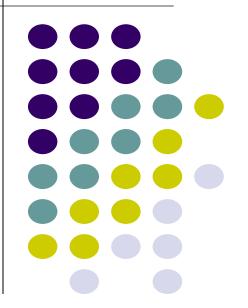
The Basics of UNIX/Linux

11-1. Arrays and Pointer. Part 1

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Lecture Outline

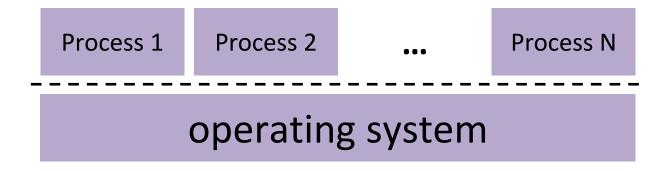


- C's Memory Model (refresher)
- Pointers (refresher)
- Arrays

OS and Processes



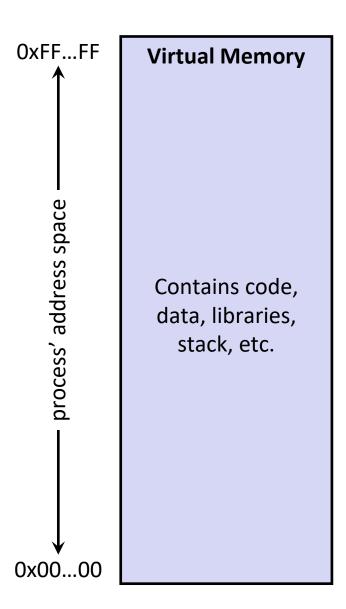
- The OS lets you run multiple applications at once
 - An application runs within an OS "process"
 - The OS timeslices each CPU between runnable processes
 - This happens very quickly: ~100 times per second



Processes and Virtual Memory



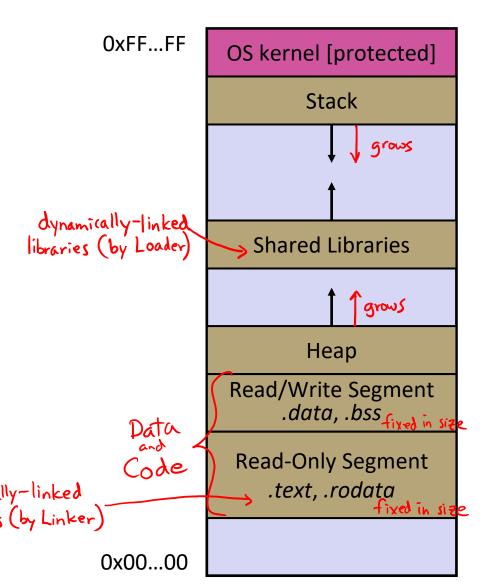
- The OS gives each process the illusion of its own private memory
 - Called the process' address space
 - Contains the process' virtual memory, visible only to it (via translation)
 - 2⁶⁴ bytes on a 64-bit machine



Loading



- When the OS loads a program it:
 - 1) Creates an address space
 - 2) Inspects the executable file to see what's in it
 - (Lazily) copies regions of the file into the right place in the address space
 - Does any final linking, relocation, or other needed preparation



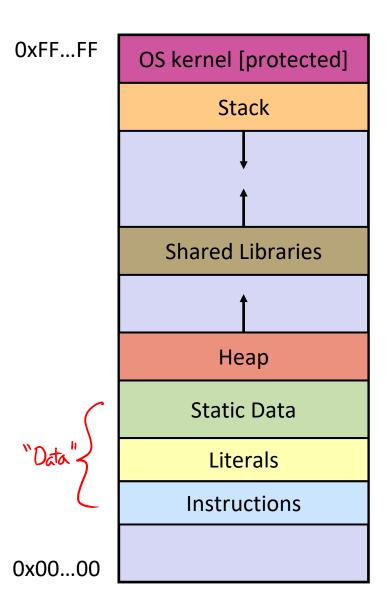
Memory Management



- Local variables on the Stack
 - Allocated and freed via calling conventions (push, pop, mov)

