact. Therefore, countable compactness is weakly hereditary.

Remark. Proof depends on the following,

¹⁶Suppose A is not a proper subset of X. Then X = A and A is countably compact.

 $^{^{17}}$ Relative topology, $\mathcal{T}_{/A} = \{G \cap A : G \in \mathcal{T}\}$ $^{18}X - A$ is open in X. If $y \not\in A$, then $y \in X - A$. If $y \in A$, then $y \in U_k$ for some k.

¹⁹Otherwise, you will have to consider two cases: $X - A \in \mathcal{V}'$ and $X - A \notin \mathcal{V}'$

- 1. There is an extension map, $\psi: P(A) \to P(X)$ that preserve open subsets (and closed subsets). This ψ is an open map which not a true inverse of the restriction, $r: P(X) \to P(A)$, defined by $r(G) = G \cap A$ for every subset G of X.
- 2. Also we rely on the subset A being closed. Suppose X have many countable open covers, but X has only uncountable open covers corresponding to a particular countable open cover of A. In such a case, X being countably compact is insufficient for A to be countably compact.

The behaviour of countinous functions

We will now study the nature of continuous functions defined on countably compact spaces. Suppose X, Y are topological space and function $f: X \to Y$ is continuous. If X is countably compact, then f(X) is also countably compact. Continuous images of countably compact spaces are countably compact. In other words, countable compactness is preserved under continuous functions. (\star^{20})

Theorem 7.5.2. Countable compactness is preserved under continuous functions.[?, 11.1.2]

Synopsis. Let X be countably compact and $f: X \to Y$ be continuous. Suppose \mathcal{U} is a countable cover of f(X), then X has a countable cover \mathcal{V} obtained by taking inverse images. Since X is countably compact, \mathcal{V} has a finite subcover \mathcal{V}' . Now taking images of members of \mathcal{V}' , we get a finite subcover \mathcal{U}' of f(X).

Proof. Suppose X is a countably compact space, Y is a topological space and $f: X \to Y$ is a continuous function. Let $\mathcal{U} = \{U_1, U_2, \cdots\}$ be a countable cover of f(X) by set open in f(X). We have to show that \mathcal{U} has a finite subcover.

Define $\mathcal{V} = \{f^{-1}(U_1), f^{-1}(U_2), \cdots\}$. Then \mathcal{V} is a countable open cover of X, since $f^{-1}(U_k)$ are open subsets of X and,

$$\bigcup_{k=1}^{\infty} U_k = f(X) \implies f^{-1} \left(\bigcup_{k=1}^{\infty} U_k \right) = X$$

$$\implies \bigcup_{k=1}^{\infty} f^{-1}(U_k) = X$$

We have, \mathcal{V} is a countable open cover of X, which is a countably compact space. Therefore \mathcal{V} has a finite subcover $\mathcal{V}' = \{f^{-1}(U_{n_1}), f^{-1}(U_{n_2}), \dots, f^{-1}(U_{n_k})\}$.

$$\bigcup_{j=1}^{k} f^{-1}(U_{n_j}) = X \implies f^{-1}\left(\bigcup_{j=1}^{k} U_{n_j}\right) = X$$

$$\implies \bigcup_{j=1}^{k} U_{n_j} = f(X)$$

 $^{^{20}}$ A topological property is preserved under continuous functions if whenever a space has that property so does every continuous image of it.[?, 6.1.9]

Clearly $\mathcal{U}' = \{U_{n_1}, U_{n_2}, \cdots, U_{n_k}\}$ is a finite subcover of \mathcal{U} . Thus every countable open cover of f(X) by sets open in f(X) has a finite subcover. Therefore, continuous images of countably compact spaces are countably compact.

Remark. 1. For a continuous function, $f: X \to Y$ the inverse images of open subsets are open in X. The relation $f^{-1} \subset f(X) \times X$ is not a function. However, we may consider a function, $\psi: P(Y) \to P(X)$ such that $\psi(U) = f^{-1}(U)$ for any subset U of Y. This ψ is an open map which maps open subsets of Y to open subsets of X.

Theorem 7.5.3. Every continuous, real-valued function on a countably compact, metric space is bounded and attains its extrema.[?, 11.1.7]

Synopsis. Let X be a countably compact space and function $f: X \to \mathbb{R}$ be continuous. Then $f(X) \subset \mathbb{R}$ is countably compact. Real line \mathbb{R} is metrisable (\star^{21}) . Then f(X) is countably compact, metric space. Therefore f(X) compact. (\star^{22}) . The subset f(X) of \mathbb{R} is bounded and closed, since every compact subset of \mathbb{R} is bounded and closed. Thus f(X) contains its supremum and infimum. Therefore, f is bounded and attains its extrema.

Proof. Let X be a countably compact space and $f: X \to \mathbb{R}$ be continuous, real-valued function on the countably compact space, X. We have to show that f is bounded and attains its extrema.

Since countable compactness is preserved under continuous functions, f(X) is countably compact subset of \mathbb{R} . Since, f(X) is a subset of the metric space, \mathbb{R} and metrisability is hereditary, f(X) is again metrisable. (suppose) We have, every countably compact, metric space is compact. Then f(X) is a compact subset of \mathbb{R} .

Since every compact subset of \mathbb{R} is bounded and closed, f(X) is bounded and closed. Since every closed subset of \mathbb{R} contains supremem and infimum, f(X) contains its extrema. Therefore, every continuous, real-valued function on a countably compact space is bounded and attains its extrema.

We have assumed that every countably compact, metric space is compact. This result will be proved in the last section of this chapter. \Box

Remark. Since countably compact, metric spaces are compact. The above theorem can be used to prove that continuous, real-valued functions on a compact, metric space attains its extrema.

Due to the Lebesgue covering lemma, next result is quite simple.*

Theorem 7.5.4. Every continuous, real-valued function on a countably compact, metric space is uniformly continuous.

Proposition 7.5.5. Let X be a first countable, Hausdorff space. Then every countably compact subset A of X is closed. ?, Exercises 11.1.7

²¹[?, 4.2 Example 4], \mathbb{R} with usual metric $d: R \to R$, d(x,y) = |x-y|

 $^{^{22}}$ [?, 11.1.11] On metric spaces, countable compactness \implies compactness.

7.5.3 Sequential Compactness

Weakly hereditary property

Theorem 7.5.6. Sequential compactness is weakly hereditary. [?, Exercises 11.1.3]

The behaviour of countinous functions

Theorem 7.5.7. Sequential compactness is preserved under continuous functions.[?, Exercises 11.1.4]

Synopsis. Let X be sequentially compact and function $f: X \to Y$ be continuous. Then any sequence, $\{y_k\}$ in f(X) has a sequence, $\{x_k\}$ in X such that $f(x_k) = y_k$. Sequence $\{x_k\}$ has a subsequence $\{x_{n_k}\}$ converging to x, then sequence $\{f(x_n)\}$ in f(X) has the subsequence $\{f(x_{n_k})\}$ converging to f(x).

Proof. Let X be a sequentially compact space, function $f: X \to Y$ be continuous and $\{y_n\}$ be a sequence in f(X) subset of Y. Construct a sequence $\{x_n\}$ such that $f(x_k) = y_k$, $\forall k$.

Every sequence in X has a convergent subsequence. Thus $\{x_n\}$ has a subsequence $\{x_{n_k}\}$ converging to $x \in X$. The image of this subsequence $\{f(x_{n_k})\}$ is a subsequence of $\{y_k\}$. We claim that, $\{f(x_{n_k})\}$ converges to $f(x) \in f(X)$.

Let U be an open subset containing f(x), then $f^{-1}(U)$ is an open subset containing x. Since $\{x_{n_k}\}$ converges to x. There exists an integer n such that for every $k \geq n$, $x_k \in f^{-1}(U)$. Clearly, for each $k \geq n$, $f(x_k) \in U$. Since U is arbitrary, $\{f(x_{n_k})\}$ converges to f(x). Therefore, the image of any sequentially compact space is sequentially compact. In other words, sequentially compactness is preserved under continuous functions.

Remark. 1. Given a sequence $\{y_n\}$ in f(X), there is a sequence of subsets $\{U_n\}$ in P(Y) such that $U_n = f^{-1}(y_n)$. Since each U_n is non-empty, we can construct a sequence $\{x_n\}$ in X using a choice function. The convergent subsequence of $\{y_n\}$ depends on the selection of this choice function.

Given every sequentially compact, metric space is countably compact. We may assert the properties of countably compact, metric spaces on sequentially compact, metric spaces.

Theorem 7.5.8. Every continuous, real-valued function on a sequentially compact, metric space is bounded and attains its extrema.

Theorem 7.5.9. Every continuous, real-valued function on a sequentially compact, metric space is uniformly continuous. [?, Exercises 11.1.6]

7.5.4 Countable Compactness on T_1 spaces

In this section, we are going to see four different characterisations of countable compactness in T_1 spaces. The first two characterisations doesn't have anything to do with the T_1 axiom.

- T_1 Space A topological space X satisfy T_1 axiom if for any two distinct points $x, y \in X$, there exists an open subset $U \subset X$ containing x but not y.[?, 7.1.2]
- **countable compactness** A topological space is countably compact if every countable open cover has a finite subcover.[?, 11.1.1]
- finite intersection property A family \mathcal{F} of subsets of X has finite intersection property(f.i.p.) if every finite subfamily of \mathcal{F} has a non-empty intersection.[?, 10.2.6]
- **accumulation point** A point $x \in X$ is accumulation point of a subset $A \subset X$ if every open subset containing x has at least one point of A other than x.[?, 5.2.7]
- **limit point** A point $x \in X$ is a limit point of a sequence $\langle x_k \rangle$ in X if for every open subset U containing x, there exists an integer $N \in \mathbb{N}$ such that $x_k \in U$ for every $k \geq N$.[?, 4.1.7]
- cluster point A point $x \in X$ is a cluster point of a sequence $\langle x_k \rangle$ in X if for any neighbourhood V of x, the sequence $\langle x_k \rangle$ assumes a point in V infinitely many times.(\star^{23})

Countable compactness in T_1 spaces

Theorem 7.5.10. In a T_1 space X, following statements are equivalent,

- 1. X is countably compact
- 2. Every countably family of closed subsets of X with finite intersection property have non-empty intersection.
- 3. Every infinite subset $A \subset X$ has an accumulation point. (\star^{24})
- 4. Every sequence $\langle x_k \rangle$ in X has a cluster point.
- 5. Every infinite open cover of X has a proper subcover.[Arens-Dugundji]

Proof. $1 \implies 2$

Suppose X is countably compact. Let $C = \{C_1, C_2, \dots\}$ be a countable family of closed subsets of X with empty intersection. Define $\mathcal{U} = \{X - C_1, X - C_2, \dots\}$ is a family of open subsets of X. By de Morgan's law, (\star^{25})

$$\bigcap_{k=1}^{\infty} C_k = \phi, \text{ then } X = X - \left(\bigcap_{k=1}^{\infty} C_k\right) = \bigcup_{k=1}^{\infty} (X - C_k)$$

We have \mathcal{U} is a countable cover of X and X is countably compact space. Thus \mathcal{U} has a finite subcover $\mathcal{U}' = \{X - C_{n_1}, X - C_{n_2}, \cdots, X - C_{n_k}\}.$

$$\mathcal{U}'$$
 is a cover of X , then $X = \bigcup_{j=1}^{k} (X - C_{n_j})$

 $^{^{23}}x$ is a cluster point of $\langle x_k \rangle$ if for every integer N, there exists k > N such that $x_k \in V$. In other words, $\langle x_k \rangle$ is frequently in V.[?, 10.1.9]

 $^{^{24}\}mathrm{Every}$ infinite subset of $\mathbb R$ has a limit point is equivalent to the completeness axiom.

²⁵Complement of Intersection = Union of complements, $X - (C \cap D) = (X - C) \cup (X - D)$,

$$X - \bigcup_{j=1}^{k} (X - C_{n_j}) = \bigcap_{j=1}^{k} (X - (X - C_{n_j})) = \bigcap_{j=1}^{k} C_{n_j} = \phi$$

Now $C' = \{C_{n_1}, C_{n_2}, \dots, C_{n_k}\}$ has empty intersection. This is a contradiction to the finite intersection property of C. Thus C has non-empty intersection. Therefore, every countably family of closed subsets of X have non-empty intersection.

$$2 \implies 1$$

Let $\mathcal{U} = \{U_1, U_2, \dots\}$ be a countable cover of X. Then $\mathcal{C} = \{X - U_1, X - U_2, \dots\}$ is a countable family of closed subsets of X.

Let $\mathcal{U}' = \{U_{n_1}, U_{n_2}, \dots, U_{n_k}\}$ be any finite subfamily of \mathcal{U} . Suppose X is not countably compact, then \mathcal{U} doesn't have a finite subcover. Therefore, \mathcal{U}' is not a cover of X. And \mathcal{C} is a family of closed subsets with finite intersection property.

Therefore by assumption, the countable family of closed subsets $\mathcal C$ has a non-empty intersection.

$$\bigcap_{k=1}^{\infty} C_k \neq \phi, \text{ then } \bigcap_{k=1}^{\infty} C_k = \bigcap_{k=1}^{\infty} (X - U_k) = X - \left(\bigcup_{k=1}^{\infty} U_k\right) \neq \phi$$

Then \mathcal{U} is not a cover of X as well. This is a contradiction, therefore X is countably compact.

$$1 \implies 3$$

Suppose X is countably compact. Let A be an infinite subset of X. Suppose A doesn't have an accumulation point.

Let B be a countably infinite subset of A. Then B also doesn't have any accumulation point. Therefore, the derived set B' is empty. Thus B is a closed subset of X. Since countable compactness is weakly hereditary, subspace B is again countably compact.

For each point $b \in B$, there is an open subset V_b such that $V_b \cap B = \{b\}$, since $b \in B$ is not an accumulation point. Thus $\mathcal{U} = \{V_b \cap B : b \in B\}$ is a countable open cover of B. Clearly, \mathcal{U} doesn't have any finite subcover.

This is a contradiction to B being countably compact. Therefore, A has an accumulation point.

7.5.5 Variations of Compactness on Metric Spaces

In this document, we will see that from metric space point of view these two notions were equivalent to the compactness and were used alternatively. For example: in functional analysis (semester 3), you will find definitions like 'a normed space is compact iff every sequence in it has a convergent subsequence', which is clearly sequential compactness for a topologist.

Lindeloff A topological space is Lindeloff iff every open cover has a countable subcover.

First countable A topological space is first countable iff every point in it has a countable local base.

Second countable A topological space is second countable iff it has a countable base.

Base A family of subsets \mathcal{B} of X is a base of a topological space if every open subset can be expressed as union of some members of \mathcal{B}

Base Characterisation A family of subsets \mathcal{B} of X is a base of a topological space iff for every $x \in X$, and for every neighbourhood U of x, there is a member $B \in \mathcal{B}$ such that $x \in B \subset U$.

Local Base A family of subsets \mathcal{L} of X is a local base at point $x \in X$ if for every neighbourhood U of x, there is a member $L \in \mathcal{L}$ such that $x \in L \subset U$.

Equivalence

We are going to see when these three notions: compactness, countable compactness and sequentially compactness are equivalent.

Theorem 7.5.11. Countably compact, metric spaces are second countable.

Synopsis. For every positive real number r, there exists a non-empty maximal subsets A_r with every pair of points at least r distance apart. A_r are finite. The union of maximal subsets $A_{\frac{1}{r}}$ for each natural number n is a countable, dense subset D of X. Thus countably compact, metric spaces are separable. The family \mathcal{B} of all open balls with center at $d \in D$ and rational radius is a countable, base for X. Thus countably compact, metric spaces are second countable.

Proof. Let (X;d) be a countably compact,, metric space. For each positive real number $r \in \mathbb{R}$, r > 0 construct a family of subsets $A_r \subset X$ such that it is a maximal set of points which are at least r distances apart.

Then A_r is finite for every positive real number r. Suppose A_r is infinite for some real number r > 0, then A_r has a accumulation point, say x by the Characterisation of countable compactness of X.

Then every neighbourhood of x must intersect A_r at infinitely many points, since every metric space is a T_1 space. Consider $B(x, \frac{r}{2})$. Since any two points of $B(x, \frac{r}{2})$ are less than r distances apart, the intersection $B(x, \frac{r}{2}) \cap A_r$ can have atmost one point in it. Thus for every positive real number r, A_r is finite.

Define $D = \bigcup_{n=1}^{\infty} A_{\frac{1}{n}}$. We claim that D is a countable, dense subset of X.

Let $x \in X$ and B(x, r) be an open ball containing x, then there exists integer $n \in \mathbb{N}$ such that $\frac{1}{n} < r.(\star^{26})$

²⁶By archimedean property of integers, we have $\forall r \in \mathbb{R}, r > 0, \exists n \in \mathbb{N} \text{ such that } nr > 1.$

Then $B(x,r) \cap D \neq \phi$, since $B(x,r) \cap A_{\frac{1}{n}} \neq \phi$. Suppose $B(x,r) \cap A_{\frac{1}{n}} = \phi$, then $A_{\frac{1}{n}}$ is not maximal. Since, x is at least $r > \frac{1}{n}$ distance apart from each points of $A_{\frac{1}{n}}$. Therefore, D intersects with every open subset and thus dense in X.

We have a countable, dense subset D of X. Therefore, X is separable. Now define $\mathcal{B} = \{B(x, r) : r \in \mathbb{Q}, x \in D\}$. Clearly, \mathcal{B} is a countable base for X. By the construction of \mathcal{B} , X is second countable. (\star^{27})

Countable Compactness, Lindeloff \iff Compactness

Theorem 7.5.12. A topological space X is compact iff it is countably compact, Lindeloff space.

Proof. Let X be a compact space. Let \mathcal{U} be a countable open cover of X, then \mathcal{U} has a finite subcover \mathcal{U}' . Therefore, every compact space is countably compact.(\star^{28})

Conversely, suppose X is a countably compact, Lindeloff space. Since X is Lindeloff, every open cover \mathcal{U} has a countable subcover \mathcal{U}' . Since X countably compact, every countable open cover \mathcal{U}' has a finite subcover \mathcal{U}'' . Thus every open cover \mathcal{U} has a finite subcover \mathcal{U}'' . Therefore every countably compact, Lindeloff space is compact.

Countable Compactness, First Countable \implies Seq. Compactness

Theorem 7.5.13. Every countably compact, first countable space is Sequentially compact.

Proof. Let X be a countably compact, first countable space. Let $\{x_n\}$ be a sequence in X. By, equivalent conditions(\star^{29}) of countably compact spaces, every sequence in countably compact space X has a cluster point, say x. We have, X is first countable. Therefore, X has a countable local base \mathcal{L} at $x \in X$. How to construct a subsequence of $\{x_n\}$ converging to x?(\star^{30})

Remark. Every sequentially compact space is countably compact.★

Theorem 7.5.14. In a second countable space, all the three forms of compactness are equivalent. [?, 11.1.10]

Proof. Every second countable space is both first countable and Lindeloff. Every countably compact, Lindeloff space is countably compact. Therefore every countably compact, second countable space compact. Again, every countably compact, first countable space is sequentially compact. Therefore every countably compact, second countable space is sequentially compact. Conversely, every compact space is countably compact and every sequentially compact space is countably compact. (\star^{31})

 $^{^{\}rm 27}{\rm Every}$ separable, metric space is second countable.

²⁸Countable compactness is a weaker notion than compactness.

 $[\]frac{29}{20}$ [?, 11.1] Conditions 1,2, and 4 are equivalent. $2 \implies 4$ without T_1 axiom is out of scope.

³⁰[?, Exercises 10.1.11]

³¹Countable compactness is a weaker notion than sequential compactness as well.

Theorem 7.5.15. In a metric space, all the three forms of compactness are equivalent.[?, 11.1.11]

Proof. In a metric space each form of compactness implies second countability. And in second countable spaces, they are all equivalent. \Box

7.6 Homotopy of Paths

Definitions 7.6.1. Let X, Y be topological spaces and $f: X \to Y$, $f': X \to Y$ be continuous functions. Then f, f' are homotopic if there exists a continuous function $F: X \times I \to Y$ such that for every $x \in X$, F(x,0) = f(x) and F(x,1) = f'(x). And we write, $f \simeq f'$.

Definitions 7.6.2. Let X be a topological space and $f: I \to X$ and $f': I \to X$ be two paths. Then f, f' are path-homotopic if they have same initial point x_0 (ie, $x_0 = f(0) = f'(0)$), same final point x_1 (ie, $x_1 = f(1) = f'(1)$) and they are homotopic (ie, $\exists F: I \times I \to X$ such that $\forall x \in I$, F(x,0) = f(x) and F(x,1) = f'(x) also fixed at the end points x_0 and x_1 (ie, $\forall t \in I$, $F(0,t) = x_0$ and $F(1,t) = x_1$). And we write $f \simeq_p f'$.

Remark. If two paths f, f' are homotopic, then they have the same end points and there exists a (topologically) continuous deformation from one path into another.

Proposition 7.6.3. The relations \simeq , \simeq_p are equivalence relations.

Proof. Homotopy: Let f, f' be continuous functions from X into Y. Then f and f' are homotopic, $f \simeq f' \iff \exists F: X \times I \to Y$ such that F is continuous, F(x,0) = f(x), and F(x,1) = f'(x)

- 1. $f \simeq f$ We have $f: X \to Y$ is continuous. Define $F: X \times I \to Y$ such that F(x,t) = f(x). Clearly, F is continuous, F(x,0) = f(x) and F(x,1) = f(x). And $\exists F: X \times I \to Y \implies f \simeq f$.
- 2. $f \simeq f' \Longrightarrow f' \simeq f$ We have, $f \simeq f'$. Thus there exists a continuous function $F: X \times I \to Y$ such that F(x,0) = f(x) and F(x,1) = f'(x). Consider $F': X \times I \to Y$ defined by F'(x,t) = F(x,1-t). Clearly, F'is continuous, F'(x,0) = F(x,1) = f'(x), and F'(x,1) = F(x,0) = f(x). Thus, $\exists F'(x,t): X \times I \to Y \Longrightarrow f' \simeq_p f$
- 3. $f \simeq f', \ f' \simeq f'' \implies f \simeq f''$ We have, $f \simeq f' \iff \exists F : X \times I \to Y \text{ such that } F \text{ is continuous,}$ F(x,0) = f(x) and F(x,1) = f'(x).

Similarly, $f' \simeq f'' \iff \exists F': X \times I \to Y \text{ such that } F' \text{ is continuous,}$ F'(x,0) = f'(x) and F'(x,1) = f''(x).

Consider $G: X \times I \to Y$ defined by

$$G(x,t) = \begin{cases} F(x,2t) & ,t \in [0,\frac{1}{2}] \\ F'(x,2t-1) & ,t \in [\frac{1}{2},1] \end{cases}$$

We have, $G(x, \frac{1}{2}) = F(x, 1) = F'(x, 0) = f'(x)$. Thus, G is continuous by pasting lemma since $[0, \frac{1}{2}] \cap [\frac{1}{2}, 1] = \{\frac{1}{2}\}$. Also G(x, 0) = F(x, 0) = f(x) and G(x, 1) = F'(x, 1) = f''(x). Thus, $\exists G : X \times I \to Y \implies f \simeq f''$

Path Homotopy: Let f, f', f'' be paths in X. Then f and f' are path homotopic, $f \simeq_p f' \iff \exists F: I \times I \to X$ such that F is continuous, $\forall s \in [0,1], \ F(s,0) = f(s), \ F(s,1) = f'(s)$ and $\forall t \in [0,1], \ F(0,t) = f(0) = f'(0), \ F(1,t) = f(1) = f'(1)$

- 1. $f \simeq_p f$ We have $f: I \to X$ is continuous. Define $F: I \times I \to X$ such that $\forall s, t \in [0,1]$, F(s,t) = f(s). Clearly, F is continuous, $\forall s \in [0,1]$, F(s,0) = f(s), F(s,1) = f(s) and $\forall t \in [0,1]$, F(0,t) = f(0), F(1,t) = f(1). Thus, $\exists F: I \times I \to X \implies f \simeq_p f$.
- 2. $f \simeq_p f' \Longrightarrow f' \simeq_p f$ We have, $f \simeq_p f'$. Thus there exists a continuous function $F: I \times I \to X$ such that $\forall s \in [0,1], \ F(s,0) = f(s), \ F(s,1) = f'(s)$ and $\forall t \in [0,1], \ F(0,t) = f(0) = f'(0), \ F(1,t) = f(1) = f'(1)$.

Consider $F': I \times I \to X$ defined by F'(s,t) = F(s,1-t). Clearly, F' is continuous. And F'(s,0) = F(s,1) = f'(s), and F'(s,1) = F(s,0) = f(s). Also, F'(0,t) = F(0,1-t) = f(0) = f'(0) and F'(1,t) = F(1,1-t) = f(1) = f'(1). Thus, $\exists F'(s,t): I \times I \to X \implies f' \simeq_p f$

3. $f \simeq f'$, $f' \simeq f'' \Longrightarrow f \simeq f''$ We have, $f \simeq f' \iff \exists F: I \times I \to X$ such that F is continuous, $\forall s \in [0,1], \ F(s,0) = f(s), \ F(s,1) = f'(s)$ and $\forall t \in [0,1], \ F(0,t) = f(0) = f'(0), \ F(1,t) = f(1) = f'(1)$

Similarly, $f' \simeq f'' \iff \exists F' : I \times I \to X$ such that F' is continuous, $\forall s \in [0,1], \ F'(s,0) = f'(s), \ F'(s,1) = f''(s)$ and $\forall t \in [0,1], \ F'(0,t) = f'(0) = f''(0), \ F'(1,t) = f'(1) = f''(1)$

Consider $G: I \times I \to X$ defined by

$$G(s,t) = \begin{cases} F(s,2t) & ,t \in [0,\frac{1}{2}] \\ F'(s,2t-1) & ,t \in [\frac{1}{2},1] \end{cases}$$

We have, $G(s, \frac{1}{2}) = F(s, 1) = F'(s, 0) = f'(s)$. Thus, G is continuous by pasting lemma[?, §18.3 pp. 106], since $[0, \frac{1}{2}] \cap [\frac{1}{2}, 1] = {\frac{1}{2}}$.

Also G(s,0) = F(s,0) = f(s) and G(s,1) = F'(s,1) = f''(s).

Again, $\forall t \in [0, \frac{1}{2}]$, G(0,t) = F(0,2t) = f(0) = f'(0) = f''(0) and $\forall t \in [\frac{1}{2}, 1]$, G(0,t) = F'(0,2t-1) = f(0) = f'(0) = f''(0). Therefore, $\forall t \in [0, 1]$, G(0,t) = f(0) = f''(0).

Similarly, $\forall t \in [0, \frac{1}{2}], G(1,t) = F(1,2t) = f(1) = f'(1) = f''(1)$ and $\forall t \in [\frac{1}{2}, 1], G(1,t) = F'(1,2t-1) = f(1) = f'(1) = f''(1)$. Therefore, $\forall t \in [0, 1], G(1,t) = f(1) = f''(1)$. Thus, $\exists G : I \times I \to X \implies f \simeq_p f''$

Definitions 7.6.4. Let f be a path in X (ie, $f: I \to X$), then [f] is the equivalence class of all paths homotopic to f in X. (ie, $g \in [f] \iff f \simeq_p g$)

Remark. The set of homotopy classes of functions from X into Y is denoted by [X, Y]. And, the set of all path-homotopic classes on X is denoted by [I, X].

Remark (Straight-line homotopy). [?, §51 Example 1 pp. 320] Let X be a topological space, and f, g be continuous functions from X into a euclidean space, say \mathbb{R}^2 . Then f, g are straight line homotopic if there exists a continuous function F from $X \times I$ such that F deforms f into g along straight line segments joining them.

For example, F(x,t) = (1-t)f(x) + tg(x).

Remark. Let A be a convex subspace of \mathbb{R}^n . Then any two paths in A from x_0 to x_1 are path homotopic in A.

Proof. —continue page 321—

Remark. [?, §51 Example 2 pp. 321] This demonstrates that the straight-line homotopy is very sensitive to the holes in the space.

Definitions 7.6.5. Let f, g be two paths in X (ie, $f: I \to X$ and $g: I \to X$) such that $f(0) = x_0$, $f(1) = g(0) = x_1$ and $g(1) = x_2$. Then the prduct h = f * g is given by $h: I \to X$ and

$$h(s) = \begin{cases} f(2s) & , s \in [0, \frac{1}{2}] \\ g(2s-1) & , s \in [\frac{1}{2}, 1] \end{cases}$$

This h is well-defined, and continuous by pasting lemma. (\star^{32})

Definitions 7.6.6. The product operation on path-homotopy classes is defined by [f] * [g] = [f * g].

Remark. The product of path-homotopic classes is well-defined.

Proof. Let F be a path-homotopy between $f, f' \in [f]$ and G be a path-homotopy between $g, g' \in [g]$. Then $H: I \times I \to X$ defined by

$$H(s,t) = \begin{cases} F(2s,t) & s \in [0,\frac{1}{2}] \\ G(2s-1,t) & s \in [\frac{1}{2},1] \end{cases}$$

Then H is well-defined, and continuous by pasting lemma.

 $^{3^2}$ Pasting Lemma: Let $X = A \cup B$, where A and B are closed in A. Let $f: A \to Y$ and $g: B \to Y$ be continuous. If f(x) = g(x) for every $x \in A \cup B$, then f and g combine to give a continuous function $h: X \to Y$, defined by setting h(x) = f(x) if $x \in A$ and h(x) = g(x) if $x \in B$.

$$\forall s \in [0, \frac{1}{2}], \ H(s, 0) = F(2s, 0) = f(2s) \text{ and }$$

$$\forall s \in [\frac{1}{2}, 1], \ H(s, 0) = G(2s - 1, 0) = g(2s - 1).$$

$$\Longrightarrow \ H(s, 0) = (f * g)(s), \text{ by the definition of } f * g$$

$$\forall s \in [0, \frac{1}{2}], \ H(s, 1) = F(2s, 1) = f'(2s) \text{ and}$$

 $\forall s \in [\frac{1}{2}, 1], \ H(s, 1) = G(2s - 1, 1) = g'(2s - 1).$
 $\implies H(s, 1) = (f' * g')(s), \text{ by the definition of } f' * g'$

$$H(0,t) = F(0,t) = f(0) = x_0 = (f * g)(0)$$
, and $H(1,t) = G(1,t) = g'(1) = x_2 = (f' * g')(1)$

Then $H: I \times I \to X$ is a path-homotopy between f * g and f' * g'.

