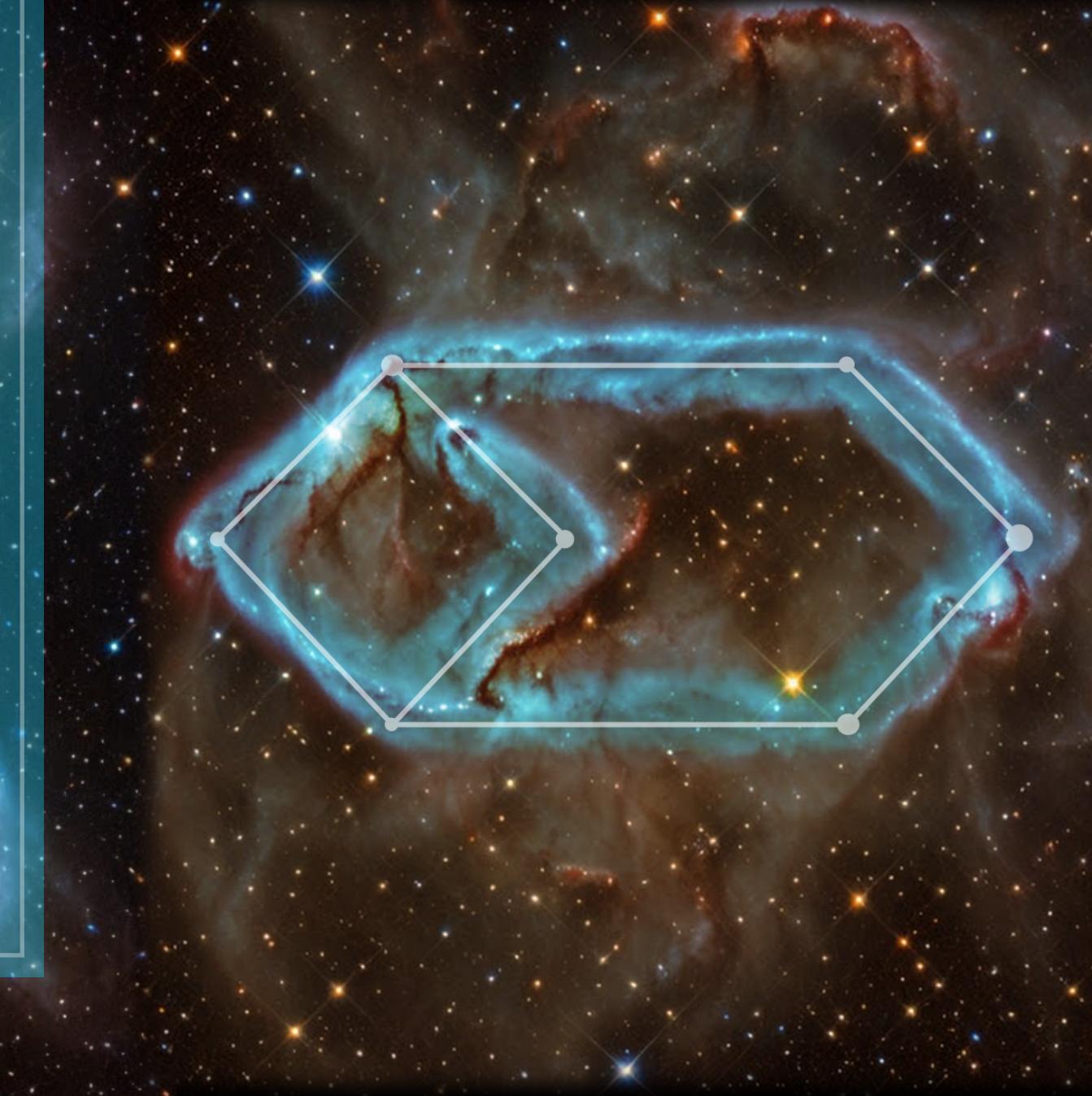


Paranatural Category Theory

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CT Octoberfest
28 October 2023



Question: What is the appropriate notion of *transformation* between difunctors?