- Global and static variables in Data
 - Allocated/freed when the process starts/exits

- Dynamically-allocated data on the Heap
 - malloc() to request; free() to free, otherwise memory leak



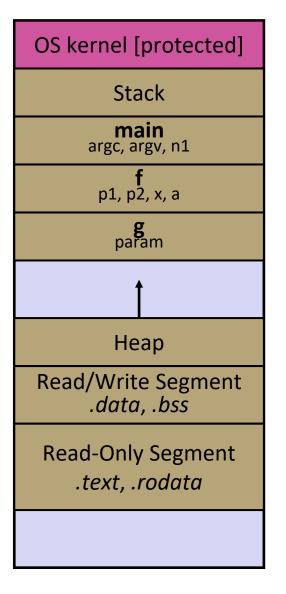
Stack in Action (1/4)

Note: arrow points to *next* instruction to be executed (like in gdb).



stack.c

```
#include <stdint.h>
int f(int, int);
int g(int);
int main(int argc, char** argv) {
  int n1 = f(3, -5);
  n1 = q(n1);
int f(int p1, int p2) {
  int x;
  int a[3];
  x = g(a[2]);
  return x;
int g(int param) {
  return param * 2;
```



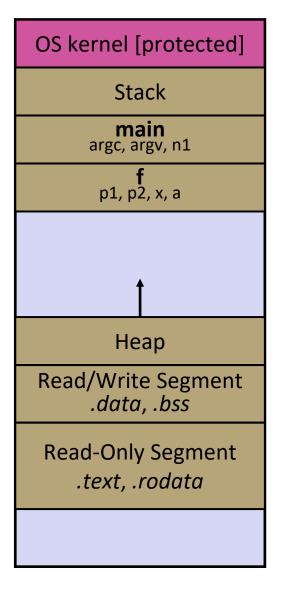
Stack in Action (2/4)

Note: arrow points to *next* instruction to be executed (like in gdb).



stack.c

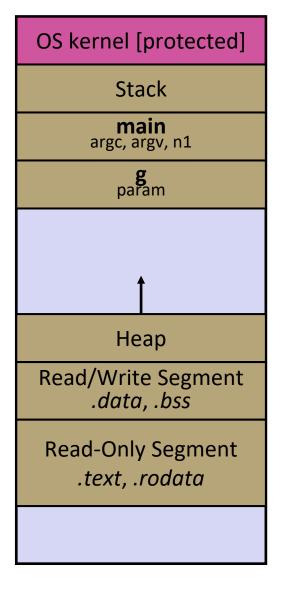
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  return x;
int g(int param) {
  return param * 2;
```



Stack in Action (3/4)

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  return param * 2;
```



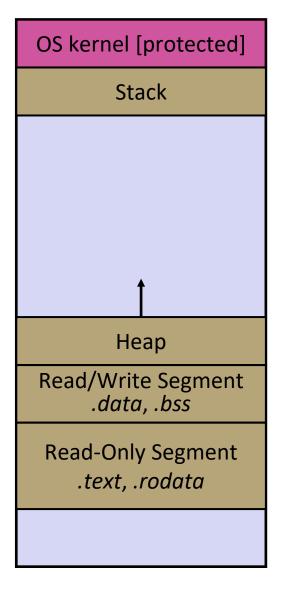
Stack in Action (4/4)

Note: arrow points to *next* instruction to be executed (like in gdb).



stack.c

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#include <stdint.h>
int f(int, int);
int g(int);
int main(int argc, char** argv) {
  int n1 = f(3, -5);
  n1 = q(n1);
int f(int p1, int p2) {
  int x;
  int a[3];
  x = g(a[2]);
  return x;
int g(int param) {
  return param * 2;
```



Lecture Outlinec



- C's Memory Model (refresher)
- Pointers (refresher)
- Arrays

Pointers



- Variables that store addresses
 - It points to somewhere in the process' virtual address space
 - &foo produces the virtual address of foo
- Generic definition:

```
type* name;
```

type *name;

Recommended to not define multiple pointers on same line:

```
int *p1, p2;
```

not the same as

Instead, use:

- Dereference a pointer using the unary * operator
 - Access the memory referred to by a pointer

Pointer Example



pointy.c

```
#include <stdio.h>
#include <stdint.h>
int main(int argc, char** argv) {
 int x = 351;
 int* p;  // p is a pointer to a int
 p = &x; // p now contains the addr of x
 printf("&x is %p\n", &x);
 printf(" p is %p\n", p);
 printf(" x is %d\n", x);
  *p = 333; // change value of x
 printf(" x is %d\n", x);
  return 0;
```

Something Curious



What happens if we run pointy.c several times?

Run 1: bash\$./pointy &x is 0x7ffff9e28524 p is 0x7ffff9e28524 x is 351 x is 333

```
Run 3: bash$ ./pointy
&x is 0x7fffe7b14644

p is 0x7fffe7b14644

x is 351

x is 333
```

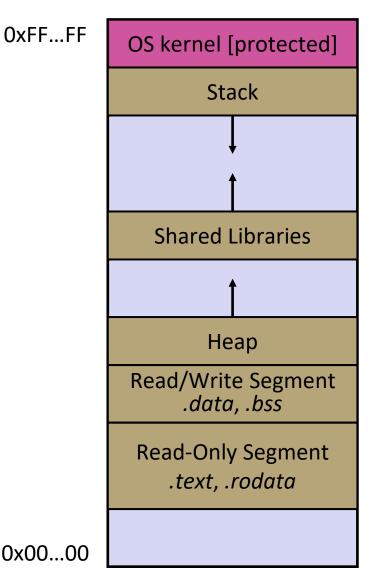
```
Run 2: bash$ ./pointy
&x is 0x7fffe847be34
p is 0x7fffe847be34
x is 351
x is 333
```

```
Run 4: bash$ ./pointy
&x is 0x7fffff0dfe54
p is 0x7fffff0dfe54
x is 351
x is 333
```

Address Space Layout Randomization



- Linux uses address space layout randomization (ASLR) for added security
 - Randomizes:
 - Base of stack
 - Shared library (mmap) location
 - Makes Stack-based buffer overflow attacks tougher
 - Makes debugging tougher
 - Can be disabled (gdb does this by default); Google if curious



0x00...00

Lecture Outline

- C's Memory Model (refresher)
- Pointers (refresher)
- Arrays

Arrays



• Definition:

```
type name[size]
```

- Allocates size*sizeof (type) bytes of contiguous memory
- Normal usage is a compile-time constant for size (e.g. int scores[175];)
- Initially, array values are "garbage"

Size of an array



- Size of an array
 - Not stored anywhere array does not know its own size!
 - sizeof (array) only works in variable scope of array definition
 - Recent versions of C allow for variable-length arrays
 - Uncommon and can be considered bad practice [we won't use]
 - https://gcc.gnu.org/onlinedocs/gcc/Variable-Length.html

```
int n = 175;
int scores[n]; // OK in C99
```

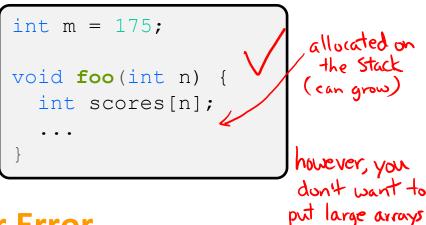
Challenge Question





• The code snippets both use a variable-length array. What will happen when we compile with C99?

```
allocated in Static Data (int m = 175; int scores[m]; (can't change size) void foo(int n) {
...
}
```



A. Compiler Error

Compiler Error

B. Compiler Error

No Error

C. No Error

Compiler Error

D. No Error

No Error

E. We're lost...

Using Arrays (1/2)



Initialization

```
type name[size] = {val0,...,valN};
```

- { } initialization can *only* be used at time of definition
- If no size supplied, infers from length of array initializer

Using Arrays (2/2)



- Array name used as identifier for "collection of data"
 - name [index] specifies an element of the array and can be used as an assignment target or as a value in an expression
 - Array name (by itself) produces the address of the start of the array
 - Cannot be assigned to / changed