Definitions 7.6.7 (Groupoid). Let G be a set and * be a binary operation on G. Then (G, *) is a groupoid if it satisfies the following axioms

- g1 Associativity $\forall x, y, z \in G$, (x * y) * z = x * (y * z)
- g2 Existence of left and right identies There exist unique elements e_L and e_R such that $\forall x \in G$, $x * e_R = x$ and $e_L * x = x$.
- g3 Existence of inverse $\forall x \in G, \exists x^{-1} \in G \text{ such that } x * x^{-1} = e_L \text{ and } x^{-1} * x = e_R$

Definitions 7.6.8 (Positive Linear Map). A positive liear map $p:[a,b] \to [c,d]$ is the unique map of the form p(x) = mx + k such that p(a) = c and p(b) = d. Clearly, scaling factor, $m = \frac{d-c}{b-a}$ as we want to transform an interval of length b-a into an interval of length d-c. And offset k is given by,

$$p(a) = \frac{d-c}{b-a}a + k = c \implies k = c - \frac{a(d-c)}{b-a} = \frac{bc-ad}{b-a}$$

But, we won't fix m and k in p(x) = mx + k, instead we will focus on the unique map with graph of positive slope and passing through required end points. The graph of a positive linear map from [a,b] to [c,d] is always a straight-line with positive slope.

Remark. The inverse of a positive linear map is also a positive linear map. Given $p:[a,b]\to [c,d], p(x)=mx+k$, where $m=\frac{d-c}{b-a}$, $k=\frac{bc-ad}{b-a}$. Then it's inverse, $\bar{p}:[c,d]\to [a,b]$ is given by p(y):m'y+k', where $m'=\frac{b-a}{d-c}=\frac{1}{m}$, $k'=\frac{ad-bc}{d-c}=\frac{-k}{m}$. Clearly $m>0 \implies m'=\frac{1}{m}>0$.

Remark. The composite of two positive linear maps is also a (piece-wise) positive linear map. Let f, g be two positive linear maps. Then their composite map f * g is given by

$$(f*g)(x) = \begin{cases} f(2x) & x \in [a, \frac{a+b}{2}] \\ g(2\left(x - \frac{b-a}{2}\right)) & x \in \left[\frac{a+b}{2}, b\right] \end{cases}$$

Remember f * g exists only if f(b) = g(a). Therefore, f * g is a well-defined, continuous (by pasting lemma) and (piecewise) positive linear map.

Lemma 7.6.9. Let f, f' be two paths in X and $k: X \to Y$ be a continuous function. Let F be the path homotopy in X between the paths f and f'. Then $k \circ F$ is a path homotopy in Y between that paths $k \circ f$ and $k \circ f'$ That is, path homotopy is preserved under a continuous function.

Lemma 7.6.10. Let f, g be two paths in X with f(1) = g(0) and $k : X \to Y$ be a continous function. Then $k \circ (f * g) = (k * f) \circ (k * g)$

Theorem 7.6.11. Let f, g, h be paths in a topological space X, and [f], [g], [h] be respective path-homotopy classes. Suppose the operation product, * is defined by

$$[f] * [g] = [f *'g] where (f *'g)(s) = \begin{cases} f(2s) & s \in [0, \frac{1}{2}] \\ g(2s-1) & s \in [\frac{1}{2}, 1] \end{cases}$$

Then the product, * has the following properties:

1. Associativity

$$\forall [f], [g], [h] \in [I, X], \ ([f] * [g]) * [h] = [f] * ([g] * [h])$$

- 2. Existence of left and right identies Let $e_x: I \to X$ defined by $\forall s \in [0,1], \ e_x(s) = x$. Let f be a path from x_0 to x_1 , then there exist unique paths e_{x_0} and e_{x_1} such that $[f] * [e_{x_1}] = [f]$ and $[e_{x_0}] * [f] = [f]$. That is, e_{x_0} , e_{x_1} are respectively the left and right path-homotopy-identities.
- 3. Existence of inverse Let f be a path in X. The path, \bar{f} , defined by $\bar{f}(s) = f(1-s)$ is the reverse path of f. Then $[f] * [\bar{f}] = [e_{x_0}]$ and $[\bar{f}] * [f] = [e_{x_1}]$. That is, the inverse of class of f is the class of reverse path of f.

In other words, Set of all path-homotopy classes together with binary operation product, * is a groupoid. ie, ([I, X], *) is a groupoid.

Proof. Step 1: Properties 2&3

Let $e_0: I \to I$ such that $e_0(t) = 0$, $\forall t \in I$. And $i: I \to I$ such that i(t) = t, $\forall t \in I$. Then $e_0 * i$ is also a path. Since I is convex, there is a path homotopy(\star^{33}) G between i and $e_0 * i$. Let $f: I \to X$ be continuous path in X from x_0 to x_1 . Then $f \circ G$ is a path homotopy (by Lemma 2) in X between $f \circ i$ and $f \circ e_0 * i$ where $f \circ i$ and $f \circ e_0$ are paths from x_0 to x_1 in X.

$$\begin{split} f\circ(e_0*i)=&(f\circ e_0)*(f\circ i), \text{ by Lemma 1}\\ =&e_{x_0}*f, \text{ since } \forall s\in I,\ f(e_0(s))=x_0=e_{x_0}(s) \text{ and } f*i\simeq_p f \end{split}$$

Therefore $[e_{x_0}] * [f] \simeq_p [f]$, since $e_0 * i \simeq_p i$, and $f \circ (e_0 * i) \simeq_p f \circ i = f$.

Similarly, $e_1: I \to I$ such that $e_1(t) = 1$. Let H be a path homotopy (\star^{34}) between $i * e_1$ and i. Thus, $f \circ H$ is a path homotopy in X from $f \circ (i * e_1)$ and $f \circ i$.

$$\begin{split} f\circ (i*e_1) = & (f\circ i)*(f\circ e_1), \text{ by Lemma 1} \\ = & f*e_{x_1}, \text{ since } f*i \simeq_p f, \ i*e_1 \simeq_p e_1, \ \forall s \in I, \ (f(e_1(s)) = x_1 = e_{x_1}(s) \end{split}$$

 $^{^{33}}G:I\times I\to I,\ G(s,0)=i(s),\ G(s,1)=(e_0*i)(s),\ G(0,t)=0,\ G(1,t)=1.$ $^{34}H:I\times I\to I,\ H(s,0)=(i*e_1)(s),\ H(s,1)=i(s),\ H(0,t)=0,\ H(1,t)=1$

Since $i * e_1 \simeq_p i$, we have $f \circ (i * e_1) \simeq_p f \circ i = f$. Therefore $[f] * [e_{x_1}] \simeq_p [f]$. Thus, $[f] * [e_{x_1}] \simeq_p [f] \simeq_p [e_{x_0}] * [f]$. Therefore, [f] has left and right inverses ie, property 2 holds.

Consider inverse path $\bar{i}:I\to I,\ \bar{i}(s)=1-s.$ Then $i*\bar{i}$ is a path in I with both end points at 0. We have, $e_0:I\to I,\ e_0(s)=0$ is also a path with both end points at 0. Since I is convex, there is a path homotopy H in I between e_0 and $i*\bar{i}$. Then $f\circ H$ is a path homotopy between $f\circ e_0=e_{x_0}$ and $f\circ (i*\bar{i})=(f\circ i)*(f\circ \bar{i})=f*\bar{f}$. Therefore, $[e_{x_0}]\simeq_p [f]*[\bar{f}]$.

Similarly $\bar{i}*i$ and e_1 are paths with both end points at 1. Since I is convex, there is a path homotopy G in I between $\bar{i}*i$ and e_1 . Then $f\circ G$ is a path homotopy between $f\circ (\bar{i}*i)=(f\circ \bar{i})*(f\circ i)=\bar{f}*f$ and $f\circ e_1=e_{x_1}$. Therefore, $[\bar{f}]*[f]\simeq_p [e_{x_1}]$. Thus the path $\bar{f}:I\to X,\ \bar{f}(s)=f(1-s),\ \forall s\in I$ is reverse of f. Also $[f]*[\bar{f}]=[e_{x_0}]$ and $[\bar{f}]*[f]=[e_{x_1}]$. ie, property 3 holds. Step 2: Property 1

Let f, g, h be three paths in X and $f(1) = g(0) = x_1$ and $g(1) = h(0) = x_2$. Then f * (g * h) is defined by

$$(g*h)(s) = \begin{cases} g(2s) & s \in [0, \frac{1}{2}] \\ h(2s-1) & s \in [\frac{1}{2}, 1] \end{cases}$$

$$(f * (g * h))(s) = \begin{cases} f(2s) & s \in [0, \frac{1}{2}] \\ (g * h)(2s - 1) & s \in [\frac{1}{2}, 1] \end{cases}$$
$$= \begin{cases} f(2s) & s \in [0, \frac{1}{2}] \\ g(2(2s - 1)) & s \in [\frac{1}{2}, \frac{3}{4}] \\ h(2(2s - 1) - 1) & s \in [\frac{3}{4}, 1] \end{cases}$$

Similarly, (f * g) * h is defined by,

$$(f*g)(s) = \begin{cases} f(2s) & s \in [0, \frac{1}{2}] \\ g(2s-1) & s \in [\frac{1}{2}, 1] \end{cases}$$

$$((f * g) * h)(s) = \begin{cases} (f * g)(2s) & s \in [0, \frac{1}{2}] \\ h(2s - 1) & s \in [\frac{1}{2}, 1] \end{cases}$$
$$= \begin{cases} f(2(2s)) & s \in [0, \frac{1}{4}] \\ g(2(2s - 1)) & s \in [\frac{1}{4}, \frac{1}{2}] \\ h(2s - 1) & s \in [\frac{1}{2}, \frac{3}{4}] \end{cases}$$

Clearly, f*(g*h) and (f*g)*h are distinct path with common endpoints. ie, (f*(g*h))(0) = f(0) = ((f*g)*h)(0). And (f*(g*h)(1) = h(1) = ((f*g)*h)(1).

We need to define a path homotopy G between f*(g*h) and (f*g)*h. Let $[a,b],[c,d]\subset I$. Consider the path $p:I\to I$ defined by the following

three unique(\star^{35}) positive linear maps $p_1:[0,a]\to[0,c],\ p_2:[a,b]\to[c,d]$ and $p_3:[b,1]\to[d,1].$

$$p(t) = \begin{cases} p_1(t) & t \in [0, a] \\ p_2(t) & t \in [a, b] \\ p_3(t) & t \in [b, 1] \end{cases}$$

We can easily construct, a path homotopy P between identity map $i: I \to I$, i(s) = s and p as follows:

$$P(s,t) = \begin{cases} t + (p_1(t) - t) \frac{s}{a} & s \in [0, a] \\ t + (p_2(t) - t) \frac{(s-a)}{(b-a)} & s \in [a, b] \\ t + (p_3(t) - t) \frac{(s-b)}{(1-b)} & s \in [b, 1] \end{cases}$$

Therefore, we have $f*(g*h) \simeq_p i$ since there exists a path homotopy P corresponding to $[a,b]=[\frac{1}{2},\frac{3}{4}]$ and $[c,d]=[x_1,x_2]$. Similarly $(f*g)*h\simeq_p i$ since there exists a path homotopy P where $[a,b]=[\frac{1}{4},\frac{1}{2}]$ and $[c,d]=[x_1,x_2]$. ie, $[f*(g*h)]\simeq_p [(f*g)*h]$.

Theorem 7.6.12. Let f be a path in X, and a_0, a_1, \dots, a_n be numbers such that $0 = a_0 < a_1 < \dots < a_n = 1$. Let $f_i : I \to X$ be the path that equals the positive linear map of I onto $[a_{i-1}, a_i]$ followed by f. Then $[f] = [f_1] * [f_2] * \dots [f_n]$. In other words, every path is path-homotopic to a piecewise-linear path.

Proof. Let f be a piece-wise positive linear map such that

$$f(t) = \begin{cases} f_1(t) & t \in [0 = a_0, a_1] \\ f_2(t) & t \in [a_1, a_2] \\ \vdots & \vdots \\ f_n(t) & t \in [a_{n-1}, a_n] \end{cases}$$

where $f_i: I \to [a_{i-1}, a_i]$ such that $f_i(t)$ are a positive linear maps.

Consider the path $p:I \to I$ defined by the unique positive linear maps on the subintervals of any partition $\{0=x_0,x_1,\cdots,x_n\}$ of I. ie, $0=x_0 < x_1 < \cdots < x_n = 1$

$$p_{1}: [x_{0}, x_{1}] \to [a_{0}, a_{1}]$$

$$p_{2}: [x_{1}, x_{2}] \to [a_{1}, a_{2}]$$

$$\vdots$$

$$p_{n}: [x_{n-1}, x_{n}] \to [a_{n-1}, a_{n}]$$

$$p_{n}: [x_{n-1}, x_{n}] \to [a_{n-1}, a_{n}]$$

$$p_{n}: [x_{n-1}, x_{n}] \to [a_{n-1}, a_{n}]$$

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$$p_{n}: [x_{n}, x_{n}] \to [x_{n}, x_{n}]$$

$$p_{n}: [x_{n}, x_{n}] \to [x_{n}, x_{n}]$$

 $^{^{35}}p_1(t) = \frac{ct}{a}, \ p_2(t) = \frac{(d-c)t}{b-a} + \frac{bc-a}{b-a}, \ p_3(t) = \frac{(1-d)t}{1-b} + \frac{d-b}{1-b}$

Then there exists a path homotopy P between identity map $i: I \to I$, i(t) = t and p given by

$$P(s,t) = \begin{cases} t + (p_1(t) - t) \frac{a_1}{x_1} & s \in [0, x_1] \\ t + (p_2(t) - t) \frac{s - x_1}{x_2 - x_1} & s \in [x_1, x_2] \\ \vdots \\ t + (p_n(t) - t) \frac{s - x_{n-1}}{x_n - x_{n-1}} & s \in [x_{n-1}, x_n] \end{cases}$$

Since any product of f_1, f_2, \dots, f_n is a path p for some partition decided by the order of associativity. This partition can be constructed as follows: Let the last product operation (by associtivity) corresponds to $\frac{1}{2}$. The expression on its left corresponds to $[0, \frac{1}{2}]$ and expression on the right corresponds to $[\frac{1}{2}, 1]$. If there are any operations on any of these parts, the last operation (by associtivity) in them corresponds to the midpoint the respective subinterval and so on.

For examples: Consider, $(f_1*(f_2*f_3))*(f_4*f_5)$. Suppose we number the operations, $(f_1*_1(f_2*_2f_3))*_3(f_4*_4f_5)$. Then we have, $*_3 \to \frac{1}{2} \Longrightarrow *_1 \to \frac{1}{4} \Longrightarrow *_2 \to \frac{3}{8}$. Again $*_3 \to \frac{1}{2} \Longrightarrow *_4 \to \frac{3}{4}$. Thus, we have $\{0, \frac{1}{4}, \frac{3}{8}, \frac{1}{2}, \frac{3}{4}, 1\}$.

Thus given two paths f and f' with distinct order of associativity of n paths : f_1, f_2, \dots, f_n . We have path homotopy P, P' given by the P(s,t) for the respective partition constructed according to the order of associativity. Then, we have $f \simeq_p i$ and $f' \simeq_p i$. Thus, irrespective of the order of associativity all these paths are path homotopic. ie, $[f] = [f_1] * [f_2] \cdots [f_n]$

Subject 8

ME010203 Numerical Analysis with Python

8.1 Interpolation and Curve Fitting

Definitions 8.1.1. Given (n+1) data points (x_k, y_k) , $k = 0, 1, \dots, n$, the problem of estimating y(x) using a function $y : \mathbb{R} \to \mathbb{R}$ that satisfy the data points is the interpolation problem. ie, $y(x_k) = y_k$, $k = 0, 1, \dots, n$.

Definitions 8.1.2. Given (n+1) data points (x_k, y_k) , $k = 0, 1, \dots, n$, the problem of estimating y(x) using a function $y : \mathbb{R} \to \mathbb{R}$ that is sufficiently close to the data points is the curve-fitting problem. ie, Given $\epsilon > 0$, $|y(x_k) - y_k| < \epsilon$, $k = 0, 1, \dots, n$.

Remark. The data could be from scientific experiments or computations on mathematical models. The interpolation problem assumes that the data is accurate. But, curve-fitting problem assumes that there are some errors involved which are sufficiently small.

Definitions 8.1.3. Given (n + 1) data points (x_k, y_k) , $k = 0, 1, \dots, n$, the problem of estimating y(x) using a polynomial function of degree n that satisfy the data points is the polynomial interpolation problem.

Remark. Polynomial is the 'simplest' interpolant. [?, 3.2]

8.2 Polynomial Interpolation

There exists a unique polynomial of degree n that satisfy (n+1) distinct data points. There are a few methods to find this polynomial: 1. Lagrange's method 2. Newton's method.

8.2.1 Lagrange's Method

Interpolation polynomial¹ is given by,

$$P(x) = \sum_{i=0}^{n} y_i l_i(x), \text{ where } l_i(x) = \prod_{j=0, j \neq i}^{n} \frac{x - x_i}{x_j - x_i}$$
(8.1)

Remark. Lagrange's cardinal functions l_i , are polynomials of degree n and

$$l_i(x_j) = \delta_{ij} = \begin{cases} 0, & i = j \\ 1, & i \neq j \end{cases}$$

Proposition 8.2.1. Error in polynomial interpolation is given by

$$f(x) - P(x) = \frac{(x - x_0)(x - x_1) \cdots (x - x_n)}{(n+1)!} f^{(n+1)}(\xi)$$
 (8.2)

where $\xi \in (x_0, x_n)$

Remark. The error increases as x moves away from the unknown value ξ .

8.2.2 Newton's Method

The interpolation polynomial is given by,

$$P(x) = a_0 + a_1(x - x_0) + \dots + a_n(x - x_0)(x - x_1) \cdots (x - x_{n-1})$$
(8.3)

where $a_i = \nabla^i y_i$, $i = 0, 1, \dots, n$.

Remark. For Newton's Method, usually it is assumed that $x_0 < x_1 < \cdots < x_n$.

Remark. Lagrange's method is conceptually simple. But, Newton's method is computationaly more efficient than Lagrange's method.

Computing coefficients a_i of the polynomial

The coefficients are given by,

$$a_0 = y_0, \ a_1 = \nabla y_1, \ a_2 = \nabla^2 y_2, \ a_3 = \nabla^3 y_3, \cdots, a_n = \nabla^n y_n$$
 (8.4)

Remark. The divided difference $\nabla^i y_i$ are computed as follows:

$$\nabla y_1 = \frac{y_1 - y_0}{x_1 - x_0}$$

$$\nabla y_2 = \frac{y_2 - y_1}{x_2 - x_1} \qquad \nabla^2 y_2 = \frac{\nabla y_2 - \nabla y_1}{x_2 - x_1}$$

$$\nabla y_3 = \frac{y_3 - y_2}{x_3 - x_2} \qquad \nabla^2 y_3 = \frac{\nabla y_3 - \nabla y_2}{x_3 - x_2} \qquad \nabla^3 y_3 = \frac{\nabla^2 y_3 - \nabla^2 y_2}{x_3 - x_2}$$

Remark. Practise Problems

Find interpolation polynomial for the following data points:

¹Using P_n to represent some polynomial of degree n. It is quite a confusing a notation when it comes to Newton's method as author construct a psuedo-recursive definition.

```
1. \{(0,7), (2,11), (3,28)\}

Ans: 5x^2 - 8x + 7

[?, Example 3.1]

2. \{(-2,-1), (1,2), (4,59), (-1,4), (3,24), (-4,-53)\}

Ans: x^3 - 2x + 3

[?, Example 3.2]

3. \{(-1.2, -5.76), (0.3, -5.61), (1.1, -3.69)\}

Ans: x^2 + x - 6

[?, Problem Set 3.1.1]

4. \{(-3,0), (2,5), (-1,-4), (3,12), (1,0)\}

Ans: x^2 + 2x - 3

[?, Problem Set 3.1.7]

5. \{(0,1.225), (3,0.905), (6,0.652)\}

Ans: 0.0037x^2 - 0.1178x + 1.225

[?, Problem Set 3.1.9]
```

Remark. In Lagrange's Method, we can interpolate at the given point even without computing the polynomial. In Newton's method, we have to compute polynomial and then interpolate for the given point.

That is, evaluate the value of cardinal polynomials at the point and substitute in Equation ?? as shown in Section 3.2./?, Example 3.1/

8.2.3 Implementation of Newton's Method

Program 8.2.2. Computing Coefficients

```
def coefficients(xData,yData):
    m = len(xData)
    a = yData.copy()
    for k in range(1,m):
        a[k:m] = (a[k:m]-a[k-1])/(xData[k:m]-xData[k-1])
    return a
```

Line 1 def coefficients(xData,yData):

Defines a function which takes two arguments/parameters, named xData and yData. In [?, 3.2], you will find coeffts which I have changed to coefficients. xData, yData are numpy array objects. xData is a array

x_0	y_0				
x_1	y_1	∇y_1			
x_2	y_2	∇y_2	$\nabla^2 y_2$		
				٠.	
x_n	y_n	∇y_n	$\nabla^2 y_n$		$\nabla^n y_n$

Table 8.1: The $\nabla^i y_i$ Computation Table

with values x_0, x_1, \dots, x_n . And yData is array with values y_0, y_1, \dots, y_n . For example, the value of x_3 can be accessed as xData[3].

Line 2 m = len(xData)

The function len() is extended by numpy to give the length of array objects. In this context, len(xData) will return the value n+1, since there are n+1 values in xData array.

Line 3 a = yData.copy()

We need a copy of yData to work with. Unlike other programming languages like java, in python a = yData will assign a new label a to the same memory location and manipulating a will corrupt the original data in yData as well. In order to avoid this, we are **making a copy of the array object using the array method provided by the numpy library.**

Line 4 for k in range(1,m):

This is a python loop statement. This ask python interpreter to repeat the following sub-block m-1 times.² In this context, Line 5 will be executed n times, since the range(1,m) object is a list-type object with values $1, 2, \dots, m-1$. And interpreter executes Line 5 for each values in the range() object, ie, $k=1,2,\dots,m-1$ before interpreting Line 6.

$\label{eq:line 5 a[k:m] = (a[k:m]-a[k-1])/(xData[k:m]-xData[k-1])} \\$

This is very nice feature available in python. This statement, evaluates m-k values in a single step.ie, $a[k], a[k+1], \dots, a[m]$. This calculation corresponds to subsequent columns of the divided difference table, that we are familiar with. For example, executing Line 5 with k=3 is same as evaluating the $\nabla^3 y_j$ column. Note that the value a[0] is never updated and similarly a[2] changes when Line 5 is executed with k=1,2. From column 3 onward, a[2] is not updated. Therefore, after completing nth executing of the Line 5, we have $a[0] = y_0, \ a[1] = \nabla y_1, \ a[2] = \nabla^2 y_2, \dots, \ a[n] = \nabla^n y_n$.

Line 6 return a

This returns the array a which is the array of coefficients.

The logic of this program is in Line 4 and Line 5. So they need more explanation/understanding than anything else.

Program 8.2.3. Interpolating using Newton's Method

```
def interpolate(a, xData, x):
    n = len(xData)-1
    p = a[n]
    for k in range(1, n+1):
        p = a[n-k]+(x-xData[n-k])*p
    return p
```

The logic this program is in Line 3, Line 4 and Line 5.

²Python block is a group of statement with same level of indentation. A sub-block is a block with an additional indentation.

Line 3: We initialize the polynomial with the coefficient $a[n] = \nabla^n y_n = a_n$.

Line 4: We are going define the polynomial recursively. This takes exactly n steps further. So we use a loop which repeats n times.

Line 5: The value of p and k changes each time Line 5 is executed. Let P_j be the value in p after executing Line 5 with k = j. Then,

```
P_0 = p = a[n]
P_1 = a[n-1] + (x - x_{n-1})P_0
P_2 = a[n-2] + (x - x_{n-2})P_1
\vdots
P_n = a[0] + (x - x_0)P_{n-1}.
```

Clearly, P_n is the unique n degree polynomial given by the Newton's method.

Program 8.2.4. How to interpolate?

```
from numpy import array
xData = array([-2,1,4,-1,3,-4])
yData = array([-1,2,59,4,24,-53])
a = coefficients(xData,yData)
print(interpolate(a,xData,2))
```

You will have to define both the functions (coefficients, interpolate) before doing this.

Line 1 from numpy import array

For defining array objects, we need to import them from numpy library.

Line 2 xData = array([-2,1,4,-1,3,-4])

You can change this line according to the first component of the given data points.

Line 3 yData = array([-1,2,59,4,24,-53])

You can change this line according to the second component of the given data points.

Line 4 a = coefficients(xData,yData)

Call function coefficients and store the array returned into a

Line 5 print(interpolate(a,xData,2))

Call function interpolate to interpolate at x=2 and print the value P(2)

Program 8.2.5 (Just for Fun). We can do more using sympy!

```
from numpy import array}
from sympy import Symbol
xData = array([-2,1,4,-1,3,-4])
yData = array([-1,2,59,4,24,-53])
a = coefficients(xData,yData)
x = Symbol('x')
p = interpolate(a,xData,x)
p.subs({x:2})
```

Remark. Programming Problems

- 1. {(0.15, 4.79867), (2.30, 4.49013), (3.15, 4.2243), (4, 85, 3.47313), (6.25, 2.66674), (7.95, 1.51909)} /?, Example 3.4/
- 2. $\{(0, -0.7854), (0.5, 0.6529), (1, 1.7390), (1.5, 2.2071), (2, 1.9425)\}$??, Problem Set 3.1.5]

8.2.4 Limitations of Polynomial Interpolation

- 1. Inaccuracy The error in interpolation increases as the point moves away from most of the data points.
- Oscilation As the number of data points considered for polynomial interpolation increases, the degree of the polynomial increases. And the graph of the interpolant tend to oscilate excessively. In such cases, the error in interpolation is quite high.
- 3. The best practice is to consider four to six data points nearest to the point of interest and ignore the rest of them.

Remark. The interpolant obtained by joining cubic polynomials corresponding to four nearest data points each, is a cubic spline³.

8.3 Roots of a Function

Definitions 8.3.1. Let $f: \mathbb{R} \to \mathbb{R}$, then $x \in \mathbb{R}$ is a root of f if f(x) = 0.

Remark. Suppose a < b and f(a), f(b) are nonzero and are of different signs. If f is continuous in [a, b], then there is a point $c \in [a, b]$ such that f(c) = 0.

Thus given a < b and f(a), f(b) are nonzero values of different sign, then there may be a bracketed root in [a,b].

Note: There is no guarantee that there exists a root in [a,b] as we are not sure about the continuity of f.

Remark. Given a bracketed root, we can find it using

- 1. Bisection Method or
- 2. Newton-Raphson Method

8.3.1 Bisection Method

Suppose a < b and f(a), f(b) are nonzero values of different signs. We evaluate f(c) where $c = \frac{a+b}{2}$. If f(c) is a nonzero value, then at least one of the pairs f(a), f(c) or f(c), f(b) are of different signs. WLOG suppose that f(a), f(c) are of different signs, then set b = c and $c = \frac{a+b}{2}$. And continue this process until we get sufficiently accurate value of a root.

³Cubic spline is a function, the graph of which is piece-wise cubic

```
Remark. Suppose f(x) = x^5 - 2. Then f(0) = -2, f(1) = -1, f(2) = 30. Since f is known to be continuous, there is a bracketed root in [1,2]. Now f(1.5) > 0 \implies [1,1.5] f(1.25) > 0 \implies [1,1.25] f(1.125) < 0 \implies [1.125,1.25] f(1.1825) > 0 \implies [1.125,1.1825] f(1.15375) > 0 \implies [1.125,1.15375] f(1.139375) < 0 \implies [1.1465625,1.15375] f(1.1465625) < 0 \implies [1.1465625,1.15015625] f(1.148359375) < 0 \implies [1.1483594,1.15015625] f(1.149257825) > 0 \implies [1.1483594,1.14925783]
```

Thus, we have 1.14 is a root of f with accuracy upto two decimal points.

8.3.2 Newton-Raphson Method

Suppose f is differentiable at $x \in \mathbb{R}$ and $f(x) \neq 0$. Then compute $x = x - \frac{f(x)}{df(x)}$ and evaluate f(x). Repeat this process to get more accurate value of a root near x.

Remark. Suppose $f(x) = x^5 - 2$. Then $df(x) = 5x^4$. Let x = 2. Then

$$x = 2 - \frac{30}{80} \implies f(1.625) = 9.330$$

$$x = 1.625 - \frac{9.330}{34.86} \implies f(1.35735) = 2.6074$$

$$x = 1.35735 - \frac{2.6074}{16.9721} \implies f(1.20373) = 0.52733$$

$$x = 1.20373 - \frac{0.52733}{10.4975} \implies f(1.15351) = 0.04224$$

$$x = 1.15351 - \frac{0.042245}{8.85225} \implies f(1.148738) = 0.00034312$$

Thus we have 1.1487 is quite close to a root of f.

8.4 Matrix Operations

Consider a system of n linear, simultaneous equations in n unknowns,

$$A_{11}x_1 + A_{12}x_2 + \dots + A_{1n}x_n = b_1$$

$$A_{21}x_1 + A_{22}x_2 + \dots + A_{2n}x_n = b_2$$

$$\vdots$$

$$A_{n1}x_1 + A_{n2}x_2 + \dots + A_{nn}x_n = b_n$$

We may represent them using matrices as Ax = b. That is,

$$\begin{bmatrix} A_{11} & A_{12} & \cdots & A_{1n} \\ A_{21} & A_{22} & \cdots & A_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ A_{n1} & A_{n2} & \cdots & A_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$

Gauss Elimination and Doolittle Decomposition are two methods to solve system of equaitons, Ax = b.

8.5 Gauss elimination method

Gauss elimination method has of two phases 1. elimination and 2. back substitution. In elimination phase, system Ax = b is transformed into an equivalent system Ux = c where U is an upper-triangular⁴ matrix. And in back substitution phase, Ux = c is solved. Since Ax = b and Ux = c are equivalent, they have the same solution x.

