# ${\bf BI}$ bliotec ${\bf A}$ de pro ${\bf GR}$ amación científic ${\bf A}$

José Angel de Bustos Pérez {jadebustos@gmail.com>

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# Contents

$\mathbf{W}\mathbf{h}$	at is $B.I.A.G.R.A$ ?	11
1.1	C language?	11
1.2	Some general ideas about $B.I.A.G.R.A$	12
1.3	How to install $B.I.A.G.R.A$ under LiNUX	12
B.I.	A.G.R.A constants (const.h)	<b>15</b>
2.1	Introduction	15
2.2	Mathematical constants	15
2.3	Logical constants	15
2.4	Error constants	16
Me	mory allocation (resmern.h)	17
3.1	Introduction	17
3.2	Vector's memory allocaten	17
	3.2.1 intptMemAll vec vaction	17
	3.2.2 dbitMemAllower function	17
3.3	Matrix memory allocation	18
	3.3.1 dblPtMek llscMat function	18
	3.3.2 dblPtMem 1 ocUpperTrMat function	18
	3.3.3 dblPtMemAllocLowerTrMat function	19
3.4	Freeing memory	20
	3.4.1 freeMemDblMat function	20
Pse	eudo random numbers (random.h)	21
4.1	Introduction	21
4.2	Pseudo random integer numbers	21
	4.2.1 intRandom function	21
	4.2.2 uintRandom function	22
4.3	Pseudo random floating point numbers	22
	4.3.1 dblRandom function	22
	4.3.2 udblRandom function	23
	1.1 1.2 1.3  B.I. 2.1 2.2 2.3 2.4  Me. 3.1 3.2  3.3	1.2 Some general ideas about B.I.A.G.R.A  1.3 How to install B.I.A.G.R.A under LiNUX  B.I.A.G.R.A constants (const.h)  2.1 Introduction  2.2 Mathematical constants  2.3 Logical constants  2.4 Error constants  Memory allocation (resmem.h)  3.1 Introduction  3.2 Vector's memory allocation  3.2.1 intPtMemAllictVect action  3.2.2 dbirtMemAllictVect function  3.3 Matrix memory allocation  3.3.1 dblPtMemAllictMeat function  3.3.2 dblPtMemAllocUpperTrMat function  3.3.3 dblPtMemAllocLowerTrMat function  3.4 Freeing memory  3.4.1 freeMemDblMat function  Pseudo random numbers (random.h)  4.1 Introduction  4.2 Pseudo random integer numbers  4.2.1 intRandom function  4.2.2 uintRandom function  4.3 Pseudo random floating point numbers  4.3.1 dblRandom function

<b>5</b>	Cor	nplex numbers (complejo.h)	25
	5.1	Introduction	25
	5.2		25
		5.2.1 biaComplex data structure	25
			26
	5.3	Arithmetical operations using complex numbers	26
		5.3.1 addComplex function	26
		5.3.2 subtractComplex function	27
		5.3.3 multiplyComplex function	27
		5.3.4 divideComplex function	27
		5.3.5 invSumComplex function	28
		5.3.6 invMulComplex function	28
	5.4	Complex number operations	29
		5.4.1 dblComplexModulus function	29
			29
		5.4.3 conjugateComplex function	29
		5.4.4 complex2Polar function	30
			30
•	<b>.</b>		3-1
6			31
	6.1		31
	6.2		31
	0.0		31
	6.3		31
			31
	0.4		$\frac{32}{32}$
	6.4		32
		6.4.1 ul factorial function	32
7	Pol	ynomial (polynomials.h)	33
	7.1		33
	7.2	Data structures	33
			33
	7.3		34
			34
			$\frac{1}{35}$
			35
			36
		1 0	37
		<del>-</del>	37
			- •

8	Mat	trix (matrix.h)		39
	8.1	Introduction		. 39
	8.2	Data structures		
		8.2.1 biaMatrix data structure		. 39
	8.3	Matrix creation		. 40
		8.3.1 identityMatrix function		. 40
		8.3.2 scalingMatrix function		. 4
		8.3.3 nullMatrix function		. 4
	8.4	Matrix operations		. 4
		8.4.1 transposeMatrix function		. 42
	8.5	Matrix checks		
		8.5.1 isIdentityMatrix function		. 4
		8.5.2 isNullMatrix function		. 43
		8.5.3 isSymmetricMatrix function	<b>)</b>	. 43
9	Roo	ots approximation (roots.h)	,	45
	9.1	Introduction		. 4.
	9.2	Data structures		. 4
		9.2.1 biaRealRoot data structure		. 4
	9.3			
		Function roots approximation		. 40
		9.3.2 newtonMeth 1 function		
<b>10</b>	Run	nge-Kutta methods (raykutta.n)		49
		Introduction		. 49
		Data structures		
		10.2.1 biaBu sherArray data structure		. 4
		10.2.2 biaDatan dzia structure		. 50
	10.3	Node number calculations		. 5
		10.3.1 intNodeNumber function		. 5
	10.4	Explicit Runge-Kutta methods (scalar problems)		. 5
		10.4.1 ExplicitRungeKutta function		
		10.4.2 classicRungeKuttaCoefs function		
		10.4.3 heunRungeKuttaCoefs function		
		10.4.4 kuttaRungeKuttaCoefs function		. 54
		10.4.5 modifiedEulerRungeKuttaCoefs function		. 5
		10.4.6 improvedEulerRungeKuttaCoefs function		. 50
11	Nur	merical integration (numintegr.h)		5'
		Introduction		. 5
		The midpoint rule (threaded version)		

		11.2.1	threadedMidPointRule function	. 57
<b>12</b>	$\pi$ co	mputa	tion (pi.h)	59
	12.1	Introd	uction	. 59
	12.2	The m	idpoint rule method	. 59
		12.2.1	threadedPiMidPointRule function	. 59
	12.3	Chudn	ovsky algorithm	. 60
			chudnovskyPi function	
Ι	Ap	pend	ices	63
A	Run	ıge-Ku	tta methods introduction	65
	A.1	What	is a Runge-Kutta method?	. 65
		A.1.1	What is a I.V.P.?	. 65
	A.2	Runge	-Kutta's method notation	. 66
		A.2.1	General formulation	. 66
		A.2.2	Matricial notation (Butche 's)	. 66
	A.3		-Kutta types	. 67
		A.3.1	Implicit Ry. Kutta	. 67
		A.3.2	Semi-implicit Rung Kutta	. 67
		A.3.3	Explicit Runge-Kutta	. 68

## **Preface**

This library was created when I was studying my degree to ease my academic tasks when I was not forced to use Mathematica, Matlab, SPSS, Linpro, ANSYS, Statgraphics, . . .