Church numerals as endo-transforms of Hom

Consider the difunctor $\text{Hom}: \mathbb{C}^{\text{op}} \times \mathbb{C} \rightarrow \text{Set}$. Then, for each $n \in \mathbb{N}$, we can define the family of maps

$$\lambda f.f^n : \text{Hom}(J, J) \rightarrow \text{Hom}(J, J)$$

indexed over objects J of \mathbb{C} .

I'll use the term **difunctor** to refer to functors of the form
 $\mathbb{C}^{\text{op}} \times \mathbb{C} \rightarrow \text{Set}$

Notion of “transformation” between difunctors Γ, Δ :

- **natural transformation**: for every $I, J \in \mathbb{C}$, a function $\alpha_{I,J}: \Gamma(I, J) \rightarrow \Delta(I, J)$, satisfying naturality
 - ▶ Doesn't capture ‘diagonal’ transformations, e.g. the Church numerals
- **dinatural transformation**: for every $J \in \mathbb{C}$, a function $\alpha_J: \Gamma(J, J) \rightarrow \Delta(J, J)$, satisfying a “dinaturality condition”.
 - ▶ Too weak—dinaturals don't compose (in general)
- **paranatural transformation**: for every $J \in \mathbb{C}$, a function $\alpha_J: \Gamma(J, J) \rightarrow \Delta(J, J)$, satisfying a “paranaturality condition”.

Dinatural Transformations

For every $i_2 \in \text{Hom}(I_0, I_1)$

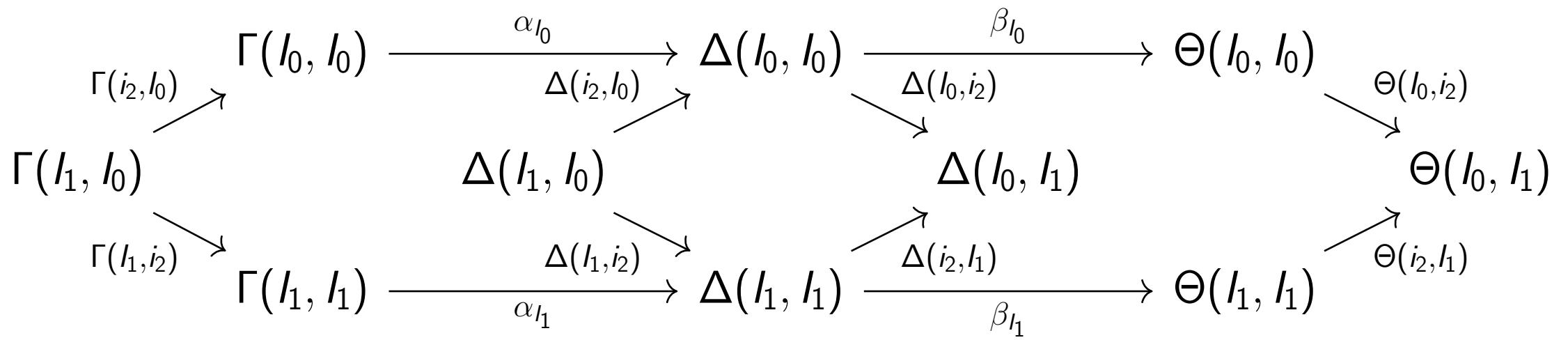
$$\begin{array}{ccccc} & \Gamma(I_0, I_0) & \xrightarrow{\alpha_{I_0}} & \Delta(I_0, I_0) & \\ \Gamma(i_2, I_0) \nearrow & & & & \searrow \Delta(I_0, i_2) \\ \Gamma(I_1, I_0) & & & & \Delta(I_0, I_1) \\ & \searrow \Gamma(I_1, i_2) & & & \nearrow \Delta(i_2, I_1) \\ & & \Gamma(I_1, I_1) & \xrightarrow{\alpha_{I_1}} & \Delta(I_1, I_1) \end{array}$$