8.5.1 Elimination Phase

We can eliminate unknowns from an equation by adding a scalar multiple of an equation to another equation of the system. In matrices, this is equivalent to adding a scalar multiple of one row to another row, say $R_i \leftarrow R_i + \lambda R_k$.

$$A_{k1}x_1 + A_{k2}x_2 + \dots + A_{kn}x_n = b_k + \lambda (A_{i1}x_1 + A_{i2}x_2 + \dots + A_{in}x_n = b_i)$$

$$(A_{k1} + \lambda A_{i1})x_1 + (A_{k2} + \lambda A_{i2})x_2 + \dots + (A_{kn} + \lambda A_{in})x_n = b_k + \lambda b_i$$

$$\begin{bmatrix} \vdots & \vdots & \vdots & \vdots \\ A_{k1} & A_{k2} & \cdots & A_{kn} \\ \vdots & \vdots & \ddots & \vdots \\ A_{i1} & A_{i2} & \cdots & A_{in} \\ \vdots & \vdots & \vdots & \vdots \end{bmatrix} \xrightarrow{R_i \leftarrow R_i + \lambda R_k} \begin{bmatrix} \vdots & \vdots & \vdots & \vdots \\ A_{k1} & A_{k2} & \cdots & A_{kn} \\ \vdots & \vdots & \ddots & \vdots \\ A_{i1} + \lambda A_{k1} & A_{i2} + \lambda A_{k2} & \cdots & A_{in} + \lambda A_{kn} \\ \vdots & \vdots & \vdots & \vdots & \vdots \end{bmatrix}$$

$$\begin{bmatrix} \vdots \\ b_k \\ \vdots \\ b_i \\ \vdots \end{bmatrix} \xrightarrow{R_i \leftarrow R_i + \lambda R_k} \begin{bmatrix} \vdots \\ b_k \\ \vdots \\ b_i + \lambda b_k \\ \vdots \end{bmatrix}$$

⁴upper triangular - all the entries below the main diagonal are zero. ie $U_{ij} = 0$, if i < j

8.5.2 Back substitution

Let Ux = c be a system of n linear equations in n unknowns and U is an upper triangular matrix. Then we can solve the system of equations from the back.

$$\begin{bmatrix} U_{1,1} & U_{1,2} & \cdots & U_{1,n-1} & U_{1,n} \\ 0 & U_{2,2} & \cdots & U_{2,n-1} & U_{2,n} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & U_{n-1,n-1} & U_{n-1,n} \\ 0 & 0 & \cdots & 0 & U_{n,n} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_{n-1} \\ x_n \end{bmatrix} = \begin{bmatrix} c_1 \\ c_2 \\ \vdots \\ c_{n-1} \\ c_n \end{bmatrix}$$

$$U_{n,n} x_n = c_n \implies x_n = \frac{c_n}{U_{n,n}}$$

$$\sum_{i=n-1}^n U_{n-1,i} x_i = c_{n-1} \implies x_{n-1} = \frac{c_{n-1} - U_{n-1,n} x_n}{U_{n-1,n-1}}$$

 $\sum_{i=1}^{n} U_{1,i} \ x_i = c_1 \implies x_1 = \frac{c_1 - \sum_{i=2}^{n} U_{1,i} x_i}{U_{1,1}}$

8.5.3 Illustrative example

Consider the following system of linear equations,

$$4x_1 - 2x_2 + x_3 = 11$$
$$-2x_1 + 4x_2 - 2x_3 = -16$$
$$x_1 - 2x_2 + 4x_3 = 17$$

We may represent the above system of linear equations using matrices,

$$\begin{bmatrix} 4 & -2 & 1 \\ -2 & 4 & -2 \\ 1 & -2 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 11 \\ -16 \\ 17 \end{bmatrix}$$

Phase 1: Elimination Process

Using eq.1, the unknown x_1 is eliminated from all subsequent equations. An equivalent operation can be performed on both the matrices A and b by adding a suitable scalar multiples of row R_1 to row R_2 and R_3 .

$$\begin{bmatrix} 4 & -2 & 1 \\ -2 & 4 & -2 \\ 1 & -2 & 4 \end{bmatrix} \xrightarrow{R_2 \leftarrow R_2 + 0.5R_1} \begin{bmatrix} 4 & -2 & 1 \\ 0 & 3 & -1.5 \\ 0 & -1.5 & 3.75 \end{bmatrix}$$
$$\begin{bmatrix} 11 \\ -16 \\ 17 \end{bmatrix} \xrightarrow{R_2 \leftarrow R_2 + 0.5R_1} \begin{bmatrix} 11 \\ -10.5 \\ 14.25 \end{bmatrix}$$

And using eq.2, x_2 is eliminated from all subsequent equations (only those rows below it). Again, we perform this by adding suitable scalar multiples of row 2 to row R_3 .

$$\begin{bmatrix} 4 & -2 & 1 \\ 0 & 3 & -1.5 \\ 0 & -1.5 & 3.75 \end{bmatrix} \xrightarrow{R_3 \leftarrow R_3 + 0.5R_2} \begin{bmatrix} 4 & -2 & 1 \\ 0 & 3 & -1.5 \\ 0 & 0 & 3 \end{bmatrix}$$
$$\begin{bmatrix} 11 \\ -16 \\ 17 \end{bmatrix} \xrightarrow{R_3 \leftarrow R_3 + 0.5R_2} \begin{bmatrix} 11 \\ -10.5 \\ 9 \end{bmatrix}$$

The elimination process is complete when all entries below the diagonal elements are reduced to zero. ie, upper triangular matrix.

Phase 2: Back substitution Process

The unknowns are easily found from the equations by solving them in the reverse order. The unknowns are solved from the bottom and solved variables are used to solve the remain unknowns.

$$\begin{bmatrix} 4 & -2 & 1 & 11 \\ 0 & 3 & -1.5 & -10.5 \\ 0 & 0 & 3 & 9 \end{bmatrix} \rightarrow \begin{cases} 4x_1 - 2x_2 + & x_3 = 11 \\ 3x_2 - & 1.5x_3 = -10.5 \\ 3x_3 = 9 \end{cases}$$
$$x_3 = \frac{9}{3} = 3$$
$$x_2 = \frac{-10.5 + 1.5x_3}{3} = -2$$
$$x_1 = \frac{11 - x_3 + 2x_2}{4} = 1$$

Remark. Why don't they use row-reduced echelon matrix of A to simplify the back substitution phase?

This doesn't have much advantage from algorithmic point of view. That is, the time complexity (number of steps for computation) is unaffected. And algorithms always prefer methods even with slight advantage in time or memory. And they won't consider complications in the manual execution of the method. Therefore, programmers won't consider alternate algorithm for the sake of computational simplicity.

8.5.4 Python: Gauss elimination method

Program 8.5.1 (Gauss elimination).

return b

Line 1 from numpy import dot

Imports the "dot()" function for numpy arrays which takes two 'numpy arrays' as input arguments and returns the dot product of them.

Line 2 def gaussElimination(a,b):

it defines "gaussElimination()" as a function which takes two arguments (inputs). First argument is the coefficient matrix A and second argument is the constant matrix b of the linear system of the form Ax = b.

Line $3 \ n = len(b)$

it assigns the length of the list b into variable n which is obviously the number of equations.

Line 4 for k in range (0, n-1):

it is a loop construct. Five instructions following it are part of this loop body, which are executed for each values of k ie, $k = 0, 1, \dots, n - 1$. For each value of k, the unknown x_{k+1} is selected for elimination process.

Line 5 for i in range(k+1,n):

it is a loop inside another loop. Four instructions following it are part of this loop body, which are executed for each values of i, ie, $i = k + 1, k + 2, \dots, n$. This eliminates x_{k+1} from all the equations after the k + 1th equation of the system. Value of i + 1 is the equation⁵ from which x_{k+1} is eliminated.

Line 6 if a[i, k]! = 0.0:

If $a[i,k] = A_{i+1,k+1} \neq 0$, then those three instruction following it are executed. Otherwise, it skips the execution of those three statements. If $x_{k+1} = x[k]$ is not there in the *i*th equation, it doesn't need to be eliminated.

Line 7 lam = a[i,k]/a[k,k]

In this step, λ is computed so that equ.(i+1) - λ equ.(k+1) doesn't have x_{k+1} term in it.

Line 8 $a[i, k+1:n] = a[i, k+1:n] - lam \times a[k, k+1:n]$

Coefficients of (i + 1)th equation are updated.

Equivalent to $a[i, 0:n] = a[i, 0:n] - lam \times a[k, 0:n]$, since zeroes need not be substracted. This is same as equ. $(i+1) \leftarrow \text{equ.}(i+1) - \lambda \text{ equ.}(k+1)$

Line 9 b[i] = b[i] - lam * b[k]

The same row operations are performed on the matrix b instead of using an augmented matrix.

Line 10 for k in range(n-1,-1,-1):

This is another loop construct. The following statement is executed n times for values of $k = n - 1, n - 2, \dots, 0$. Value of k + 1 gives the unknown x_{k+1} which is solved by the back substitution process.

⁵Python starts counting from zero. For example: $A_{11} = a[0,0], x_1 = x[0]$ and $b_1 = b[0]$

Line 11 x[k] = (b[k] - dot(a[k, k+1:n], x[k+1:n]))/a[k, k]

This is the back substitution process. After elimination phase we have k equation in the form $A_{k,k}x_k + A_{k,k+1}x_{k+1} + \cdots + A_{k,n}x_n = b_k$. And we already have values of $x_{k+1}, x_{k+2}, \cdots, x_n$. Then

$$x_k = \frac{b_k - (A_{k,k+1}x_{k+1} + A_{k,k+2}x_{k+2} + \cdots)}{A_{k,k}}$$

This is equivalent to

$$b_k - \begin{bmatrix} A_{k,k+1} & A_{k,k+2} & \cdots & A_{k,n} \end{bmatrix} \begin{bmatrix} x_{k+1} \\ x_{k+2} \\ \vdots \\ x_n \end{bmatrix}$$

$$b_k \leftarrow \frac{A_{k,k}}{A_{k,k}}$$

Remember: The values of x_k are updated into b_k as they are computed. Thus x_k, x_{k+1}, \dots, x_n are stored in b for next back substitution ie, for evaluating x_{k-1} . We start with x_{n-1} , as $x_n = b_n$ is already solved.

Line 12 return b

It returns the new b matrix as output of the "gaussElimination()" function where $x_k = b_k$, $\forall k$.

8.6 LU Decomposition Method: Doolittle

Let Ax = b be a linear system of n equations in n unknowns and let A = LU for some lower triangular matrix L and upper triangular matrix U. Then we have LUx = Ly = b where y = Ux. There are two phases for this method : 1. LU decomposition and 2. substitution.

First, we compute L and U such that A = LU using Gauss elimination. Then We can solve Ly = b using forward substitution process and then solve Ux = y using back substitution process.

For Doolittle decomposition, we prefer to write A as a product LU as shown below:

$$A = \begin{bmatrix} 1 & 0 & \cdots & 0 \\ L_{2,1} & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ L_{n,1} & L_{n,2} & \cdots & 1 \end{bmatrix} \begin{bmatrix} U_{1,1} & U_{1,2} & \cdots & U_{1,n} \\ 0 & U_{2,2} & \cdots & U_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & U_{n,n} \end{bmatrix}$$

$$A = \begin{bmatrix} U_{1,1} & U_{1,2} & \cdots & U_{1,n} \\ L_{2,1}U_{1,1} & L_{2,1}U_{1,2} + U_{2,2} & \cdots & L_{2,1}U_{1,n} + U_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ L_{n,1}U_{1,1} & L_{n,1}U_{1,2} + L_{n,2}U_{2,2} & \cdots & \sum_{k=1}^{n-1} L_{n,k}U_{k,n} + U_{n,n} \end{bmatrix}$$

Note that in Doolittle's decomposition method, the diagonal entries of the lower triangular matrix L are all 1. ie, $L_{ii} = 1$, $\forall i$. Thus, we can use an $n \times n$ matrix to represent both L and U by overwriting trivial entries (zeroes and ones) of both the matrices. And this matrix is represented by $[L \setminus U]$.

$$[L \backslash U] = \begin{bmatrix} U_{1,1} & U_{1,2} & \cdots & U_{1,n} \\ L_{2,1} & U_{2,2} & \cdots & U_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ L_{n,1} & L_{n,2} & \cdots & U_{n,n} \end{bmatrix}$$

is the combined matrix made from both the triangular matrices L and U.

The triangular matrices L and U such that LU = A can be computed the variables in the Gauss elimination method.

$$A = \begin{bmatrix} 1 & 0 & \cdots & 0 \\ L_{2,1} & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ L_{n,1} & L_{n,2} & \cdots & 1 \end{bmatrix} \begin{bmatrix} U_{1,1} & U_{1,2} & \cdots & U_{1,n} \\ 0 & U_{2,2} & \cdots & U_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & U_{n,n} \end{bmatrix}$$

We can break down this matrix multiplication into the following row operations on the rows of the upper triangular matrix^7

$$U_{R1} \leftarrow U_{R1}$$

$$U_{R2} \leftarrow L_{2,1} \cdot U_{R1} + U_{R2}$$

$$\cdots$$

$$U_{Rn} \leftarrow L_{n,1} \cdot U_{R1} + L_{n,2} \cdot U_{R2} + \cdots + L_{n,n-1} \cdot U_{R(n-1)} + U_{Rn}$$

Clearly, λ we use to eliminate x_k from row i are $L_{i,k}$. And the matrix obtained after Gauss elimination is the upper triangular matrix U.

8.6.1 Illustrative example

Solve
$$\begin{bmatrix} -3 & 6 & -4 \\ 9 & -8 & 24 \\ -12 & 24 & -26 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} -3 \\ 65 \\ -42 \end{bmatrix}$$

Phase 1: LU Decomposition

Suppose, we have a system of three linear equations, then

$$A = LU = \begin{bmatrix} 1 & 0 & 0 \\ L_{21} & 1 & 0 \\ L_{31} & L_{32} & 1 \end{bmatrix} \begin{bmatrix} U_{1,1} & U_{1,2} & U_{1,3} \\ 0 & U_{2,2} & U_{2,3} \\ 0 & 0 & U_{3,3} \end{bmatrix}$$

$$A = \begin{bmatrix} U_{1,1} & U_{1,2} & U_{1,3} \\ L_{2,1}U_{1,1} & L_{2,1}U_{1,2} + U_{2,2} & L_{2,1}U_{1,3} + U_{2,3} \\ L_{3,1}U_{1,1} & L_{3,1}U_{1,2} + L_{3,2}U_{2,2} & L_{3,1}U_{1,3} + L_{3,1}U_{2,3} + U_{3,3} \end{bmatrix}$$

 $^{^6}$ algorithmic implementation all decomposition algorithms prefer to use a combined matrix $\overline{}$

 $^{^{7}}U_{Rk}$: kth row of the matrix U

We can compute L and U using the Gauss elimination process⁸. The matrix obtained after Gauss elimination on A is U and the values of the variable lam used in Gauss elimination are the entries in L. That is, in order to eliminate x_k from row i, we use $lam = L_{i,k}$.

Given,
$$A = \begin{bmatrix} -3 & 6 & -4\\ 9 & -8 & 24\\ -12 & 24 & -26 \end{bmatrix}$$

$$\begin{bmatrix} -3 & 6 & -4 \\ 9 & -8 & 24 \\ -12 & 24 & -26 \end{bmatrix} \xrightarrow{R_2 \leftarrow R_2 + 3R_1 \atop R_3 \leftarrow R_3 - 4R_1} \begin{bmatrix} -3 & 6 & -4 \\ 0 & 10 & 12 \\ 0 & 0 & -10 \end{bmatrix} \implies L_{2,1} = 3, \ L_{3,1} = -4$$

We store these non-trivial entries of L into A itself. That is, $A_{2,1} = L_{2,1}, \ A_{3,1} = L_{3,1}$.

"In this case, $A_{3,2}$ became zero (this is not a trivial zero yet), and we won't eliminate x_2 from row 3 to save computation time. Thus, we are not computing $L_{3,2}=0$ or storing it. However, the variable representing $L_{3,2}$ is $A_{3,2}$, which is already zero after Gauss elimination and we are quite happy with that."

Clearly, $L_{3,2} = 0$. Therefore, we have

$$U = \begin{bmatrix} -3 & 6 & -4 \\ 0 & 10 & 12 \\ 0 & 0 & -10 \end{bmatrix}, \ L = \begin{bmatrix} 1 & 0 & 0 \\ 3 & 1 & 0 \\ -4 & 0 & 1 \end{bmatrix}$$

Since we are already stored those two non-trivial entries of L into A. We get,

$$[L \backslash U] = \begin{bmatrix} -3 & 6 & -4 \\ 3 & 10 & 12 \\ -4 & 0 & -10 \end{bmatrix}$$

Phase 2: Substitution

Suppose Ly = b,

$$\begin{bmatrix} 1 & 0 & 0 \\ L_{2,1} & 1 & 0 \\ L_{3,1} & L_{3,2} & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix} \rightarrow \begin{cases} y_1 & = b_1 \\ L_{2,1}y_1 + y_2 & = b_2 \\ L_{3,1}y_1 + L_{3,2}y_2 + y_3 & = b_3 \end{cases}$$

Now, we can find the values of y_k and store them into the matrix b itself.

$$\begin{aligned} b_1 &\leftarrow y_1 = b_1 \\ b_2 &\leftarrow y_2 = b_2 - \begin{bmatrix} L_{2,1} \end{bmatrix} \begin{bmatrix} b_1 \end{bmatrix}, \text{ since } b_1 = y_1 \\ b_3 &\leftarrow y_3 = b_3 - \begin{bmatrix} L_{3,1} & L_{3,2} \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}, \text{ since } b_1 = y_1, \ b_2 = y_2 \end{aligned}$$

⁸We usually need a proof for such a strong statement. In this paper, they are more focussed on the application side and therefore we will don't present any vigorous proof.

In general,

$$b_k \leftarrow y_k = b_k - \begin{bmatrix} L_{k,1} & L_{k,2} & \cdots & L_{k,k-1} \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_k - 1 \end{bmatrix}$$
, since $b_j = y_j$, $j = 1, 2, \cdots, (k-1)$

We have $A = LU \implies LUx = b$. Suppose Ux = y, then we get Ly = b. First of all, we will solve Ly = b using forward substitution.

$$\begin{bmatrix} 1 & 0 & 0 \\ 3 & 1 & 0 \\ -4 & 0 & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} -3 \\ 65 \\ -42 \end{bmatrix} \rightarrow \begin{cases} y_1 & = -3 \\ 3y_1 + y_2 & = 65 \\ -4y_1 + y_3 & = -42 \end{cases}$$

$$y_1 = -3$$

 $y_2 = 65 - 3y_1 = 74$
 $y_3 = -42 + 4y_1 = 54$

Suppose Ux = y,

$$\begin{bmatrix} U_{1,1} & U_{1,2} & U_{1,3} \\ 0 & U_{2,2} & U_{2,3} \\ 0 & 0 & U_{3,3} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix}$$

Now, we can find the values of x_k and store them into the matrix y itself.

$$y_{3} \leftarrow x_{3} = \frac{y_{3}}{U_{3,3}}$$

$$y_{2} \leftarrow x_{2} = \frac{y_{2} - [y_{3}][U_{2,3}]}{U_{2,2}}, \text{ since } y_{3} = x_{3}$$

$$y_{1} \leftarrow x_{1} = \frac{y_{1} - \begin{bmatrix} y_{2} \\ y_{3} \end{bmatrix}[U_{1,2} \quad U_{1,3}]}{U_{1,1}}, \text{ since } y_{2} = x_{2}, y_{3} = x_{3}$$

In general,

$$x_{k} = \frac{y_{k+1}}{y_{k+2}} \begin{bmatrix} y_{k+1} & & & & & & & & \\ y_{k+2} & & & & & & & \\ \vdots & & & & & & & \\ y_{n} & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & \\ & & & \\ & &$$

8.6.2 Python: Doolittle's LU Decomposition method

```
Program 8.6.1.

from numpy import dot

def LUdecomposition(a):

n = len(a)

for k in range(0,n-1):

for i in range(k+1,n):
```

```
if a[i,k] != 0.0:
    lam = a[i,k]/a[k,k]
    a[i,k+1:n] = a[i,k+1:n] - lam*a[k,k+1:n]
    a[i,k] = lam
    return a

def LUsolve(a,b):
    n = len(a)
    for k in range(1,n):
        b[k] = b[k]-dot(a[k,0:k],b[0:k])
    b[n-1] = b[n-1]/a[n-1,n-1]
    for k in range(n-2,-1,-1):
        b[k] = (b[k] - dot(a[k,k+1:n],b[k+1:n]))/a[k,k]
    return b
```

This program mainly uses the Gauss elimination algorithm. Thus, the explanation for Lines 3-8 are not repeated here.

But remember the loop at Line 4 has inner loop at Line 5 and Line 7-9 are at same level of indentation which means they all are either executed or skipped depending on the truthness of the condition in Line 6. And Line 6-9 are executed for each instance of inner loop. Again, Line 5-9 are executed for each instance of the outer loop.

This time the gaussElimination() function which you have seen earlier is split into two functions 1. LUdecomposition() and 2. LUsolve(). And forward substitution is also added to LUsolve().

Line 2 def LUdecomposition(a):

LUdecomposition(A) computes L and U such that A = LU and combine both triangular matrices into a single matrix [L/U], by over-writting their trivial entries. And returns this combined matrix.

Line 9 a[i, k] = lam

Clearly, lam used for eliminating x_k from row i, $\lambda_{i,k} = a[i,k]/a[k,k] = L[i,k]$, $\forall k, \forall i$, (i > k). Also a[i,k] which is reduced zero by Gauss elimination process is not used anymore⁹ in Gauss elimination process. Thus L[i,k] can stored at a[i,k] straight away. And U[i,j], $j \leq i$ are already the entries of the matrix obtained from Gauss elimination. Thus for each iterations of k, the matrix a is updated (k+1th row and k+1th column) with respective entries of the combined matrix $[L \setminus U]$.

Line 10 return a

Matrix a is already $[L \setminus U]$, and thus LUdecomposition(A) returns $[L \setminus U]$ such that A = LU.

Line 11 def LUsolve(a,b):

LUsolve() function does both forward substitution and back substitution. Suppose Ax = b is the system to be solved. Then the inputs of LUsolve() are $a = [L \setminus U]$ where A = LU.

 $^{{}^9}A[i,k]$ is not used after elimination of x_k from row i - It turns out that the trivial zeroes which are ignored on the row operations in Gauss elimination not only save time, but also provide a variable to store our intermediate result $L_{i,k}$ in Doolittle method.

Line 12 n = len(a)

We have to compute this again as this function starts fresh and thus value of the variable n from LUdecomposition() is lost.

Line 13 for k in range(1,n):

This is the loop for forward substitution.

Line 14 b[k] = b[k] - dot(a[k, 0:k], b[0:k])updating b_{k^*} with y_{k^*} such that Ly = b where $k^* = k - 1$.

$$b_k \leftarrow b_k - \begin{bmatrix} L_{k,1} & L_{k,2} & \cdots & L_{k,k-1} \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_k - 1 \end{bmatrix}$$

Line 15 b[n-1] = b[n-1]/a[n-1, n-1]Computing¹⁰ y_n and storing it into b_n .

$$b_n \leftarrow \frac{b_n}{U_{n,n}}$$

Line 16 for k in range(n-2,-1,-1):

This is the loop for back substitution.

Line 17 b[k] = (b[k] - dot(a[k, k+1:n], b[k+1:n]))/a[k, k] updating b_{k^*} with x_{k^*} such that Ux = y where $k^* = k-1$.

$$b_k - \begin{bmatrix} x_{k+1} \\ x_{k+2} \\ \vdots \\ x_n \end{bmatrix} \begin{bmatrix} U_{k,k+1} & U_{k,k+2} & \cdots & U_{k,n} \end{bmatrix}$$

$$U_{k,k}$$

Line 18 return b

LUsolve($[L \setminus U], b$) returns b where $b[i+1] = x_i$.

Programmer's Tip : There are few things to remember when spliting a function into two functions.

- 1. These functions are completely independent of one another.
- 2. Variables defined inside a function are not available outside.
- 3. The best way to give/take data to/from a function is through arguements/return-value

Beginner's Tip: In any programming language, we reuse variable. Thus, same variable may represent different values at different points of time. In Doolittle LU Decomposition, the variable 'a' initially represent matrix A, this variable is passed into LUdecomposition() function. In that function, A is changed to $[L\backslash U]$ in a step-by-step fashion. The value of 'a' in updated in step 7, 8 and 9. This is bit hard to imagine this transition of 'a' from A to $[L\backslash U]$ for a beginner at programming. Simiarly, in LUsolve() function, the variable 'b' changes from matrix b to matrix y, and then to matrix x.

 $^{^{10}\}mathrm{Mathematically},$ you can define dot product of empty matrices as zero, but nympy dot function can't handle such a situation. Therefore, we have to do this step separately.

8.7 Numerical Integraion

Numerical integration/Quadrature is the numerical approximation of $\int_a^b f(x)dx$ by $\sum_{i=0}^n A_i f(x_i)$ where x_i are nodal abscissas, and A_i are weights. There are two methods to determine these nodal abscissas and suitable weights so that the sum is sufficiently accurate to the value of the integral.

- 1. Newton-Cotes forumulas
- 2. Gauss quadrature

Newton-Cotes formulas are useful when f(x) can be evaluated without much computation. And using those values f(x) can be interpolated to a piecewise-polynomial function. Then using equally spaces nodal abscissas and suitable weights $\int_a^b f(x)dx$ can be numerically approximated.

Gauss quadrature rules require lesser evalutions of f. And therefore are quite useful when evaluation of f(x) has much computational complexity. Also, this method can manage integrable singularities where as Newton-Cote formulas can't numerically integrate function with singularities.

8.7.1 Newton-Cotes formulas

We divide the interval of integral (a, b) into n subintervals of equal length, ie, h = (b - a)/n. Let $x_0 = a, x_1, x_2, \dots, x_{n-1}, x_n = b$ be the end points of these subintervals. Then we can find an n degree polymonial interpolant satisfying f at those points, using Lagrange's method.

Polynomial,
$$P(x) = \sum_{i=0}^{n} f(x_i) l_i(x)$$
 where $l_i(x) = \prod_{\substack{j=0 \ j \neq i}}^{n} \frac{x - x_i}{x_i - x_j}$

Thus the integral $I = \int_a^b f(x) dx$ can be numerically evaluated as follows:

$$I = \int_a^b P_n(x)dx = \sum_{i=0}^n \left(f(x_i) \int_a^b l_i(x)dx \right)$$
$$= \sum_{i=0}^n A_i f(x_i), \text{ where } A_i = \int_a^b l_i(x)dx$$

The simplest cases of Newton-Cotes formulas are when n = 1, 2, and 3

Trapezoidal rule $n=1 \implies A_0 = \frac{h}{2}, A_1 = \frac{h}{2}$ and

$$\int_{a}^{b} f(x)dx = A_0 f(x_0) + A_1 f(x_1) = \frac{h}{2} (f(a) + f(b))$$

Simpson's 1/3 rule $n=2 \implies A_0 = \frac{h}{3}, A_1 = \frac{4h}{3}, A_2 = \frac{h}{3}$ and

$$\int_{a}^{b} f(x)dx = \sum_{i=0}^{2} A_{i}f(x_{i}) = \frac{h}{3} \left(f(a) + 4f\left(\frac{a+b}{2}\right) + f(b) \right)$$

Simpson's 3/8 **rule** $n=3 \implies A_0 = \frac{3h}{8}, \ A_1 = \frac{9h}{8}, \ A_2 = \frac{9h}{8}, \ A_3 = \frac{3h}{8}$ and

$$\int_{a}^{b} f(x)dx = \sum_{i=0}^{3} A_{i}f(x_{i}) = \frac{3h}{8}(f(a) + 3f(a+h) + 3f(a+2h) + f(b))$$

Trapezoidal Rule: n=1

Consider interval (a, b). Since n = 1, we have $x_0 = a$ and $x_1 = b$.

$$l_0(x) = \frac{x - x_1}{x_0 - x_1}$$
$$l_1(x) = \frac{x - x_0}{x_1 - x_0}$$

$$A_0 = \int_a^b l_0(x)dx = \int_a^b \frac{x - x_1}{x_0 - x_1} dx = \frac{-1}{h} \int_a^b (x - b)dx$$

$$= \frac{-1}{h} \left(\frac{(x - b)^2}{2}\right)_a^b = \frac{-1}{h} \left(\frac{0 - (a - b)^2}{2}\right) = \frac{h}{2}$$

$$A_1 = \int_a^b l_1(x)dx = \int_a^b \frac{x - x_0}{x_1 - x_0} dx = \frac{1}{h} \int_a^b (x - a)dx$$

$$= \frac{1}{h} \left(\frac{(x - a)^2}{2}\right)_a^b = \frac{1}{h} \left(\frac{(b - a)^2 - 0}{2}\right) = \frac{h}{2}$$

Therefore,

$$\int_{a}^{b} f(x)dx = A_0 f(x_0) + A_1 f(x_1) = \frac{h}{2} (f(a) + f(b))$$

Simpon's 1/3 Rule : n=2

Consider interval (a,b) divided into two subintervals of equal length $h=\frac{a+b}{2}$. We have $x_0=a,\ x_1=\frac{a+b}{2},\ x_2=b$.

$$l_0(x) = \frac{x - x_1}{x_0 - x_1} \frac{x - x_2}{x_0 - x_2}$$
$$l_1(x) = \frac{x - x_0}{x_1 - x_0} \frac{x - x_2}{x_1 - x_2}$$
$$l_2(x) = \frac{x - x_0}{x_2 - x_0} \frac{x - x_1}{x_2 - x_1}$$

$$A_0 = \int_a^b l_0(x)dx$$

$$= \int_a^b \frac{x - x_1}{x_0 - x_1} \frac{x - x_2}{x_0 - x_2} dx$$

$$= \int_a^b \left(\frac{x - \frac{a+b}{2}}{a - \frac{a+b}{2}}\right) \left(\frac{x - b}{a - b}\right) dx$$

Changing variable of integration $y = x - \frac{a+b}{2}$

$$y = x - \frac{a+b}{2} \implies dy = dx$$

 $x = a \implies y = -h$
 $x = b \implies y = h$

Continuing with the value of A_0 ,

$$A_{0} = \int_{-h}^{h} \left(\frac{y}{-h}\right) \left(\frac{y-h}{-2h}\right) dy$$

$$= \frac{1}{2h^{2}} \int_{-h}^{h} y^{2} - \frac{1}{2h} \int_{-h}^{h} y dy$$

$$= \frac{1}{2h^{2}} \left(\frac{h^{3}}{3} - \frac{(-h)^{3}}{3}\right) - \frac{1}{2h} \left(\frac{h^{2}}{2} - \frac{(-h)^{2}}{2}\right)$$

$$= \frac{h}{3}$$

$$A_1 = \int_a^b l_1(x)dx$$

$$= \int_a^b \frac{x - x_0}{x_1 - x_0} \frac{x - x_2}{x_1 - x_2} dx$$

$$= \int_a^b \left(\frac{x - a}{h}\right) \left(\frac{x - b}{-h}\right) dx$$

Applying change of variable, $y = x - \frac{a+b}{2}$

$$= \int_{-h}^{h} \left(\frac{y+h}{h}\right) \left(\frac{y-h}{-h}\right) dy$$

$$= \frac{-1}{h^2} \int_{-h}^{h} y^2 dy + \int_{-h}^{h} 1 dy$$

$$= \frac{-1}{h^2} \left(\frac{h^3}{3} - \frac{(-h)^3}{3}\right) + (h - (-h))$$

$$= \frac{-2h^3}{3h^2} + 2h$$

$$= \frac{4h}{2}$$

$$A_2 = \int_a^b l_2(x)dx$$

$$= \int_a^b \left(\frac{x - x_0}{x_2 - x_0}\right) \left(\frac{x - x_1}{x_2 - x_1}\right) dx$$

$$= \int_a^b \left(\frac{x - a}{2h}\right) \left(\frac{x - \frac{a + b}{2}}{h}\right) dx$$

Applying change of variable, $y = x - \frac{a+b}{2}$

$$\begin{split} &= \int_{-h}^{h} \left(\frac{y+h}{2h} \right) \left(\frac{y}{h} \right) dy \\ &= \frac{1}{2h^2} \int_{-h}^{h} y^2 dy + \frac{1}{2h} \int_{-h}^{h} y dy \\ &= \frac{1}{2h^2} \left(\frac{h^3}{3} - \frac{(-h)^3}{3} \right) + \frac{1}{2h} \left(\frac{h^2}{2} - \frac{(-h)^2}{2} \right) \\ &= \frac{2h^3}{6h^2} \\ &= \frac{h}{3} \end{split}$$

Therefore,

$$\int_{a}^{b} f(x)dx = \sum_{i=0}^{2} A_{i}f(x_{i}) = \frac{h}{3} \left(f(a) + 4f\left(\frac{a+b}{2}\right) + f(b) \right)$$

8.7.2 Composite Trapezoidal Rule

Suppose an interval (a, b) is divided into n subintervals. In Composite Trapezoidal Rule, Trapezoidal Rule is applied to each subinterval. Thus we have,

$$\begin{split} I &= I_0 + I_1 + \dots + I_{n-1} \text{ where } I_k \text{ is the integral over } (x_k, x_{k+1}) \\ &= \frac{h(f(x_0) + f(x_1))}{2} + \frac{h(f(x_1) + f(x_2))}{2} + \dots + \frac{h(f(x_{n-1} + f(x_n)))}{2} \\ &= \frac{h}{2} (f(x_0) + 2f(x_1) + \dots + 2f(x_{n-1}) + f(x_n)) \end{split}$$

8.7.3 Recursive Trapezoidal Rule

Suppose an interval (a,b) is divided into 2^{k-1} subintervals. In Recursive Trapezoidal Rule, we apply Trapezoidal Rule on each subinterval. And there is a recursive forumula since the number of intervals doubles as value of k increases by one. Let H=b-a

$$k = 1 \implies 2^0 = 1$$
 and $I_1 = \frac{H}{2}(f(a) + f(b))$

$$k = 2 \implies 2^1 = 2 \text{ and}$$

$$I_2 = \frac{H}{4}(f(a) + 2f\left(a + \frac{H}{2}\right) + f(b)$$

$$= \frac{H}{4}(f(a) + f(b)) + \frac{H}{2}f\left(a + \frac{H}{2}\right)$$

$$= \frac{I_1}{2} + \frac{H}{2}f\left(a + \frac{H}{2}\right)$$

$$\begin{split} k &= 3 \implies 2^2 = 4 \text{ and} \\ I_3 &= \frac{H}{8} (f(a) + 2f\left(a + \frac{H}{4}\right) + 2f\left(a + \frac{2H}{4}\right) + 2f\left(a + \frac{3H}{4}\right) + f(b)) \\ &= \frac{H}{8} (f(a) + 2f\left(a + \frac{2H}{4}\right) + 2f\left(a + \frac{4H}{4}\right) + 2f\left(a + \frac{6H}{4}\right) + f(b)) \\ &+ \frac{H}{4} (f\left(a + \frac{H}{4}\right) + f\left(a + \frac{3H}{4}\right)) \\ &= \frac{I_2}{2} + \frac{H}{4} (f\left(a + \frac{H}{4}\right) + f\left(a + \frac{3H}{4}\right)) \end{split}$$

Consider interval (0,64). We have b-a=H=64.

```
I_1 = 32(f(0) + f(64))
I_2 = 16(f(0) + 2f(32) + f(64))
I_3 = 8(f(0) + 2f(16) + 2f(32) + 2f(48) + f(64))
I_4 = 4(f(0) + 2f(8) + 2f(16) + \dots + 2f(56) + f(64)
I_5 = 2(f(0) + 2f(4) + 2f(8) + \dots + 2f(60) + f(64)
I_6 = f(0) + 2f(2) + 2f(4) + \dots + 2f(62) + f(64)
```

The values corresponding to the intervals in I_{k-1} appear in I_k as alternate terms. Other terms, corresponds to the odd multiples of $\frac{H}{2^k}$. We separate them into two sums and represent the first sum as $\frac{I_{k-1}}{2}$.