B.I.A.G.R.A stands for **BI**bliotec**A** de pro**GB**amación científic**A** which means Scientific Programming Library.

About the name? Well, for those days a new medicine was introduced into the market, my mind . . .

First version was "published" in 1998 but it is was not widely distributed.

Some months ago I started reviewing my backups to centralize useful stuff into only one repository and I found this library. Unfortunately the version I found was one of the first versions so a lot of stuff is missing:

- Matrix operations.
- Matrix factorization
- System linear equations.
- Partial differential equations.
- . . .

This library was originally written in Spanish so I decided to translate into English and making it public available on Github.

It is important to remark that this release candidate version has not been tested and the results maybe not accurate enough. I will check and add tests to check it in a not distant future.

Today there are a lot of resources in many programming languages but in the days this library was published there was no much stuff available due to access to the internet was not as generalized as nowadays.

If you want to know how to use C language for scientific programming this code could be useful, specially how pointers are used to optimize memory usage and how function's pointers are used to abstact C functions from the mathematical functions they are going to use.

You can also check how to use threads using **OpenMP** and how to use the **GNU Multiple Precision Arithmetic Libary (GMP)**.

I will also add some code I created to test how a Fujitsu Primergy server scaled using Pi digits calculations.

I do not have any plans to increase this library. If I and the missing code I will add it though.

It is possible that I add some code from time to time due to the possibility I have to code some numerical methods in the future.

I hope this code is useful for you.

José Ángel de Bustos Pérez

# License

This code is distributed under the GPLv2 License.

The author will be deeply grateful if this code is not used to:

- Attempt World domination.
- Use as compiling massive weapon.
- Translate it into Klingon.
- Ask for /home nationalisms.
- Reverse spelling it.

Just joking!!! Of course you can reverse spelling it (do you dare?).

José Ángel de Bustos Pérez



# Chapter 1

## What is B.I.A.G.R.A?

- B.I.A.G.R.A stands for **BI**bliotec**A** de pro**GR**amación científic**A** which means Scientific Programming Library.
- B.I.A.G.R.A is enterely coded using C language.
- B.I.A.G.R.A has been developed and tested under LiNUX.
- B.I.A.G.R.A is distributed as open source and its author does not take any responibility.
- I wrote B.I.A.G.R.A in the 90s to help me with some academic tasks when studying my degree.

## 1.1 C language?

C language instead of FORTRAN?

- C is modular and structured.
- C is a general purpose language programming.
- C is a very powerful language and its code is very fast.
- C allows dynamic memory allocation.
- C code is portable.
- C is able to handle graphic modes.

#### 1.2 Some general ideas about B.I.A.G.R.A

B.I.A.G.R.A has been developed under **Linux** and some **Linux** knowledge is needed.

B.I.A.G.R.A was developed to solve general problems instead of particular ones. For instance, instead of writing a program to get the inverse of a 4x4 matrix and having to change the source code to get the inverse of a 5x5 matrix B.I.A.G.R.A was developed to allow to write programs to get the inverse of any matrix without having to change de sorce code.

To be able to do that *pointers* were used instead of using *arrays*.

When we talk about *vectors* we will be talking about a *pointer* using dynamic memory allocation. When we talk about matrices we will be talking about *pointer* to a *pointer* using dynamic memory allocation.

B.I.A.G.R.A uses some data structures to store data.

For common errors as:

- Errors in dynamic memory allocation.
- Division by zero.
- ...

B.I.A.G.R.A uses its own constants to notify these errors (Chapter 2).

#### 1.3 How to install B.I.A.G.R.A under LiNUX

To install B.I.A.G.R.A:

# make static

The following actions will take place:

- Header files will be copied to /usr/include/biagra.
- The library will be placed at /usr/lib/libbiagra.a-X.Y.Z and a symbolic link named /usr/lib/libbiagra.a will be created pointing to /usr/lib/libbiagra.a-X.Y.Z.

• Example files will be copied to /usr/share/biagra-X.Y.Z.

To uninstall B.I.A.G.R.A:

# make uninstall





# Chapter 2

# B.I.A.G.R.A constants (const.h)

#### 2.1 Introduction

B.I.A.G.R.A includes its own constants to be used it needed.

These constants are defined in const.

## 2.2 Mathematical constants

Table 2.1 shows the B.I.A.G.R.A 's mathematical constanst.

Constant	Name	Value
е	BIAE	2.71828182845904523536029
$\pi$	BIA_PI	3.14159265358979323846264

Table 2.1: B.I.A.G.R.A mathematical constants.

## 2.3 Logical constants

The following logical constants are defined:

BIA\_FALSE when a condition is not met.

**BIA\_TRUE** when a condiction is met.

## 2.4 Error constants

The following error constants are defined:

 ${\bf BIA\_ZERO\_DIV}$  division by zero.

 ${\bf BIA\_MEM\_ALLOC}$  error in memory allocation.



## Chapter 3

# Memory allocation (resmem.h)

#### 3.1 Introduction

B.I.A.G.R.A includes its own memory allocation functions which are defined in resmem.h file.

## 3.2 Vector's memory allocation

Some functions are provided to handle memory allocations for vectors.

#### 3.2.1 intPtMemAllocVec function

This functions allocates memory for a vector of int.

The definition of this function:

```
double *intPtMemAllocVec(int intElements);
```

This function has only one argument, intElements, which is the dimension of the vector and a int pointer is returned.

#### 3.2.2 dblPtMemAllocVec function

This functions allocates memory for a vector of doubles.

The definition of this function:

```
double *dblPtMemAllocVec(int intElements);
```

This function has only one argument, intElements, which is the dimension of the vector and a double pointer is returned.

#### 3.3 Matrix memory allocation

Some functions are provided to handle memory allocations for vectors.

#### 3.3.1 dblPtMemAllocMat function

This function allocates memory for a matrix of doubles.

The definition of this function:

double \*\*dblPtMemAllocMat(int intRows, int intCols);

where:

intRows number of rows.

intCols number of columns.

#### 3.3.2 dblPtMemAllogOpperTiMat function

This function allocates memory for a upper triangular square matrix.

The definition of this function

double \*\*db1PtMemAllocUpperTrMat(int intOrder);

This function has only one argument, intOrder, which is the order of the matrix and a double pointer to pointer is returned.