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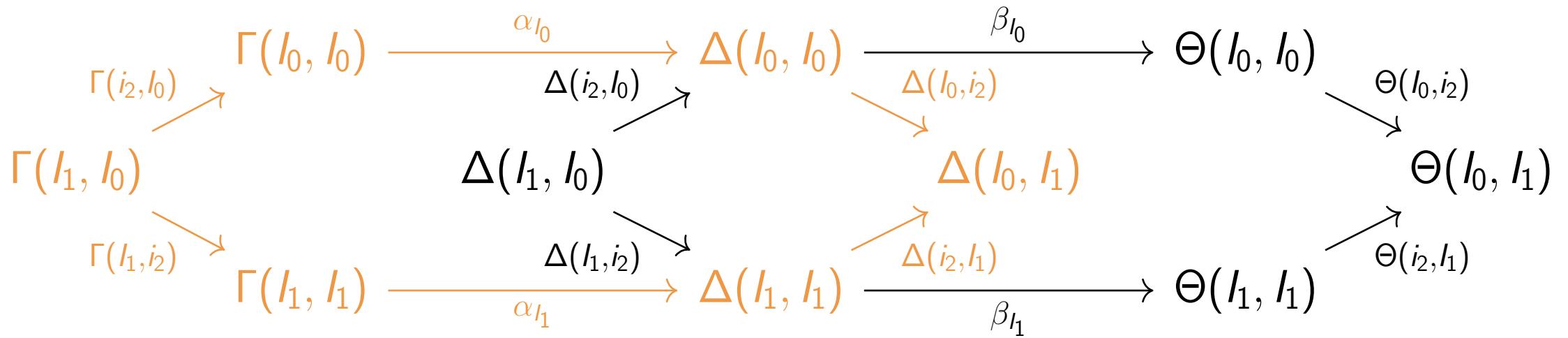
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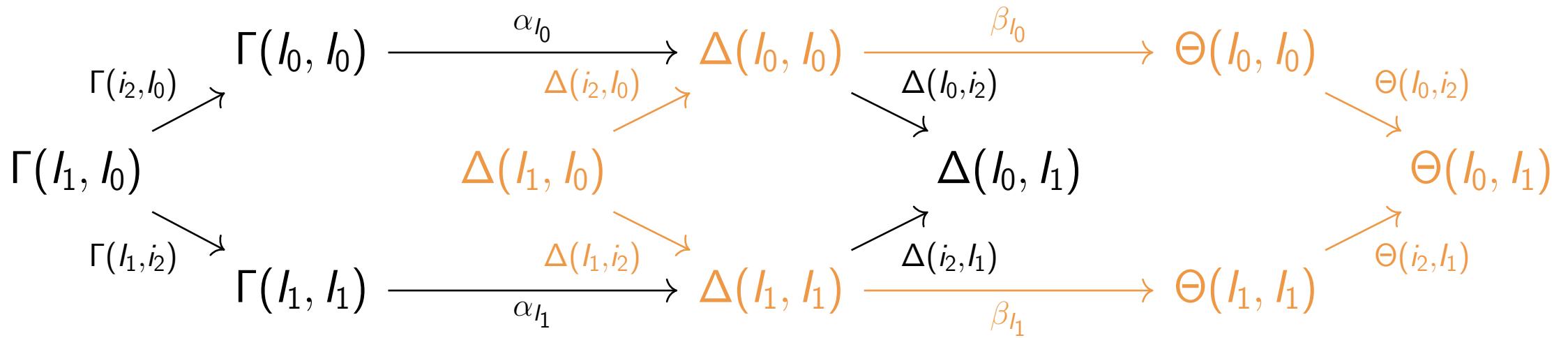
Dinaturals don't compose