Clearly,
$$I_k = \frac{H}{2^k} \sum_{i=0}^{2^{k-2}} f\left(a + \frac{2iH}{2^k}\right) + \frac{2H}{2^k} \sum_{i=1}^{2^{k-2}} f\left(a + \frac{(2i-1)H}{2^{k-1}}\right)$$
$$= \frac{I_{k-1}}{2} + \frac{H}{2^{k-1}} \sum_{i=1}^{2^{k-2}} f\left(a + \frac{(2i-1)H}{2^{k-1}}\right)$$

Python: Recursive Trapezoidal Rule

```
Program 8.7.1.
```

```
def recursiveTrapezoidalRule(f,a,b,Iold,k):
    if k == 1 :
    Inew = (f(a)+f(b))*(b-a)/2.0
    n = 2**(k-2)
    h = (b-a) * 1.0 / n

x = a + h/2.0
    sum = 0.0
    for i in range(n):
      sum = sum + f(x)
    Inew = (Iold + h * sum) / 2.0
  return Inew
```

Line 1 def recursiveTrapezoidalRule(f,a,b,Iold,k):

This function has five input arguments. (a) f is a real function (b) a,b are start and end of interval in which f is going to be integrated (c) Iold is

the value of the integral for 2^{k-1} subintevals using recursive trapezoidal method (d) k is a variable such that 2^k is the number of subintervals considered for Integration.

Line 2 if k == 1:

If k = 1, we proceed to Line 3, otherwise we go to Line 4.

Line 3 Inew = (f(a) + f(b)) * (b - a)/2.0

For k = 1, we use trapezoidal rule $I_1 = \frac{b-a}{2}(f(a) + f(b))$. When writing this in python, we are use 2.0 so that the python won't ignore the decimal part of this fraction. In python, 5/2 = 2. And 5/2.0 = 2.5

Line 4 else:

If Line 2 is false, (ie $k \neq 1$) python executes Line 5-8. These line implements the recursive formula for I_k .

- Line 5 n = 2 * *(k-2)Equivalent to $n \leftarrow 2^{k-2}$.
- Line 6 h = (b a) * 1.0/n

This the length of a subinterval when we divide (a,b) into 2^k subinterval-s/panels. Equivalent to $h \leftarrow \frac{b-a}{2^k-2}$

Line 7 x = a + h/2.0

This the parameter of f in the first term in the sum $\sum_{i=1}^{2^{k-2}} f\left(a + \frac{(2i-1)H}{2^{k-1}}\right)$ in the recursive formula for I_k . Equivalent to $x \leftarrow a + \frac{h}{2}$.

Line 8 sum = 0.0

We are going to use this variable to find that sum. To start with, we will make it 0 and will add each term to it one-by-one. Equivalent to $sum \leftarrow 0$.

Line 9 for i in range(n):

This the variable i in the recursive forumula for I_k . For each value of $i = 1, 2, \dots, 2^{k-2}$, Lines 10 and 11 are executed. That is, for each value of i, the corresponding term in the sum is computed and added to the variable sum.

Line 10 sum = sum + f(x)

Value of f at x is computed and added to the partial sum. Equivalent to $sum \leftarrow sum + f(x)$. However, the value of x is changed for each i in Line 11. Thus, for next value of i, x and sum have the new values to use.

Line 11 x = x + h

Equivalent to $x \leftarrow x + h$. For i = 1, $x = a + \frac{h}{2}$ before Line 11. At Line 11, $x \leftarrow a + \frac{3h}{2}$. And this is the value of x for i = 2 before Line 11 next time. Thus, x iterates through $a + \frac{h}{2}, a + \frac{3h}{2}, \cdots, a + \frac{(2n-1)h}{2}$. This x is updated and used for next execution of Line 10 and 11.

Line 12 Inew = (Iold + h * sum)/2.0

We reach here only after executing Line 10-11 for all values of i. That is, sum in the recursive forumula is already computed. This line, implements the recursive formula and stores that value into the variable Inew. Equivalent to $Inew \leftarrow \frac{Iold + h \times sum}{2}$.

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Line 13 return Inew

It returns the value of I_k , the intergral of f over (a,b) using recursive trapezoidal rule for 2^k subintervals.

ME010204 Complex Analysis

ME010205 Measure & Integration

Semester III

ME010301 Advanced Complex Analysis

ME010302 Partial Differential Equations

ME010303 Multivariate Calculus & Integral Transforms

13.1 Integral Transforms

13.1.1 The Weierstrass Approximation Theorem

Every continuous, real valued function on a compact interval has a polynomial approximation.[?, Theorem 11.17]

Theorem 13.1.1 (Weierstrass). Let f be a real-valued, continuous function on a compact interval [a,b]. Then for every $\epsilon > 0$, there is a polynomial p such that $|f(x) - p(x)| < \epsilon$ for every $x \in [a,b]$.

Synopsis. Given a real-valued continuous function on compact interval [a,b], we can construct a real-valued, continuous function g on \mathbb{R} which is periodic with period 2π . We have, if $f \in L(I)$ and f is bounded almost everywhere in I, then $f \in L^2(I)$. [?, Theorem 10.52]. By Fejer's theorem ([?, Theorem 11.15]), the fourier series generated by g ([?, definition 11.3]) converges to the Cesaro sum ([?, Definition 8.47]), which is g itself in this case. Thus for any $\epsilon > 0$, there is a finite sum of trignometric functions. The power series expansions of trignometric functions ([?, definition 9.27]) being uniformly convergent, there exists a polynomial p_m which approximates g. And we can construct p (polynomial approximation of g) using p_m .

Proof. Define $g: \mathbb{R} \to \mathbb{R}$,

$$g(t) = \begin{cases} f(a + (b - a)t/\pi), & t \in [0, \pi) \\ f(a + (2\pi - t)(b - a)/\pi), & t \in [\pi, 2\pi] \\ g(t - 2n\pi), & t > 2\pi, & n \in \mathbb{N} \\ g(t + 2n\pi), & t < 0, & n \in \mathbb{N} \end{cases}$$

Thus g is a continuous, real-valued, periodic function with period 2π such that

$$f(x) = g\left(\frac{\pi(x-a)}{b-a}\right), \ x \in [a,b]$$
(13.1)

The fourier series generated by g is given by,

$$g(t) \sim \frac{a_0}{2} + \sum_{k=1}^{\infty} \left(a_k \cos kt + b_k \sin kt \right)$$

where
$$a_k = \frac{1}{\pi} \int_0^{2\pi} f(t) \cos kt \ dt$$
, $b_k = \frac{1}{\pi} \int_0^{2\pi} f(t) \sin kt \ dt$

Let $\{s_n(t)\}$ be the sequence of partial sums of the fourier series generated by g. And $\{\sigma_n(t)\}$ be the sequence of averages of $s_n(t)$ given by,

$$\sigma_n(t) = \frac{1}{n} \sum_{k=1}^n s_k(t)$$
, where $s_k(t) = \frac{a_0}{2} + \sum_{j=1}^k (a_j \cos jt + b_j \sin jt)$

Function $f \in L(I)$ being real-valued continuous function on a compact interval, it is bounded and hence is Lebesgue square integrable.ie, $f \in L^2(I)$. Thus, $g \in L^2(I)$.

Since g is continuous on \mathbb{R} , the function $s: \mathbb{R} \to \mathbb{R}$ defined by,

$$s(t) = \lim_{h \to 0^+} \frac{g(t+h) - g(t-h)}{2}$$

is well-defined on \mathbb{R} and $s(t) = g(t), \ \forall t \in \mathbb{R}$.

Then by Fejer's Theorem, the sequence $\{\sigma_n(t)\}$ converges uniformly to g(t) for every $t \in \mathbb{R}$. Thus, given $\epsilon > 0$, there exists $N \in \mathbb{N}$ such that $\forall t \in \mathbb{R}$, $|g(t) - \sigma_N(t)| < \frac{\epsilon}{2}$.

We have,

$$\sigma_N(t) = \sum_{k=0}^{N} (A_k \cos kt + B_k \sin kt), \text{ where } A_k, B_k \in \mathbb{R}$$
 (13.2)

By the power series expansion of the trignometric functions about origin,

$$\cos kt = \sum_{j=1}^{\infty} \left(\frac{\cos^{(j)} 0}{j!} (kt)^j \right) = \sum_{j=1}^{\infty} A'_j t^j \text{ where } A'_j \in \mathbb{R}$$
 (13.3)

$$\sin kt = \sum_{j=1}^{\infty} \left(\frac{\sin^{(j)} 0}{j!} (kt)^j \right) = \sum_{j=1}^{\infty} B_j' t^j \text{ where } B_j' \in \mathbb{R}$$
 (13.4)

Since the above power series expansions of trignometric functions are uniformly convergent, their finite linear combination $\{\sigma_N(t)\}$ is also uniformly convergent. ie, Given $\epsilon > 0$ there exists $m \in \mathbb{N}$ such that for every $t \in \mathbb{R}$

$$\left| \sum_{k=0}^{m} C_k t^k - \sigma_N(t) \right| < \frac{\epsilon}{2} \text{ where } C_k \in \mathbb{R}$$

Therefore, $|p_m(t)-g(t)| \leq |p_m(t)-\sigma_N(t)|+|\sigma_N(t)-g(t)| < \epsilon$ where $p_m(t)=\sum_{k=0}^m C_k t^k$. Define $p:[a,b]\to\mathbb{R}$ by,

$$p(x) = p_m \left(\frac{\pi(x-a)}{b-a}\right) \tag{13.5}$$

By equations ?? and ??, $|p(x) - f(x)| < \epsilon$ for every $x \in [a, b]$.

13.1.2 Other Forms of Fourier Series

Let $f \in L([0, 2\pi])$, then the fourier series generated by f is given by,

$$f(x) \sim \frac{a_0}{2} + \sum_{n=1}^{\infty} \left(a_n \cos nx + b_n \sin nx \right)$$

where
$$a_n = \frac{1}{\pi} \int_0^{2\pi} f(t) \cos nt \ dt, \qquad b_n = \frac{1}{\pi} \int_0^{2\pi} f(t) \sin nt \ dt$$

By Euler's forumula $e^{inx}=\cos nx+i\sin nx$. We have, $\cos nx=\frac{(e^{inx}+e^{-inx})}{2}$ and $\sin nx=\frac{(e^{inx}-e^{-inx})}{2i}$

$$f(x) \sim \frac{a_0}{2} + \sum_{n=1}^{\infty} \left(\alpha_n e^{inx} + \beta_n e^{-inx} \right)$$

where
$$\alpha_n = \frac{(a_n - ib_n)}{2}$$
 $\beta_n = \frac{(a_n + ib_n)}{2}$

Therefore, by assigning $\alpha_0 = a_0/2$, $\alpha_{-n} = \beta_n$, we get the following exponential form of fourier series generated by f,

$$f(x) \sim \sum_{n=-\infty}^{\infty} \alpha_n e^{inx}$$
 where $\alpha_n = \frac{1}{2\pi} \int_0^{2\pi} f(t) e^{-int} dt$

Note: If f is periodic with period 2π , then the interval of integration $[0, 2\pi]$ can be replaced with any interval of length 2π . eg. $[-\pi, \pi]$

Periodic with period p

Let $f \in L([0,p])$ and f is periodic with period p. Then

$$f(x) \sim \frac{a_0}{2} + \sum_{n=1}^{\infty} \left(a_n \cos \frac{2\pi nx}{p} + b_n \sin \frac{2\pi nx}{p} \right)$$

where
$$a_n = \frac{2}{p} \int_0^p f(t) \cos \frac{2\pi nt}{p} dt$$
 $b_n = \frac{2}{p} \int_0^p f(t) \sin \frac{2\pi nt}{p} dt$

Therefore, we have the exponential form of the above fourier series given by,

$$f(x) \sim \sum_{n=-\infty}^{\infty} \alpha_n e^{\frac{2\pi i n x}{p}}$$
, where $\alpha_n = \frac{1}{p} \int_0^p f(t) e^{\frac{-2\pi i n t}{p}} dt$

13.1.3 Fourier Integral Theorem

Theorem 13.1.2 (Fourier Integral Theorem). Let $f \in L(-\infty, \infty)$. Suppose $x \in \mathbb{R}$ and an interval $[x - \delta, x + \delta]$ about x such that either

1. f is of bounded variation on an interval $[x - \delta, x + \delta]$ about x or

2. both limits f(x+) and f(x-) exists and both Lebesgue intergrals

$$\int_0^\delta \frac{f(x+t) - f(x+)}{t} dt \text{ and } \int_0^\delta \frac{f(x-t) - f(x-)}{t} dt$$

exists.

Then,

$$\frac{f(x+) + f(x-)}{2} = \frac{1}{\pi} \int_0^{\infty} \int_{-\infty}^{\infty} f(u) \cos v(u-x) \ du \ dv,$$

the integral \int_0^∞ being an improper Riemann integral.

Synopsis.

$$f(x+t)\frac{\sin \alpha t}{\pi t}dt \to f(u)\frac{\sin \alpha (u-x)}{\pi (u-x)} \to \frac{f(u)}{\pi} \int_0^\alpha \cos v(u-x)dv$$

By Riemann-Lebesgue lemma/?, Theorem 11.6/,

$$f \in L(I) \implies \lim_{\alpha \to +\infty} \int_I f(x) \sin \alpha t \ dt = 0$$

By Jordan's Theorem[?, Theorem 10.8], if g is of bounded variation on $[0, \delta]$, then

$$\lim_{\alpha \to +\infty} \frac{2}{\pi} \int_0^{\delta} g(t) \frac{\sin \alpha t}{t} dt = g(0+)$$

By Dini's Theorem[?, Theorem 10.9], if the limit g(x+) exists and Lebesgue integral $\int_0^\delta \frac{g(t)+g(0+)}{t} dt$ exists for some $\delta > 0$, then

$$\lim_{\alpha to + \infty} \frac{2}{\pi} \int_0^{\delta} g(t) \frac{\sin \alpha t}{t} dt = g(0+)$$

The order of Lebesgue integrals can be interchanged?, Theorem 10.40], Suppose $f \in L(X)$ and $g \in L(Y)$. Then

$$\int_X f(x) \left(\int_Y g(y)k(x,y) dy \right) dx = \int_Y g(y) \left(\int_X f(x)k(x,y) dx \right) dy$$

Proof. Consider $\int_{-\infty}^{\infty} f(x+t) \frac{\sin \alpha t}{\pi t} dt$. We prove that this integral is equal to the either sides.

$$\int_{-\infty}^{\infty} f(x+t) \frac{\sin \alpha t}{\pi t} dt = \int_{-\infty}^{-\delta} + \int_{-\delta}^{0} + \int_{0}^{-\delta} + \int_{\delta}^{\infty} f(x+t) \frac{\sin \alpha t}{\pi t} dt$$

We have, function $\frac{f(x+t)}{\pi t}$ is bounded on $(-\infty, -\delta) \cup (\delta, \infty)$, hence $\frac{f(x+t)}{\pi t}$ is Lebesgue integrable on $(-\infty, -\delta) \cup (\delta, \infty)$.

By Riemann Lebesgue lemma,

$$\frac{f(x+t)}{\pi t} \in L(-\infty, -\delta) \implies \int_{-\infty}^{-\delta} f(x+t) \frac{\sin \alpha t}{\pi t} dt = 0,$$

$$\frac{f(x+t)}{\pi t} \in L(\delta, \infty) \implies \int_{\delta}^{\infty} f(x+t) \frac{\sin \alpha t}{\pi t} dt = 0$$

Case 1 Suppose f is of bounded variation on $[x - \delta, x + \delta]$, put g(t) = f(x + t) then g is of bounded variation on $[-\delta, \delta]$. Thus g is of bounded variation on $[0, \delta]$. Then by Jordan's Theorem

$$\lim_{\alpha \to +\infty} \frac{2}{\pi} \int_0^{\delta} f(x+t) \frac{\sin \alpha t}{t} dt = \lim_{\alpha \to +\infty} \frac{2}{\pi} \int_0^{\delta} g(t) \frac{\sin \alpha t}{t} dt = g(0+) = f(x+)$$

Case 2 Suppose both the limits f(x+) and f(x-) exists and both Lebesgue integrals

$$\int_0^\delta \frac{f(x+t) - f(x+)}{t} dt \text{ and } \int_0^\delta \frac{f(x-t) - f(x-)}{t} dt$$

exists.

Thus, we have f(x+) exists and the Lebesgue integral $\int_0^\delta \frac{f(x+t)-f(x+)}{t}dt$ exists. Put g(t)=f(x+t), then g(0+)=f(x+) exists and the Lebesgue integral $\int_0^\delta \frac{g(t)-g(0+)}{t}dt$ exists, then by Dini's Theorem,

$$\lim_{\alpha \to +\infty} \frac{2}{\pi} \int_0^{\delta} f(x+t) \frac{\sin \alpha t}{t} dt = \lim_{\alpha \to +\infty} \frac{2}{\pi} \int_0^{\delta} g(t) \frac{\sin \alpha t}{t} dt = g(0+) = f(x+)$$

Similarly, f(x-) exists and the Lebesgue integral $\int_0^\delta \frac{f(x-t)-f(x-)}{t}dt$ exists. Put g(t)=f(x-t), then g(0+)=f(x-) exists and the Lebesgue integral $\int_0^\delta \frac{g(t)-g(0+)}{t}dt$ exists, then by Dini's Theorem,

$$\lim_{\alpha \to +\infty} \frac{2}{\pi} \int_{-\delta}^{0} f(x+t) \frac{\sin \alpha t}{t} dt = \lim_{\alpha \to +\infty} \frac{2}{\pi} \int_{0}^{\delta} f(x-\tau) \frac{\sin \alpha \tau}{\tau} d\tau$$
$$= \lim_{\alpha \to +\infty} \frac{2}{\pi} \int_{0}^{\delta} g(\tau) \frac{\sin \alpha \tau}{\tau} d\tau = g(0+) = f(x-)$$

Then by either cases,

$$\lim_{\alpha \to +\infty} \int_{-\infty}^{\infty} f(x+t) \frac{\sin \alpha t}{\pi t} dt = \lim_{\alpha \to +\infty} \int_{-\delta}^{0} + \int_{0}^{\delta} f(x+t) \frac{\sin \alpha t}{\pi t} dt$$
$$= \frac{f(x+) + f(x-)}{2}$$

We have, $\int_0^{\alpha} \cos v(u-x) dv = \frac{\sin v(u-x)}{u-x}$.

$$\begin{split} \lim_{\alpha \to +\infty} \int_{-\infty}^{\infty} f(x) \frac{\sin \alpha t}{\pi t} dt &= \lim_{\alpha \to +\infty} \int_{-\infty}^{\infty} f(u) \frac{\sin \alpha (u-x)}{u-x} du, \text{ (put } u = x+t) \\ &= \lim_{\alpha \to +\infty} \int_{-\infty}^{\infty} f(u) \left(\int_{0}^{\alpha} \cos v (u-x) dv \right) du \\ &= \lim_{\alpha \to +\infty} \int_{0}^{\alpha} \left(\int_{-\infty}^{\infty} f(u) \cos v (u-x) du \right) dv, \end{split}$$

since, the order of Lebesgue integrals can be reversed.

$$= \int_{0}^{\infty} \left(\int_{-\infty}^{\infty} f(u) \cos v (u - x) du \right) dv$$

where, \int_0^∞ is not a Lebesgue integral, but an improper Riemann integral

Therefore,

$$\int_0^\infty \left(\int_{-\infty}^\infty f(u) \cos v(u - x) du \right) dv = \lim_{\alpha \to +\infty} \int_{-\infty}^\infty f(x) \frac{\sin \alpha t}{\pi t} dt$$
$$= \frac{f(x+) + f(x-)}{2}$$

Remark. If a function f on $(-\infty, \infty)$ is non-periodic, then it may not have a fourier series representation. In such cases, we have fourier intergral representation.

13.1.4 Exponential form of Fourier Integral Theorem

Let $f \in L(-\infty, \infty)$. Suppose $x \in \mathbb{R}$ and an interval $[x - \delta, x + \delta]$ about x such that either

- 1. f is of bounded variation on an interval $[x \delta, x + \delta]$ about x or
- 2. both limits f(x+) and f(x-) exists and both Lebesgue intergrals

$$\int_0^{\delta} \frac{f(x+t) - f(x+)}{t} dt \text{ and } \int_0^{\delta} \frac{f(x-t) - f(x-)}{t} dt$$

exists.

Then,

$$\frac{f(x+) + f(x-)}{2} = \lim_{\alpha \to \infty} \frac{1}{2\pi} \int_{-\alpha}^{\alpha} \left(\int_{-\infty}^{\infty} f(u)e^{iv(u-x)} \ du \right) dv$$

Proof. Let $F(v) = \int_{-\infty}^{\infty} f(u) \cos v(u-x) du$. Then F(v) = F(-v) and

$$\lim_{\alpha \to \infty} \frac{1}{2\pi} \int_{-\alpha}^{\alpha} F(v)dv = \lim_{\alpha \to \infty} \frac{1}{\pi} \int_{0}^{\alpha} \int_{-\infty}^{\infty} f(u) \cos v(u - x) du dv$$
$$= \frac{f(x+) + f(x-)}{2}$$

Let $G(v) = \int_{-\infty}^{\infty} f(u) \sin v(u-x) du$. Then G(v) = -G(-v) and

$$\lim_{\alpha \to \infty} \frac{1}{2\pi} \int_{-\alpha}^{\alpha} G(v) dv = 0$$

Thus

$$\lim_{\alpha \to \infty} \frac{1}{2\pi} \int_{-\alpha}^{\alpha} F(v) + iG(v) dv = \frac{f(x+) + f(x-)}{2}$$

13.1.5 Integral Transforms

Definitions 13.1.3. Integral transform g(y) of f(x) is a Lebesgue integral or Improper Riemann integral of the form

$$g(y) = \int_{-\infty}^{\infty} K(x, y) f(x) \ dx$$

, where K is the kernal of the transform. We write $g = \mathcal{K}(f)$.

Remark. Integral transforms(operators) are linear operators. ie, $\mathcal{K}(af_1 + bf_2) = a\mathcal{K}f_1 + b\mathcal{K}f_2$

Remark. A few commonly used integral transforms,

1. Exponential Fourier Transform F,

$$\mathscr{F}f = \int_{-\infty}^{\infty} e^{-ixy} f(x) \ dx$$

2. Fourier Cosine Transform C,

$$\mathscr{C}f = \int_0^\infty \cos xy f(x) \ dx$$

3. Fourier Sine Transform \mathscr{S} ,

$$\mathscr{S}f = \int_0^\infty \sin xy f(x) \ dx$$

4. Laplace Transform \mathcal{L} ,

$$\mathscr{L}f = \int_0^\infty e^{-xy} f(x) \ dx$$

5. Mellin Transform \mathcal{M} ,

$$\mathscr{M}f = \int_0^\infty x^{y-1} f(x) \ dx$$

Remark. Suppose f(x) = 0, $\forall x < 0$.

$$\int_{-\infty}^{\infty} e^{-ixy} f(x) \, dx = \int_{0}^{\infty} e^{-ixy} f(x) \, dx = \int_{0}^{\infty} \cos xy \, f(x) \, dx + i \int_{0}^{\infty} \sin xy \, f(x) \, dx$$
$$\mathscr{F} f = \mathscr{C} f + i \mathscr{F} f$$

Therefore Fourier Cosine $\mathscr C$ and Sine $\mathscr S$ transforms are special cases of fourier integral transform, $\mathscr F$ provided f vanishes on negative real axis.

Remark. Let y = u + iv, f(x) = 0, $\forall x < 0$.

$$\int_{0}^{\infty} e^{-xy} f(x) = \int_{0}^{\infty} e^{-xu} e^{-ixv} f(x) \ dx = \int_{0}^{\infty} e^{-ixv} \phi_{u}(x) dx$$

where $\phi_u(x) = e^{-xu} f(x)$.

$$\mathcal{L}f = \mathscr{F}\phi_u$$

Therefore Laplace transform, $\mathcal L$ is a special case of Fourier integral transform, $\mathscr F$.

Remark. Let $g(y) = \mathscr{F}f(x)$.

$$g(y) = \int_{-\infty}^{\infty} e^{-ixy} f(x) \ dx$$

Suppose f is continuous at x, then by fourier integral theorem,

$$f(x) = \frac{1}{2\pi} \int_{-\infty}^{\infty} \left(\int_{-\infty}^{\infty} f(u)e^{iv(u-x)} du \right) dv$$

$$= \int_{-\infty}^{\infty} e^{-ivx} \left(\frac{1}{2\pi} \int_{-\infty}^{\infty} e^{ivu} f(u) du \right) dv$$

$$= \int_{-\infty}^{\infty} g(v)e^{-ivx} dv = \mathscr{F}g \text{ where } g(v) = \frac{1}{2\pi} \int_{-\infty}^{\infty} f(u)e^{ivu} du$$

The above function g(v) gives the inverse fourier transformation of f.

Let g be fourier transform of f, then f is uniquely determined by its fourier transform g by,

$$f(x) = \mathscr{F}^{-1}g(y) = \frac{1}{2\pi} \lim_{\alpha \to \infty} \int_{-\alpha}^{\alpha} g(y)e^{ixy}dy$$

6. Inverse Fourier Transform \mathscr{F}^{-1} ,

$$\mathscr{F}^{-1}f = \int_{-\infty}^{\infty} \frac{e^{ixy}}{2\pi} f(x) \ dx$$

13.1.6 Convolutions

Definitions 13.1.4. Let $f, g \in L(-\infty, \infty)$. Let S be the set of all points x for which the Lebesgue integral

$$h(x) = \int_{-\infty}^{\infty} f(t)g(x-t)dt$$

exists. Then the function $h: S \to \mathbb{R}$ is a convolution of f and g. And h = f * g.

Remark. Convolution operator is commutative. ie, h = f * g = g * f

Remark. Suppose f, g vanishes on negative real axis, then

$$h(x) = \int_{-\infty}^{\infty} f(t) \ g(x-t) \ dt = \int_{-\infty}^{0} + \int_{0}^{x} + \int_{x}^{\infty} f(t) \ g(x-t) \ dt = \int_{0}^{x} f(t) \ g(x-t) \ dt$$

Remark. Singularity is a point at which the convolution integral fails to exists.