In a upper triangular square matrix all elements below the diagonal are zero:

$$\left( \begin{array}{cccccc} a_{0,0} & a_{0,1} & a_{0,2} & a_{0,3} & a_{0,4} \\ 0 & a_{1,1} & a_{1,2} & a_{1,3} & a_{1,4} \\ 0 & 0 & a_{2,2} & a_{2,3} & a_{2,4} \\ 0 & 0 & 0 & a_{3,3} & a_{3,4} \\ 0 & 0 & 0 & 0 & a_{4,4} \end{array} \right)$$

For intOrder = 5:

myMatrix = dbpPtMemAllocUpperTrMat(5);

and:

Pointer	# elements	First element	Last element
myMatrix[0]	5	0	4
myMatrix[1]	4	0	3
myMatrix[2]	3	0	2
myMatrix[3]	2	0	1
myMatrix[4]	1	0	0

so:

$$myMatrix[i][j] = *(*(myMatrix + i) + j) = \begin{cases} a_{i,j+i} & \forall i \leq j \\ 0 & \forall i > j \end{cases}$$

#### 3.3.3 dblPtMemAllocLowerTrMat function

This function allowcates memory for a lower triangular square matrix.

The definition of this function:

double \*\*dblPtMemAllocLowerTrMat(int intOrder)

This function has only one argument, intorder, which is the order of the matrix and a double pointer to pointer is returned.

In a lower triangular square matrix all elements above the diagonal are zero:

$$\begin{pmatrix} a_{0,0} & 0 & 0 & 0 & 0 \\ a_{1,0} & a_{1,1} & 0 & 0 & 0 \\ a_{2,0} & a_{2,1} & a_{2,2} & 0 & 0 \\ a_{3,0} & a_{3,1} & a_{3,2} & a_{3,3} & 0 \\ d_{4,0} & a_{4,1} & a_{4,2} & a_{4,3} & a_{4,4} \end{pmatrix}$$

For intOrder = 5:

myMatrix = dbpPtMemAllocLowerTrMat(5);

and:

Pointer	# elements	First element	Last element
myMatrix[0]	1	0	0
myMatrix[1]	2	0	1
myMatrix[2]	3	0	2
myMatrix[3]	4	0	3
myMatrix[4]	5	0	4

so:

$$myMatrix[i][j] = *(*(myMatrix + i) + j) = \begin{cases} a_{i,j} & \forall i \leq j \\ 0 & \forall i < j \end{cases}$$

## 3.4 Freeing memory

B.I.A.G.R.A includes its own functions to free memory.

#### 3.4.1 freeMemDblMat function



# Chapter 4

# Pseudo random numbers (random.h)

#### 4.1 Introduction

B.I.A.G.R.A includes its own functions to pseudo random number generation and they are defined in random. file.



This functions have not been tested to produce unpredictable sequences, so be careful when use them.

## 4.2 Pseudo random integer numbers

#### 4.2.1 intRandom function

This function generates random integers.

The definition of this function:

int intRandom(int limit);

The pseudo random integer is placed in the interval (-limit, limit).



Before using this function srand must be used to initialize rand. You can use srand((unsigned)time(NULL)).

The pseudo random number is generated with the following formula:

$$\left[\frac{limit \cdot rand()}{RAND\_MAX + 1}\right] \in (-limit, limit)$$

Then randomly is choosed if the number is positive or negative using the above formula with limit = 2 and then taking modulus 2. If modulus is 1 then the number will be a negative one.

#### 4.2.2 uintRandom function

This function generates random integers.

The definition of this function:

int uintRandom(int limit);

The pseudo random integer is placed in the interval [0, limit).



Before using this function srand must be used to initialize rand. You can use srand((unsigned)time(NULL)).

The pseudo random number is generated with the following formula:

$$\begin{bmatrix} limit \cdot rand() \\ RAND\_MAX + 1 \end{bmatrix} \in [0, limit)$$

## 4.3 Pseudo random floating point numbers

#### 4.3.1 dblRandom function

This function generates random floating point numbers.

The definition of this function:

int dblRandom(int limit);

The pseudo random floating point number is placed in the interval (-limit, limit).



Before using this function **srand** must be used to initialize **rand**. You can use **srand((unsigned)time(NULL))**.

The pseudo random number is generated with the following formula:

$$\frac{limit \cdot rand()}{RAND\_MAX + 1} \in (-limit, limit)$$

Then randomly is choosed if the number is positive or negative using the above formula with limit = 2 and then taking modulus 2. If modulus is 1 then the number will be a negative one.

#### 4.3.2 udblRandom function

This function generates random floating point numbers

The definition of this function:

int udblRandom(int limit);

The pseudo random floating point number is placed in the interval [0, limit).



Before using this function srand must be used to initialize rand. You can use srand((unsigned)time(NULL)).

The pseudo random number is generated with the following formula:

$$\frac{limit \cdot rand()}{RAND MAX + 1} \in [0, limit)$$



# Chapter 5

# Complex numbers (complejo.h)

#### 5.1 Introduction

Functions to manage complex numbers are defined in complex h file.

## 5.2 Data structures

Some data structures are defined in B.I.A.G.R.M to manage complex numbers.

#### 5.2.1 biaComplex data structure

```
biaComplex data structure is defined in figure 7.1 where:
```

intDegree polynomial degree

intRealRoots number of real roots (if any).

intCompRoots number of complex roots (if any).

**dblCoef** pointer to store polynomial coeficients.

```
typedef struct {
  double dblReal,
       dblImag;
} biaComplex;
```

Figure 5.1: biaComplex data structure.

#### 5.2.2 biaPolar data structure

This data structure is used to store data for root approximation. Data structure is defined in figure 9.1 where:

intNMI maximum number of iterations to get the root with a maximum error of dblTol.

intIte iterations used to get the root.

**dblx0** initial approximation to get the root.

dblRoot root approximation.

**dblTol** maximum tolerance when calculating the root.

**dblError** error in root approximation. Difference between the las two root approximations.

Figure 5.2: hiaPolar data structure.

# 5.3 Arithmetical operations using complex numbers

#### 5.3.1 addComplex function

This function adds two complex numbers.

The definition of this function:

where:

\*ptCmplx1 first complex number to be added.

\*ptCmplx2 second complex number to be added.

\*ptRes result of the operation.

#### 5.3.2 subtractComplex function

This function subtracts two complex numbers.

The definition of this function:

where:

\*ptCmplx1 complex number.

\*ptCmplx2 complex number to be subtracted to the bove.

\*ptRes result of the operation.

## 5.3.3 multiplyComplex function

This function multiplies two complex numbers.

The definition of this function

where:

ptCmplx1 first complex number to be multiplied.

ptCmplx2 second complex number to be multiplied.

ptRes result of the operation.