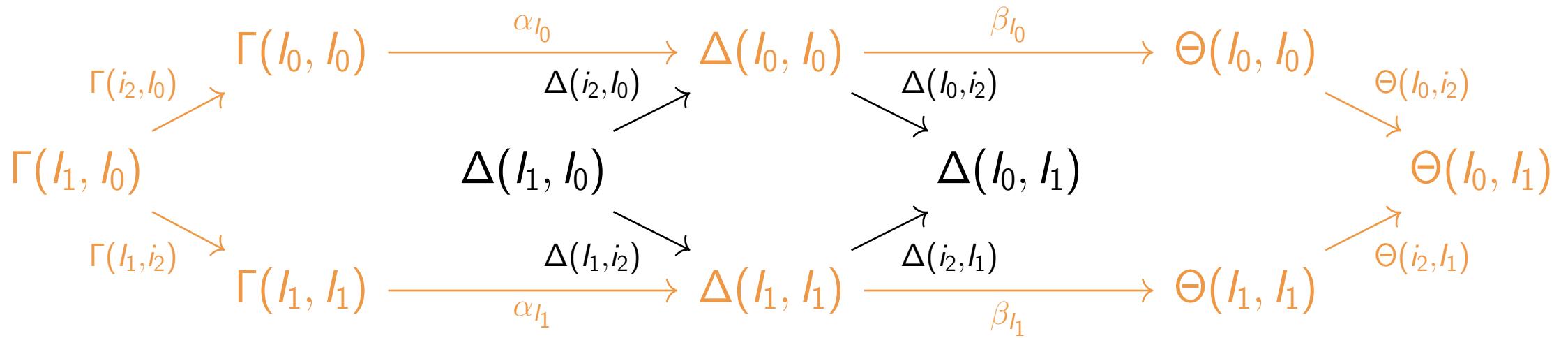
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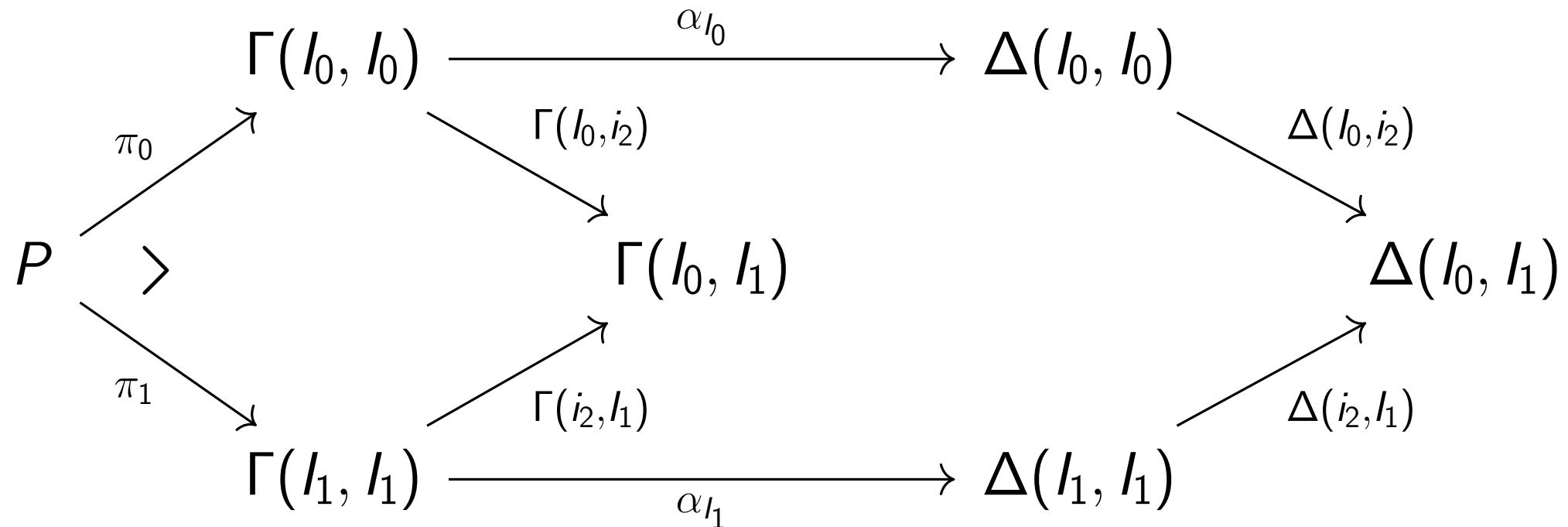