Theorem 13.1.5. Let $f, g \in L(\mathbb{R})$ and either f or g is bounded in \mathbb{R} . Then the convolution integral

$$h(x) = \int_{-\infty}^{\infty} f(t)g(x-t)dt$$

exists for every $x \in \mathbb{R}$ and the function h so defined is bouned in \mathbb{R} . In addition, if the bounded function is continuous on \mathbb{R} , then h is continuous and $h \in L(\mathbb{R})$.

Synopsis.

Proof. \Box

Remark. If f, g are both unbounded, the convolution integral may not exist.

eg:
$$f(t) = \frac{1}{\sqrt{t}}$$
, $g(t) = \frac{1}{\sqrt{1-t}}$

Theorem 13.1.6. Let $f, g \in L^2(\mathbb{R})$. Then the convolution integral f * g exists for each $x \in \mathbb{R}$ and the function $h : \mathbb{R} \to \mathbb{R}$ defined by h(x) = f * g(x) is bounded in \mathbb{R} .

Synopsis.

Proof.

13.1.7 The Convolution Theorem for Fourier Tranforms

Theorem 13.1.7. Let $f, g \in L(\mathbb{R})$ and at least one of f or g is continuous and bounded on \mathbb{R} . Let h = f * g. Then for every real u,

$$\int_{-\infty}^{\infty} h(x) e^{-ixu} dx = \left(\int_{-\infty}^{\infty} f(t) e^{-itu} dt \right) \left(\int_{-\infty}^{\infty} g(y) e^{-iyu} dy \right)$$

The integral on the left exists both as a Lebesgue integral and an improper Riemann integral.

Synopsis.

Proof.
$$\Box$$

Remark (Application of Convolution Theorem).

$$B(p,q) = \frac{\Gamma p \Gamma q}{\Gamma p + q}$$
, where $B(p,q) = \int_0^1 x^{p-1} (1-x)^{q-1} dx$, $\Gamma p = \int_0^\infty t^{p-1} e^{-t} dt$

13.2 Multivariate Differential Calculus

In this chapter, we deal with real functions of several variables. Instead of \mathbf{c} , we write $\overline{c} \in \mathbb{R}^n$, then $\overline{c} = (c_1, c_2, \dots, c_n)$ where $c_j \in \mathbb{R}$ for every $j = 1, 2, \dots, n$. Again, suppose $f : \mathbb{R}^n \to \mathbb{R}^m$ and $f(\overline{x}) = \overline{y}$, then $\overline{y} = (y_1, y_2, \dots, y_m)$ where each y_k is real. The unit co-ordinate vector, $\overline{u_k}$ is given by $u_{kj} = \delta_{j,k}$

13.2.1 Directional Derivative

Motivation: The existence of all partial derivatives of a multivariate real function f at a point \bar{c} doesn't imply the continuity of f at \bar{c} . Thus, we need a suitable generalisation for the partial derivative which could characterise continuity. And directional derivative is such an attempt.

Definitions 13.2.1 (Directional Derivative). Let $S \subset \mathbb{R}^n$ and $f: S \to \mathbb{R}^m$. Let \overline{c} be an interior points of S and $\overline{u} \in \mathbb{R}^n$, then there exists an open ball $B(\overline{c}, r)$ in S. Also for some $\delta > 0$ the line segment $\alpha: [0, \delta] \to S$ given by $\alpha(t) = \overline{c} + t\overline{u}$ lie in $B(\overline{c}, r)$.

Then the Directional derivative of f at an interior point \overline{c} in the direction \overline{u} is given by

$$f'(\overline{c}, \overline{u}) = \lim_{h \to 0} \frac{f(\overline{c} + h\overline{u}) - f(\overline{c})}{h}$$

Remark. The direction derivative of f at an interior point \bar{c} in the direction \bar{u} exists only if the above limit exists.

Remark. Example, [?, Exercise 12.2a]

Suppose $\overline{x}, \overline{a}, \overline{c}, \overline{u} \in \mathbb{R}^n$. Let $f : \mathbb{R}^n \to \mathbb{R}$ such that $f(\overline{x}) = \overline{a} \cdot \overline{x}$. Then

$$f'(\overline{c}, \overline{u}) = \lim_{h \to 0} \frac{\overline{a} \cdot (\overline{c} + h\overline{u}) - \overline{a} \cdot \overline{c}}{h} = \overline{a} \cdot \overline{u}$$

Remark (Properties). Let $f: S \to \mathbb{R}^m$, where $S \subset \mathbb{R}^n$

1. $f'(\overline{c}, \overline{0}) = \overline{0}$

Note: The zero vectors belongs to \mathbb{R}^n , \mathbb{R}^m respectively.

- 2. $f'(\overline{c}, \overline{u_k}) = \frac{\partial f}{\partial u_k}(\overline{c}) = D_k f(\overline{c})$, the k^{th} partial derivative of f.
- 3. Let $f = (f_1, f_2, \dots, f_m)$, such that $f(\overline{c}) = (f_1(\overline{c}), f_2(\overline{c}), \dots, f_m(\overline{c}))$. Then, $\exists f'(\overline{c}, \overline{u}) \iff \forall k, \exists f'_k(\overline{c}, \overline{u}) \text{ and } f'(\overline{c}, \overline{u}) = (f'_1(\overline{c}, \overline{u}), f'_2(\overline{c}, \overline{u}), \dots, f'_m(\overline{c}, \overline{u}))$
 - ie, Directional derivative of f exists iff directional derivative of each component function f_k exists. And the components of the directional derivatives of f are the directional derivatives of the components of f.

 Thus $D_k f(\overline{c}) = (D_k f_1(\overline{c}), D_k f_2(\overline{c}), \dots, D_k f_m(\overline{c}))$ holds.
- 4. Let $F(t) = f(\overline{c} + t\overline{u})$, then $F'(0) = f'(\overline{c}, \overline{u})$ and $F'(t) = f'(\overline{c} + t\overline{u}, \overline{u})$
- 5. Let $f(\overline{c}) = \overline{c} \cdot \overline{c} = \|\overline{c}\|^2$, and $F(t) = f(\overline{c} + t\overline{u})$, then $F'(t) = 2\overline{c} \cdot \overline{u} + 2t\|\overline{u}\|^2$ and $F'(0) = f'(\overline{c}, \overline{u}) = 2\overline{c} \cdot \overline{u}$
- 6. Let f be linear, then $f'(\overline{c}, \overline{u}) = f(\overline{u})$
- 7. Existence of all partial derivatives doesn't imply existence of all directional derivatives.

$$f(x,y) = \begin{cases} x+y & \text{if } x = 0 \text{ or } y = 0\\ 1 & \text{otherwise} \end{cases}$$

For above f, directional derivatives exists only along the co-ordinates (ie, partial derivatives).

8. Existence of all directional derivatives doesn't imply continuity.

$$f(x,y) = \begin{cases} xy^2(x^2 + y^4) & x \neq 0 \\ 0 & x = 0 \end{cases}$$

Above f is discontinuous at (0,0), however all directional derivatives exists and has finite value.

13.2.2 Total Derivative

We may define a total derivative $T_c(h) = hf'(c)$ in the case of real-functions of single variable as follows:-

Let
$$E_c(h) = \begin{cases} \frac{f(c+h) - f(c)}{h} - f'(c), & h \neq 0\\ 0, & h = 0 \end{cases}$$

Then, $f(c+h) = f(c) + hf'(c) + hE_c(h)$ and as $h \to 0$, $E_c(h) \to 0$. Also $T_c(h) = f'(c)h$ is a linear function of h. ie, $T_c(ah_1 + bh_2) = aT_c(h_1) + bT_c(h_2)$. Now, we will define a total derivative of multivariate function that has these two properties.

Definitions 13.2.2 (Total Derivative). The function $f: \mathbb{R}^n \to \mathbb{R}^m$ is differentiable at \overline{c} if there exists a **linear** function $T_{\overline{c}}: \mathbb{R}^n \to \mathbb{R}^m$ such that $f(\overline{c} + \overline{v}) = f(\overline{c}) + T_{\overline{c}}(\overline{v}) + \|\overline{v}\| E_{\overline{c}}(\overline{v})$ where $E_{\overline{c}}(\overline{v}) \to \overline{0}$ as $\overline{v} \to \overline{0}$.

The linear function $T_{\overline{c}}$ is the total derivative of f at \overline{c} , $T_{\overline{c}}(\overline{0}) = \overline{0}$ and the condition above gives the First Order Taylor's Formula for $f(\overline{c} + \overline{v}) - f(\overline{c})$.

Remark (Properties). Let $f: \mathbb{R}^n \to \mathbb{R}^m$ and $f'(\overline{c})(\overline{v}) = T_{\overline{c}}(\overline{v})$ be the total derivative of f at \overline{c} evaluated at \overline{v} . Then,

- 1. $f'(\overline{c})(\overline{v}) = f'(\overline{c}, \overline{u})$
- 2. If f is differentiable at \bar{c} , then f is continuous at \bar{c} .
- 3. $f'(\overline{c})(\overline{v}) = v_1 D_1 f(\overline{c}) + v_2 D_2 f(\overline{c}) + \dots + v_n D_n f(\overline{c})$

Note. The above f' is a function from \mathbb{R}^n to the set of all linear functions $\mathcal{L} = \{h : \mathbb{R}^n \to \mathbb{R}^m\}$. $f'(\overline{c})$ is a linear function (in fact, total derivative $T_{\overline{c}}$) which maps \overline{v} into the directional derivatives of f at \overline{c} in the direction \overline{v} . This notation generalises f' for univariate f as well. (put n = m = 1)

In this subject, we use the following notations,

 $D_k f(\overline{c})$ partial derivative

 $f'(\overline{c}, \overline{v})$ directional derivative

 $f'(\overline{c})(\overline{v})$ total derivative

 $\nabla f(\overline{c})$ gradient vector

Theorem 13.2.3. If f is differentiable at \overline{c} with total derivative $T_{\overline{c}}$, then for every $\overline{u} \in \mathbb{R}^n$, $T_{\overline{c}}(\overline{u}) = f'(\overline{c}, \overline{u})$. (ie, $f'(\overline{c})(\overline{v}) = f'(\overline{c}, \overline{v})$)

Proof. For $\overline{v} = \overline{0}$, we have $T_{\overline{c}}(\overline{0}) = 0 = f'(\overline{c}, \overline{0})$.

Suppose $\overline{v} \neq \overline{0}$, then put $\overline{v} = h\overline{u}$. Since f is differentiable at \overline{c} , f has total derivative at \overline{c} . That is, there exists a linear function $T_{\overline{c}}$ such that $f(\overline{c} + h\overline{u}) =$

$$f(\overline{c}) + T_{\overline{c}}(h\overline{u}) + ||h\overline{u}|| E_{\overline{c}}(h\overline{u}) \text{ where } E_{\overline{c}}(h\overline{u}) \to \overline{0} \text{ as } h\overline{u} \to \overline{0}.$$

$$\implies f(\overline{c} + h\overline{u}) = f(\overline{c}) + hT_{\overline{c}}(\overline{u}) + |h|||\overline{u}|| E_{\overline{c}}(h\overline{u}), \ E_{\overline{c}}(h\overline{u}) \to \overline{0} \text{ as } h\overline{u} \to \overline{0}$$

$$\implies \frac{f(\overline{c} + h\overline{u}) - f(\overline{c})}{h} = T_{\overline{c}}(\overline{u}) + \frac{|h|||\overline{u}|| E_{\overline{c}}(h\overline{u})}{h}, \ E_{\overline{c}}(h\overline{u}) \to \overline{0} \text{ as } h \to 0$$

$$\implies \lim_{h \to 0} \frac{f(\overline{c} + h\overline{u}) - f(\overline{c})}{h} = T_{\overline{c}}(\overline{u}) + \lim_{h \to 0} \frac{|h|||\overline{u}|| E_{\overline{c}}(h\overline{u})}{h}$$

$$\implies f'(\overline{c}, \overline{u}) = T_{\overline{c}}(\overline{u})$$

Note. $T_{\overline{c}}$ is linear, however $E_{\overline{c}}$ is not linear. Thus $E_{\overline{c}}(h\overline{u}) \neq hE_{\overline{c}}(\overline{u})$.

As $h \to 0$, $h\overline{u} \to \overline{0}$ and $E_{\overline{c}}(h\overline{u}) \to \overline{0}$. Since the order of the function $E_{\overline{c}}(h\overline{u})$ is much smaller than that of h, the limit on the right converges to 0.

Theorem 13.2.4. If f is differentiable at \bar{c} , then f is continuous at \bar{c} .

Proof. Let $\overline{v} \neq 0$, then

$$\overline{v} = v_1 \overline{u_1} + v_2 \overline{u_2} + \dots + v_n \overline{u_n},$$

$$\overline{v} \to \overline{0} \implies \forall j, \ v_j \to 0$$

$$T \text{ is linear } \implies T_{\overline{c}}(\overline{v}) = v_1 T_{\overline{c}}(\overline{u_1}) + v_2 T_{\overline{c}}(\overline{u_2}) + \dots + v_n T_{\overline{c}}(\overline{u_n})$$

$$\text{Thus, } T_{\overline{c}}(\overline{v}) \to \overline{0} \text{ as } \overline{v} \to 0$$

Since f differentiable at \bar{c} , there exists linear function $T_{\bar{c}}$ such that

$$f(\overline{c} + \overline{v}) = f(\overline{c}) + T_{\overline{c}}(\overline{v}) + \|v\| E_{\overline{c}}(\overline{v})$$

$$\implies \lim_{\overline{v} \to \overline{0}} f(\overline{c} + \overline{v}) = f(\overline{c}) + \lim_{\overline{v} \to \overline{0}} T_{\overline{c}}(\overline{v}) + \lim_{\overline{v} \to \overline{0}} \|v\| E_{\overline{c}}(\overline{v})$$

$$\implies \lim_{\overline{v} \to \overline{0}} f(\overline{c} + \overline{v}) = f(\overline{c})$$

Theorem 13.2.5. Let $S \subset \mathbb{R}^n$ and $f: S \to \mathbb{R}^m$ be differentiable at an interior point \overline{c} of S, where $S \subseteq \mathbb{R}^n$. If $\overline{v} = v_1 \overline{u_1} + v_2 \overline{u_2} + \cdots + v_n \overline{u_n}$, then

$$f'(\overline{c})(\overline{v}) = \sum_{k=1}^{n} v_k D_k f(\overline{c})$$

In particular, if f is real-valued (m=1) we have, $f'(\bar{c})(\bar{v}) = \nabla f(\bar{c}).\bar{v}$

Proof. Suppose $f: S \to \mathbb{R}^m$ is differentiable at \overline{c} , then there exists a linear function $f'(\overline{c}): S \to \mathbb{R}^m$ such that $f(\overline{c} + \overline{v}) = f(\overline{c}) + f'(\overline{c})(\overline{v}) + \|\overline{v}\| E_{\overline{c}}(\overline{c})$ where $E_{\overline{c}} \to \overline{0}$ as $\overline{v} \to \overline{0}$.

$$f'(\overline{c})(\overline{v}) = f'(\overline{c}) \left(\sum_{k=1}^{n} v_k \overline{u_k} \right)$$

$$= \sum_{k=1}^{n} v_k f'(\overline{c})(\overline{u_k}), \text{ since } f'(\overline{c}) \text{ is linear}$$

$$= \sum_{k=1}^{n} v_k D_k f(\overline{c}), \text{ since } f'(\overline{c})(\overline{u_k}) = f'(\overline{c}, \overline{u_k}) = D_k f(\overline{c})$$

Let m = 1, then $f: S \to \mathbb{R}$

$$f'(\overline{c})(\overline{v}) = \sum_{k=1}^{n} v_k D_k f(\overline{c}) = \nabla f(\overline{c}) . \overline{v}$$

since $\nabla f(\overline{c}) = (D_1 f(\overline{c}), D_2 f(\overline{c}), \dots, D_n f(\overline{c}))$

Remark. Let $f: S \to \mathbb{R}$, then $f(\overline{c} + \overline{v}) = f(\overline{c}) + \nabla f(\overline{c}).\overline{v} + o(\|\overline{v}\|)$ as $\overline{v} \to \overline{0}$.

Remark (Complex-valued Functions). Seminar - Robitha

13.2.3 Matrix of Linear Function

Let $T: \mathbb{R}^n \to \mathbb{R}^m$ be a linear function. Let $\{\overline{u_1}, \overline{u_2}, \cdots, \overline{u_n}\}$ be standard basis for \mathbb{R}^n and $\{\overline{e_1}, \overline{e_2}, \cdots, \overline{e_m}\}$ be standard basis for \mathbb{R}^m . Let $\overline{v} \in \mathbb{R}^n$, then $\overline{v} = \sum_{k=1}^n v_k \overline{u_k}$ and $T(\overline{v}) = \sum_{k=1}^n v_k T(\overline{u_k})$ and

$$\begin{split} T(\overline{v}) &= \begin{bmatrix} v_1 & v_2 & \vdots & v_n \end{bmatrix} \begin{bmatrix} T(\overline{u_1}) \\ T(\overline{u_2}) \\ \ddots \\ T(\overline{u_n}) \end{bmatrix} \\ &= \begin{bmatrix} v_1 & v_2 & \vdots & v_n \end{bmatrix} \begin{bmatrix} t_{11}\overline{e_1} + t_{21}\overline{e_2} + \dots + t_{m1}\overline{e_m} \\ t_{12}\overline{e_1} + t_{22}\overline{e_2} + \dots + t_{m2}\overline{e_m} \\ & \ddots \\ t_{1n}\overline{e_1} + t_{2n}\overline{e_2} + \dots + t_{mn}\overline{e_m} \end{bmatrix} \\ &= \begin{bmatrix} v_1 & v_2 & \cdots & v_n \end{bmatrix} \begin{bmatrix} t_{11} & t_{21} & \cdots & t_{m1} \\ t_{12} & t_{22} & \cdots & t_{m2} \\ \vdots & \vdots & \ddots & \vdots \\ t_{1n} & t_{2n} & \cdots & t_{mn} \end{bmatrix} \begin{bmatrix} \overline{e_1} \\ \overline{e_2} \\ \vdots \\ \overline{e_m} \end{bmatrix} \end{split}$$

We may take the transpose,

$$T(\overline{v}) = \begin{bmatrix} \overline{e_1} & \overline{e_2} & \cdots & \overline{e_m} \end{bmatrix} \begin{bmatrix} t_{11} & t_{12} & \cdots & t_{1n} \\ t_{21} & t_{22} & \cdots & t_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ t_{m1} & t_{m2} & \cdots & t_{mn} \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \\ \vdots \\ v_n \end{bmatrix}$$

$$T(\overline{v}) = T\left(\sum_{k=1}^{n} v_k \overline{u_k}\right) = \sum_{k=1}^{n} v_k T(\overline{u_k}) = \sum_{k=1}^{n} v_n \sum_{j=1}^{m} t_{kj} \overline{e_j}$$

Thus matrix of T is given by, $m(T) = (t_{ik})$ where $T(\overline{u_k}) = \sum_{k=1}^n t_{ik} \overline{e_i}$.

Remark (Example). Let $T: \mathbb{R}^3 \to \mathbb{R}^2$ defined by T(x, y, z) = (2x + y, y - z).

$$\begin{split} T(1,2,3) &= T((1,0,0) + 2(0,1,0) + 3(0,0,1)) \\ &= T(\overline{u_1} + 2\overline{u_2} + 3\overline{u_3}) \\ &= T(\overline{u_1}) + 2T(\overline{u_2}) + 3T(\overline{u_3}) \\ &= \begin{bmatrix} 1 & 2 & 3 \end{bmatrix} \begin{bmatrix} T(\overline{u_1}) \\ T(\overline{u_2}) \\ T(\overline{u_3}) \end{bmatrix} \\ &= \begin{bmatrix} 1 & 2 & 3 \end{bmatrix} \begin{bmatrix} (2,0) \\ (1,1) \\ (0,-1) \end{bmatrix} \\ &= \begin{bmatrix} 1 & 2 & 3 \end{bmatrix} \begin{bmatrix} 2(1,0) \\ (1,0) + (0,1) \\ -1(0,1) \end{bmatrix} \\ &= \begin{bmatrix} 1 & 2 & 3 \end{bmatrix} \begin{bmatrix} \frac{2\overline{e_1}}{\overline{e_1} + \overline{e_2}} \\ -\overline{e_2} \end{bmatrix} \\ &= \begin{bmatrix} 1 & 2 & 3 \end{bmatrix} \begin{bmatrix} 2 & 0 \\ 1 & 1 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} \overline{e_1} \\ \overline{e_2} \end{bmatrix} \\ &= 4\overline{e_1} - \overline{e_2} = 4(1,0) - 1(0,1) = (4,-1) \end{split}$$

In the above case,
$$m(T) = \begin{bmatrix} 2 & 0 \\ 1 & 1 \\ 0 & -1 \end{bmatrix}$$

Using the matrix of linear function m(T), we can compute the image of any point in \mathbb{R}^3 by matrix multiplication.

Matrix of the composition of two linear functions

Let $T: \mathbb{R}^n \to \mathbb{R}^m$ and $S: \mathbb{R}^m \to \mathbb{R}^p$ be two linear functions with domain of S containing the range of T (so that $S \circ T$ is well defined). Then $S \circ T: \mathbb{R}^n \to \mathbb{R}^p$ is defined by

$$S \circ T(\overline{x}) = S(T(\overline{x})), \ \forall \overline{x} \in \mathbb{R}^n$$

Since S, T are linear, $S \circ T$ is also linear.

$$S \circ T(a\overline{x} + b\overline{y}) = S(T(a\overline{x} + b\overline{y})) = S(aT(\overline{x}) + bT(\overline{y})) = aS(T(\overline{x})) + bS(T(\overline{y}))$$
$$= aS \circ T(\overline{x}) + bS \circ T(\overline{y}), \ \forall a, b \in \mathbb{R}, \ \forall \overline{x}, \overline{y} \in \mathbb{R}^n$$

Let $\{\overline{u_1}, \overline{u_2}, \dots, \overline{u_n}\}$ be the standards basis for \mathbb{R}^n , $\{\overline{e_1}, \overline{e_2}, \dots, \overline{e_m}\}$ be the standards basis for \mathbb{R}^m and $\{\overline{w_1}, \overline{w_2}, \dots, \overline{w_p}\}$ be the standards basis for \mathbb{R}^p . Let

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$$\overline{v} \in \mathbb{R}^n$$
, then $\overline{(v)} = \sum_{i=1}^n v_i \overline{u_i}$, and $S \circ T(\overline{v}) = \sum_{i=1}^n v_i S \circ T(\overline{u_i})$

$$\begin{split} S \circ T(\overline{v}) &= \begin{bmatrix} v_1 & v_2 & \vdots & v_n \end{bmatrix} \begin{bmatrix} S \circ T(\overline{u_1}) \\ S \circ T(\overline{u_2}) \\ \vdots \\ S \circ T(\overline{u_n}) \end{bmatrix} \\ &= \begin{bmatrix} v_1 & v_2 & \vdots & v_n \end{bmatrix} \begin{bmatrix} S(t_{11}\overline{e_1} + \dots + t_{m1}\overline{e_m}) \\ S(t_{12}\overline{e_1} + \dots + t_{m2}\overline{e_m}) \\ \vdots \\ S(t_{1n}\overline{e_1} + \dots + t_{mn}\overline{e_m}) \end{bmatrix} \\ &= \begin{bmatrix} v_1 & v_2 & \vdots & v_n \end{bmatrix} \begin{bmatrix} t_{11}S(\overline{e_1}) + \dots + t_{m1}S(\overline{e_m}) \\ t_{12}S(\overline{e_1}) + \dots + t_{m2}S(\overline{e_m}) \\ \vdots & \vdots & \ddots & \vdots \\ t_{1n}S(\overline{e_1}) + \dots + t_{mn}S(\overline{e_m}) \end{bmatrix} \\ &= \begin{bmatrix} v_1 & v_2 & \cdots & v_n \end{bmatrix} \begin{bmatrix} t_{11} & t_{21} & \cdots & t_{m1} \\ t_{12} & t_{22} & \cdots & t_{m2} \\ \vdots & \vdots & \ddots & \vdots \\ t_{1n} & t_{2n} & \cdots & t_{mn} \end{bmatrix} \begin{bmatrix} S(\overline{e_1}) \\ S(\overline{e_2}) \\ \vdots \\ S(\overline{e_m}) \end{bmatrix} \\ &= \begin{bmatrix} v_1 & v_2 & \cdots & v_n \end{bmatrix} \begin{bmatrix} t_{11} & t_{21} & \cdots & t_{m1} \\ t_{12} & t_{22} & \cdots & t_{m2} \\ \vdots & \vdots & \ddots & \vdots \\ t_{1n} & t_{2n} & \cdots & t_{mn} \end{bmatrix} \begin{bmatrix} s_{11}\overline{w_1} + s_{12}\overline{w_2} + \cdots + s_{1p}\overline{w_p} \\ s_{12}\overline{w_1} + s_{22}\overline{w_2} + \cdots + s_{p2}\overline{w_p} \\ \vdots & \ddots & \vdots \\ s_{1m}\overline{w_1} + s_{2m}\overline{w_2} + \cdots + s_{pm}\overline{w_p} \end{bmatrix} \begin{bmatrix} \overline{w_1} \\ \overline{w_2} \\ \vdots & \vdots & \ddots & \vdots \\ t_{1n} & t_{2n} & \cdots & t_{mn} \end{bmatrix} \begin{bmatrix} s_{11} & s_{21} & \cdots & s_{p1} \\ s_{12} & s_{22} & \cdots & s_{p2} \\ \vdots & \vdots & \ddots & \vdots \\ s_{1m} & s_{2m} & \cdots & s_{pm} \end{bmatrix} \begin{bmatrix} \overline{w_1} \\ \overline{w_2} \\ \vdots & \vdots & \ddots & \vdots \\ \overline{w_p} \end{bmatrix} \end{split}$$

We may take transpose,

$$S \circ T(\overline{v}) = \begin{bmatrix} \overline{w_1} & \overline{w_2} & \vdots & \overline{w_p} \end{bmatrix} \begin{bmatrix} s_{11} & s_{12} & \cdots & s_{1m} \\ s_{21} & s_{22} & \cdots & s_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ s_{p1} & s_{p2} & \cdots & s_{pm} \end{bmatrix} \begin{bmatrix} t_{11} & t_{12} & \cdots & t_{1n} \\ t_{21} & t_{22} & \cdots & t_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ t_{m1} & t_{m2} & \cdots & t_{mn} \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \\ \cdots \\ v_n \end{bmatrix}$$

Remember: Given $T: \mathbb{R}^n \to \mathbb{R}^m$, then we may take m(T) either as $m \times n$ matrix or $n \times m$ matrix. Since, we chose $m \times n$, $m(S \circ T) = m(S)m(T)$. Otherwise, $m(S \circ T) = m(T)m(S)$. This may change for different authors.

Suppose $m(S) = (s_{ij})$ and $m(T) = (t_{ij})$ respectively. Then

$$S(e_k) = \sum_{i=1}^p s_{ik} \overline{w_i}, \ k = 1, 2, \cdots, m \text{ and}$$

$$T(u_j) = \sum_{k=1}^{m} t_{kj} \overline{e_k}, \ j = 1, 2, \cdots, n$$

$$(S \circ T)(\overline{u_j}) = S(T(\overline{u_j})) = S\left(\sum_{k=1}^m t_{kj}\overline{e_k}\right) = \sum_{k=1}^m t_{kj}S(\overline{e_k})$$
$$= \sum_{k=1}^m t_{kj}\left(\sum_{i=1}^p s_{ik}\overline{w_i}\right) = \sum_{i=1}^p \left(\sum_{k=1}^m s_{ik}t_{kj}\right)\overline{w_i}$$

Therefore, $m(S \circ T) = \sum_{k=1}^{m} s_{ik} t_{kj} = (s_{ik})(t_{kj}) = m(S)m(T)$.

13.2.4 The Jacobian Matrix

Let $\overline{u_1}, \overline{u_2}, \dots, \overline{u_n}$ be the unit co-ordinate vectors in \mathbb{R}^n and $\overline{e_1}, \overline{e_2}, \dots, \overline{e_m}$ be the unit co-ordinate vectors in \mathbb{R}^m . Let function $f: \mathbb{R}^n \to \mathbb{R}^m$ be differentiable at $\overline{c} \in \mathbb{R}^n$. Then there exists a linear function $T = f'(\overline{c}) : \mathbb{R}^n \to \mathbb{R}^m$ such that $f(\overline{c} + \overline{v}) = f(\overline{c}) + f'(\overline{c})(\overline{v}) + \|\overline{v}\| + E_{\overline{c}}(\overline{v})$. We have, $T(\overline{u_k}) = f'(\overline{c})(\overline{u_k}) = f'(\overline{c}, \overline{u_k}) = D_k f(\overline{c}) = D_k \sum_{i=1}^m f_i(\overline{c}) \overline{e_i}$.

Clearly, the matrix of total derivative T, $m(T) = (t_{ik}) = (D_k f_i(\overline{c}))$. This matrix is called Jacobian matrix of f at \overline{c} and is denoted by $Df(\overline{c})$.

$$Df(\overline{c}) = \begin{bmatrix} D_1 f_1(\overline{c}) & D_2 f_1(\overline{c}) & \cdots & D_n f_1(\overline{c}) \\ D_1 f_2(\overline{c}) & D_2 f_2(\overline{c}) & \cdots & D_n f_2(\overline{c}) \\ \vdots & \vdots & \ddots & \vdots \\ D_1 f_m(\overline{c}) & D_2 f_m(\overline{c}) & \cdots & D_n f_m(\overline{c}) \end{bmatrix}$$

Properties of Jacobian matrix

1. kth row of $Df(\bar{c})$ is gradient vector of f_k

$$\nabla f_k(\overline{c}) = (D_1 f_k(\overline{c}), D_2 f_k(\overline{c}), \cdots, D_n f_k(\overline{c}))$$

- 2. When m = 1, $Df(\overline{c}) = \nabla f(\overline{c})$.
- 3. $f'(\overline{c})(\overline{v}) = \sum_{k=1}^{m} (\nabla f_k(\overline{c}) \cdot \overline{v}) \overline{e_k}$
- 4. $||f'(\overline{c})(\overline{v})|| \le M||v||$ where $M = \sum_{k=1}^m ||\nabla f_k(\overline{c})||$, by property (3)
- 5. $f'(\overline{c})(\overline{v}) \to \overline{0}$ as $\overline{v} \to \overline{0}$, by property (4)

Chain Rule

Chain Rule for real function :
$$\frac{dF \circ G}{dx}(x) = \frac{d}{dy}F(y)\frac{d}{dx}G(x) = F'(y)\ G'(x)$$

For example : $\frac{d}{dx}(ax+3)^3 = \frac{d}{dy}y^3\frac{d}{dx}(ax+3) = 3ay^2 = 3a(ax+3)^2$

Theorem 13.2.6. Let g be differentiable at \overline{a} , with total derivative $g'(\overline{a})$ and $\overline{b} = g(\overline{a})$. Let f is differentiable at \overline{b} , with total derivative $f'(\overline{b})$. Then $h = f \circ g$ is differentiable at \overline{a} with total derivative $h'(\overline{a}) = f'(\overline{b}) \circ g'(\overline{a})$. Try to read $h'(\overline{a}) = H$, $f'(\overline{b}) = F$, $g'(\overline{a}) = G$, then $H = F \circ G \Longrightarrow H(x) = F(G(x))$ In other words, $h'(\overline{a})(\overline{v}) = f'(\overline{b}) \circ g'(\overline{a})$ (\overline{v}) = $f'(\overline{b})(g'(\overline{a})(\overline{v}))$.

