

Lec5_Scheduling and Binary Search Trees

Scheduling Problem

- Runway Reservation System
 - Airport with single runway
 - “Reservations” for future landing
 - When plane lands, it is removed from set of pending events
 - Reserve req specify “requested landing time” t
 - Add t to the set if no other landings are scheduled within k mins either way

Assume that k can vary

⇒ Goal : Run this System efficiently in $O(\lg n)$ time

- How to solve
 - Sorted List
 - Appending and sorting : $\Theta(n \lg n)$
 - Insertion : $\Theta(n)$
 - Check : $O(1)$
 - Sorted array
 - Find position to insert : $\Theta(\lg n)$
 - Insertion : $\Theta(n)$
 - Unsorted list/array
 - Check : $O(n)$
 - Min-Heap
 - Insertion : $O(\lg n)$
 - Check : $O(n)$
 - Dictionary
 - Insertion : $O(1)$
 - Check : $\Omega(n)$

Binary Search Trees (BST)

- Properties
 - Each node x in the binary tree has a key $key(x)$. Nodes other than the root have a parent $p(x)$. Nodes may have a left child $left(x)$ and/or a right child $right(x)$. These are pointers unlike in a heap
 - The invariant is : for any node x , for all nodes y in the left subtree of x , $key(y) \leq key(x)$. For all nodes y in the right subtree of x $key(y) \geq key(x)$
- Functions
 - Insertion : $insert(val)$

- Finding a value in the BST if it exists : find(val)
- Finding the minimum element in a BST : findmin()
- Complexity
 - All operations are $O(h) \rightarrow h$: height of the BST
 - Function
 - Next-larger

```
if right child not NIL, return minimum(right)
else y = parent(x)
```

```
while y not NIL && x = right(y)
    x = y; y = parent(y)
```

```
return (y)
```

- Height h of the tree should be $O(\lg n)$
→ **Balancing is needed!!**