#### 5.3.4 divideComplex function

This function divides one complex number by other:

$$\frac{\mathbf{a} + \mathbf{b} \cdot i}{\mathbf{c} + \mathbf{d} \cdot i} = (\mathbf{a} + \mathbf{b} \cdot i) \cdot (\mathbf{c} + \mathbf{d} \cdot i)^{-1}$$

The definition of this function:

where:

\*ptCmplx1 complex number.

\*ptCmplx2 complex number used as divisor.

\*ptRes result of the operation.

The following codes are returned:

BIA_ZERO_DIV	Division by zero
BIA_TRUE	Success

#### 5.3.5 invSumComplex function

This function gets the additive inverse of a complex number

$$\forall z_1 \in \mathbb{C} \quad \exists z_2 \in \mathbb{C} \mid z_1 + z_2 = 0$$

The definition of this function:

void invSumComplex(biaComplex \*ptCmplx, biaComplex \*ptRes);
where:

\*ptCmplx complex number to get its additive inverse.

\*ptRes where the additive inverse will be stored.

#### 5.3.6 invMulComplex function

This function gets the multiplicative inverse of a complex number:

$$\forall z \in \mathbb{C} - \{0\} = \mathbb{C}^* \quad \exists z_2 \in \mathbb{C} \mid z_1 \cdot z_2 = 1$$

The definition of this function:

int invMulComplex(biaComplex \*ptCmplx, biaComplex \*ptRes) ;
where:

\*ptCmplx complex number to get its multiplicative inverse.

\*ptRes where the additive multiplicative will be stored.

The following codes are returned:

BIA_ZERO_DIV	Division by zero
$\mathbf{BIA\_TRUE}$	Success

### 5.4 Complex number operations

#### 5.4.1 dblComplexModulus function

This function gets the modulus of a complex number.

The definition of this function:

double dblComplexModule(biaComplex \*ptCmplx);

where:

\*ptCmplx complex number to get its modulus.

This function returns the complex number modulus.

#### 5.4.2 dblComplexArg function

This function gets the argument of a complex number.

The definition of this function:

double dblComplexArg(biaComplex \*ptCmplx);

where:

\*ptCmplx complex number to get its argument.

This function returns the complex number argument (radians).

#### 5.4.3 conjugateComplex function

This function gets the conjugate complex of a complex number:

$$z = a + b \cdot i \in \mathbb{C} \Rightarrow \overline{z} = a - b \cdot i \in \mathbb{C}$$

The definition of this function:

void conjugateComplex(biaComplex \*ptCmplx, biaComplex \*ptRes);

where:

\*ptCmplx complex number to get its conjugate.

\*ptRes complex conjugate.

#### 5.4.4 complex2Polar function

This function gets the polar coordinates of a complex number.

The definition of this function:

```
void complex2Polar(biaComplex *ptCmplx, biaPolar *ptRes);
where:
```

 ${
m *ptCmplx}$  complex number to calculate polar coordinates.

## 5.4.5 polar2Complex function

This function gets the cartesian coordinates of a polar coordinates for a complex number.

The definition of this function:

```
void polar2Complex(biaPolar *ptPolar, biaComplex *ptRes);
where:
```

<sup>\*</sup>ptRes complex number in cartesian coordinates.



Argument is supposed to be in radians.

<sup>\*</sup>ptRes polar coordinates.

<sup>\*</sup>ptPolar polar coordinates.

# Chapter 6

# Integer numbers (integers.h)

#### 6.1 Introduction

B.I.A.G.R.A includes functions about integer numbers in integers.h file.

## 6.2 Sum integers

#### 6.2.1 uintSumFirstNIntegers function

This function gets the sum of the first n integers.

The definition of this function unsigned uintSumFirstNIntegers(int n);
If the sum is bigger than an unsigned int 0 is returned.

### 6.3 Prime numbers

#### 6.3.1 isPrime function

This function checks if a number is a prime number.

The definition of this function:

int isPrime(int intN);

The following codes are returned:

	intN is not a prime number
BIA_TRUE	intN is a prime number

#### 6.3.2 getFirstNPrimes function

This function checks if a number is a prime number.

The definition of this function:

```
void getFirstNPrimes(unsigned int *ptPrimes, int intNumber, int *ptCalc);
where:
```

\*ptPrimes array where primes will be stored. Memory allocation for this array has to be initialized before using this function.

intNumber number of primes to be computed.

\*ptCalc in this variable the total amount of computed primes will be stored.

#### 6.4 General

#### 6.4.1 uliFactorial function

This function returns a integer's factorial.

The definition of this function:

```
unsigned long int uliFactorial(int intN);
```

where:

intN is the integer the factorial is going to be computed.



Factorial es returned as an unsigned long int.

# Chapter 7

# Polynomial (polynomials.h)

#### 7.1 Introduction

Functions to manage polynomials are defined in polynomial. It file.

A polynomial used to be represented as shown in equation 7.1.

$$p(x) = a_0 + a_1 \cdot x + \dots + a_n \cdot x^n = \sum_{i=0}^n a_i \cdot x^i \quad \text{where } a_i \in \mathbb{R}$$
 (7.1)

## 7.2 Data structures

Some data structures are defined in B.I.A.G.R.A to manage polynomials.

#### 7.2.1 biaRealPol data structure

This data structure is used to handle polinomials  $p(x) \in \mathbb{R}[x]$ . biaPol data structure is defined in figure 7.1 where:

intDegree polynomial degree.

intRealRoots number of real roots (if any).

intCompRoots number of complex roots (if any).

**dblCoefs** pointer to store polynomial coeficients.

**dblRoots** pointer to store real polynomial roots.

```
typedef struct {
  int intDegree = 0,
     intRealRoots = 0,
     intCompRoots = 0;

double *dblCoefs,
     *dblRoots;
} biaRealPol;
```

Figure 7.1: biaRealPol data structure.

Polynomial coeficients are stored in dblCoefs pointer which has to be previously initialized:

```
\begin{array}{cccc}
\text{dblCoefs}[0] & = & a_0 \\
\text{dblCoefs}[1] & = & a_1 \\
& & & & & \\
\text{dblCoefs}[n] & = & a_n
\end{array}
```



dbtRoots are stored in the same way as dblCoefs.

## 7.3 Polynomials operations

#### 7.3.1 addPolyfunction

This function adds two polynomials.