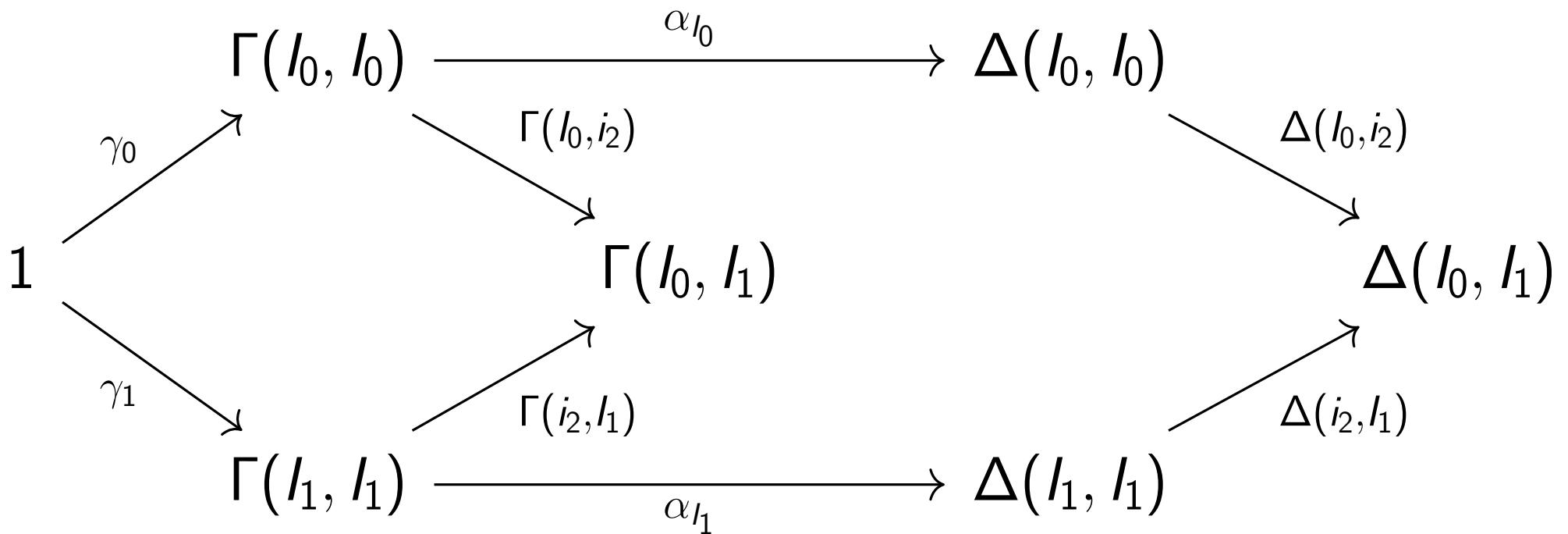
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Defn. Given difunctors Γ, Δ , we say a family of maps $\alpha_J: \Gamma(J, J) \rightarrow \Delta(J, J)$ is a **paranatural transformation** from Γ to Δ and write $\alpha: \Gamma \xrightarrow{\diamond} \Delta$ if, for every $i_2 \in \text{Hom}(I_0, I_1)$, the following hexagon commutes.





“if the diamond commutes, so does the hexagon”

Some Results

Prop. The Church numerals are paranatural transformations

$\text{Hom} \xrightarrow{\diamond} \text{Hom}$

Prop. If $\alpha: \Gamma \xrightarrow{\diamond} \Delta$ and $\beta: \Delta \xrightarrow{\diamond} \Theta$, then the pointwise-defined

composite $(\beta \circ \alpha)_I := \beta_I \circ \alpha_I$ is a paranatural transformation $\Gamma \xrightarrow{\diamond} \Theta$

Defn. Write $\hat{\mathcal{C}}$ for the category whose objects are difunctors and whose morphisms are paranatural transformations.

Prop. $\hat{\mathcal{C}}$ has all finite products

Prop. $\hat{\mathcal{C}}$ is cartesian closed

Conj. $\hat{\mathcal{C}}$ is an elementary topos

Conj. Hom^{Hom} is a natural numbers object in $\hat{\mathcal{C}}$

Defn. The **diYoneda embedding** $\mathbf{yy} : \mathbb{C}^{\text{op}} \times \mathbb{C} \rightarrow \hat{\mathbb{C}}$ is the functor whose object part is given by

$$\mathbf{yy}(I_0, I_1)(J_0, J_1) := \text{Hom}(I_0, J_1) \times \text{Hom}(J_0, I_1)$$

and whose four morphism parts are given by appropriate pre- and post-compositions.

Lemma For any difunctor $\Delta : \mathbb{C}^{\text{op}} \times \mathbb{C} \rightarrow \text{Set}$, there is a bijection

$$\Delta(I, J) \cong \mathbf{yy}(J, I) \xrightarrow{\diamond} \Delta$$

paranatural in I, J .

