Traversing the k-mer Landscape of NGS Read Datasets for Quality Score Sparsification

Jacopo Notarstefano jacopo.notarstefano [at] gmail.com

May 19, 2014

Next-generation sequencing (NGS)

NGS is the name given to several different sequencing technologies that were developed in the past decade.

These methods have in common a much higher throughput and a much lower cost compared to Sanger sequencing, which was used to sequence the Human Genome [LLL+12].

Sequencers	HiSeq 2000	SOLiDv4	Sanger 3730xl
Output data/run	600 Gb	120 Gb	1.9~84 Kb
Cost/million bases	\$0.07	\$0.13	\$2400

The problem with NGS

"In the past two decades, genomic sequencing capabilities have increased exponentially, outstripping advances in computing power and storage."

- Moore's law predicts that the number of transistors on integrated circuits doubles every 24 months [Moo65].
- Kryder's law predicts that hard drive density doubles approximately every 13 months [Wal05].
- Sequencing output has doubled every 9 months [Kah11].

How to deal with storage

	Lossless	Lossy
General purpose	GZIP BZIP2 7zip	
Specific for NGS		RQS

An example of the FASTQ format

```
@EAS54_6_R1_2_1_443_348
GTTGCTTCTGGCGTGGGTGGGGGGG
+EAS54_6_R1_2_1_443_348
;;;;;;;;;;;9;7;;.7;393333
```

- First and third rows contain an identifier for the read.
- Second row is the read itself.
- Fourth element represents the error probability of each base call using Phred quality scores.

Phred quality score

Definition (Phred quality score)

Let P_e be the estimated probability that a base call is incorrect. Then we define its *Phred quality score* Q as

$$Q = -10\log_{10}P_e.$$

This quantity is encoded in the Sanger FASTQ format as a single byte, where the character '!' represents the lowest quality while '~' is the highest ($|\Sigma| = 94$).

The RQS algorithm in brief

Yu, Yorukoglu and Berger [YYB14] propose the following:

- Generate a dictionary D of all *k*-mers that appear with high multiplicity in a representative collection of reads.
- ② Let γ be a read. Its k-mers close to D in Hamming distance have nearly all of their quality score discarded.
- Ompress the resulting high redundancy sequence with a general purpose lossless compression algorithm.

The DICT algorithm

Algorithm 1 Dict

```
Input: C, k, r
Output: D
 1: D ← {}
 2: A \leftarrow [0, \dots, 0] \in \mathbb{N}^{4^k}
 3: for x \in C_k do
 4: A[x]++
 5: for x \in [4^k] do
 6: if A[x] \ge r then
    D.append(x)
 8: return D
```

The MARKKMER algorithm

Algorithm 2 MarkKmer **Input:** x, D Output: M 1: if $\Delta(x, D) > 1$ then $M \leftarrow [false, \dots, false] \in \{true, false\}^k$ 3: **else** $M \leftarrow [true, ..., true] \in \{true, false\}^k$ for $y \in D \mid \Delta(x, y) = 1$ do 5: for $i \in [k]$ do 6: if $x_i \neq y_i$ then 7: $M_i \leftarrow false$ 8: 9: **return** M

The MARKREAD algorithm

Algorithm 3 MarkRead

```
Input: \gamma, D
Output: \mathcal{M}
  1: // Let x^a be the k-mer in \gamma starting at a.
 2: // Cover \gamma by k-mers \{x^{a_1}, \ldots, x^{a_n}\}.
  3: for i \in [n] do
 4: M^i \leftarrow MARKKMER(x^{a_i}, D)
  5: \overline{\mathsf{M}}^{\mathsf{i}} \leftarrow [\mathsf{false}, \dots, \mathsf{false}] \in \{\mathsf{true}, \mathsf{false}\}^{\mathsf{length}(\gamma)}
  6: for j \in [k] do
  7: \overline{M}_{j+a_i-1}^i \leftarrow M_i^i
  8. \mathcal{M} \leftarrow \overline{M}^1 \setminus \cdots \setminus \overline{M}^n
  9: return \mathcal{M}
```

The $\operatorname{SparsifyRQ}$ algorithm

Algorithm 4 SPARSIFYRQ Input: γ , Q, D, Q_{threshold} Output: Q' 1: Q' \leftarrow Q 2: $\mathcal{M} \leftarrow \text{MARKREAD}(\gamma, D)$ 3: for $i \in [\text{length}(\gamma)]$ do 4: if $(Q_i > Q_{\text{threshold}})$ or $(\mathcal{M}_i = \text{true})$ then 5: $Q'_i \leftarrow Q_{\text{threshold}}$ 6: return Q'

An example

Results

Results of compressing HG02215, chromosome 21 with r=50, k=32 and $Q_{threshold}=40$, along with the F1 score of downstream variant calling:

Method	Size	Bits/Q	F1
Uncompressed	273 MiB	8.0000	1
GZIP	143 MiB	4.1923	1
BZIP2	133 MiB	3.8791	1
7zip (PPMd)	124 MiB	3.6269	1
RQS + GZIP	14 MiB	0.3825	0.9914
RQS + BZIP2	8.7 MiB	0.2540	0.9914
RQS + 7zip (PPMd)	11 MiB	0.2953	0.9914

The implementation of MARKKMER

The search for k-mers in the dictionary within Hamming Distance 1 from D is accomplished by Yu et al. with a Boost multi-index hash table with 4 keys and a tricky algorithm described in the appendix of their paper.

A variant of the SPELLMODELS algorithm [Sag98] could be used here instead:

- The dictionary of k-mers D is mantained in memory as its generalized suffix tree.
- Given a *k*-mer, we search in *D* for the set of *k*-mers at Hamming Distance 1 from it.

Bibliography I

- Scott D. Kahn, *On the future of genomic data*, Science **331** (2011), 728–729.
- Lin Lui, Yinhu Li, Siliang Li, Ni Hu, Yimin He, Ray Pong, Danni Lin, Lihua Lu, and Maggie Law, *Comparison of next-generation sequencing systems*, Journal of Biomedicine and Biotechnology **2012** (2012).
- Gordon E. Moore, *Cramming more components onto integrated circuits*, Electronics (1965), 114–117.
- Marie-France Sagot, *Spelling approximate repeated or common motifs using a suffix tree*, LATIN '98: Theoretical Informatics, Lecture Notes in Computer Science, vol. 1380, 19998, pp. 374–390.

Bibliography II

- Chip Walter, Kryder's law, Scientific American (2005).
- Y. William Yu, Deniz Yorukoglu, and Bonnie Berger, Traversing the k-mer Landscape of NGS Read Datasets for Quality Score Sparsification, Research in Computational Molecular Biology, Lecture Notes in Computer Science, vol. 8394, 2014, pp. 385–399.