Proof. Given $\epsilon > 0$, let $y \in \mathbb{R}^p$ such that $||y|| < \epsilon$. Let $f : \mathbb{R}^n \to \mathbb{R}^m$ and $g : \mathbb{R}^p \to \mathbb{R}^n$, then $h = f \circ g : \mathbb{R}^p \to \mathbb{R}^m$.

We have,
$$h(\overline{a} + \overline{y}) - h(\overline{a}) = f(g(\overline{a} + \overline{y})) - f(g(\overline{a})) = f(\overline{b} + \overline{v}) - f(\overline{b})$$
 where $\overline{b} = g(\overline{a})$, and $\overline{v} = g(\overline{a} + \overline{y}) - g(\overline{a})$.

Since g is differentiable at \overline{a} , g satisfies first-order Taylor's formula.

$$g(\overline{a} + \overline{y}) = g(\overline{a}) + g'(\overline{a})(\overline{y}) + \|\overline{y}\| E_{\overline{a}}(\overline{y}) \text{ where } E_{\overline{a}} \to \overline{0} \text{ as } \overline{y} \to \overline{0}$$

$$\Longrightarrow \overline{v} = g(\overline{a} + \overline{y}) - g(\overline{a}) = g'(\overline{a})(\overline{y}) + \|\overline{y}\| E_{\overline{a}}(\overline{y})$$

Clearly, as $\overline{y} \to \overline{0} \implies \overline{v} \to g'(\overline{a})(\overline{0}) = \overline{0}$. Again, we have f is differentiable at \overline{b} , thus f satisfies first-order Taylor's formula.

$$f(\overline{b} + \overline{v}) = f(\overline{b}) + f'(\overline{b})(\overline{v}) + \|\overline{v}\| E_{\overline{b}}(\overline{v})$$
 where $E_{\overline{b}} \to \overline{0}$ as $\overline{v} \to \overline{0}$

$$\implies f(\overline{b} + \overline{v}) - f(\overline{b}) = f'(\overline{b})(\overline{v}) + \|\overline{v}\| E_{\overline{b}}(\overline{v})$$

$$= f'(\overline{b}) (g'(\overline{a})(\overline{y}) + \|\overline{y}\| E_{\overline{a}}(\overline{y})) + \|\overline{v}\| E_{\overline{b}}(\overline{v})$$

$$= f'(\overline{b})(g'(\overline{a})(\overline{y})) + \|\overline{y}\| E(\overline{y})$$
where $E(\overline{y}) = f'(\overline{b})(E_{\overline{a}}(\overline{y})) + \frac{\|\overline{v}\|}{\|\overline{y}\|} E_{\overline{b}}(\overline{v}), \ \overline{y} \neq \overline{0}$

$$\implies h(\overline{a} + \overline{y}) - h(\overline{a}) = f(\overline{b} + \overline{v}) - f(\overline{b}) = f'(\overline{b})(g'(\overline{a})(\overline{y})) + \|\overline{y}\|E(\overline{y})$$

Since $f'(\overline{b})$ and $g'(\overline{a})$ are linear, their composition is also linear. Therefore, h is differentiable at \overline{a} with a linear, total derivative $h'(\overline{a}) = f'(\overline{b}) \circ g'(\overline{a})$ as it satisfies first-order Taylor's formula if $E_{\overline{y}} \to \overline{0}$ as $\overline{y} \to \overline{0}$.

We have, $\|\overline{v}\| \leq \|g'(\overline{a})(\overline{y})\| + \|\overline{y}\| \|E_{\overline{a}}(\overline{y})\| \leq M\|y\| + \|E_{\overline{a}}(\overline{y})\| \|\overline{y}\|.$

$$\implies \frac{\|\overline{v}\|}{\|\overline{y}\|} \le M + \|E_{\overline{a}}(\overline{y})\|$$

Thus, $\overline{v} \to \overline{0}$ as $\overline{y} \to \overline{0}$. Then $f'(\overline{b})(\overline{v}) \to f'(\overline{b})(\overline{0}) = \overline{0}$. And $E_{\overline{a}}(\overline{y}) \to \overline{0}$. Therefore, $E(\overline{y}) \to \overline{0} + M\overline{0} = \overline{0}$ as $\overline{y} \to \overline{0}$.

Matrix form of the chain rule

Let $f: \mathbb{R}^n \to \mathbb{R}^m$, $g: \mathbb{R}^p \to \mathbb{R}^n$. And $h = f \circ g: \mathbb{R}^p \to \mathbb{R}^m$. Suppose g is differentiable at $\overline{a} \in \mathbb{R}^p$ and f is differentiable at $g(\overline{a}) = \overline{b} \in \mathbb{R}^n$. Then h is differentiable at \overline{a} and the Jacobian matrix of h is given by the chain rule,

$$Dh(\overline{a}) = Df(\overline{b})Dg(\overline{a})$$
 where $h = f \circ g, \ \overline{b} = g(\overline{a})$

In other words,

$$D_{j}h_{i}(\overline{a}) = \sum_{k=1}^{n} D_{k}f_{i}(\overline{b})D_{j}g_{k}(\overline{a}), \ i = 1, 2, \dots, m, \ j = 1, 2, \dots, p$$

For
$$m = 1$$
, $D_j h(\overline{a}) = \sum_{k=1}^n D_k f(\overline{b}) D_j g_k(\overline{a})$

For
$$m=1$$
 and $p=1$, $h'(\overline{a})=\sum_{k=1}^n Df(\overline{b})g'_k(\overline{a})=\nabla f(\overline{b})\cdot Dg(\overline{a})$

Theorem 13.2.7. Let f and D_2f be continuous functions on a rectangle $[a,b] \times [c,d]$. Let p and q be differentiable on [c,d], where $p(y) \in [a,b]$ and $q(y) \in [c,d]$ for each $y \in [c,d]$. Define F by the equation,

$$F(y) = \int_{p(y)}^{q(y)} f(x, y) dx, \ y \in [c, d]$$

Then F'(y) exists for each $y \in (c,d)$ and is given by,

$$F'(y) = \int_{p(y)}^{q(y)} D_2 f(x, y) dx + f((q, y), y) q'(y) - f(p(y), y) p'(y)$$

The following two theorems are required for proving the theorem on differentiating an integral.

Theorem 13.2.8. Let α be of bounded variation on [a,b] and assume that $f \in \mathcal{R}(\alpha)$ on [a,b].

Define
$$F(x) = \int_{a}^{x} f \ d\alpha, \ x \in [a, b]$$

Then F is of bounded variation on [a,b] and F is continuous at x if α is continuous at x. If α is increasing on [a,b], then the derivative F'(x) exists at each $x \in (a,b)$ where $\alpha'(x)$ exists and where f is continuous. And

$$F'(x) = f(x)\alpha'(x)$$

Theorem 13.2.9. Let $Q = \{(x,y) : a \le x \le b, c \le y \le d\}$. Assume that α is of bounded variation on [a,b] and for each $y \in [c,d]$, assume that the integral

$$F(y) = \int_a^b f(x,y) \ d\alpha(x)$$

exists. If the partial derivative D_2f is continuous on Q, the derivative F'(y) exists for each $y \in (c,d)$ and is given by

$$F'(y) = \int_a^b D_2 f(x, y) \ d\alpha(x)$$

Proof. Let $G(x_1, x_2, x_3) = \int_{x_1}^{x_2} f(t, x_3) dt$. Then we may write F(y) in terms of G. That is, F(y) = G(p(y), q(y), y).

By 1-dimensional chain rule, we have

$$F'(y) = \frac{dF}{dy} = \frac{\partial G}{\partial p} \frac{dp}{dy} + \frac{\partial G}{\partial q} \frac{dq}{dy} + \frac{\partial G}{\partial y}$$
$$= D_1 G p'(y) + D_2 G q'(y) + D_3 G$$

D₁G

Since the variable of differentition is present in the limit of the integral, we use theorem ?? to compute the derivative of the integral. We may write, $G(p(y),q(y),y)=-\int_{q(y)}^{p(y)}f(t,y)\,dt$. We are differentiating (partially) with respect to p(y). Thus q(y),y are constants for this differentiation. Suppose $G(x,a,y)=-H(x)=-\int_a^x f(t,y)\,dt \implies D_1G=-H'(x)=-f(x,y)$. Thus, $D_1G=-f(p(y),y)$.

D₂G

Again, variable of differentiation is persent in the limit of the integral. Thus, we write, $G(p(y), q(y), y) = \int_{p(y)}^{q(y)} f(t, y) dt$. Now we are differentiating (partially) with respect to the the second component of G which is q(y). Clearly, p(y) and y are treated as constants. $G(a, x, y) = H(x) = \int_a^x f(t, y) dt \implies D_2G = H'(x) = f(q(y), y)$.

D₃G

Now the variable of integration is not affecting the limits of the integral. Also it is given that D_2f is continuous on $[a,b] \times [c,d]$. We write $G(a,b,x) = H(x) = \int_a^b f(t,x) \ dt \implies D_3G = H'(x) = \int_a^b D_2f(t,x) \ dt$. Thus $D_3G = \int_{p(y)}^{q(y)} f(t,y) \ dt$.

The mean-value theorem for differentiable functions

Theorem 13.2.10 (Mean-Value). Let S be an open subset of \mathbb{R}^n . Assume $f: S \to \mathbb{R}^m$ is differentiable at each point of S. Let \overline{x} , \overline{y} be two points in S such that $L(\overline{x}, \overline{y}) = \{t\overline{x} + (1-t)\overline{y}: t \in [0,1]\}$ is subset of S. Then for every $\overline{a} \in \mathbb{R}^m$, there exists a point $\overline{z} \in L(\overline{x}, \overline{y})$ such that

$$\overline{a}.(f(\overline{y}) - f(\overline{x})) = \overline{a}.f'(\overline{z})(\overline{y} - \overline{x})$$

Proof. Let $\overline{u} = \overline{y} - \overline{x}$. We have S is open subset and $L(\overline{x}, \overline{y}) \subset S$, thus there exists $\delta > 0$ such that $\overline{x} + t\overline{u} \in S, \forall t \in (-\delta, 1 + \delta)$. In other words, the 'Line $L(\overline{x}, \overline{y})$ ' is properly contained in S, in such a way that extending the Line from \overline{x} to \overline{y} a little bit extra one either sides is still contained in S.

Let $\overline{a} \in \mathbb{R}^m$ and $F: (-\delta, 1+\delta) \to \mathbb{R}$ defined by $F(t) = \overline{a}.f(\overline{x}+t\overline{u})$. Then F is differentiable at each $t \in (-\delta, 1+\delta)$ and the derivative $F'(t) = \overline{a}.f'(\overline{x}+t\overline{u},\overline{u})$, the directional derivative of $f(\overline{x}+t\overline{u})$ with respect to \overline{u} .

$$f'(\overline{x} + t\overline{u}, \overline{u}) = f'(\overline{x} + t\overline{u})(\overline{u}) \implies F'(t) = \overline{a}.f'(\overline{x} + t\overline{u})(\overline{u})$$

By 1-dimensional mean-value theorem, we have

$$\exists \theta \in (0,1) \text{ such that } F(1) - F(0) = F'(\theta)$$

By definition of F, $F(1) = \overline{a} \cdot f(\overline{x} + \overline{u}) = \overline{a} \cdot f(\overline{y})$. And $F(0) = \overline{a} \cdot f(\overline{x})$. Therefore,

$$F'(\theta) = F(1) - F(0) = \overline{a} \cdot f(\overline{y}) - \overline{a} \cdot f(\overline{x}) = \overline{a} \cdot (f(\overline{y}) - f(\overline{x}))$$

We also have,

$$F'(\theta) = \overline{a}.f'(\overline{x} + \theta \overline{u})(\overline{u}) = \overline{a}.f'(\overline{z})(\overline{y} - \overline{x}), \text{ where } \overline{z} = \overline{x} + \theta \overline{u} \in L(\overline{x}, \overline{y})$$

Remark. Suppose S is convex in \mathbb{R}^m . Then for every pair of points $\overline{x}, \overline{y} \in S$, $L(\overline{x}, \overline{y}) \subset S$. Thus Mean-value theorem holds for all $\overline{x}, \overline{y} \in S$.

13.3 Multivariate Calculus

13.3.1 A sufficient condition for differentiability

Theorem 13.3.1. Suppose one of the partial derivatives $D_1 f, D_2 f, \dots, D_n f$ exists at \overline{c} . And the remaining n-1 partial derivatives exists in some n-ball $B(\overline{c})$ and are continuous at \overline{c} . Then f is differentiable at \overline{c} .

Proof. Step 1 : We claim that the function $f: \mathbb{R}^n \to \mathbb{R}^m$ is differentiable at \overline{c} iff each component f_k is differentiable at \overline{c} .

Suppose f is differentiable at \overline{c} , then there exists a linear, total derivative function $f'(\overline{c})$ satisfying first-order Taylor's formula at \overline{c} . ie, $f(\overline{c} + \overline{v}) = f(\overline{c}) + f'(\overline{c})(\overline{v}) + ||\overline{v}|| E_{\overline{c}}(\overline{v})$ where $E_{\overline{c}}(\overline{v}) \to \overline{0}$ as $\overline{v} \to \overline{0}$.

$$f(\overline{c} + \overline{v}) = (f_1(\overline{c} + \overline{v}), f_2(\overline{c} + \overline{v}), \cdots, f_m(\overline{c} + \overline{v}))$$

$$f(\overline{c}) = (f_1(\overline{c}), f_2(\overline{c}), \cdots, f_m(\overline{c}))$$

$$f'(\overline{c})(\overline{v}) = (f'_1(\overline{v}), f'_2(\overline{v}), \cdots, f'_m(\overline{v}))$$

$$E_{\overline{c}}(\overline{v}) = (E_1(\overline{v}), E_2(\overline{v}), \cdots, E_m(\overline{v}))$$

where each component of the error function $E_k(\overline{v}) \to 0$ as $\overline{v} \to \overline{0}$. Also since $f'(\overline{c})$ is linear, each of its components $f'_k : \mathbb{R}^n \to \mathbb{R}$ are linear.

$$\begin{split} f(\overline{c}+\overline{v}) &= (f_1(\overline{c}+\overline{v}), f_2(\overline{c}+\overline{v}), \cdots, f_m(\overline{c}+\overline{v})) \\ &= (f_1(\overline{c}), f_2(\overline{c}), \cdots, f_m(\overline{c})) + (f_1'(\overline{v}), f_2'(\overline{v}), \cdots, f_m'(\overline{v})) \\ &+ \|\overline{v}\| \left(E_1(\overline{v}), E_2(\overline{v}), \cdots, E_m(\overline{v}) \right) \text{ where } E_k(\overline{v}) \to 0 \text{ as } \overline{v} \to \overline{0} \\ &= (f_1(\overline{c}), f_2(\overline{c}), \cdots, f_m(\overline{c})) + (f_1'(\overline{v}), f_2'(\overline{v}), \cdots, f_m'(\overline{v})) \\ &+ (\|\overline{v}\| E_1(\overline{v}), \|\overline{v}\| E_2(\overline{v}), \cdots, \|\overline{v}\| E_m(\overline{v})) \text{ where } E_k(\overline{v}) \to 0 \text{ as } \overline{v} \to \overline{0} \\ &= (f_1(\overline{c}) + f_1'(\overline{v}) + \|\overline{v}\| E_1(\overline{v}), \cdots, f_m(\overline{c}) + f_m'(\overline{v}) + \|\overline{v}\| E_m(\overline{v})) \end{split}$$

Thus first-order Taylor's forumula for f at \overline{c} gives first-order Taylor's forumula for each of its components f_k . ie, $f_k(\overline{c}+\overline{v})=f_k(\overline{c})+f'_k(\overline{v}+\|\overline{v}\|E_k(\overline{v})$ where $E_k(\overline{v})\to 0$ as $\overline{v}\to \overline{0}$. Therefore, f_k are differentiable at \overline{c} for $k=1,2,\cdots,m$.

Suppose each component f_k of f are differentiable at \overline{c} . Then there exists linear, total derivative functions f'_k satisfying first-order Taylor's formula at \overline{c} . ie, $f_k(\overline{c} + \overline{v}) = f_k(\overline{v}) + f'_k(\overline{v}) + ||\overline{v}|| E_k(\overline{v})$ where $E_k(\overline{v}) \to 0$ as $\overline{v} \to \overline{0}$.

Define $E_{\overline{c}}(\overline{v}) = (E_1(\overline{v}), E_2(\overline{v}), \dots, E_k(\overline{v}))$. Then $E_{\overline{c}}(\overline{v}) \to \overline{0}$ as $\overline{v} \to \overline{0}$. Therefore, there exists a linear, total derivative function $f'(\overline{c}) = (f'_1, f'_2, \dots, f'_m)$

satisfying first-order Taylor's formula at \overline{c} .

Thus, if each (real-valued) component function f_k are differentiable, then f is also differentiable. Therefore, it is sufficient to prove the theorem for a real-valued function.

Step 2: Telescopic Sum

Assume (without loss of generality) that D_1f exists at \overline{c} and D_2f, D_3f, \dots, D_nf exist and continuous in some n-ball $B(\overline{c})$. Suppose D_rf exists at \overline{c} and all partial derivatives execept D_rf are continuous. Then $v_0 = \overline{0}$, $v_1 = y_r\overline{u_r}$, $v_2 = y_r\overline{u_r} + y_1\overline{u_1}$, Then the following proof can be applied without any loss of generality.

Let $\overline{v} = \lambda \overline{y}$ where $\overline{y} = \frac{\overline{v}}{\|\overline{v}\|}$. Clearly, $\|\overline{y}\| = 1$ and $\lambda = \|\overline{v}\|$. Choose $\lambda > 0$ such that $\overline{c} + \overline{v} \in B(\overline{c})$ and all the partial derivatives $D_2 f, D_3 f, \dots, D_n f$ exists and are continuous in $B(\overline{c})$.

We have,
$$\overline{y} = (y_1, y_2, \dots, y_n) = y_1 \overline{u_1} + y_2 \overline{u_2} + \dots + y_n \overline{u_n}$$
.

Define
$$\overline{v_0} = \overline{0}$$
, $\overline{v_1} = y_1 \overline{u_1}$, \cdots , $\overline{v_n} = y_1 \overline{u_1} + y_2 \overline{u_2} + \cdots + y_n \overline{u_n}$.

$$f(\overline{c} + \overline{v}) - f(\overline{c}) = (f(\overline{c} + \lambda \overline{v_n}) - f(\overline{c} + \lambda \overline{v_{n-1}}))$$

$$+ (f(\overline{c} + \lambda \overline{v_{n-1}}) - f(\overline{c} + \lambda \overline{v_{n-2}}))$$

$$+ \dots + (f(\overline{c} + \lambda \overline{v_1}) - f(\overline{c} + \lambda \overline{v_0}))$$

$$= \sum_{k=1}^{n} f(\overline{c} + \lambda \overline{v_k}) - f(\overline{c} + \lambda \overline{v_{k-1}})$$

$$= \sum_{k=1}^{n} f(\overline{c} + \lambda \overline{v_{k-1}} + \lambda y_k \overline{u_k}) - f(\overline{c} + \lambda \overline{v_{k-1}})$$

Step 3: Mean-value theorem

Define $\overline{b_k} = \overline{c} + \lambda \overline{v_{k-1}}$. Then we have

$$f(\overline{c} + \overline{v}) - f(\overline{c}) = \sum_{k=1}^{n} f(\overline{b_k} + \lambda y_k \overline{u_k}) - f(\overline{b_k})$$
 (13.6)

We know that all partial derivatives exists in $B(\overline{c})$. Thus f is continuous in $B(\overline{c})$. Therefore by mean-value theorem we have,

$$f(\overline{b_k} + \lambda y_k \overline{u_k}) - f(\overline{b_k}) = \lambda y_k D_k f(\overline{a_k}) \text{ where } \overline{a_k} \in L(\overline{b_k}, \overline{b_k} + \lambda y_k \overline{u_k})$$

$$f(\overline{c} + \overline{v}) - f(\overline{c}) = \lambda \sum_{k=1}^{n} y_k D_k f(\overline{a_k}) \text{ where } \overline{a_k} \in L(\overline{b_k}, \overline{b_k} + \lambda y_k \overline{u_k})$$
 (13.7)

Step 4: Continuity of partial derivatives in $B(\overline{c})$

As $\lambda \to 0$, $\overline{v} \to \overline{0}$. And both $\overline{b_k}$, $\overline{b_k} + \lambda y_k \overline{u_k} \to \overline{c}$. Clearly, $\overline{a_k}$ in the line between $\overline{b_k}$ and $\overline{b_k} + \lambda y_k \overline{u_k}$ also converges to \overline{c} .

For $k \geq 2$, $D_k f$ are continuous in the *n*-ball $B(\overline{c})$. Thus $D_k f(\overline{a_k}) \to D_k f(\overline{c})$. We may write, $D_k f(\overline{a_k}) = D_k f(\overline{c}) + E_k(\lambda)$ where $E_k(\lambda) \to \overline{0}$ as $\lambda \to 0$. Also, since $D_1 f$ exists, $D_1 f(\overline{c} + \lambda y_1 \overline{u_1}) \to D_1 f(\overline{c})$ as $\lambda \to 0$.

Remember:
$$D_1 f(\overline{c}) = \lim_{h \to 0} \frac{f(\overline{c} + h\overline{u_1}) - f(\overline{c})}{h}$$

$$f(\overline{c} + \overline{v}) - f(\overline{c}) = \lambda \sum_{k=1}^{n} y_k D_k f(\overline{c}) + \lambda \sum_{k=1}^{n} y_k E_k f(\lambda)$$
$$= \nabla f(\overline{c}) \cdot \overline{v} + \|\overline{v}\| E(\lambda)$$
where $E(\lambda) = \sum_{k=1}^{n} y_k E_k(\lambda) \to \overline{0}$ as $\overline{v} \to \overline{0}$

That is, we have a linear function which satisfies first-order Taylor's formula at \bar{c} . Therefore, f is differentiable.

13.3.2 Sufficient conditions for the equality of mixed partial derivatives

Let $f: \mathbb{R}^n \to \mathbb{R}^m$. Then $D_r f$ and $D_k f$ are two partial derivatives of f. And $D_{r,k} f = D_r(D_k f)$ and $D_{k,r} f = D_k(D_r f)$.

$$D_{r,k}f = \frac{\partial^2 f}{\partial x_r \partial x_k} = \frac{\partial}{\partial x_r} \frac{\partial f}{\partial x_k} \text{ and } D_{k,r}f = \frac{\partial^2 f}{\partial x_k \partial x_r} = \frac{\partial}{\partial x_k} \frac{\partial f}{\partial x_r}$$

There are two sufficient conditions for the equality of these mixed partial derivatives in our scope. 1. differentiability of $D_k f$ or 2. continuity of $D_{r,k} f$ and $D_{k,r}$ at \bar{c} where the mixed partial derivatives are to be equal.

Differentiability

Theorem 13.3.2. Suppose $D_r f$ and $D_k f$ exists in an n-ball about \overline{c} and are both differentiable at \overline{c} . Then $D_{r,k} f = D_{k,r} f$.

Proof. Step 1: It is sufficient to prove the theorem for real-valued functions.

Let $f: \mathbb{R}^n \to \mathbb{R}^m$, then $f(\overline{c}) = (f_1(\overline{c}), f_2(\overline{c}), \cdots, f_m(\overline{c}))$. And $D_k f(\overline{c}) = (D_k f_1(\overline{c}), D_k f_2(\overline{c}), \cdots, D_k f_m(\overline{c}))$. Thus it is sufficient to prove that $D_{r,k} f_j(\overline{c}) = D_{k,r} f_j(\overline{c})$, $j = 1, 2, \cdots, m$. That is, it is sufficient to prove equality of mixed partial derivatives of a real-valued function $f_j: \mathbb{R}^n \to \mathbb{R}$. Also, we will prove it for n = 2 and $\overline{c} = (0, 0)$. From Step 2 onwards, we will write f instead of f_j for ease of notation.

Step 2: $\nabla(h)$

Let $f: \mathbb{R}^n \to \mathbb{R}$. Suppose that the partial derivatives $D_k f$, $D_r f$ exist in the n-ball B(n). And let h > 0 such that the rectangle with vertices (0,0), (0,h), (h,0), (h,h) lies in B(n). Suppose n=3, c=(x,y,z), and we want to prove equality of $D_{2,3}f$ and $D_{3,2}f$. Then we will consider the rectangle with vertices (x,y,z), (x,y,z+h), (x,y+h,z), (x,y+h,z+h). Again, we are taking n=2 and $\overline{c}=(0,0)$, only for the ease of notation as the same proof is applicable for any finite natural number, n and any vector $\overline{c} \in \mathbb{R}^n$.

Define
$$\nabla(h) = f(h,h) - f(h,0) - f(0,h) + f(0,0)$$
.

Step 3:
$$D_{1,2}f = \frac{\nabla(h)}{h^2} = D_{2,1}f$$

Define G(x) = f(x,h) - f(x,0). Then we have, $\nabla(h) = G(h) - G(0)$ and $G'(x) = D_1 f(x,h) - D_1 f(x,0)$. By 1-dimensional mean value theorem,

$$G(h) - G(0) = hG'(x_1)$$
 where $x_1 \in (0, h)$
= $h(D_1 f(x_1, h) - D_1 f(x_1, 0))$

We have D_1f is differentiable at (0,0). There exists linear, total derivative function $(D_1f)'(0,0)$ where $(D_1f)'(0,0)(x,y) = \nabla D_1f(0,0) \cdot (x,y)$ satisfying first-order Taylor's formula at (0,0).

Remember :
$$f'(\overline{c})(\overline{v}) = \sum_{k=1}^{n} v_k D_k f(\overline{c}) = \nabla f(\overline{c}) \cdot \overline{v}$$

$$D_1 f((0,0) + (x_1,h)) = D_1 f(0,0) + \nabla D_1 f(0,0) \cdot (x_1,h) + \|(x_1,h)\| E_1(h)$$

where $E_1(h) \to 0$ as $h \to 0$

$$D_1 f(x_1, h) = D_1 f(0, 0) + x_1 D_{1,1} f(0, 0) + h D_{2,1} f(0, 0) + \left| \sqrt{x_1^2 + h^2} \right| E_1(h)$$

Similarly,

$$D_1 f((0,0) + (x_1,0)) = D_1 f(0,0) + \nabla D_1 f(0,0) \cdot (x_1,0) + \|(x_1,0)\| E_2(h)$$
where $E_2(h) \to 0$ as $h \to 0$

$$D_1 f(x_1,0) = D_1 f(0,0) + x_1 D_{1,1} f(0,0) + |x_1| E_2(h)$$

Therefore $\nabla(h) = h(D_1 f(x_1, h) - D_1 f(x_1, 0)) = h^2 D_{2,1} f(0, 0) + E(h)$ where $E(h) = h |\sqrt{x_1^2 + h^2}|E_1(h) - h|x_1|E_2(h)$ and $E(h) \to 0$ as $h \to 0$.

$$\lim_{h \to 0} \frac{\nabla(h)}{h^2} = D_{2,1} f(0,0) \tag{13.8}$$

Define H(y) = f(h, y) - f(0, y). Then we have, $\nabla(h) = H(h) - H(0)$ and $H'(y) = D_2 f(h, y) - D_2 f(0, y)$. By 1-dimensional mean value theorem,

$$H(h) - H(0) = hH'(y_1)$$
 where $y_1 \in (0, h)$
= $h(D_2f(h, y_1) - D_2f(0, y_1))$

We have $D_2 f$ is differentiable at (0,0). Thus there exists a linear, total derivative function $(D_2 f)'(0,0)$ where $(D_2 f)'(0,0)(x,y) = \nabla D_2 f(0,0) \cdot (x,y)$ satisfying first-order Taylor's formula at (0,0). That is,

$$D_2 f((0,0) + (h,y_1)) = D_2 f(0,0) + \nabla D_2 f(0,0) \cdots (h,y_1) + ||(h,y_1)|| E_3(h)$$
where $E_3(h) \to 0$ as $h \to 0$

$$D_2 f(h,y_1) = D_2 f(0,0) + h D_{1,2} f(0,0) + y_1 D_{2,2} f(0,0) + \left| \sqrt{h^2 + y_1^2} \right| E_3(h)$$

Again,

$$D_2 f((0,0) + (0,y_1)) = D_2 f(0,0) + \nabla D_2 f(0,0) \cdots (0,y_1) + |y_1| E_4(h)$$
where $E_4(h) \to 0$ as $h \to 0$

$$D_2 f(0,y_1) = D_2 f(0,0) + y_1 D_{2,2} f(0,0) + |y_1| E_4(h)$$

Therefore, $\nabla(h) = h(D_2 f(h, y_1) - D_2 f(0, y_1)) = h^2 D_{1,2} f(0, 0) + E'(h)$ where $E'(h) = \left| \sqrt{h^2 + y_1^2} \right| E_3(h) - |y_1| E_4(h)$ and $E'(h) \to 0$ as $h \to 0$.

$$\lim_{h \to 0} \frac{\nabla(h)}{h^2} = D_{1,2} f(0,0) \tag{13.9}$$

Therefore, $D_{1,2}f(0,0) = D_{2,1}f(0,0)$.

Continuity

Theorem 13.3.3. Suppose $D_r f$ and $D_k f$ exists in an n-ball about \overline{c} . And $D_{r,k} f$ and $D_{k,r} f$ are continuous at \overline{c} . Then $D_{r,k} f = D_{k,r} f$.

Proof. We have $D_r f = (D_r f_1, D_r f_2, \cdots, D_r f_m)$. Therefore, it is sufficient to prove the theorem for real-valued functions. Suppose n = 2, $\bar{c} = (0,0)$ and the partial derivatives $D_1 f$ and $D_2 f$ exist and are continuous in some 2-ball about (0,0). Suppose (h,h) lies in that 2-ball, then $D_1 f(h,h) \to D_1 f(0,0)$ as $h \to 0$.

