

Spring 2023 - ICSI 526

Homework 4

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May 11, 2023

Contents

1	Question 1	1
1.1	Hash function A ($n = p \times q$)	1
1.2	Hash function B ($h2(m) = M1 \oplus M2 \dots \oplus Mn.$)	1
2	Question 2	2
2.1	Design and implement a PRNG using AES (OFB mode)	2
2.2	Calculate the Fraction of One Bits	2

1 Question 1

1.1 Hash function A ($n = p \times q$)

Properties:

- One Way Property: This does satisfy the One Way Property even if we know n . We because of the Modulus, we don't know how many times the original function has 'wrapped around' so to speak. We can't predict the output.
- Weak Collision Resistance: This is NOT Weak Collision Resistant, it is possible to generate the same hash value for two different inputs. We can use -2 and 2 as our input x values and $p = 107$ and $q = 127$ for our primes. Because of the x^2 , they will be the same value $N = 13589$ when being fed into the mod function and will output the same hash. Because for any two inputs $x1$ and $x2$, if $x1 \neq x2$, then there exists a possibility that $h1(x1) = h1(x2)$ which means that this is false.
- Strong Collision Resistance: This is NOT Strong Collision Resistant. The reason is that for any two inputs $x1$ and $x2$, the hash function $h1(x)$ produces the same output if $x1 = -x2 \pmod n$. That is, $h1(x1) = h1(x2)$ if $x1 \equiv -x2 \pmod n$. An attacker who knows the primes p and q can easily find two inputs $x1$ and $x2$ such that $x1 \equiv -x2 \pmod n$ and then generate a collision by computing $h1(x1)$ and $h1(x2)$.

1.2 Hash function B ($h_2(m) = M_1 \oplus M_2 \dots \oplus M_n$)

Properties:

- One Way Property: This does satisfy the One Way Property. We divide a message
- Weak Collision Resistance: This is Weak Collision Resistant.
- Strong Collision Resistance: This is Strong Collision Resistant.

2 Question 2

2.1 Design and implement a PRNG using AES (OFB mode)

Design Algo...

2.2 Calculate the Fraction of One Bits

Calculate FOOB...