The definition of this function:

```
int addPol(biaPol *ptPol1, biaPol *ptPol2, biaPol *ptRes);
where:
```

**ptPol1** pointer to a biaRealPol struct with the first polynomial to be added.

**ptPol2** pointer to a biaRealPol struct with the second polynomial to be added.

**ptRes** pointer to a biaRealPol struct where the add operation will be stored.

The following codes are returned:

BIA_MEM_ALLOC	Memory allocation error
$oxed{BIA\_TRUE}$	Success



ptRes will be released and memory allocation will be carried out to store the derivative.

ptRes member dblCoefs has to be initialized to a NULL pointer to avoid a Segment Fault error if it was not previously initialized.

#### 7.3.2 dblEvaluatePol function

This function returns the polynomial value in a given point.

The definition of this function:

double dblEvaluatePol(blaRealPol \*ptPol, double dblX);

where:

ptPol pointer to a biaRealPol struct with the polynomial to be evaluated.

**dblX** value in which the polynomial will be evaluated.



Polynomial is defined over  $\mathbb{R}[x]$ .

#### 7.3.3 derivativePol function

This function gets the n-th derivative of a polynomial.

The definition of this function:

int derivativePol(biaPol \*ptPol, biaPol \*ptDer, int intN);

where:

**ptPol** pointer to a biaRealPol struct with the polynomial to get its derivative is stored.

ptDer pointer to a biaRealPol struct where the derivative will be stored.

intN order of the derivative to get.

The following codes are returned:

BIA_MEM_ALLOC	Memory allocation error
$\mathbf{BIA\_TRUE}$	Success



ptDer will be released and memory allocation will be carried out to store the derivative



ptDer member dblGoefs has to be initialized to a NULL pointer to avoid a Segment Fault error if it was not previously initialized.

#### 7.3.4 multiply ol function

This functions multiplies two polynomials.

The definition of this function:

int multiplyPol(biaPol \*ptPol1, biaPol \*ptPol2, biaPol \*ptRes);
where:

ptPol1 pointer to a biaRealPol struct with the first polynomial.

ptPol2 pointer to a biaRealPol struct with the second polynomial.

**ptRes** pointer to a biaRealPol struct where the multiplication operation will be stored.

The following codes are returned:

BIA_MEM_ALLOC	Memory allocation error
BIA_TRUE	Success



ptRes will be released and memory allocation will be carried out to store the derivative.

ptRes member dblCoefs has to be initialized to a NULL pointer to avoid a Segment Fault error if it was not previously initialized.

#### 7.3.5 polFromRealRoots function

This function get the polynomial expression from its real roots.

The definition of this function:

int polFromRealRoots(double \*dblRoots, int intNumber, biaRealPol \*ptPol);
where:

dblRoots pointer to a double with the polynomial real roots.

intNumber number or real roots

ptPol pointer to a biaRealPol struct where the polynomial will be stored.

The following codes are returned:

BIA_MEM_AL	OC	A	Memory allocation error
BIA_TRUE	Ĭ	S	Success



intNumber includes roots multiplicity so intNumber is equal to polynomial degree.



Each polynomial root must be included as many times as its multiplicity.

#### 7.3.6 subtractPol function

This function subtracts two polynomials.

The definition of this function:

int subtractPol(biaPol \*ptPol1, biaPol \*ptPol2, biaPol \*ptRes);

where:

ptPol1 pointer to a biaRealPol struct with the first polynomial.

**ptPol2** pointer to a biaRealPol struct with the polynomial to be subtracted from the above.

**ptRes** pointer to a biaRealPol struct where the subtract operation will be stored.

The following codes are returned:

BIA_MEM_ALLOC	Memory allocation error
BIA_TRUE	Success



ptRes will be released and memory allocation will be carried out to store the derivative.



ptRes member dblCoefs has to be initialized to a NULL pointer to avoid a Segment Fault error if it was not previously initialized.

# Chapter 8

# Matrix (matrix.h)

## 8.1 Introduction

Functions to manage matrices are defined in matrix in file.

#### 8.2 Data structures

Some data structures are defined in B.I.A.G.R.A to manage matrices.

#### 8.2.1 biaMatrix data structure

This data structure is used to store a matrix. **biaMatrix** data structure is defined in figure 8.1 where:

intRows number of rows.

intCols number of columns.

dblCoefs pointer to store matrix coeficients.

```
typedef struct {
  int intRows,
      intCols;

double **dblCoefs;
} biaMatrix;
```

Figure 8.1: biaMatrix data structure.

#### 8.3 Matrix creation

B.I.A.G.R.A includes functions to create some kind of matrices.

#### 8.3.1 identityMatrix function

This function stores the identity matrix with order taken from intRows member of ptMatrix:

The definition of this function:

```
void identityMatrix(biaMatrix *ptMatrix);
```

where:

\*ptMatrix matrix that has to be created before using this function. Memory allocation for dblCoefs must be done before using this function.



intRows is used to get the matrix order.

#### 8.3.2 scalingMatrix function

This function stores the scaling matrix with factor  $\lambda$  and order taken from intRows member of ptMatrix:

$$\begin{pmatrix} \lambda & 0 & 0 & 0 & 0 \\ 0 & \lambda & \ddots & 0 & 0 \\ 0 & \ddots & \ddots & \ddots & 0 \\ 0 & 0 & \ddots & \lambda & 0 \\ 0 & 0 & 0 & 0 & \lambda \end{pmatrix}$$

The definition of this function:

void scalingMatrix(biaMatrix \*ptMatrix, double dolFactor)

where:

\*ptMatrix matrix that has to be created before using this function. Memory allocation for dblCoefs must be done before using this function.



intRows is used to get the matrix order.

## 8.3.3 nullMatrix function

This function stores the null matrix with order taken from intRows member of ptMatrix:

$$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \ddots & 0 & 0 \\ 0 & \ddots & \ddots & \ddots & 0 \\ 0 & 0 & \ddots & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

The definition of this function:

void nullMatrix(biaMatrix \*ptMatrix);

where:

\*ptMatrix matrix that has to be created before using this function. Memory allocation for dblCoefs must be done before using this function.



intRows and intCols is used to get the matrix order.

## 8.4 Matrix operations

#### 8.4.1 transposeMatrix function

This function stores the transpose matrix of a given matrix.

The definition of this function:

void transposeMatrix(biaMatrix \*ptMatrix, biaMatrix \*ptRes);
where:

<sup>\*</sup>ptRes matrix to store the transpose matrix. Memory has to be preallocated before using this function.



intRows and IntCols is used to get the matrix order.

#### 8.5 Matrix checks

#### 8.5.1 isIdentityMatrix function

This function checks if a matrix is the identity matrix.

The definition of this function:

int isIdentityMatrix(biaMatrix \*ptMatrix);

where:

\*ptMatrix matrix to check.

<sup>\*</sup>ptMatrix matrix to get its transpose matrix

#### 8.5.2 isNullMatrix function

This function checks if a matrix is a null matrix.