```
int primes[6] = {2, 3, 5, 6, 11, 13};
primes[3] = 7;
primes[100] = 0; // memory smash!
```

Multi-dimensional Arrays



Generic 2D format:

```
type name[rows][cols] ={{values},...,{values}};
```

- Still allocates a single, contiguous chunk of memory
- C is row-major

Parameters: reference vs. value (1/2)



- There are two fundamental parameter-passing schemes in programming languages
- Call-by-value / Pass by value
 - Parameter is a local variable initialized when the function is called and gets a copy of the calling argument
 - manipulating the parameter only changes copy, not the calling argument
 - C, Java, C++ primitives

Parameters: reference vs. value (2/2)



- Call-by-reference / Pass-by-reference
 - Parameter is an alias for the supplied argument
 - manipulating the parameter manipulates the calling argument
 - C++ references (you'll learn later)

Arrays as Parameters



- It's tricky to use arrays as parameters
 - What happens when you use an array name as an argument?
 - Arrays do not know their own size

```
int sumAll(int a[]); // prototype

int main(int argc, char** argv) {
   int numbers[] = {9, 8, 1, 9, 5};
   int sum = sumAll(numbers);
   return 0;
}

int sumAll(int a[]) {
   int i, sum = 0;
   for (i = 0; i < ...???
}</pre>
```

Solution 1: Declare Array Size



```
int sumAll(int a[5]); // prototype
int main(int argc, char** argv) {
  int numbers[] = \{9, 8, 1, 9, 5\};
  int sum = sumAll(numbers);
 printf("sum is: %d\n", sum);
  return 0;
int sumAll(int a[5]) {
  int i, sum = 0;
  for (i = 0; i < 5; i++) {
    sum += a[i];
  return sum;
```

Problem: loss of generality/flexibility!

Solution 2: Pass Size as Parameter



This is the standard idiom in C programs

```
int sumAll(int a[], int size); //
prototype
int main(int argc, char** argv) {
  int numbers[] = \{9, 8, 1, 9, 5\};
  int sum = sumAll(numbers, 5);
  printf("sum is: %d\n", sum);
 return 0;
int sumAll(int a[], int size) {
  int i, sum = 0;
  for (i = 0; i < size; i++) {</pre>
    sum += a[i];
  return sum;
```

Returning an Array



- Local variables, including arrays, are allocated on the Stack
 - They "disappear" when a function returns!
 - Can't safely return local arrays from functions
 - Can't return an array as a return value why not?

Solution: Output Parameter



- Create the "returned" array in the caller
 - Pass it as an output parameter to copyarray ()
 - A pointer parameter that allows the callee to leave values for the caller to use
 - Works because arrays are "passed" as pointers
 - "Feels" like call-by-reference, but it's not

```
void copyArray(int src[], int dst[], int size) {
  int i;

for (i = 0; i < size; i++) {
   dst[i] = src[i];
  }
}</pre>
```

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Arrays: Call-By-Value or Call-By-Reference?



Technical answer:

- a T[] array parameter is "promoted" to a pointer of type T^* , and the *pointer* is passed by value
 - So it acts like a call-by-reference array (if callee changes the array parameter elements it changes the caller's array)
 - But it's really a call-by-value pointer (the callee can change the pointer parameter to point to something else(!))

```
void copyArray(int src[], int dst[], int size) {
  int i;
  dst = src;
  for (i = 0; i < size; i++) {
    dst[i] = src[i]; // copies source array to itself!
  }
}</pre>
```

Output Parameters



- Output parameters are common in library functions
 - long int strtol(char* str, char** endptr, int base);
 - int sscanf(char* str, char* format, ...);

```
int num, i;
char* pEnd, str1 = "333 rocks";
char str2[10];

// converts "333 rocks" into long -- pEnd is
conversion end
num = (int) strtol(str1, &pEnd, 10);

// reads string into arguments based on format string
num = sscanf("3 blind mice", "%d %s", &i, str2);
```

outparam.c

Q&A