Remark. A function f such that $D_{1,2}f \neq D_{2,1}f$.

Let,
$$f(x,y) = \begin{cases} \frac{xy(x^2 - y^2)}{x^2 + y^2} & (x,y) \neq (0,0) \\ 0 & (x,y) = (0,0) \end{cases}$$

$$D_1 f(x,y) = \frac{\partial}{\partial x} \frac{x^3 y - xy^3}{x^2 + y^2}$$

$$= \frac{(3x^2 y - y^3)(x^2 + y^2) - 2x(x^3 y - xy^3)}{(x^2 + y^2)^2}$$

$$= \frac{x^4 y + 4x^2 y^3 - y^5}{(x^2 + y^2)^2}$$

$$D_1 f_{(x=0)} = -y \implies D_{2,1} f_{(x=0)} = \frac{\partial}{\partial y} - y = -1$$

$$D_2 f(x,y) = \frac{\partial}{\partial y} \frac{x^3 y - xy^3}{x^2 + y^2}$$

$$= \frac{(x^3 - 3xy^2)(x^2 + y^2) - 2y(x^3 y - xy^3)}{(x^2 + y^2)^2}$$

$$= \frac{x^5 - 4x^3 y^2 - xy^4}{(x^2 + y^2)^2}$$

$$D_2 f_{(y=0)} = x \implies D_{1,2} f_{(y=0)} = \frac{\partial}{\partial x} x = 1$$

Therefore, $D_{1,2}f \neq D_{2,1}f$ in the neighbourhood of (0,0). This treatment save a lot of time. After D_1f , we are planning to perform $D_{2,1}f = D_2(D_1f)$ in which the value of x is going to be treated as a constant. Therefore, we can simplify the expression by substituting x = 0 at this stage. If you are not confident enough to substitute that "early". You may take partial derivative with respect to y and then substitute x = 0 and y = 0. Why don't we substitute y = 0 before D_2f is something you should know already!

13.3.3 Implicit Functions and Extremum Problems

Definitions 13.3.4 (Implicit function). Let f be a function. Consider the equation, $f(\overline{x}, \overline{y}) = 0$. If there exists a function g such that $\overline{x} = g(\overline{y})$, then g is an implicit form of f or g is defined implicitly by f.

For example, a linear system of equations Ax - b = 0 implicitly defines $x = A^{-1}b$ provided A has non-zero determinant.

Definitions 13.3.5 (Jacobian Determinant). Let $f : \mathbb{R}^n \to \mathbb{R}^n$, then determinant of the Jacobian matrix $Df(\overline{x})$ is the Jacobian determinant of f, $J_f(\overline{x})$.

Theorem 13.3.6. Let $f: \mathbb{C} \to \mathbb{C}$. Then $J_f(z) = |f'(z)|^2$.

Proof. Suppose $f: \mathbb{C} \to \mathbb{C}$ where f(z) = u(z) + iv(z) where $u: \mathbb{C} \to \mathbb{R}$ and $v: \mathbb{C} \to \mathbb{R}$. These real-valued functions u, v have respective u^*, v^* multivariate real functions such that $u^*: \mathbb{R}^2 \to \mathbb{R}$, where $u(z) = u^*(x, y)$ and z = x + iy. Let $f(z) = z^2 + 1$. Then $u^*(x, y) = x^2 - y^2 + 1$ and $v^*(x, y) = -2xy$. And theoretically we use derivatives of u^* when we mention derivatives of u.

Then f has a derivative at z only if the partial derivatives D_1u , D_2u , D_1v , D_2v exists at z and satisfies Cauchy-Riemann equations. ie $D_1u(z) = D_2v(z)$ and $D_2v(z) = -D_1v(z)$.[?, Theorem 5.22].

Thus we have $f'(z) = D_1 u + i D_1 v$ [?, Theorem 12.6]¹.

$$f'(z) = D_1 u(z) + i D_1 v(z)$$
$$|f'(z)|^2 = (D_1 u(z))^2 + (D_1 v(z))^2$$

For ease of representation, we write D_1u instead of $D_1u(z)$

$$|f'(z)|^2 = (D_1 u)^2 + (D_1 v)^2$$

¹Prove using first-order Taylor's formula

We also have

$$J_f(z) = |Df(z)| = \begin{vmatrix} D_1 u & D_2 u \\ D_1 v & D_2 v \end{vmatrix} = D_1 u D_2 v - D_1 v D_2 u = (D_1 u)^2 + (D_1 v)^2$$
Therefore, $J_f(z) = |f'(z)|^2$.

Functions with non-zero Jacobian determinant

That is, $f: \mathbb{R}^n \to \mathbb{R}^n$ such that $J_f \neq 0$ in an *n*-ball. In other words, we have an *n*-ball $B(\overline{x})$ such that $J_f(\overline{y}) \neq 0$, $\forall \overline{y} \in B(\overline{x})$.

Theorem 13.3.7. Let B be an n-ball about \overline{a} in \mathbb{R}^n , ∂B be its boundary and $\overline{B} = B \cup \partial B$ be its closure. Let $f: \mathbb{R}^n \to \mathbb{R}^n$ be continuous in \overline{B} and all partial derivatives, $D_j f_i(\overline{x})$ exists for every $\overline{x} \in B$. Let $f(\overline{x}) \neq f(\overline{a})$ for every $\overline{x} \in \partial B$ and $J_f(\overline{x}) \neq 0$ for every $\overline{x} \in B$. Then f(B) contains an n-ball about $f(\overline{a})$.

$$B = \{ \overline{x} : ||\overline{x} - \overline{a}|| < r \}$$
$$\partial B = \{ \overline{x} : ||\overline{x} - \overline{a}|| = r \}$$
$$\overline{B} = \{ \overline{x} : ||\overline{x} - \overline{a}|| \le r \}$$

Proof. Define $g: \partial B \to \mathbb{R}$ where $g(\overline{x}) = \|f(\overline{x}) - f(\overline{a})\|$. We have, $f(\overline{x}) \neq f(\overline{a})$ for every $\overline{x} \in \partial B$, thus $g(\overline{x}) > 0$ for every $\overline{x} \in \partial B$. Function f is continuous on \overline{B} , thus g is continuous on \overline{B} and thus g is continuous on its subset ∂B . Since ∂B is compact, every continuous function on ∂B attains its extrema³ and thus g attains its minimum value m > 0 somewhere on ∂B .

Consider *n*-ball T about $f(\overline{a})$ with radius $\frac{m}{2}$,

$$T = B\left(f(\overline{a}), \frac{m}{2}\right) = \left\{\overline{y} \in \mathbb{R}^n: \|f(\overline{a}) - \overline{y}\| < \frac{m}{2}\right\}$$

Therefore, it is sufficient to prove that $T \subset f(B)$.

Let $\overline{y} \in T$. Define $h: \overline{B} \to \mathbb{R}$ where $h(\overline{x}) = \|f(\overline{x}) - \overline{y}\|$. Again this continuous function h on compact set \overline{B} attains its extrema somewhere on \overline{B} . Since $\overline{y} \in T$, $h(\overline{a}) = \|f(\overline{a}) - \overline{y}\| < \frac{m}{2}$. Thus, the minimum of h on \overline{B} is less than $\frac{m}{2}$, since $\overline{a} \in \overline{B}$.

Let $\overline{x} \in \partial B$, then

$$\begin{split} h(\overline{x}) = & \| f(\overline{x}) - \overline{y} \| \\ = & \| f(\overline{x}) - f(\overline{a}) + f(\overline{a}) - \overline{y} \| \\ \geq & \| f(\overline{x}) - f(\overline{a}) \| + \| f(\overline{a}) - \overline{y} \| \\ = & g(\overline{x}) - h(\overline{a}) \\ > & \frac{m}{2} \text{ since } g(\overline{x}) \geq m \text{ and } h(\overline{a}) < \frac{m}{2} \end{split}$$

 $^{{}^2\}overline{B}$: The line above B has a different meaning compare to \overline{a} (situations like this are an abuse of language).

³ "Every continuous function on a compact set attains its extrema"

Thus h doesn't attain its minimum on ∂B , but at an interior point $\bar{c} \in B$. Consider

$$h^{2}(\overline{x}) = ||f(\overline{x}) - \overline{y}||^{2} = \sum_{r=1}^{n} (f_{r}(\overline{x}) - y_{r})^{2}$$

The function h^2 also has minimum at the same point \bar{c} . Thus all partial derivatives of h^2 at \bar{c} are zero. ie,

$$D_k h^2(\overline{c}) = \sum_{r=1}^n (f_r(\overline{c}) - y_r) D_k f_r(\overline{c}) = 0$$

This is a system of linear equations with non-zero determinant since $\overline{c} \in B$ and we have $J_f(\overline{c}) \neq 0$. Therefore, $f_r(\overline{c}) = y_r$. That ie, $f(\overline{c}) = y \in f(B)$. Since $y \in T$ is arbitrary, $T \subset f(B)$.

Theorem 13.3.8. Let A be an open subset of \mathbb{R}^n and $f: A \to \mathbb{R}^n$ is continuous and has continuous partial derivatives $D_j f_i$ on A. If f is one-to-one on A and $J_f(\overline{x}) \neq 0$, $\forall \overline{x} \in A$, then f(A) is open.

Proof. Let $\overline{b} \in f(A)$. Then $\overline{b} = f(\overline{a})$ for some $\overline{a} \in A$. Since f is continuous, f has continuous partial derivatives on A and $J_f(\overline{x}) \neq 0$ for every $\overline{x} \in A$, there exists an open ball $B \subset A$ containing \overline{a} such that $f(B) \subset f(A)$ contains an n-ball about $f(\overline{a})$. Since $\overline{b} \in f(A)$ is arbitary, every point in f(A) has an n-ball containing it in f(A). Therefore, f(A) is open.

Two assumption in above theorem are trivial. 1. f is continuous in the closed ball, \overline{B} . Set B so chosen that $\overline{B} \subset A$ and f is continuous in A. Thus, f is continuous in \overline{B} . 2. f has different value at boundary compared to center. ie, $f(\overline{a}) \neq f(\overline{x})$, $\forall x \in \partial B$. We have, f is injective on A, and $\overline{B} \subset A$. Thus f has different values for any two distinct points in it. Thus, $\forall \overline{x}, \overline{y} \in \overline{B}, \ \overline{x} \neq \overline{y} \Longrightarrow \overline{x}, \overline{y} \in A$, and $\overline{x} \neq \overline{y} \Longrightarrow f(\overline{x}) \neq f(\overline{y})$

Theorem 13.3.9. Let S be an open subset of \mathbb{R}^n and $f: S \to \mathbb{R}^n$. Let components of f has continuous partial derivatives on S, $D_j f_i$ and $J_f(\overline{a}) \neq 0$ for some point $\overline{a} \in S$. Then there is an n-ball B about \overline{a} on which f is injective.

Proof. Let $\overline{z} = (\overline{z_1}, \overline{z_2}, \dots, \overline{z_n})$ where $\overline{z_i} \in \mathbb{R}^n$. ie, $\overline{z} \in \mathbb{R}^{n^2}$. Define function $h : \mathbb{R}^{n^2} \to \mathbb{R}$ by $h(\overline{z}) = \det[D_j f_i(\overline{z_i})]$. Since f has continuous partial derivatives on S, each component of f has continuous partial derivatives in S and thus h is continuous on S^n which is a subset of \mathbb{R}^{n^2} since S is an open subset of \mathbb{R}^n .

Let $\overline{a} \in S$ such that $J_f(\overline{a}) \neq 0$. Existence of such a point in S is assumed. Consider, $\overline{z_i} = \overline{a}$, $\forall i$. Then $\overline{z} = (\overline{a}, \overline{a}, \dots, \overline{a})$. And $h(\overline{z}) = \det[D_j f_i(\overline{a})] = J_f(\overline{a}) \neq 0$.

Since h is continous and $h(\overline{z}) \neq 0$. There exists an n-ball B about \overline{a} in S such that $h(\overline{z}) \neq 0$ for $\overline{z_i} \in B$. We claim that f is injective on B.

Theorem 13.3.10.

Proof. \Box

Remark (Properties). Functions with non-zero Jacobian determinant has following properties:

- 1. If $J_f \neq 0$ in n-ball B about \overline{a} which has different values at its boundaries, then f(B) has an n-ball about $f(\overline{a})$.
- 2. If $J_f \neq 0$, f has continuous partial derivatives in S, and f is injective in an open set A, then f(A) is open.
- 3. Let S be an open set in \mathbb{R}^n , f has continuous partial derivatives in S, and $J_f(\overline{a}) \neq 0$ for some $\overline{a} \in S$, then f is injective on an n-ball $B(\overline{a})$ in S.
- 4. Let A be an open set in \mathbb{R}^n , f has continuous partial derivatives in A, and $J_f \neq 0$ in A, then f is an open mapping.

Inverse function Theorem

Theorem 13.3.11 (Inverse function). Let S be an open subset of \mathbb{R}^n and f be a continuously differentiable function $f: S \to \mathbb{R}^n$. If $J_f(\overline{a}) \neq 0$ for some $\overline{a} \in S$, then there are two open sets $X \subset S$, and $Y \subset f(S)$ such that $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are function $f(\overline{a}) \in Y$ and $f(\overline{a}) \in Y$ are

Proof. \Box

Implicit function Theorem

Theorem 13.3.12 (Implicit function). Let S be an open subset of \mathbb{R}^{n+k} and f be a function $f: S \to \mathbb{R}^n$. Suppose f is continuously differentiable on S. Let $(\overline{x_0}, \overline{t_0}) \in \mathbb{R}^n \times \mathbb{R}^k$ such that $f(\overline{x_0}, \overline{t_0}) = \overline{0}$ and $J_f(\overline{x_0}) \neq 0$. Then there exists an open set T_0 containing $\overline{t_0}$ in \mathbb{R}^k and a unique function $g: \mathbb{R}^k \to \mathbb{R}^n$ such that 1. g is continuously differentiable on T_0 2. $g(\overline{t_0}) = \overline{x_0}$ 3. $f(g(\overline{t}), \overline{t}) = \overline{0}$, $\forall \overline{t} \in T_0$

Proof.

Extrema of function of one variable

Theorem 13.3.13 (Extrema of functions of one variable). Let $n \geq 1$ and function f has nth partial derivative in open interval (a,b). Suppose for some $c \in (a,b)$,

$$f'(c) = f''(c) = \dots = f^{(n-1)}(c) = 0$$
 and $f^{(n)}(c) \neq 0$

If n is even, 1. f has a local minimum at c if $f^{(n)}(c) > 0$ 2. f has a local maximum at c if $f^{(n)}(c) < 0$ and If n is odd, there is neither a local minimum nor a local maximum at c.

Proof.

 $^{^4} f \in C'(S)$: f is continuously differentiable on S

Extrema of function of several variables

Definitions 13.3.14 (stationary point). If function f is differentiable at \overline{a} and $\nabla f(\overline{a}) = \overline{0}$, the point \overline{a} is a stationary point of f.

Definitions 13.3.15 (saddle point). A stationary point is a saddle point if every n-ball $B(\overline{a})$ contains points \overline{x} such that $f(\overline{x}) > f(\overline{a})$ and other points such that $f(\overline{x}) < f(\overline{a})$.

Definitions 13.3.16 (quadratic form). A function $Q : \mathbb{R}^n \to \mathbb{R}$ defined $Q(\overline{x}) = \sum_{i=1}^n \sum_{j=1}^n a_{ij} x_i x_j$ where $a_{ij} \in \mathbb{R}$ is a function of the quadratic form or simply a quadratic form.

symmetric quadratic form $a_{ij} = a_{ji}, \forall i, j$.

positive definite quadratic form $Q(\overline{x}) > 0, \forall \overline{x} \neq \overline{0}$.

negative definite quadratic form $Q(\overline{x}) < 0, \forall \overline{x} \neq \overline{0}$.

Theorem 13.3.17 (Second derivative test for extrema). Let f be a function $f: \mathbb{R}^n \to \mathbb{R}^m$. Suppose that the second order partial derivatives $D_{i,j}f$ exists in an n-ball $B(\overline{a})$ and are continuous at \overline{a} where \overline{a} is a stationary point of f.

Let
$$Q(\overline{t}) = \frac{1}{2}f''(\overline{a}, \overline{t}) = \frac{1}{2}\sum_{i=1}^{n}\sum_{j=1}^{n}D_{i,j}f(\overline{a})t_it_j$$

- 1. If $Q(\overline{t}) > 0$ for all $\overline{t} \neq \overline{0}$, f has a relative minimum at \overline{a} .
- 2. If $Q(\overline{t}) < 0$ for all $\overline{t} \neq \overline{0}$, f has a relative maximum at \overline{a} .
- 3. If $Q(\overline{t})$ takes both positive and negative values, then f has a saddle point at \overline{a} .

Proof.

Theorem 13.3.18 (Second derivative test: simplified). Let f be a real-valued function $f: \mathbb{R}^2 \to \mathbb{R}$ with continuous second order partial derivatives at a stationary point $\overline{a} \in \mathbb{R}^2$. Let $A = D_{1,1}f(\overline{a})$, $B = D_{1,2}f(\overline{a}) = D_{2,1}f(\overline{a})$, and $C = D_{2,2}f(\overline{a})$. And let $\Delta = \begin{vmatrix} A & B \\ B & C \end{vmatrix} = AC - B^2$. Then we have,

- 1. If $\Delta > 0$ and A > 0, then f has a relative minimum at \overline{a} .
- 2. If $\Delta > 0$ and A < 0, then f has a relative maximum at \overline{a} .
- 3. if $\Delta < 0$, then f has a saddle point at \overline{a} .

Proof. \Box

ME010304 Functional Analysis

ME010305 Optimization Technique

Semester IV

ME010401 Spectral Theory

ME010402 Analytic Number Theory

ME800401 Differential Geometry

ME800402 Algorithmic Graph Theory

19.1 Networks

19.1.1 An Introduction to Networks

Definitions 19.1.1. A network N is a digraph D with two special vertices source s and sink t together with a capacity function $c: E(D) \to \mathbb{Z}$ such that for every arc a = (u, v) of the digraph, c(u, v) is non-negative.

Remark. Mathematical Modeling using Network,

- 1. There is no restriction on indegree/outdegree of source/sink vertices of the digraph D of a network N.
- 2. Applications of Network: Transportation problem.

c(u,v) is the capacity of the arc (u,v) of D

 $N^+(x) = \{y \in V(D) : (x,y) \in E(D)\}$ is the out-neighbourhood of x.

$$N^-(x) = \{y \in V(D) : (y, x) \in E(D)\}$$
 is the in-neighbourhood of x.

Definitions 19.1.2. A flow f in a network N is function $f: E(D) \to \mathbb{Z}$ such that 1. each edge satisfies capacity constraint and 2. each vertex except source and sink satisfies conservation equation.

capacity constraint

$$0 \le f(a) \le c(a)$$
 for every arc $a \in V(D)$ (19.1)

conservation equation

$$\sum_{y \in N^{+}(x)} f(x,y) = \sum_{y \in N^{-}(x)} f(y,x), \text{ Vertex } x \in V(D) - \{s,t\}$$
 (19.2)

net flow out of x

$$\sum_{y \in N^+(x)} f(x,y) - \sum_{y \in N^-(x)} f(y,x)$$

net flow into x

$$\sum_{y \in N^-(x)} f(y,x) - \sum_{y \in N^+(x)} f(x,y)$$

Definitions 19.1.3. The flow f in a network N is the net flow out of source s.

Remark. 1. net flow out of/into $x \in V(D) - \{s, t\}$ is zero.

2. Without loss of generality¹, underlying digraph is always assymetric.

$$(X,Y) = \{(x,y) \in E(D) : x \in X, y \in Y\}.$$

Let X Y be non-empty subsets of $V(D)$ such that X Y

Let X, Y be non-empty subsets of V(D) such that X, Y are disjoint. Then (X, Y) is the set of all arcs from X to Y.

flow from X **to** Y is the sum of flow on each arc in (X,Y)

$$f(X,Y) = \sum_{(x,y)\in(X,Y)} f(x,y)$$
 (19.3)

capacity of the partition (X,Y) is the total capacity of arcs in (X,Y)

$$c(X,Y) = \sum_{(x,y)\in(X,Y)} c(x,y)$$
 (19.4)

cut Let $P \subset V(D)$ such that $s \in P$ and $t \notin P$ and $\bar{P} = V(D) - P$, then (P, \bar{P}) is a cut.

flow from P to \bar{P} is the sum of flow on each arc in (P, \bar{P}) .

$$f(P,\bar{P}) = \sum_{(x,y)\in(P,\bar{P})} f(x,y)$$
 (19.5)

flow from \bar{P} to P is the sum of flow on each arc in (\bar{P}, P)

$$f(\bar{P}, P) = \sum_{(x,y)\in(\bar{P}, P)} f(x,y)$$
 (19.6)

capacity of the cut (P, \bar{P}) is the total capacity of the arcs in (P, \bar{P})

$$c(P, \bar{P}) = \sum_{(x,y)\in(P,\bar{P})} c(x,y)$$
 (19.7)

Theorem 19.1.4. For any cut (P, \bar{P}) , the flow in N is $f(N) = f(P, \bar{P}) - f(\bar{P}, P)$.

Synopsis. The net flow out of source s is the flow f(N) in the network N. Let (P, \bar{P}) be a cut of N, then $s \in P$ and $t \notin P$. Suppose $P = \{s\}$, then the theorem is true. Suppose P is not singleton, then for each vertex $x \in P$, $x \neq s$, the net flow out of x is zero by flow conservation equation. And flow between vertices in P cancels out each other. Thus adding net flow out of each vertex in P, will be same as the net flow out of source which is the flow in the network, f(N).

¹If underlying digraph of a network is symmetric, then by replacing an arc (u, v) with a new vertex w and two arcs (u, w), (w, v) gives an assymetric digraph. [?] pp.131

Proof.

Flow,
$$f = \sum_{y \in N^+(s)} f(s, y) - \sum_{y \in N^-(s)} f(y, s)$$
 (19.8)

By conservation equation, we have $\forall x \in P, x \neq s$,

$$\sum_{y \in N^{+}(x)} f(x,y) - \sum_{y \in N^{-}(x)} f(y,x) = 0$$
 (19.9)

By above equations,

Flow,
$$f = \sum_{x \in P} \sum_{y \in N^+(x)} f(x, y) - \sum_{x \in P} \sum_{y \in N^-(x)} f(y, x)$$

$$= \sum_{(x,y) \in (P,\bar{P})} f(x,y) - \sum_{(y,x) \in (\bar{P},P)} f(y,x)$$
(19.10)

Corollary 19.1.4.1. Flow cannot exceed the capacity of any cut (P, \bar{P}) . Further, $f(N) \leq \min c(P, \bar{P})$.

Synopsis. Let (P, \bar{P}) be a cut in network N, then by theorem the flow f(N) = flow from P to \bar{P} - flow from \bar{P} to P. Since the flow from \bar{P} to P is nonnegative, $f(N) \leq flow$ from P to \bar{P} . Clearly, $f(x,y) \leq c(x,y)$ by the capacity constaint. Thus $f(N) \leq f(P, \bar{P}) \leq c(P, \bar{P}) \leq \min c(P, \bar{P})$.

Proof.

$$\begin{split} f(N) &= \sum_{(x,y) \in (P,\bar{P})} f(x,y) - \sum_{(y,x) \in (\bar{P},P)} f(y,x) \\ &\leq \sum_{(x,y) \in (P,\bar{P})} f(x,y) = f(P,\bar{P}) \\ &\leq \sum_{(x,y) \in (P,\bar{P})} c(x,y) = c(P,\bar{P}), \quad \because \forall x,y \in V(D), \ f(x,y) \leq c(x,y) \\ &\leq \min c(P,\bar{P}) \end{split}$$

Corollary 19.1.4.2. In a network N flow is the net flow into the sink of N.

Synopsis. Let $\bar{P} = \{t\}$, then by theorem f(N) is the net flow into the sink.

Proof. Suppose $P = V(D) - \{t\}$. Then by theorem, we have

$$f(N) = \sum_{(x,y)\in(P,\bar{P})} f(x,y) - \sum_{(y,x)\in(\bar{P},P)} f(y,x)$$
$$= \sum_{x\in N^{-}(t)} f(x,t) - \sum_{x\in N^{+}(t)} f(t,x)$$

Remark. Exercise 5.1

4. Let N be a network with underlying digraph D which has a vertex $v \in V(D) - \{s,t\}$ with zero indegree. Clearly the flow into v is zero. Thus flow out of v is also zero by flow conservation equation. Let N' be the network obtained from N by deleting the vertex v. Then f(N) = f(N').

19.1.2 The Max-Flow Min-Cut Theorem

maximum flow A flow f in network N is maximum flow in N, if $f(N) \ge f'(N)$ for each flow f' in N.

minimum cut A cut (P, \bar{P}) in network N is minimum cut of N, if $c(P, \bar{P}) \leq c(X, \bar{X})$ for each cut (X, \bar{X}) in N.

f-unsaturated Let f be a flow in network N with underlying digraph D, and $Q = u_0, a_1, u_1, a_2, \dots, u_{n-1}, a_n, u_n$ be a semipath in D such that every forward arc $a_i = (u_{i-1}, u_i)$ has flow not upto its capacity, $f(a_i) < c(a_i)$ and every reverse arc $a_i = (u_i, u_{i-1})$ has some positive flow in it, $f(a_i) > 0$

f-augmenting semipath Let f be a flow in a network N with underlying digraph D. Suppose semipath $Q = s, a_1, u_1, a_2, \cdots, u_{n-1}, a_n, t$ (from source to sink) is f-unsaturated, then Q is an f-augmenting semipath.

Theorem 19.1.5. Let f be a flow in a network N with underlying digraph D. The flow f is maximum in N iff there is no f-augmenting semipath in D.

Synopsis. Suppose Q is an f-augmenting semipath in D, then there exists a flow f^* in N such that $f(N) + \Delta = f^*(N)$. Therefore, f is not a maximum flow in N. Suppose there is no f-augmenting semipath in D, then there exists a cut (P, \bar{P}) such that $f(a) = c(a) \ \forall a \in (P, \bar{P})$ and $f(a) = 0 \ \forall a \in (\bar{P}, P)$. Suppose f^* in a maximum flow in N, then $f(N) \leq f^*(N) \leq c(P, \bar{P}) = f(N)$.

Proof. Let f be a flow in a network N with underlying digraph D and $Q = s, a_1, u_1, a_2, u_2, \cdots, u_{n-1}, a_n, t$ be an f-augmenting semipath in D.

define
$$\Delta_i = \begin{cases} c(a_i) - f(a_i) \text{ for every forward arc } a_i \in Q, \\ f(a_i) \text{ for every reverse arc } a_i \in Q, \end{cases}$$

Define $\Delta = \min\{\Delta_i\}$. Also define $f^* : E(D) \to \mathbb{Z}$ such that

$$f^*(a_i) = \begin{cases} f(a) + \Delta, \text{ for every forward arc } a_i \in Q, \\ f(a) - \Delta, \text{ for every reverse arc } a_i \in Q, \\ f(a_i), \text{ for every arc of } D \text{ which are not in } Q. \end{cases}$$

Since Q is an f-augmenting semipath in D, $\Delta > 0$ and $f(N) + \Delta = f^*(N)$.

Clearly $f(N) < f^*(N)$, and it is enough to show that f^* is a flow in N. f^* is a flow if it satisfies 1. capacity constraint and 2. conservation equation. For any arc $a_i \notin Q$, $f^*(a_i) = f(a_i) \le c(a_i)$. Suppose $a_i \in Q$. If $a_i = (u_{i-1}, u_i)$, a_i is a forward arc and we have $f^*(a_i) = f(a_i) + \Delta \le f(a_i) + \Delta_i = f(a_i) + c(a_i) - f(a_i) = c(a_i)$. If $a_i = (u_i, u_{i-1})$, then a_i is a reverse arc and we have $f^*(a_i) = \Delta \le \min\{\Delta_i\} = \Delta_i = c(a_i)$. Thus f^* satisfies capacity constraint on every arc of D.

19.1. NETWORKS

Let $x \in V(D) - \{s, t\}$. Suppose $x \notin Q$,

Net flow out of
$$x=\sum_{y\in N^+(x)}f^*(x,y)-\sum_{y\in N^-(x)}f^*(y,x)$$

$$=\sum_{y\in N^+(x)}f(x,y)-\sum_{y\in N^-(x)}f(y,x)$$

$$=0$$

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Suppose $x = u_i \in Q$, then Q has two arc having vertex x say, a_{i-1} , and a_i . There are four possibilities for these two arcs,

- 1. Both a_{i-1} , a_i are forward arcs.
- 2. Arc a_{i-1} is forward, but arc a_i is reverse.
- 3. Arc a_{i-1} is reverse, but arc a_i is forward.
- 4. Both a_{i-1} , a_i are reverse arcs.

Case 1 $a_{i-1} = (u_{i-1}, u_i)$ and $a_i = (u_i, u_{i+1})$.

Net flow out of
$$x = \sum_{y \in N^+(x)} f^*(x,y) - \sum_{y \in N^-(x)} f^*(y,x)$$

$$= \sum_{\substack{y \in N^+(x) \\ y \neq u_{i+1}}} f^*(x,y) + f^*(u_i,u_{i+1}) - \left(\sum_{\substack{y \in N^-(x) \\ y \neq u_{i-1}}} f^*(y,x) + f^*(u_{i-1},u_i)\right)$$

$$= \sum_{\substack{y \in N^+(x) \\ y \neq u_{i+1}}} f(x,y) + f(u_i,u_{i+1}) + \Delta - \left(\sum_{\substack{y \in N^-(x) \\ y \neq u_{i-1}}} f(y,x) + f(u_{i-1},u_i)\right) - \Delta$$

$$= \sum_{\substack{y \in N^+(x) \\ y \neq u_{i+1}}} f(x,y) - \sum_{\substack{y \in N^-(x) \\ y \in N^-(x)}} f(y,x)$$

Case 2 $a_{i-1} = (u_{i-1}, u_i)$ and $a_i = (u_{i+1}, u_i)$.

Net flow out of
$$x = \sum_{y \in N^+(x)} f^*(x,y) - \sum_{y \in N^-(x)} f^*(y,x)$$

$$= \sum_{\substack{y \in N^+(x) \\ y \neq u_{i+1}, u_{i-1}}} f^*(x,y) + f^*(u_i, u_{i+1}) + f^*(u_i, u_{i-1}) - \sum_{y \in N^-(x)} f^*(y,x)$$

$$= \sum_{\substack{y \in N^+(x) \\ y \neq u_{i+1}, u_{i-1}}} f(x,y) + f(u_i, u_{i+1}) + \Delta + f(u_i, u_{i-1}) - \Delta - \sum_{y \in N^-(x)} f(y,x)$$

$$= \sum_{y \in N^+(x)} f(x,y) - \sum_{y \in N^-(x)} f(y,x)$$

$$= 0$$

Case 3 $a_{i-1} = (u_i, u_{i-1})$ and $a_i = (u_i, u_{i+1})$.

