



#lang primal-form

Jacob Adley and Lalo Viduarri



Fundamental Theorem of Arithmetic

- ∫ Every integer $n \geq 2$ can be factored into a product of primes in exactly one way (aside from rearranging the factors)
- ∫ $\mathbb{Z} \rightarrow$ Prime Factorization
- ∫ Default form of integers
- ∫ Still have access to base 10 integers
- ∫ Make working easy in this form
- ∫ User does not have to know implementation to do basic work



Why

- ∫ By representing number in their primal-form it becomes very simple to do some normally 'complex' operations.
 - ∫ Multiplication -> Set union
 - ∫ Division -> Set subtraction
- ∫ Facilitate reasoning and working with numbers in their prime factorization
 - ∫ Euler's Totient
 - ∫ Co-primes (\equiv greatest common denominator)



Consequences

∫ Big numbers = slow

∫∫ We have ideas

∫ Simple is hard?

∫∫ add1

∫∫ sub1

∫∫ +

∫∫ -

∫ Going between representations of integers



Data Type

- ∫ Normal stuff
 - ∫∫ Intuitive
 - ∫∫ Clean
 - ∫∫ Dank



Data Type: Syntax

∫ User input

∫∫ Raw numbers

∫∫∫ \mathbb{Z}

∫∫ Primal numbers: **Literals:** $^$: () **Variables:** x ,

$w \in \mathbb{Z}; y, z \in \mathbb{N}$

∫∫∫ (x)

∫∫∫ $(x\ y)$

∫∫∫ $(x : y)$

∫∫∫ $(x \wedge y)$

∫∫∫ $(x \wedge y\ w)$

∫∫∫ $(x \wedge y : w)$

∫∫∫ $(x \wedge y : w \wedge z)$

∫∫∫ $(x \wedge y\ w \wedge z)$

∫∫ Negative primal numbers

∫∫∫ $(\text{neg } \dots)$

∫∫∫ $(\neg \dots)$



Data Type: Syntax II

∫ User input

∫∫ Normal integers, $n \in \mathbb{N}$

∫∫∫ $(\text{int } n) \rightarrow n$

∫∫ The no kiddin operator !t

∫∫∫ Trust on user

∫∫∫ No checks performed on input values

∫∫∫ $(!t \dots) \rightarrow (\dots)$



Data Type: Implementation

∫ #%datum

∫ #%app

∫ parse-primal

∫ !-parse-primal

∫ base-normalize



Match

- ∫ Needed to provide ways to allow the programmer to use numbers without having to dig into our implementation of them. (lists)
- ∫ Adding match cases was a simple way to this.
 - ∫∫ Allows programmers to access and describe properties that they want numbers to have.
 - ∫∫ Since we are using match expander then programmers can use the match they are used to



Match: Syntax

Match on a Factor

```
(mach num  
  [(fac 7 -> a b) bod])
```

Looks to see if there is a factor
of 7

Match on a power

```
(mach num  
  [(pow 4 -> a b) bod])
```

Looks for anything raised to the
4th power

Match on both

```
(match num  
  [(n-to-the 3 6 -> a b) bod])
```

Looks for that specific factor
raised to that power



Match: Implementation

```
(define-match-expander fac
  (λ (stx)
    (syntax-case stx (>> >>))
    [(_ num >> n m) #`(app (factor num) n m)]
    [(_ num >> n m) #`(app (factor num) n m))]))
```

- ∫ Write a function to traverse a primal
- ∫ Use matches app syntax to call it and bind the variables.
- ∫ Function always succeeds and needs to be wrapped with guards since it always succeeds.



Dank Examples

\int φ

\int gcd

\int lcm

\int prime?



Going →: Changes

∫ Efficiency

- ∩∩ Table with previous computed primes

∫ Printing

- ∩∩ Current print shows implementation

- ∩∩ Want to print similar to construction

- ∩∩∩ Use structs

∫ Normalize output of implemented functions

- ∩∩ Can have -1^x s.t. $x > 1$

- ∩∩ ****Bug****



Going →: Future Implementation

∫ Variables in number construction

∫∫ (define x 5)

∫∫ (define y 2)

∫∫ (x ^ y)

∫∫ (!t x ^ y)

∫ Modular Arithmetic

∫∫ Define at the top of the file what space to work in: $\mathbb{Z} \pmod{n}$