The definition of this function:

```
int isNullMatrix(biaMatrix *ptMatrix, double dblTol);
where:
```

\*ptMatrix matrix to check.

**dblTol** if a matrix element is minor than this value it is assumed it is a null element.

#### 8.5.3 isSymmetricMatrix function

This function checks if a matrix is a symmetric matrix

The definition of this function:

```
int isSymmetricMatrix(biaMatrix *ptMatrix);
where:
```

<sup>\*</sup>ptMatrix matrix to check.



## Chapter 9

# Roots approximation (roots.h)

#### 9.1 Introduction

Functions to compute function's roots approximation are defined in roots.h file.

#### 9.2 Data structures

Some data structures are defined in B.I.A.G.R.A to manage roots.

#### 9.2.1 biaRealRoot data structure

This data structure is used to store data for root approximation.

Data structure is defined in figure 9.1 where:

intNMI maximum number of iterations to get the root with a maximum error of dblTol.

intIte iterations used to get the root.

**dblx0** initial approximation to get the root.

**dblRoot** root approximation.

**dblTol** maximum tolerance when calculating the root.

**dblError** error in root approximation. Difference between the las two root approximations.

```
typedef struct {
  int intNMI,
    intIte;

double dblx0,
    dblRoot,
    dblTol,
    dblError;
} biaRealRoot;
```

Figure 9.1: biaRealRoot data structure.

## 9.3 Function roots approximation

#### 9.3.1 newtonPol function

This function approaches a polynomial root using the Newton method.

The definition of this function:

```
int newtonPol(biaPol *ptPol, biaRealRoot *ptRoot);
```

The following codes are returned:

BIA_MEM_ALLO	Memory allocation error
BIA_ZÉRO DIV	Division by zero
BIA_TRUE	the root was computed satisfying the problem
	conditions (intMNI and dblTol);
BIA_FALSE	Root approximation could not be calculated
	satisfying the requirements (intMNI and dblTol).

The following values in \*ptRoot need to be initialized:

intMNI maximum number of iterations to compute the root.

dblx0 initial approximation.

dblTol tolerance to compute de root.

The following data will be stored:

intIte iterations used to compute the root.

dblRoot approximation of the root.

dblError error in the approximation.



When two consecutive approximations are close enough, dblTol, last approximation will be considered as good and will be stored in \*biaRealRoot \*ptRoot in dblRoot.

#### 9.3.2 newtonMethod function

This function approaches a function's root using the **Newton** method.

The definition of this function:

```
int newtonMethod(biaRealRoot *ptRoot,
    int (*func)(double dblx0, double *ptRes),
    int (*der)(double dblx0, double *ptRes));
```

Function's pointers are used to avoid having to recode the C function every time a root function need to be approximated for different mathematical functions.

The following codes are returned:

BIA_ZERO_DIV	Division by zero
BIA_TRUE	the root was computed satisfying the problem
	conditions (intMNI and dblTol);
BIA_FALSE	Root approximation could not be calculated
	satisfying the requirements (intMNI and dblTol).

where:

ptRoot is a pointer to a biaRealRoot variable.

**func** pointer to a function implementing the function which root is going to be computed.

**der** pointer to a function implementing the derivative of the function which root is going to be computed.

The following values in \*ptRoot need to be initialized:

intMNI maximum number of iterations to compute the root.

dblx0 initial approximation.

dblTol tolerance to compute de root.

The following data is stored:

intIte iterations used to compute the root.

**dblRoot** approximation of the root.

dblError error in the approximation.



When two consecutive approximations are close enough, dblTol, last approximation will be considered as good and will be stored in \*biaRealRoot \*ptRoot in dblRoot.

#### Usage example

To approximate a root for the function  $f(x) = \sqrt{2}$  to functions need to be created. A C function for the mathematical function implementation:

```
/* f(x) = x^2 - 2 */
int myfunc(double x0, double *fx0) {
  *fx0 = (double)(x0 * x0 2.);
  return BIA_TRUE;
}
```

A C function for the mathematical derivative function implementation:

```
/* f'(x) = 2*x */
int myfuncder(double x0, double *fx0) {
  *fx0 = 2.*x0;
  return BIA_TRUE;
}
```

Both functions must meet the following requirements:

- An integer value is returned, **BIA\_TRUE** when function is evaluated in x0 and **BIA\_ZERO\_DIV** if a division by zero takes place.
- x0 value to evaluate the function.
- \*fx0 pointer to a double to store the function's value in x0.

So to approximate function's root using Newton Method:

```
i = newtonMethod(&myRoot, &myfunc, &myfuncder);
```

# Chapter 10

# Runge-Kutta methods (rngkutta.h)

#### 10.1 Introduction

Runge-Kutta are a family of implicit and explicit teracive methods used to approximate solutions of ordinary differential equations or ODE.

Butcher matricial notation is used in this implementation.

## 10.2 Data structures

#### 10.2.1 biaButcherArray data structure

This structure is used to store the Butcher matricial notation.

Data structure is defined in figure 10.1 where:

intStages method stages.

\*dblC  $c_i$  coefficients stored in an array with size intStages.

\*dblB  $b_i$  coefficients stored in an array with size intStages.

\*\*dblMatrix matrix to store  $a_{i,j}$  method's coeficients.

```
typedef struct {
  double *dblC,
          *dblB,
          **dblMatrix;
  int intStages;
} biaButcherArray;
```

Figure 10.1: biaButcherArray data structure.



See appendix A if you need information about Butcher matricial notation.

#### 10.2.2 biaDataRK data structure

This structure is used to store all the data needed to apply a Runge-Kutta method.

Data structure is defined in figure 10.2 where:

intNumApprox number of approximations to be done (size of the array dblPoints).

intImplicit when the Runge Kutta method is implicit or not. The following constants are defined in the header file:

	Name	Value
BI	A_IMPLICIT_RK_TRUE	0
BI	A_IMPLICIT_RK_FALSE	1

\*dblPoints array with dimension intNumApprox and its elements will be the approximations in  $x_i$  where:

```
x_i = dblFirst + i \cdot dblStepSize where 0 \le i < intNumApprox
```

dblStepSize method's step-size.

**dblFirst** first point used to compute all the approximations. The value of the function in this point is known (initial condition).

dblLast last point in which approximations will be computed.

**strCoefs** variable of type biaButcherArray (section 10.2.1) storing Butcher matricial notation.

```
typedef struct {
  int intNumApprox,
    intImplicit;

double *dblPoints,
    dblStepSize,
    dblFirst,
    dblLast;

biaButcherArray strCoefs;
} biaDataRK;
```

Figure 10.2: biaDataRK data structure

#### 10.3 Node number calculations

#### 10.3.1 intNodeNumber function

This function gets the number of nodes that can be placed in an interval. All nodes are equidistant.

