

Instructions

- You have one hour and fifty minutes of “in-class time” followed by three days of “take-home” time to complete this test.
- The exam consists of five questions, each of equal value.
- You may choose any **four of five** problems to complete during the in-class time. Please indicate clearly which problems you wish to have graded. You must answer **all five** questions during the take-home time.
- Take-home due date: Friday, October 29, 2021, 1:30pm (in class). Please return a *complete solution* for the take-home portion, even if you believe that you answered the question correctly in class.
- Your grade will be computed as a weighted average of your “in-class” grade and your “take-home” grade. (Weights to be determined later.)
- The exam is open book, and all aids are permitted, but **all work is to be done independently**. Consultation with others is **not** permitted. Good luck!

1. Let X be a memoryless Bernoulli source over $\mathcal{X} = \{0, 1\}$, where $p_X(0) = 2/3$ and $p_X(1) = 1/3$.
 - (a) Compute the entropy of this source.
 - (b) Determine the rate of a binary Huffman code (with 8 codewords, possibly of variable length) for the third extension of this source.
 - (c) Determine the rate of a binary Tunstall code (with 8 codewords of fixed length) for this source.
 - (d) Determine the probability of each of the eight phrases produced by the Tunstall parsing tree of (c). Is the fixed length code optimal for this distribution?
2. Let X be a discrete memoryless source with alphabet \mathcal{X} , and entropy $H(X)$ in bits. For a positive integer n and any $\epsilon > 0$, let

$$A_\epsilon^{(n)} = \left\{ (x_1, \dots, x_n) \in \mathcal{X}^n : \left| -\frac{1}{n} \log p(x_1, \dots, x_n) - H(X) \right| \leq \epsilon \right\}$$

be the typical set as defined in class, where $p(x_1, \dots, x_n) = \prod_{i=1}^n p(x_i)$. For $0 < \delta \leq 1$, let

$$B_\delta^{(n)} = \left\{ (x_1, \dots, x_n) \in \mathcal{X}^n : \left| -\frac{1}{n} \log p(x_1, \dots, x_n) - H(X) \right| \leq \delta H(X) \right\}.$$

Prove the following.

- (a) If $(x_1, \dots, x_n) \in B_\delta^{(n)}$, then

$$2^{-n(1+\delta)H(X)} \leq p(x_1, \dots, x_n) \leq 2^{-n(1-\delta)H(X)}.$$

- (b) $\Pr(B_\delta^{(n)}) \geq 1 - \delta$ when n is sufficiently large.
- (c) $|B_\delta^{(n)}| \leq 2^{n(1+\delta)H(X)}$.
- (d) $|B_\delta^{(n)}| \geq (1 - \delta)2^{n(1-\delta)H(X)}$ when n is sufficiently large.
- (e) Show that either $A_\epsilon^{(n)} \subseteq B_\delta^{(n)}$ or $B_\delta^{(n)} \subseteq A_\epsilon^{(n)}$. Under what conditions are the two sets equal?

3. Let X be a random variable with range $\mathcal{X} = \{1, 2, 3, \dots\}$ and expected value $E(X)$.

- (a) Show that $H(X) \leq E(X)$ when $H(X)$ is measured in bits.
- (b) When does equality occur?

Now let Y be a random variable with range $\mathcal{Y} = \{0, 1, 2, 3, \dots\}$ and expected value $E(Y)$.

- (c) Show that $H(Y) \leq 1 + E(Y)$ when $H(X)$ is measured in bits.
- (d) When does equality occur?

4. Random variables X and Y are jointly distributed over the finite alphabet $\mathcal{X} \times \mathcal{Y}$ according to the joint probability mass function $p_{X,Y}(x, y) = p_X(x)p_{Y|X}(y | x)$, where $p_X(x)$ denotes, for any $x \in \mathcal{X}$, the marginal probability mass function for X , and $p_{Y|X}(y | x)$ denotes, for any $x \in \mathcal{X}$ and any $y \in \mathcal{Y}$, the conditional probability mass function of Y given that $X = x$.

The mutual information $I(X; Y)$ between X and Y , which is a function of p_X and $p_{Y|X}$, will be denoted as $I(X; Y) = I(p_X; p_{Y|X})$.

- (a) Show that $I(p_X; p_{Y|X})$ is **concave** in p_X for fixed $p_{Y|X}$, i.e., for any $\lambda \in [0, 1]$, and any probability mass functions p'_X and p''_X , the mutual information satisfies

$$I(p_X^*; p_{Y|X}) \geq \lambda I(p'_X; p_{Y|X}) + (1 - \lambda) I(p''_X; p_{Y|X}),$$

where $p_X^* = \lambda p'_X + (1 - \lambda) p''_X$. *Hint:* The inequality $A \geq B$ is equivalent to $A - B \geq 0$.

- (b) Show that $I(p_X; p_{Y|X})$ is **convex** in $p_{Y|X}$ for fixed p_X , i.e., for any $\lambda \in [0, 1]$, and any conditional probability mass functions $p'_{Y|X}$ and $p''_{Y|X}$, the mutual information satisfies

$$I(p_X; p_{Y|X}^*) \leq \lambda I(p_X; p'_{Y|X}) + (1 - \lambda) I(p_X; p''_{Y|X}),$$

where $p_{Y|X}^* = \lambda p'_{Y|X} + (1 - \lambda) p''_{Y|X}$. *Hint:* The inequality $A \leq B$ is equivalent to $B - A \geq 0$.

5. In a coin-tossing game, a player must try to guess the outcome of a fair coin toss. The guess must be made while the coin is in the air, i.e., after it is flipped, but before it is caught. A one dollar bet returns one additional dollar if the guess is correct. In other words, if C_0 denotes a player's total capital *before* the bet, and a player bets B dollars, then the player's total capital C_1 *after* the bet is given by

$$C_1 = \begin{cases} C_0 + B & \text{if the guess is correct;} \\ C_0 - B & \text{if the guess is incorrect.} \end{cases}$$

After years of patient observation, a player has learned to partially predict the outcome of the coin toss. By using this information, they hope to gain a tidy sum for their retirement.

Let X denote the outcome of a given coin toss, and let Y denote the player's prediction. Assume that $P[X = Y] = p > 1/2$, and that successive coin tosses are independent.

- (a) Calculate $I(X; Y)$.
- (b) The player has an initial capital C_0 . They have decided to wager a fraction q of their total capital on each successive coin toss, and withhold $1 - q$. After N trials, show that their total capital is the random variable

$$C_N = C_0 \prod_{n=1}^N (1 + q)^{Z_n} (1 - q)^{1 - Z_n}$$

where $Z_n = 1$ if the n th prediction is correct and 0 otherwise. Define the *growth rate* as

$$R_N = \frac{1}{N} \log_2 \frac{C_N}{C_0}.$$

Find the value of q that maximizes the expected growth rate $E(R_N)$. Also find the value of q that maximizes $E(C_N)$. Compare the maximum expected growth rate with $I(X; Y)$.

- (c) If you were the player, which value of q would you use and why? *Hint:* does the law of large numbers apply to either R_N or C_N as $N \rightarrow \infty$?