## Review

- □ The waitpid() System Call
- □ The system() System Call
- Concept of Signals
  - Linux Signals
  - Signal() System Call
  - The kill() and raise() System Call
  - The alarm() System Call
  - The pause() System Call

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#### Preview

- □ The pause() System call
- □ The signal() system call
- Signal set
- □ The sigprocmask() system call
- □ The sigaction() system call
- □ Interprocess Communication
- lacktriangledown The sigsuspend() system call
- □ The abort() system call

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# The pause() System Call

□ The pause() system call suspends the calling process until a signal is caught

```
#include <unistd.h>
int pause (void);

Return -1 with error
```

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# The pause() System Call

```
/* alarml.c: demonstrate a signal system call*/
#include <ayof.vpes.h>
#include <ayof.vpes.h>
#include <atof.oh>
#include <atof.oh>
#include <atof.oh>
#include <atof.oh>
woid ding (int sig)
{
   printf("alarm fired signal number = %d\n", sig);
}
int main()
{
   printf("set alram 5 second for a process\n");
   alarm(5);
   signal (SIGALEM, ding);
   pause();
   pause();
}
```

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# The signal() System Call

```
#include <signal.h>
void (*signal(int signo, void (*func(int)))(int);
signo: name of signal SIG...
func: pointer to siglal handler function
```

- □ Function **signa**l <u>accept two arguments</u> and <u>return a pointer</u> to a function that returns nothing.
- Second argument is pointer to a function <u>that take a single</u> <u>integer argument</u> and return nothing.

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```
/* sig tablet of demonstrate a signal system call*/
sinclode csignal.b>
signal (sinclode, ding);
signal (sinclo
```

```
The signal() System Call

//*signal.c send signal to a process */
#include <agrail.b?
#include <agrail.c.
```

# 

# Signal Set Distinct Sigemptyset() or sigfillset() must be called for every object of type sigset\_t before any other use of the object. The sigemptyset() function initializes a signal set to be empty. The sigfillset() function initializes a signal set to contain all signals. The sigaddset() function adds the specified signal signo to the signal set. The sigdelset() function deletes the specified signal signo from the sig- nal set. The sigsmember() function returns whether a specified signal signo is contained in the signal set.

# The sigprocmask() System Call Signal mask of a process is the set of signals currently blocked from delivery to that process. The sigprocmask() system call examines, or changes, or both examines and changes the signal mask of the calling process. Sinclude <signal.h> int sigprocmask( int how, const sigset\_t \*set, sigset\_t \*oset ); return 0 for ok, -1 for error

# The sigaction() System Call

■ We can modify or examine the action associated with a particular signal by using system call sigaction().

```
#include <signal.h>
int sigaction ( int signo, const struct sigaction *act,
struct sigaction *oact);
 truct segaction(
void (*aa handler)(int)/*pointer to signal handler */
signet_ taa_mask; /*additional set of signals to block */
int sa_flag; /*signal options */
void (*aa_sigaction)(int signifo_t *, void *);/* alternate handler */
```

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# The sigaction() System Call