The definition of this function:

int intNodeNumber(double dblLong, double dblStepSize)
where:

dblLong interval length.

dblStepSize distance between two nodes.

The function returns the number of nodes that can be placed.



Arguments are supposed to be different from zero.



Arguments are supposed to be positive.

# 10.4 Explicit Runge-Kutta methods (scalar problems)

#### 10.4.1 ExplicitRungeKutta function

This function solves an I.V.P. using an explicit Runge-Kutta method.

The definition of this function:

where:

ptData pointer to a biaDataRK variable. This variable contains all the necessary data to solve the *I.V.P.* 

**IVP** C function's pointer to a function implementing the O.D.E.. This functions needs to have two double arguments and returns the value of the I.V.P.:

**dblX** point where we want to evaluate the O.D.E. **dblY** O.D.E. value in dblX  $(y_i \approx y(x_i)).$ 

The following codes are returned:

	Memory error allocation
BIA_TRUE	Success

#### Usage example

For instance, to so we this I.V.P.:

$$\begin{cases} y' = y(x) * \frac{x - y(x)}{x^2} \\ y(1) = 2 \end{cases}$$

the implementation of the IVP would be:

A biaDataRK variable has to be initialized. B.I.A.G.R.A provides several functions to initialize the biaButcherArray member.

To solve the I.V.P. problem:

i = explicitRungeKutta(&varDataRK, IVP);

#### 10.4.2 classicRungeKuttaCoefs function

This function initialize the Butcher array for the classic Runge-Kutta method (four-stage method and order two).

The definition of this function:

where:

**ptData** pointer to a **biaDataRK** variable<sup>2</sup>. In this variable the following members will be initialized:

intStages will be initialized to 4.

intImplicit will be initialized to BIA\_IMPLICIT\_RK\_FALSE.

strCoefs will be initialized with the Butcher array.

The Butcher array for this method:

The following codes are returned:

BIA_MEM_ALLOC	Memory error allocation
$oxed{BIA\_TRUE}$	Success

 $<sup>^{2}</sup>$ Section (10.2.2).

#### 10.4.3 heunRungeKuttaCoefs function

This function initialize the Butcher array for the Heun Runge-Kutta method (three-stage method and order three).

The definition of this function:

int heunRungeKuttaCoefs(biaDataRK \*ptData);

where:

ptData pointer to a biaDataRK variable<sup>3</sup>. In this variable the following members will be initialized:

intStages will be initialized to 3.

intImplicit will be initialized to BIA\_MPLICIT\_RK\_FALSE.

strCoefs will be initialized with the Butcher array.

The Butcher array for this method.



The following codes are returned:

BIA_ME	$\mathbf{M}_{-}\mathbf{A}$	LOC	Memory error allocation
$BIA_{-}TR$	JE		Success

## 10.4.4 kutta ungeKutta Coefs function

This function initialize the Butcher array for the Kutta Runge-Kutta method (three-stage method and order three).

The definition of this function:

int kuttaRungeKuttaCoefs(biaDataRK \*ptData);

where:

**ptData** pointer to a **biaDataRK** variable<sup>4</sup>. In this variable the following members will be initialized:

 $<sup>^{3}</sup>$ Section (10.2.2).

<sup>&</sup>lt;sup>4</sup>Section (10.2.2).

intStages will be initialized to 3.

intImplicit will be initialized to BIA\_IMPLICIT\_RK\_FALSE.

strCoefs will be initialized with the Butcher array.

The Butcher array for this method:

$$\begin{array}{c|cccc}
0 & 0 \\
\frac{1}{2} & \frac{1}{2} & 0 \\
1 & -1 & 2 & 0 \\
\hline
& \frac{1}{6} & \frac{2}{3} & \frac{1}{6}
\end{array}$$

The following codes are returned:

BIA_MEM_ALLOC	Memory error allocation
BIA_TRUE	Success

#### 10.4.5 modifiedEulerRungeKuttaCoefs hunckion

This function initialize the Butcher array for the modified Euler Kutta Runge-Kutta method (two-stage method and order two).

The definition of this function.

int modifiedEulerRungeKuttaCoefs(biaDataRK \*ptData);

where:

ptData pointer to a biaDataRK variable<sup>5</sup>. In this variable the following members will be initialized:

intStages will be initialized to 2.

intImplicit will be initialized to BIA\_IMPLICIT\_RK\_FALSE.

strCoefs will be initialized with the Butcher array.

The Butcher array for this method:

$$\begin{array}{c|cccc}
0 & 0 \\
\frac{1}{2} & \frac{1}{2} & 0 \\
\hline
& 0 & 1 \\
\end{array}$$

The following codes are returned:

BIA_MEM_ALLOC	Memory error allocation
${f BIA\_TRUE}$	Success

<sup>&</sup>lt;sup>5</sup>Section (10.2.2).

#### 10.4.6 improvedEulerRungeKuttaCoefs function

This function initialize the Butcher array for the improved Euler Runge-Kutta method (two-stage method and order two).

The definition of this function:

int improvedRungeKuttaCoefs(biaDataRK \*ptData);

where:

ptData pointer to a biaDataRK variable<sup>6</sup>. In this variable the following members will be initialized:

intStages will be initialized to 3.

intImplicit will be initialized to BIA\_IMPLICIT\_RK\_FALSE

strCoefs will be initialized with the Butcher array.

The Butcher array for this method:



The following codes are returned:

BIA_ME	MALLOC	Memory error allocation
${ m BIA\_TR}$	JE 🖊	Success

 $<sup>^{6}</sup>$ Section (10.2.2).

## Chapter 11

# Numerical integration (numintegr.h)

#### 11.1 Introduction

In this chapter numerical methods to compute defined integrals will be defined.

The following libraries will be used:

• OpenMP the Open Multi Processing API which provides multiplatform shared memory capabilities for parallel programming.

## 11.2 The midpoint rule (threaded version)

This implementation of the proposition rule is a multithreaded implementation using OpenMP.

Using n subintervals the midpoint rule:

$$\int_{a}^{b} f(x)dx \approx \sum_{i=0}^{n-1} f(\frac{1}{2} \cdot (x_i + x_{i+i})) \triangle x \quad \text{where} \quad \triangle x = \frac{b-a}{n}$$

#### 11.2.1 threadedMidPointRule function

This function uses a threaded version of the midpoint rule to compute an integral over an inteval.

