CSE215: Lecture 07
Foundations of Computer Science

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Things

- Attendance check on a weekly basis. Usually done on a day where content is not heavy

- Solutions to online classes will be provided by the end of the lecture day

- Homework grades will be available before the recitation after the due week

- Still seeking a wireless mic. Waiting for IT Department.

Previous lectures

Truth table

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

Equivalence laws

Laws	Formula	Formula
Commutative laws	$p \wedge q \equiv q \wedge p$	$p \vee q \equiv q \vee p$
Associative laws	$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	$(p \vee q) \vee r \equiv p \vee (q \vee r)$
Distributive laws	$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$

Propositional logic: A formal language to express facts and argue about them

Valid arguments

Premise$_1$
Premise$_2$
$\vdots$
Premise$_m$
$\therefore$ Conclusion

Inference

Name	Rule	Name	Rule	
Modus Ponens	$p \rightarrow q$ p $\therefore q$	Elimination	$p \vee q$ $\sim q$ $\therefore p$	$p \vee q$ $\sim p$ $\therefore q$
Modus Tollens	$p \rightarrow q$ $\sim q$ $\therefore \sim p$	Transitivity	$p \rightarrow q$ $q \rightarrow r$ $\therefore p \rightarrow r$	

Today

- Study propositions with quantifiers: for all integers a, b, and c, (a + b) +c = a + (b + c)?

- Predicates & quantifiers

- Negation

Predicate & Quantifiers

Predicate

- A **propositional function** or **predicate** is a sentence that contains **one or more variables**
- A predicate is neither true nor false
- A predicate becomes a proposition when the variables are substituted with specific values
- The **domain** of a predicate variable is the set of all values that may be substituted for the variable

Examples

Symbol	Predicate	Domain	Propositions
$p(x)$	$x > 5$	$x \in \mathbb{R}$	$p(6), p(-3.6), p(0), \ldots$
$p(x,y)$	$x + y$ is odd	$x \in \mathbb{Z}, y \in \mathbb{Z}$	$p(4,5), p(-4,-4), \ldots$

Universal quantifier ($\forall$)

Definition

- Let $p(x)$ be a predicate and D be the domain of x
- A universal statement is a statement of the form

$$\forall x \in D, p(x)$$

- Forms:
  - "$p(x)$ is true for all values of x"
  - "For all x, $p(x)$"
  - "For each x, $p(x)$"
  - "For every x, $p(x)$"
  - "Given any x, $p(x)$"

Existential quantifier ($\exists$)

Definition

- Let $p(x)$ be a predicate and D be the domain of x
- An existential statement is a statement of the form

$$\exists x \in D, p(x)$$

- Forms:
  - "There exists an x such that $p(x)$"
  - "For some x, $p(x)$"
  - "We can find an x, such that $p(x)$"
  - "There is some x such that $p(x)$"
  - "There is at least one x such that $p(x)$"

Examples

Universal st.s	Domain
$\exists x \in D, x^2 \geq x$	$D = \{1, 2, 3\}$
$\exists x \in \mathbb{R}, x^2 \geq x$	$\mathbb{R}$

Universal conditional statement $(\forall, \rightarrow)$

Definition

- A **universal conditional statement** is of the form

$$\forall x, \text{ if } p(x) \text{ then } q(x)$$

Examples

- $\forall x \in \mathbb{R}$, if $x > 2$ then $x^2 > 4$
- $\forall$ real number x, if x is an integer then x is rational
  $\forall$ integer x, x is rational
- $\forall x$, if x is a square then x is a rectangle
  $\forall$ square x, x is a rectangle

Can be extended to **existential conditional statement** $(\exists, \rightarrow)$

Implicit quantifiers

Examples

- If a number is an integer, then it is a rational number
  Implicit meaning: $\forall$ number x, if x is an integer, x is rational
- The number 10 can be written as a sum of two prime numbers
  Implicit meaning: $\exists$ prime numbers p and q such that $10 = p+q$
- If $x > 2$, then $x^2 > 4$
  Implicit meaning: $\forall$ real x, if $x > 2$, then $x^2 > 4$

Definition

- Let $p(x)$ and $q(x)$ be predicates and D be the common domain of x. Then implicit quant. symbols $\Rightarrow, \Leftrightarrow$ are defined as:

$$p(x) \Rightarrow q(x) \equiv \forall x, p(x) \to q(x)$$
$$p(x) \Leftrightarrow q(x) \equiv \forall x, p(x) \leftrightarrow q(x)$$

Propositions with multiple quantifiers

$$\forall x \in D, \exists y \in E, \text{such that } p(x, y)$$

$$\exists x \in D, \forall y \in E, \text{such that } p(x, y)$$

$$\forall x \in D, \forall y \in E, \text{such that } p(x, y)$$

$$\exists x \in D, \exists y \in E, \text{such that } p(x, y)$$

Note:
the order of quantifiers matter

- Every lock has a key

- For any lock L, there exists a key K, such that K can unlock L.

- There is a key for every lock

- There exists a key K, such that for any lock L, K can unlock L.

Exercise:
translate to formal logic

Every nonzero real number has a reciprocal

The reciprocal of 4 is 1/4 (namely 0.25)

- for any nonzero real number r, r has a reciprocal

- for any nonzero real number r, there exists a real number s, such that r * s = 1

Exercise:
translate to formal logic

- "There is a person supervising every detail of the production process."