- Note that I and J are flipped on the right
- To prove this, we construct an $\alpha_d : \mathbf{yy}(I, I) \xrightarrow{\diamond} \Delta$ for each $d : \Delta(I, I)$ and vice-versa.

Claim The category of difunctors $\hat{\mathcal{C}}$ has exponential objects.

Proof. By “diYoneda reasoning”: for difunctors Γ, Δ , suppose their exponential Δ^Γ existed. Then

$$\begin{aligned}\Delta^\Gamma(I, J) &\cong \mathbf{yy}(J, I) \xrightarrow{\diamond} \Delta^\Gamma \\ &\cong \mathbf{yy}(J, I) \times \Gamma \xrightarrow{\diamond} \Delta\end{aligned}$$

diYoneda Lemma
(desired property)

so now define $\Delta^\Gamma(I, J)$ to be $\mathbf{yy}(J, I) \times \Gamma \xrightarrow{\diamond} \Delta$, and verify this satisfies all the necessary properties.

Have we actually done
anything new here?

Is paranaturality an instance
of naturality?

Question Given difunctors $\Gamma, \Delta: \mathbb{C}^{\text{op}} \times \mathbb{C} \rightarrow \text{Set}$, can we define functors

$$\overline{\Gamma}, \overline{\Delta}: \mathbb{C}' \rightarrow \mathbb{D}$$

(for some appropriately-picked \mathbb{C}', \mathbb{D}) such that *paranatural transformations* $\Gamma \xrightarrow{\diamond} \Delta$ are the same thing as *natural transformations* $\overline{\Gamma} \rightarrow \overline{\Delta}$?

- **Positive**: No need to develop paranatural category theory separately, diYoneda is just an instance of Yoneda, $\hat{\mathbb{C}}$ “is” a presheaf category
- **Negative**: Paranatural category theory is indeed a novel branch of category theory, diYoneda is a distinct result from Yoneda, difunctor categories may be differently-behaved than presheaf categories

Further topics in paranatural category theory

- Equivalent formulations of paranaturality
- Categories of diagonal elements (“ Γ -structures”)
- “Splice categories”
- Strong (Co)End calculus
 - ▶ Structural ends
 - ▶ Initial algebras and [Uus10]’s Yoneda-like lemma
 - ▶ Structural coends, terminal coalgebras, and bisimulations
- Dependent paranatural transformations (and maybe a dependent diYoneda Lemma?)

- Mathematical framework for categorical semantics of (co)inductive types (generalizing and dualizing [AFS18])
- Parametricity: Paranatural transformations encode impredicative ‘universal’ and ‘existential’ types (e.g. from System F)
 - ▶ Paranatural transformations correspond to parametrically polymorphic functions, with the paranaturality condition matching the ‘free theorems’ of [Wad89]
 - ▶ Structural coends encode ‘abstract data structures’
- Difunctor models of type theory
 - ▶ Uses diYoneda for “lifting Grothendieck universes”, à la [HS99]

- Collection of links: jacobneu.github.io/research/paranat
- arXiv preprint: arxiv.org/abs/2307.09289
- HoTTEST talk:
 - ▶ Video: youtube.com/watch?v=X4v5HnnF2-o
 - ▶ Slides: research/slides/HoTTEST-2022.pdf
- Midlands Graduate School talk: research/slides/MGS-2023.pdf
- CMU HoTT Seminar Talk: research/slides/CMU-2023.pdf
- Lean formalization (in progress) will be made public soon!

- [AFS18] Steve Awodey, Jonas Frey, and Sam Speight.
Impredicative encodings of (higher) inductive types.
In *Proceedings of the 33rd Annual ACM/IEEE Symposium on Logic in Computer Science*, pages 76–85, 2018.
- [HS99] Martin Hofmann and Thomas Streicher.
Lifting grothendieck universes.
Unpublished note, 199:3, 1999.
- [Uus10] Tarmo Uustalu.
A note on strong dinaturality, initial algebras and uniform parameterized fixpoint operators.
In *FICS*, pages 77–82, 2010.

[Wad89] Philip Wadler.

Theorems for free!

In *Proceedings of the fourth international conference on Functional programming languages and computer architecture*, pages 347–359, 1989.

Thank you!