

1. (a)  $total = 0 - \Theta(1)$   
 for  $i$  in  $range(n)$ :  
      $total += i - \Theta(1)$   $\left[ \Theta(n) \right]$   
     for  $j$  in  $range(1, n+1)$ :  
          $total += j - \Theta(1)$   $\left[ \Theta(n) \right]$   $\left[ \begin{array}{l} \Theta(2n) \\ \parallel \\ \Theta(n) \end{array} \right]$

$$\therefore T_{worst}(n) = \Theta(n)$$

(b)  $total = 0 - \Theta(1)$   
 for  $i$  in  $range(n)$ :  
      $total += i - \Theta(1)$   
     for  $j$  in  $range(i, n)$ :  
          $total += j - \Theta(1)$   $\left[ \Theta(n) \right]$   $\left[ \Theta(n^2) \right]$

$$\therefore T_{worst}(n) = \Theta(n^2)$$

(c)  $total = 0 - \Theta(1)$   
 while  $n > 1$ :  $-\Theta(n)$   
      $total += n \% 2 - \Theta(1)$   $\left[ \right]$   
      $n = n // 2$   $\downarrow$   $\left[ \Theta(\log_2 n) \right]$

$n \rightarrow \frac{n}{2} \rightarrow \frac{n}{4} \rightarrow \frac{n}{8} \rightarrow \dots : size$

$$\therefore T_{worst}(n) = \Theta(\log_2 n)$$

$$2(a) \quad 3n^4 + 8n^3 - 3n = \Theta(n^4)$$

$$C_1 n^4 \leq 3n^4 + 8n^3 - 3n \leq C_2 n^4$$

Upper bound:

$$3n^4 + 8n^3 - 3n \leq 3n^4 + 8n^4 + 3n^4$$

$$3n^4 + 8n^3 - 3n \leq 14n^4 \quad \text{for } n \geq 0$$

lower bound:

$$3n^4 \leq 3n^4 + 8n^3 - 3n \quad \text{for } n \geq 0$$

$$3n^4 \leq 3n^4 + 8n^3 - 3n \leq 14n^4 \quad (n \geq 0)$$

$$\therefore C_1 = 3, C_2 = 14, n_0 = 0$$

$$(b) \quad \sqrt{17n^2 + 4n - 7} = \Theta(n)$$

$$C_1 n \leq \sqrt{17n^2 + 4n - 7} \leq C_2 n$$

$$C_1^2 n^2 \leq 17n^2 + 4n - 7 \leq C_2^2 n^2$$

Upper bound:

$$17n^2 + 4n - 7 \leq 17n^2 + 4n^2 + 7n^2 = 28n^2 \quad \text{for } n \geq 0$$

lower bound:

$$17n^2 - 4n^2 - 7n^2 \leq 17n^2 + 4n - 7$$

$$6n^2 \leq 17n^2 + 4n - 7 \quad \text{for } n \geq 1$$

$$6n^2 \leq 17n^2 + 4n - 7 \leq 28n^2 \quad \begin{matrix} C_1^2 = 6 \\ C_2^2 = 28 \end{matrix}$$

$$C_1 = \sqrt{6}$$

$$C_2 = \sqrt{28}$$

$$\therefore C_1 = \sqrt{6}, C_2 = \sqrt{28}, n_0 = 1$$

$$(c) \quad f(n) = O(g(n)) \Rightarrow f(n) \text{ is upper bound.}$$

$$f(n) = 4n \quad g(n) = 2n$$

$$g(n) = O(h(n)) \Rightarrow g(n) \text{ is upper bound.}$$

$$g(n) = 2n \quad h(n) = n$$

$$\therefore n \leq 4n \Rightarrow h(n) \leq f(n) \quad \text{for } n \geq 0$$

$$f(n) = O(h(n))$$

3. def find\_primes(n):

prime\_list = []  $- \Theta(1)$

for i in range(1, n+1):  $- \Theta(n)$

if is\_prime(i):  $- \Theta(\sqrt{n})$

prime\_list.append(i)  $- \Theta(1)$  Amortized

return prime\_list  $- \Theta(1)$

$\Theta(n\sqrt{n})$

def is\_prime(n):

if n == 1:  $- \Theta(1)$

return False

for k in range(2, int(math.sqrt(n)+1)):

if n % k == 0:  $- \Theta(1)$

return False

return True

$\Theta(\sqrt{n})$

$$\therefore T(n) = \Theta(n\sqrt{n})$$

4. (a)

def reverse\_l(lst):

rev\_list = []  $- \Theta(1)$

for i in range(len(lst)):  $\rightarrow \Theta(n)$

rev\_list.insert(0, lst[i])  $\rightarrow \Theta(n)$

return rev\_list  $- \Theta(1)$

$\Theta(n^2)$

$$1 + 2 + 3 + 4 + 5 + \dots + n = \frac{n(n+1)}{2} = \Theta(n^2)$$

$$\therefore T_{\text{worst}}(n) = \Theta(n^2)$$

4. (b)

def reverse2(lst):

rev\_lst = []  $\rightarrow O(1)$

for i in range(~~len~~ len(lst)-1, -1, -1):

rev\_lst.append(lst[i])  $\rightarrow O(1)$  Amortized  $\rightarrow O(n)$

return rev\_lst  $\rightarrow O(1)$

$$1 + 1 + 1 + 1 + 1 + \dots + 1 = O(n)$$

$$\therefore T_{\text{worst}}(n) = O(n)$$