```
#include <signal.h>
int sigaction (int signo, const struct sigaction *act,
struct sigaction *oact);
```

- □ **signo** (Input) A signal from the list
- \*act (Input) A pointer to the sigaction structure that describes the action to be taken for the signal. Can be NULL. If act is a NULL pointer, signal handling is unchanged. **sigaction()** can be used to inquire about the current handling of signal sig. If act is not NULL, the action specified in the sigaction structure becomes the new action associated with sig.
- \*oact (Output) A pointer to a storage location where sigaction() can store a sigaction structure. This structure contains the action currently associated with signo. Can be NULL. If oact is a NULL pointer, sigaction() does not store this information.

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```
// SIGINT is blocked
if (sigprocmaskSIG_BLOCK, &mask, &orig_mask) < 0) {
    perror (*sigprocmask*);
    return 1;</pre>
 int count =0;
while (1)
           if (count == 10)
```

```
nt main( int argc, char *argv[] )
    sigset_t new_set, old_set; // signal set
time t start, finish;
             cmask( SIG BLOCK, &new set, &old set); /*two signals are blocked */
   time( &start );
printf( "SIGTSTP and SIGINT are blocked at %s\n", ctime(&start));
    sime( &finish );

printf( "SIGTSTP and SIGINT are unblocked at %s\n", ctime(&finish));

sigprocmask( SIG_SETMASK, &old_set, NULL ); // unblock two signals
```

```
signmetyment (smect.em_mask ); /* Clear signmention structure mask */
smect.em_mask = (smect.em_mask ); /* Clear signmention structure mask */
smect.em_mask = (smect.em_mask ); * set signmin serion hander as function eather */
signeticing stilled, smect, NULL); /* change the action change for signmin SIGAIUM */
signetizet( smec.em_t); /* clear signmin sets: block all signmin*/
signetizet( smec.em_t); /* clear signmin sets: block all signmin*/
signetizet( smec.em_t); /* clear signmin sets: block all signmin signmi
          alarm( 1 ); /* SIGALRN will be sent in 1 second */
printf( "SIGALRN is fired but it is blocked.\n");
   do {
   time( sfinish );
   printf( "Current time is %s\n", ctime(sfinish) );
   diff = difftime( finish, start );
} while (diff < 10);</pre>
signrocmask ( \ SiG_SETWASK, \ told\_set, \ NULL \ ); \ /* \ signal set are reset with old set */printf( "SIGALRM signals unblocked at &s\n", ctime(sfinish) ); \\
                                                                                                                                                                                                                                                                     COSC350 System Software, Fall 2020
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```

# Interprocess Communication

- Three issues in interprocess communication
  - How one process can pass information to another
  - How to make sure two or more processes do not get into the **critical section** (mutual exclusion)
  - Proper sequencing when dependencies are present (ex. Producer-Consumer problem, Dinning Philosopher problem)

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# Interprocess Communication

(the Producer-Consumer Problem)

#### Description

- Two processes (or threads) share a common, fixed-sized buffer.
- □ Producer puts information into the buffer, and consumer takes it out.

#### Troubles arises

- □ When the producer wants to put a new item in the buffer, but it is already full.
- When the consumer tries to take a item from the buffer, but buffer is already empty.

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### Interprocess Communication

(the Producer-Consumer Problem)

- When the producer wants to put a new item in the buffer, but it is already full.
  - Solution producer is go to sleep, awakened by consumer when consumer has removed on or more items.
- When the consumer tries to take a item from the buffer, but buffer is already empty.
  - Solution consumer is go to sleep, awakened by the producer when producer puts one or more information into the buffer.

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# Interprocess Communication (Dining Philosophers Problem)



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# Interprocess Communication

(Dining Philosophers Problem)

- Five silent philosophers sit at a round table with bowls of spaghetti. Chopsticks are placed between each pair of adjacent philosophers.
- Each philosopher must alternately think and eat. However, a philosopher can only eat spaghetti when they have both left and right chopsticks.
- Each chopstick can be held by only one philosopher and so a philosopher can use the chopstick only if it is not being used by another philosopher.
- After an individual philosopher finishes eating, they need to put down both chopsticks so that the chopstick s become available to others. A philosopher can take the chopstick on their right or the one on their left as they become available, but cannot start eating before getting both chopsticks.

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# Interprocess Communication

- critical section (critical region) The part of program where the shared memory is accessed.
- If we could arrange matters such that no two processes were ever in their critical sections at the same time, we can avoid races condition.

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## Interprocess Communication

- We can change the signal mask for a process to block and unblock selected signal by sequence of system calls.
- ☐ It might be possible to use this technique to protect critical region (critical section).

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# Interprocess Communication

```
sigset_t newmask, oldmask;
sigemptyset(&newmask);
sigaddset(&newmask, SIGINT);
/*block SIGINT and save current signal mask */
if (sigprocmask(SIG_BLOCK, &newmask, &oldmask)<0)
    error_sys(" SIG_BLOCK ERROR ");
/**** Critical Region of code *******/
/*reset signal mask, which unblocks SIGINT */
if (sigprocmask(SIG_SETMASK, &oldmask, NULL)<0)
    error_sys(" SIG_BLOCK ERROR ");
/* hole */
pause();</pre>
```

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# Interprocess Communication

- If a signal is sent to the process while it is blocked, the signal delivery will be differed until the signal is unblocked.
- If a signal does occur <u>between the</u> <u>unblocking and the pause</u>, the signal can be lost.
- □ The result is pause forever!!!
- □ The sigsuspend() system guarantee both reset and put a process to sleep.

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# The sigsuspend() System Call

#include <signal.h>
int sigsuspend(const sigset\_t \*sigmask);

- The sigsuspend() function replaces the current signal mask of a process with the signal set given by \*sigmask and then suspends processing of the calling process.
- □ The process does not resume running until a signal is delivered

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# The sigsuspend() System Call

```
sigset_t newmask, oldmask;
sigemptyset(&newmask);
sigaddset(&newmask, SIGINT);
/*block SIGINT and save current signal mask */
if (sigprocmask(SIG_BLOCK, &newmask, &oldmask)<0)
    error_sys("SIG_BLOCK ERROR");

    /**** Critical Region of code *******/
/*reset signal mask, which unblocks SIGINT */
if (sigsuspend(&oldmask)<0)
    error_sys("SIG_SUSPEND_ERROR");</pre>
```

COSC350 System Software, Fall 2020 Dr. Sang-Eon Park #include (signal.b)
#include (signal.b)
#include (stdio.b)

void timestamp( char "str ) (

time t t;

time (t t);

printf( "%s the time is %s\n", str, ctime(£t) );

j int main( int argo, char "argv[ ) (

struct signation signat; /\* for signation call \*/

signet; block, sst;

signifilate( f&llock set ) / /\*set all signals are included \*/

signet; block set;

signifilate( f&llock set; ) /\*set all signals are included \*/

signation (stdio.b) \*/ \*signation (stdio.b) \*/ \*set signation (stdio.b) \*/ \*signation (stdio.b) \*/ \*signation (stdio.b) \*/ \*signation (stdio.b) \*/ \*set signation (stdio.b) \*/ \*signation (st

# The abort() System Call

- □ The abort() system call cause abnormal program termination.
- ☐ The abort() system call send SIGABRT signal to caller process

#include <stdlib.h>
void abort(void);

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