The definition of this function:

where:

intThreads number of threads to be used.

intN number of subintervals to be used.

- a Interval left endpoint.
- **b** Interval right endpoint.

func C function's pointer to a function implementing the function we want to compute the integral. This function needs to have one long double argument and returns the function's value in the argument point (as long double).

#### Usage example

To compute:



the implementation of the func would be:

```
long double func(double x) {
  return (1/(1+(x*x)));
}
```

and to compute it using 4 threads an 15000 subintervals:

```
value = threadedMidPointRule(4, 15000, 0., 1., &func);
```

## Chapter 12

# $\pi$ computation (pi.h)

#### 12.1 Introduction

In this chapter several methods will be shown to compute x digits.

The following libraries will be used:

- OpenMP the Open Multi-Processing API which provides multiplatform shared memory capabilities for parallel programming.
- GMP the GNU Multiple Prediction Arithmetic Library.

## 12.2 The midpoint rive method

To compute  $\pi$  value using the midpoint rule the following is used:

$$\int_0^1 \frac{1}{1+x^2} dx = \frac{\pi}{4}$$

So:

$$\pi = 4 \cdot \int_0^1 \frac{1}{1 + x^2} dx$$

#### 12.2.1 threadedPiMidPointRule function

This function approaches  $\pi$  using the midpoint rule.

The definition of this function:

long double threadedPiMidPointRule(int intThreads, int intN);

where:

intThreads number of threads to be used.

intN number of subintervals to be used.



To compute the integral the midpoint rule will be used. The implementation used for the midpoint rule used is the one provided by B.I.A.G.R.A in section 11.2.1.

### 12.3 Chudnovsky algorithm

The Chudnovsky algorithm uses the following to compute  $\pi$ :

$$\pi \cdot \sum_{i=0}^{\infty} \frac{(6 \cdot i)! \cdot (13591409 + 545140134 \cdot i)}{(3 \cdot i)! \cdot (i!)^3 \cdot (-640320)^{3 \cdot i}} = 126880 \cdot \sqrt{10005}$$

For high performance computation.

$$\pi \cdot \sum_{i=0}^{\infty} \frac{(6 \cdot i)! \cdot (13591409 + 545140134 \cdot i)}{(3 \cdot i)! \cdot (i!)^3 \cdot (-262537412640768000)^i} = 426880 \cdot \sqrt{10005}$$

## 12.3.1 chudnovskyP function

This function approaches a using the Chudnovsky algorithm.

The definition of this function:

int chudnovskyPi(int intThreads, int intPiDigits, char \*ptPi);
where:

intThreads number of threads to be used.

**intPiDigits** number of  $\pi$  digits to compute.

**ptPi** char point to store  $\pi$  as a string.



This algorithm uses the GNU Multiple Precision Library (GMP) to be able to perform all the computations. C default number types do not provide enough precision.



This algorithm uses  ${\it OpenMP}$  to be able to parallelize computations.



ptPi must be initialized to store  $\pi$  as a string. For instance to store 100  $\pi$  digits the size to allocate memory is 102 to include the "." character and the "null character" at the end of the string.









# Appendix A

# Runge-Kutta methods introduction

This appendix is intended to help to know how Runge-Kutta methods are implemented and used in this library.

## A.1 What is a Runge-Kutta method?

Runge-Kutta methods are a family of numerical methods to approach solutions of ordinary differential equations (O.D.E). These methods are iterative methods used to solve "initial problem value" (I.P.V) or "Cauchy problem".

These methods are only-one-step methods with a fixed size for the method step<sup>1</sup>.

#### A.1.1 What is a I.V.P.?

An I.V.P. is:

$$\begin{cases} y' = f(x, y(x)) \\ y(x_0) = y_0 \end{cases}$$
 (A.1)

So y' is a function depending on the variable x, and the function y(x). y(x) is the solution of the equation A.1 and the point  $(x_0, y_0)$  belongs to the curve y(x).

<sup>&</sup>lt;sup>1</sup>It is also possible to implement methods with a variable step known as *embedding*.

Solving the I.V.P. A.1 is finding a function y(x) such as the equation A.1 is met.

An example of a I.V.P.:

$$\begin{cases} y' = \frac{x * y(x) - y(x)^2}{x^2} \\ y(1) = 2 \end{cases}$$
 (A.2)

The solution of the A.2 will be:

$$y(x) = \frac{x}{\frac{1}{2} + \ln x} \tag{A.3}$$

## A.2 Runge-Kutta's method notation

 $y(x_i)$  will be the exact value of the function y(x) evaluated in  $x_i$ .  $y_i$  will be the approximation of the function y(x) in the point  $x_i$ . h is the step used by the method in each iteration.

## A.2.1 General formulation

A s-stages Runge-Kutta's method formulation is:

$$y_{n+1} = y_n + h \cdot \sum_{i=0}^{s-1} b_i \cdot k_i$$
 (A.4)

where:

$$k_i = f(x_n + c_i \cdot h, y_n + h \cdot \sum_{j=0}^{s-1} a_{i,j} \cdot k_j)$$
 (A.5)

satisfying:

$$\sum_{j=0}^{s-1} a_{i,j} = c_i \tag{A.6}$$

### A.2.2 Matricial notation (Butcher's)

Matricial notation is used to represent method's coeficients using a matrix.

For a s-stages Runge-Kutta method the matricial notation will be:



In section 10.2.1 is shown a data structure used to store the Butcher array.

## A.3 Runge-Kutta types

There are several types of Runge-Kutta methods.

## A.3.1 Implicit Runge-Kutta

A Runge-Kutta method is said to be implicit when the  $a_{i,j} \neq 0$  for some j > i.

The 2-stages Gauss method is an implicit Runge-Kutta method of 2-stages:

## A.3.2 Semi-implicit Runge-Kutta

A **Runge-Kutta** method is said to be semi-implicit when the  $a_{i,j} = 0$  when j > i.

A 2-stages semi-implicit Runge-Kutta method:

$$\begin{array}{c|c}
\frac{3+\sqrt{3}}{6} & \frac{3+\sqrt{3}}{6} & 0 \\
\frac{3-\sqrt{3}}{6} & \frac{-\sqrt{3}}{3} & \frac{3+\sqrt{3}}{6} \\
\hline
 & \frac{1}{2} & \frac{1}{2}
\end{array}$$

## A.3.3 Explicit Runge-Kutta

A Runge-Kutta method is said to be explicit when the  $a_{i,j}=0$  when  $j\geq i$ .

A 4-stages explicit  $\bf Runge\text{-}Kutta$  method also known as "classic Runge-Kutta":