Net flow out of
$$x = \sum_{y \in N^+(x)} f^*(x, y) - \sum_{y \in N^-(x)} f^*(y, x)$$

$$= \sum_{y \in N^+(x)} f^*(x, y) - \left(\sum_{\frac{y \in N^-(x)}{y \neq u_{i-1}, u_{i+1}}} f^*(y, x) + f^*(u_{i-1}, u_i) + f^*(u_{i+1}, u_i) \right)$$

$$= \sum_{y \in N^+(x)} f^*(x, y) - \left(\sum_{\frac{y \in N^-(x)}{y \neq u_{i-1}, u_{i+1}}} f(y, x) + f(u_{i-1}, u_i) + \Delta + f(u_{i+1}, u_i) - \Delta \right)$$

$$= \sum_{y \in N^+(x)} f(x, y) - \sum_{y \in N^-(x)} f(y, x)$$

$$= 0$$

Case 4 $a_{i-1} = (u_i, u_{i-1})$ and $a_i = (u_{i+1}, u_i)$.

Net flow out of
$$x = \sum_{y \in N^+(x)} f^*(x,y) - \sum_{y \in N^-(x)} f^*(y,x)$$

$$= \sum_{\substack{y \in N^+(x) \\ y \neq u_{i-1}}} f^*(x,y) + f^*(u_i,u_{i-1}) - \left(\sum_{\substack{y \in N^-(x) \\ y \neq u_{i+1}}} f^*(y,x) + f^*(u_{i+1},u_i)\right)$$

$$= \sum_{\substack{y \in N^+(x) \\ y \neq u_{i-1}}} f(x,y) + f(u_i,u_{i-1}) - \Delta - \left(\sum_{\substack{y \in N^-(x) \\ y \neq u_{i+1}}} f(y,x) + f(u_{i+1},u_i)\right) + \Delta$$

$$= \sum_{y \in N^+(x)} f(x,y) - \sum_{y \in N^-(x)} f(y,x)$$

Therefore, f^* is a flow on N. We have $f(N) < f^*(N)$. Thus f is not maximum flow in N due to the existence of an f-augmenting semipath in D.

Conversely, assume that there is no f-augmenting semipath in D. Now, we construct a cut (P, \bar{P}) of N. Let P be the set of all vertices $x \in V(D)$ such that there is an f-unsaturated s-x semipath in D. Trivially, $s \in P$. And $t \notin P$ since there are no f-augmenting semipath in D.² Clearly, (P, \bar{P}) is a cut of the network N.

We claim that $c(P, \bar{P}) = f(N)$. Suppose there is a forward arc $(x, y) \in (P, \bar{P})$, then flow in it is saturated. If f(x, y) < c(x, y), then there is an f-unsaturated s - y semipath in D. ie, s - x semipath + arc(x, y). Thus every forward arc $(x, y) \in (P, \bar{P})$ is saturated. Suppose there is a reverse arc $(y, x) \in (P, \bar{P})$

²An f-augmenting semipath is an f-unsaturated s-t semipath in D.

 (\bar{P}, P) , then there is no flow in it(saturated reversed arc). If f(y, x) > 0, then there is an f-unsaturated s - y semipath in D. ie, s - x semipath + arc (y, x). Thus every reverse arc $(y, x) \in (\bar{P}, P)$ is saturated. And we have,

$$\begin{split} \sum_{(x,y)\in(P,\bar{P})} f(x,y) &= \sum_{(x,y)\in(P,\bar{P})} c(x,y) \\ \sum_{(y,x)\in(\bar{P},P)} f(y,x) &= 0 \\ f(N) &= \sum_{(x,y)\in(P,\bar{P})} f(x,y) - \sum_{(y,x)\in(\bar{P},P)} f(y,x) \\ &= \sum_{(x,y)\in(P,\bar{P})} c(x,y) \\ &= c(P,\bar{P}) \end{split}$$

Suppose f^* is maximum flow in network N and (X, \bar{X}) is minimum cut of N. Then $f(N) \leq f^*(N)$. Thus we have, $f(N) \leq f^*(N) \leq c(X, \bar{X}) \leq c(P, \bar{P}) = f(N)$. Therefore, $f(N) = f^*(N)$. ie, the flow f is maximum in network N if there are no f-augmenting semipaths in D.

Theorem 19.1.6 (maximum-flow, min-cut). In every network, the value of maximum flow equals capacity of minimum cut.

Proof. Suppose flow f in network N in maximum, then by previous theorem there is no f-augmenting semipath in D. And $f(N) \leq c(X, \bar{X})$ for any cut (X, \bar{X}) in N. We can construct a cut (P, \bar{P}) in N such that $f(N) = c(P, \bar{P})$. Let P be the set of all vertices x in D such that there is an f-unsaturated s-x semipath in D. Clearly $s \in P$ and $t \notin P$. Also $f(P, \bar{P}) = c(P, \bar{P})$ and $f(\bar{P}, P) = 0$. Then the cut (P, \bar{P}) is minimum cut of N. Suppose there is a cut (X, \bar{X}) such that $c(X, \bar{X}) < c(P, \bar{P})$. Then $f(N) = f(P, \bar{P}) - f(\bar{P}, P) = c(P, \bar{P}) < c(X, \bar{X})$ which is a contradiction. Therefore, the value of maximum flow equals capacity of minimum cut.

Remark. Exercise 5.2

- 1. Suppose (X, \bar{X}) is a cut of N such that f(a) = c(a), $\forall a \in (X, \bar{X})$ and f(a) = 0, $\forall a \in (\bar{X}, X)$. By the definition of cut, $s \in X$ and $t \in \bar{X}$. Thus there is no f-augmenting semipath in D. Suppose there is an f-augmenting semipath Q in D, then there is either (a) a forward arc $(x,y) \in (X,\bar{X})$ such that f(x,y) < c(x,y) or (b) a reverse arc $(y,x) \in (\bar{X},X)$ such that f(y,x) > 0 which is a contradition. Therefore, the flow f(N) is maximum and the given cut (X,\bar{X}) is minimum as shown in the proof of the maximum-flow min-cut theorem.
- 3. The algorithm suggested in the hint of this exercise won't work if two subnetworks have a common arc such that the direction of flow in which is not consistent. Suppose, the generalized network is not supposed to have any common arcs. Then construct subnetworks for each pair (s,t) with all those arcs which are on some s-t semipath. Define subnetwork capacity function c'(a) = c(a) for every arc in N'.

Let N be a generalized network with set of sources S and set of sinks T. A flow in N is maximum if there is not f-augmenting s-t semipath for each pair $(s,t) \in S \times T$.

19.1.3 A max-flow min-cut algorithm

Theorem 19.1.7. Let N be a network with underlying digraph D, source s, sink t, capacity function c and flow f. Let D' be the digraph with same vertex set as D and arc set defined by $E(D') = \{(x,y) : (x,y) \in E(D), c(x,y) > f(x,y) \text{ or } (y,x) \in E(D), f(y,x) > 0\}$. ie, D' has only the unsaturated arcs of D. Then D' has an s-t directed path iff D has an f-augmenting semipath. Moreover, shortest s-t path in D' has the same length as shortest f-augmenting semipath in D.

Synopsis. Each directed s-t path in D' has respective f-augmenting semipath in D and vice versa. Clearly, they have the same length.

Proof. Let N be a network with underlying digraph D, capacity c and flow f. Let D' be the digraph with vertex set V(D') = V(D) and arc set $E(D') = \{(x,y) : \text{either } (x,y) \text{ or } (y,x) \text{ is unsaturated in } N\}$.

Suppose D' has a directed s-t path $Q': s, u_1, u_2, \dots, u_{n-1}, t$. Then by the construction of D', for each $u_i \in Q$, there exists an f unsaturated arc a_i in D. ie, either forward arc $a_i = (u_{k-1}, u_k)$ such that $f(u_{k-1}, u_k) < c(u_{k-1}, u_k)$ or reverse arc $a_i = (u_k, u_{k-1})$ such that $f(u_k, u_{k-1}) > 0$. Therefore, we have an s-t semipath $Q: s, a_1, u_1, a_2, \dots, u_{n-1}, a_n, t$ in D such that Q is an f-augmenting semipath since every arc in Q is f-unsaturated. Clearly, Q, Q' are of the same length.

Conversely, suppose that the digraph D has an f-augmenting semipath Q: $s, a_1, u_1, a_2, \dots, u_{n-1}, a_n, t$. Then each arc $a_i \in Q$ are f-unsaturated and by the construction of D', there exists a directed s-t path $Q'=s, u_1, u_2, \dots, u_{n-1}, t$ in D'. And Q, Q' are of the same length.

There is a one-one correspondence between the directed s-t paths in D' and f-augmenting semipaths in D. Clearly, they have the same length. Thus shortest directed s-t path in D and shortest f-augmenting semipath in D' are of the same length.

saturation arc of N with respect to the flow f is an arc a_j in an f-augmenting semipath Q with $\Delta_j = \Delta$.

augmentation path is an f-augmenting semipath Q in D.

Algorithm 19.1.8 (max-flow min-cut). An algorithm to find maximum flow and minimum cut of a network N with underlying digraph D, source s, sink t, capacity function c and initial flow f.

1. Construct digraph D' with vertex set V(D') = V(D) and arc set $E(D') = \{(x,y): (x,y) \in E(D) \& f(x,y) < c(x,y) \text{ or } (y,x) \in E(D) \& f(y,x) > 0\}$

2. Find (shortest) s-t directed path in D' using Moore's breadth first search (BFS) algorithm. If D' doesn't have an s-t path, then proceed to step 5. Otherwise, let $Q': s, u_1, u_2, \cdots, u_{n-1}, t$ be a (shortest) s-t path in D'.

- 3. Let $Q: s, a_1, u_1, a_2, \dots, u_{n-1}, a_n, t$ be the respective semipath in D such that $f(a_j) < c(a_j)$ for forward arcs and $f(a_i) > 0$ for reverse arcs. Let $\Delta_j = c(a_j) f(a_j)$ for forward arcs and $\Delta_j = f(a_j)$ for reverse arcs. And let $\Delta = \min\{\Delta_j\}$. And augment flow f by Δ ie, $f(a_j) \leftarrow f(a_j) + \Delta$ for forward arcs and $f(a_j) \leftarrow f(a_j) \Delta$ for reverse arcs.
- 4. Goto step 1 (Proceed with new flow f and find whether there are any directed s-t paths in D'. If any, augment the flow along the new augmentation path Q by saturating the flow along the saturation arc.)
- 5. There is no s-t directed path in D'. Thus there is no f-augmenting semipath in D. Therefore the flow f in N is maximum. Let P be the set of all vertices in D' with non-zero breadth first index(bfi) from Moore's BFS algorithm applied in step 2. (P, \bar{P}) is minimum cut of N.

Remark. Validity of the algorithm is proved in the previous theorem.

19.1.5 Connectivity and Edge-Connectivity

edge cutset is the set U subset of E(G) such that G-U is disconnected.

vertex cutset is the set S subset of V(G) such that G-S is disconnected.

edge connectivity $\lambda(G)$ is the minimum cardinality of all edge cutsets of G.

connectivity $\kappa(G)$ is the minimum cardinality of all vertex cutsets of G.

Theorem 19.1.9. For every graph G, $\kappa(G) \leq \lambda(G) \leq \delta(G)$

Proof. Suppose graph G is disconnected then $\kappa(G) = \lambda(G) = \delta(G) = 0$. Let G be a connected graph. Then G has at least one vertex v with degree $\delta(G)$. Therefore $\lambda(G) \leq \delta(G)$ since edges incident with v form an edge cutset of G and $\lambda(G)$ is the cardinality of all edge cutsets.

Let G be a graph with edge connectivity $\lambda(G) = c$. Let U be a edge cutset with cardinality c and let edge $uv \in U$. Construct a set of vertices $S \subset V(G)$ such that (S) is of minimal cardinality and) for each edge in U other uv, S has a vertex incident with it. Cardinality of S is at most c-1, since we can select one vertex each for each edge in U other than u, v. If G - S is a disconnected graph, then $\kappa(G) < \lambda(G)$. Suppose G - S is a connected graph, then delete a non-pendent vertex u or v from G - S, say v. Since G - S is a connected graph with a singleton edge cutset, $\{uv\}$. We have a vertex cutset $S \cup \{v\}$ of G. Therefore, $\kappa(G) \leq c = \lambda(G)$.

Theorem 19.1.10. If G is a graph of diameter 2, then $\lambda(G) = \delta(G)$

n-edge connected G is n-edge connected if $\lambda(G) \geq n$.

n connected G is n-connected if $\kappa(G) \geq n$.

Theorem 19.1.11. Let G be a graph of order p and n be an integer such that $1 \le n \le p-1$. If $\delta(G) \ge \frac{p+n-2}{2}$, then G is n-connected.

connection number c(G) is the smallest integer such that $2 \le c(G) \le p$ and every subgraph of order n in G is connected.

l-connectivity $\kappa_l(G)$ is minimum number of vertices whose removal will produce a disconnected graph with at least l components or a graph with fewer than l vertices.

(n, l)-connected A graph G is (n, l)-connected if $\kappa_l(G) \geq n$.

Remark. Exercises 5.5

1.
$$\lambda(K_{m,n}) = \kappa(K_{m,n}) = m$$

8. $c(K_p) = 2$, $c(K_{m,n}) = n+1$, $c(C_p) = p-1$ Every two vertices of complete graph of order p are adjacent. For complete bi-partitie graph $K_{m,n}$ such that $1 \le m \le n$, there exists a totally disconnected subgraph of order n. Therefore $c(K_{m,n}) \ge n+1$. And with n+1 vertices, both partitions have at least two vertices each and therefore the graph is connected and $c(K_{m,n}) \le n+1$. For cycle C_p , any subgraph is disconnected if two non-adjacent vertices are deleted. Therefore $c(C_p) \ge p-1$. And C_p remains connected even after deletion of any vertex, therefore $c(C_p) \le p-1$.

9.

$$\delta(G) \ge \frac{p + (l-1)(n-2)}{l} \implies \kappa_l(G) \ge n$$

19.1.6 Menger's Theorem

Theorem 19.1.12. For a non-trivial graph G, $\lambda(u,v) = M'(u,v)$ for every pair (u,v) of vertices of G.

Corollary 19.1.12.1. Graph G is n-edge connected iff every two vertices of G are connected by at least n edge disjoint paths.

Theorem 19.1.13. For every pair of non-adjacent vertices u, v in graph G, $\kappa(u, v) = M(u, v)$.

Corollary 19.1.13.1. Graph G is n-connected iff every pair of vertices of G are connected by at least n internally disjoint paths.

Algorithm 19.1.14 (connectivity $\kappa(G)$). .

- 1. If degree of every vertex is p-1, then output $\kappa = p-1$ and stop. Otherwise, continue.
- 2. If G is disconnected, output $\kappa = 0$ and stop. Otherwise, continue.
- 3. $\kappa \leftarrow p$
- 4. $i \leftarrow 0$
- 5. If $i \leq \kappa$, then $i \leftarrow i+1$ and continue. Otherwise, output κ and stop.

- 6. $j \leftarrow i + 1$
- 7. (1) If j = p + 1, then return to step 5. Otherwise continue.
 - (2) If v_iv_j ∉ E(G), construct network N with digraph D as follows: for each vertex v ∈ V(G), there are two vertices v', v" ∈ V(D) and an arc (v', v") ∈ E(D). And for each edge uv ∈ E(G), there are two arcs (u", v), (v", u) ∈ E(D). The capacity function is given by, c(v', v") = 1 for every v ∈ V(G) and c(a) = ∞ for every other arc in D. Set source s = v_i" and sink t = v_j and find maximum flow in N using max-flow min-cut algorithm. Otherwise proceed to step 7d
 - (3) If $f(N) < \kappa$, then $\kappa \leftarrow f(N)$. Otherwise, continue.
 - (4) $j \leftarrow j + 1$ and return to step 7a

19.2 Matchings and Factorizations

19.2.1 An Introduction to Matching

Marriage Problem Given a collection of men and women, where each womon knows some of the men. Can every women marry a man she knows?

Assignment Problem Given several job openings and applicants for one or more of these positions. Find an assignment so that maximum positions are filled?

Optimal Assignment Problem Given several job openings and applicants for one or more of these positions. The benefits of employing these applicants on those positions are also given. Find an assignment of maximum benefit to the company?

matching in G is a 1-regular³ subgraph of G.

maximum matching in G is a matching of G with maximum cardinality.

perfect matching in G is a matching of cardinality p/2. ie, p/2 edges.

 $\mathbf{maximum}$ weight $\mathbf{matching}$ in a weighted graph G is a matching with maximum weight.

Definitions 19.2.1. Let M be a matching in a graph G,

matched edge is an edge in subgraph M of G.

unmatched edge is an edge of G that doesn't belong to M.

matched vertex with respect to M is a vertex incident with an edge of M.

single vertex is a vertex that is not incident with any edge of M.

alternating path in G is a path with edges alternately matched and unmatched.

augmenting path in G is a non-trivial alternating path with single vertices as end vertices.

³A graph G is k-regular, if every vertex of G has degree k.

Theorem 19.2.2. Let M_1, M_2 be two matchings in G such that there is a spanning subgraph H of G with edges that are either in M_1 or M_2 , but not both. Then the components of H are either 1. isolated vertex 2. even cycle with edge alternately from M_1 and M_2 3. a non-trivial path with edges alternately from M_1 and M_2 such that each end vertex is single with respect to either M_1 or M_2 , but not both.

Synopsis. $\Delta(H) \leq 2$ by Pigeonhole principle. Any component of H is either a path or a cycle. A cycle with edge alternately from M_1 and M_2 is even. If an end vertex of a non-trivial path is matched with respect to $M_1(WLOG)$, then it is there in $M_1 - M_2$ ie, it is not there in M_2 . If there is another edge in M_2 incident with it, then it has to be in H and it will cease to be an end vertex of path component. Therefore, it is unmatched with respect to M_2 .

Theorem 19.2.3. A matching M in a graph G is maximum iff there is no augmenting path with respect to M in G.

Synopsis. If M is maximum matching and P an M-augmenting path. Since both end-vertices are single, length of P is odd. Let M', M'' be edges of P which are in M and not in M respectively. Then M-M'+M'' is a matching of cardinality one greater than that of M which is a contradiction since M is maximum. Conversely, suppose M be a matching such that there no M-augmenting paths in G. Let M' be a maximum matching in G. Then a nontrivial path component of the graph induced by $M\Delta M'$ is of even length otherwise both end-vertices are matched with respect to one of the matching M or M' which is a contradiction. Again every cycle components are even. Therefore |M| = |M'|, since $M\Delta M'$ doesn't have a nontrivial component of another kind.

Definitions 19.2.4. Let U_1, U_2 be two nonempty, disjoint, subsets of the vertex set of a graph G. Then U_1 is matched to U_2 if there exists a matching M in G such that every edge in M incident with a vertex in U_1 and a vertex in U_2 . And every vertex of U_1 (or U_2) is incident with some edge in M. Suppose M^* be a matching such that $M \subset M^*$, then U_1 is matched under M^* to U_2 .

Definitions 19.2.5. Let U be a nonempty set of vertices of a graph G. U is **nondeficient**, 4 if $|N(S)| \ge |S|$ for every nonempty subset S of U.

Theorem 19.2.6. Let G be a bipartite graph with partite sets V_1, V_2 . The set V_1 can be matched to a subset of V_2 iff V_1 is nondeficient.

Corollary 19.2.6.1. Every r-regular bipartite multigraph has a perfect matching.

Theorem 19.2.7. A collection S_1, S_2, \dots, S_n of finite non-empty sets has a system of distinct representatives iff for each k, $0 \le k \le n$, the union of any k of these sets contains at least k elements.

Remark (Hall's Marriage Theorem). Suppose there are n women. Then every women can marry a man she knows iff each subset of k women $(1 \le k \le n)$ colletively knows atleast k men.

Remark. Let W be a set of n women. Then there are $2^n - 1$ nonempty subset for W. Thus, Hall's Marriage Theorem suggests that we ensure $|N(S)| \ge |S|$ for every nonempty subset S of W. This method has complexity $O(2^n)$.

⁴N(S) is the neighbourhood set of all vertices adjacent to some vertex in S

19.2.2 Maximum Matching in Bipartite Graphs

Definitions 19.2.8. Let M be a matching in a graph G and P is an augmenting path with respect to M. Let M' be set of edge in P and M. And M'' be the set of edges in P and not in M. Then $M_1 = (M - M') \cup M''$ is the **matching** obtained by augmenting M along path P.

Remark. $|M_1| = |M| + 1$

Theorem 19.2.9. Let M be a a matching of a graph G that is not maximum, and let v be a single vertex with respect to M. Let M_1 denote the mathing obtained by augmenting M along some augmenting path. If G contains an augmenting path with respect to M_1 that has v and an end-vertex, then G contains an augmenting path with respect to M that has v as an end-vertex

Corollary 19.2.9.1. Let M be a matching of a graph G. Suppose that $M = M_1, M_2, \cdots, M_k$ is a finite sequence of matchings of G such that M_i $(2 \le i \le k)$ is obtained by augmenting M_{i-1} along some augmenting path. Suppose v is a single vertex with respect to M for which there exists no augmenting path starting at v. Then G does not contain an augmenting path with respect to M_i $(2 \le i \le k)$ that has v as an end-vertex.

Definitions 19.2.10. An alternating tree with respect to a matching M is a tree such that every path from it's root are alternating path with respect to M.

Algorithm 19.2.11 (Maximum Matching Algorithm for Bipartite Graphs). .

- 1. $i \leftarrow 1$ and $M \leftarrow M_1$
- 2. If i < p, then continue; otherwise stop.
- 3. If v_i is matched, then $i \leftarrow i+1$ and return to Step 2; otherwise, $v \leftarrow v_i$ and Q is initialized to contain v only.
- 4. (1) For $j = 1, 2, \dots, p$ and $j \neq i$, let $TREE(v_j) \leftarrow F$. Also, $TREE(v_j) \leftarrow T$.
 - (2) If $Q = \phi$, then $i \leftarrow i + 1$ and return to Step 2; otherwise, delete a vertex x from Q and continue.
 - (3) (1) Suppose that $N(x) = \{y_1, y_2, \dots, y_k\}$. Let $j \leftarrow 1$.
 - (2) If $j \leq k$, then $y \leftarrow y_j$; otherwise return to Step 4.2
 - (3) If TREE(y) = T, then $j \leftarrow j+1$ and return to Step 4.3.2. Otherwise, continue.
 - (4) If y is incident with a matching edge yz, then $TREE(y) \leftarrow T$, $TREE(z) \leftarrow T$, $PARENT(y) \leftarrow x$, $PARENT(z) \leftarrow y$ and add z to Q, $j \leftarrow j+1$ and return to Step 4.3.2. Otherwise, y is a single vertex and continue.
 - (5) Use PARENT to determine the alternating v-x path P' in the alternating tree. Let P be the augmenting path obtained from P' by adding the path x, y. Proceed to Step 5
- 5. Augment M along P to obtain a new matching M'. Let $M \leftarrow M'$, $i \leftarrow i+1$, and return to Step 2.

Definitions 19.2.12. Let G be a weighted complete bipartite graph with partite sets V_1 and V_2 . A **feasible vertex labeling** is a real function $l:V(G) \to \mathbb{R}$ on vertex set of G such that $l(v) + l(u) \ge w(vu)$ where $v \in V_1$ and $u \in V_2$.

Definitions 19.2.13. Consider the function $l: V(G) \to \mathbb{R}$ such that $\forall v \in V_1, \ l(v) = \max\{w(vu) : u \in V_2\}$ and $\forall u \in V_2, \ l(u) = 0$. Then l is a feasible vertex labeling on V(G). And,

 E_l is the set of all edge of the weighted complete bipartite graph G such that l(v) + l(u) = w(vu).

 H_l is the spanning subgraph of G induced by the edge set E_l .

Theorem 19.2.14. Let l be a feasible vertex labeling of a weighted complete bipartite graph G. If H_l contains a perfect matching M', then M' is a maximum weight matching of G.

Algorithm 19.2.15 (Kuhn-Munkres). .

- 1. (1) For each $v \in V_1$, let $l(v) \leftarrow \max\{w(uv) : v \in V_2\}$.
 - (2) For each $u \in V_2$, let $l(u) \leftarrow 0$.
 - (3) Let H_l be the spanning subgraph of G with edge set E_l .
 - (4) Let G_l be the underlying graph of H_l .
- 2. Apply Algorithm ?? to determine a maximum matching M in G_l .
- 3. (1) If every vertex v of V_1 is matching with respect to M, output M and stop. Otherwise, continue.
 - (2) Let x be the first single vertex of V_1 .
 - (3) Construct an alternating tree with respect to M that is rooted at x. If an augmenting path P is discovered, then augmenting M along P and return to Step 3.1. Otherwise, let T be the alternating tree with respect to M and rooted at x that cannot be expanded further in G_l.
- 4. Compute $m_l \leftarrow \min\{l(v) + l(u) w(vu) : v \in V_1 \cap V(T), u \in V_2 V(T)\}$. Let

$$l'(v) = \begin{cases} l(v) - m_l \text{ for } v \in V_1 \cap V(T) \\ l(v) + m_l \text{ for } v \in V_2 \cap V(T) \\ l(v) \text{ otherwise} \end{cases}$$

5. Let $l \leftarrow l'$, construct G_l and return to Step 3.3.

19.2.4 Factorizations

Definitions 19.2.16. A factor of a graph G is a spanning⁵ subgraph of G.

Definitions 19.2.17. Let G_1, G_2, \dots, G_n be edge-disjoint factors of G such that $E(G) = \bigcup_{i=1}^n E(G_i)$. Then G is **factorable** and $G = G_1 \oplus G_2 \oplus \dots \oplus G_n$.

Definitions 19.2.18. An r-regular factor of G is an r-factor of G.

 $^{^5\}mathrm{Spanning}$ subgraph of a graph G has every vertex of G

Definitions 19.2.19. If G has a factorisation to r-factors, then G is **r-factorable**.

Remark. $K_{3,3}$ is 1-factorable. K_5 is 2-factorable.

Definitions 19.2.20. An odd component of G is a component of G with odd number of vertices. And an even component of G is a component of G of with even number of vertices.

Theorem 19.2.21 (Tutte). A nontrivial graph G has a 1-factor iff for every proper subset S of V(G), the number of odd components of G-S does not exceed |S|.

Remark. There exist cubic graphs that doesn't have a 1-factor.

Theorem 19.2.22 (Petersen). Every bridgeless cubic graph contains a 1-factor.

Remark. Every brideless cubic graphs has a 1-factor. Let G be a bridgeless cubic graph. Consider every pair of factors G_1, G_2 such that $G = G_1 \oplus G_2$ where G_1 is a 1-factor and G_2 is a 2-factor. G is not 1-factorable only if every such G_2 doesn't have a 1-factor.

Theorem 19.2.23. Petersen graph is not 1-factorable.

Theorem 19.2.24. Every r-regular bipartite multigraph $(r \ge 1)$ is 1-factorable.

Remark (Application of 1-factorisation). For even number p, a 1-factorisation of K_p corresponds to the schedule of a round of the round robin tournament among p teams. If p is odd, consider K_{p+1} where v_{p+1} is an imaginary team called by eteam. A game with by eteam is a by e.

Definitions 19.2.25. A hamiltonian cycle is a spanning cycle. And, Hamiltonian graph is a graph containing a hamiltonian cycle.

Theorem 19.2.26. Complete graph K_{2n+1} can be factored into n hamiltonian cycles.

Remark. For n = 3, K_7 can be factored into three hamiltonian cycles.

Theorem 19.2.27. Let $0 \le r < p$. Then there exists an r-regular graph of order p iff pr is even.

Definitions 19.2.28. Let $\{E_1, E_2, \dots, E_n\}$ be partition of E(G). And let H_i be subgraph of G induced by the edge set E_i . A **decomposition** of a graph G is a colletion of these subgraphs H_1, H_2, \dots, H_n . And $G = H_1 \oplus H_2 \oplus \dots \oplus H_n$.

Definitions 19.2.29. Let $G = H_1 \oplus H_2 \oplus \cdots \oplus H_n$ be a decomposition of G such that $H \cong H_i$. Then G is H-decomposible.

Remark. $K_{3,3}$ is $3K_2$ -decomposible. K_5 is C_5 -decomposible. K_{2n} is nK_2 -decomposible. K_{2n+1} is C_{2n+1} -decomposible.

Every graph is K_2 -decomposible. Every complete bipartite graph $K_{m,n}$ is $K_{1,m}$ -decomposible and $K_{1,n}$ -decomposible.

19.2.5 Block Designs

Definitions 19.2.30. A block design on a set V is a collection of k-element subsets of V such that each element of V appears exactly in r subsets.

variety The elements of V are called varieties.

block k-element subsets of V are called blocks.

balanced design If each variety appears in exactly r blocks and each pair of varieties appears in exactly λ blocks.

incomplete design If blocks are proper subsets of V. ie, k < v.

Definitions 19.2.31. A balanced incomplete block design of v varities in b blocks of cardinality k such that each variety appears in exactly r blocks and each pair of varities appears in exactly λ blocks is a (b, v, r, k, λ) -design.

Theorem 19.2.32. bk = vr

Theorem 19.2.33. $\lambda(v-1) = r(k-1)$

Corollary 19.2.33.1. $\lambda < r$

Theorem 19.2.34 (Fisher's Inequality). $b \ge v$

symmetric design If b = v

Theorem 19.2.35. In a symmetric (b, v, r, k, λ) -design with even $v, r - \lambda$ is a perfect square.

steiner triple system (b, v, r, k, λ) -design with $k = 3, \lambda = 1$.

Remark. (b, v, r, k, λ) -designs are incomplete. However, a complete block design (ie, v = 3) is also included as a Steiner triple system.

Theorem 19.2.36. Steiner triple system with v varieties exists iff v = 6n + 1 or v = 6n + 3 or v = 3.

Definitions 19.2.37 (Kirkman's Schoolgirls Problem). A class of 15 girls. Parade 15 girls in five rows (3 girls in a row). Is it possible to plan 7 days parade so that two girls are together in a row exactly once?

kirkman triple system Steiner triple system with v = 6n + 3.

Remark. It is proved that kirkman triple system exists with v = 6n + 3 for every $n \ge 0$.

Theorem 19.2.38. The code consisting of the rows of the incidence matrix of $a(b, v, r, k, \lambda)$ -design (b = v, r = k) is t-error correcting, where $t = k - \lambda - 1$.

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