- $\exists$ person p such that $\forall$ detail d, p supervises d

Negation

Negation of quantified statements ($\sim$)

Definition

- Formally,

$$\sim (\forall x \in D, p(x)) \equiv \exists x \in D, \sim p(x)$$
$$\sim (\exists x \in D, p(x)) \equiv \forall x \in D, \sim p(x)$$

- Negation of a **universal** statement ("all are") is logically equivalent to an **existential** statement ("there is at least one that is not")

  Negation of an **existential** statement ("some are") is logically equivalent to a **universal** statement ("all are not")

Negation of quantified statements ($\sim$)

Examples

- All mathematicians wear glasses
  Negation (incorrect): No mathematician wears glasses

  Negation (correct): There is at least one mathematician who does not wear glasses
- Some snowflakes are the same
  Negation (incorrect):: Some snowflakes are different
  Negation (correct):: All snowflakes are different

Negation of quantified statements ($\sim$)

Examples

- $\forall$ primes p, p is odd
  Negation: $\exists$ primes p, p is even
- $\exists$ triangle T, sum of angles of T equals $200°$
  $\forall$ triangles T, sum of angles of T does not equal $200°$
- No politicians are honest
  Formal statement: $\forall$ politicians x, x is not honest
  Formal negation: $\exists$ politician x, x is honest
  Informal negation: Some politicians are honest
- 1357 is not divisible by any integer between 1 and 37
  Formal statement: $\forall n \in [1, 37]$, 1357 is not divisible by n
  Formal negation: $\exists n \in [1, 37]$, 1357 is divisible by n
  Informal negation: 1357 is divisible by some integer between 1 and 37

Negation of universal conditional statements

Definition

- Formally,

$$\sim (\forall x, p(x) \rightarrow q(x)) \equiv \exists x, \sim (p(x) \rightarrow q(x))$$
$$\equiv \exists x, (p(x) \wedge \sim q(x))$$

Examples

- $\forall$ real x, if $x > 10$, then $x^2 > 100$.
  Negation: $\exists$ real x such that $x > 10$ and $x^2 \leq 100$.
- If a computer program has more than 100,000 lines, then it contains a bug.
  Negation: There is at least one computer program that has more than 100,000 lines and does not contain a bug.

Negation of propositions with multiple quantifiers

$\sim (\forall x \text{ in } D, \exists y \text{ in } E \text{ such that } P(x,y))$

$\equiv \exists x \text{ in } D \text{ such that } \sim (\exists y \text{ in } E \text{ such that } P(x,y))$

$\equiv \exists x \text{ in } D \text{ such that } \forall y \text{ in } E, \sim P(x,y)$

$\sim (\exists x \text{ in } D \text{ such that } \forall y \text{ in } E, P(x,y))$

$\equiv \forall x \text{ in } D, \sim (\forall y \text{ in } E, P(x,y))$

$\equiv \forall x \text{ in } D, \exists y \text{ in } E \text{ such that } \sim P(x,y)$

Break if time is good;

Adapted exam exercises

To finish by 1h50pm

Exercise: 2021 midterm-1

Problem 3. **[10 points]**

Give negations of the following statements. Reasoning is not required.

(a) [1 point] $p \wedge q$

(b) [1 point] $p \vee q$

(c) [1 point] $p \oplus q$

(d) [1 point] $p \rightarrow q$

(e) [1 point] $p \leftrightarrow q$

(f) [1 point] $\forall x, \forall y$ such that $p(x, y)$

(g) [1 point] $\forall x, \exists y$ such that $p(x, y)$

(h) [1 point] $\exists x, \forall y$ such that $p(x, y)$

(i) [1 point] $\exists x, \exists y$ such that $p(x, y)$

(j) [1 point] $\exists x, \forall y, \exists z$ such that $p(x, y, z)$

Solution

(a) [1 point] $p \wedge q$

(b) [1 point] $p \vee q$

(c) [1 point] $p \oplus q$

(d) [1 point] $p \rightarrow q$

(e) [1 point] $p \leftrightarrow q$

(f) [1 point] $\forall x, \forall y$ such that $p(x, y)$

(g) [1 point] $\forall x, \exists y$ such that $p(x, y)$

(h) [1 point] $\exists x, \forall y$ such that $p(x, y)$

(i) [1 point] $\exists x, \exists y$ such that $p(x, y)$

(j) [1 point] $\exists x, \forall y, \exists z$ such that $p(x, y, z)$

- a. ~p $\vee$ ~q

- b. ~p $\wedge$ ~q

- c. (~p $\vee$q) $\wedge$ (p $\vee$ ~q), namely, p<->q

- d. p $\wedge$ ~q

- e. (p $\wedge$~q) $\vee$ (q $\wedge$ ~p), namely p xor q

- f-j. change \forall to \exists, \exists to \forall, p to ~p

Final 2021

Prove that $n^2 + 9n + 27$ is odd for all natural numbers n. You can use any proof technique.

- Students: Express the propositions we need to prove here

- Instructor: Prove

Solution

- for all natural number n, n^2 + 9n + 27 is odd.

Final 2021

Prove that if $n^2 + 8n + 20$ is odd, then n is odd for natural numbers n.

- Students: Express the propositions we need to prove here

- Instructor: Prove

Solution

- For all natural number n, n^2 + 8n + 20 is odd -> n is odd

Final 2021

Prove using contradiction that the cube root of an irrational number is irrational.

- Students: Express the propositions we need to prove here

- Instructor: Prove

Solution

- For all irrational number i, cubic root of i is irrational

Final 2021

Use mathematical induction to prove the following identities.

(a) [5 points] For all natural numbers n,

$$1^2 \times 2 + 2^2 \times 3 + 3^2 \times 4 + \cdots + n^2 \times (n+1) = \frac{n(n+1)(n+2)(3n+1)}{12}$$

- Students: Express the propositions we need to prove here

- Instructor: Prove

Summary

- Predicates

- Universal and existential quantifiers

- Negations

Thank you!