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ELEC 5200/6200 (Fall 2021)
Homework 2
Assigned 09/01/21, due 09/08/21

Note: This homework will be easier if you have completed the assigned reading for Chapter 2. You also might find the RISC-V reference data sheet useful (located on the green insert inside your textbook).

Question 1: Show how the value 0x EFBEADDA would be arranged in the memory of a big-endian machine. Assume that the data is stored starting at address 0. Show your answer in hexadecimal. (2 pts)

Address	0x03	0x02	0x01	0x00
Byte	0xDA	0xAD	0xBE	0xEF

Question 2: Show how the binary value 1111 1010 1100 1110 1011 1110 1010 1101 would be arranged in the memory of a little-endian machine. Assume that the data is stored starting at address 0. Show your answer in hexadecimal. (2 pts)

Address	0x03	0x02	0x01	0x00
Byte	0xFA	0xCE	0xBE	0xAD

Question 3: For the following C statement, assume that variable **f** is already stored in register x7, **g** in register x6, and **h** in x5.

```
f = g + (h - 11)
```

a) Write the corresponding RISC-V assembly code. (4 pts)

```
ADD x5, x5, -11
ADD x7, x6, x5
```

b) Write the machine language for the assembly code you came up with in part a. (4 pts)
Note: write the binary in groups of 4 for legibility. Ex: 0000 1111 0000 1111

1. 1111 1111 0101 0010 1000 0010 1001 0011

2. 0000 0000 0101 0011 0000 0011 1011 0011

Question 4: Consider the following assembly code written using the RISC-V instruction set:

```
        addi   x29, x0, 100
        add    x18, x0, x0
        addi   x28, x18, 1
        add    x19, x18, x28
        sw     x18, 0(x11)
        sw     x19, 4(x11)
        addi   x11, x11, 8
        addi   x10, x0, 1

LOOP:   add    x10, x10, x28 // ++n
        lw     x5, -4(x11)    // x5 = result [n-1]
        lw     x6, -8(x11)   //x6 = result [n-2]
        add    x7, x5, x6
        sw     x7, 0(x11)     // x7 = result[n]
        addi   x11, x11, 4   // result++
        blt    x10, x29, LOOP
```

Translate the code above into C. Assume that the integer `n` is the loop variable and is located in register x10. Register x11 is initialized to the base address of the integer array named `Result`. Further assume that 32-bit integers are used. (5 pts)

```
int result[100], n;    // define array and loop var
result[0]=0;           // sd x18, 0(x11)
result[1] = 1          // sd x19, 8(x11)

for (n=2; n<100; n++){
  result [n] = result [n-1]+ result[n-2];
}
```

- Sets first element in array to 0 and second element to 1
- loop starts from 3rd element in array and goes to 100th
- for each element from 3rd on, the progam would set the element to the sum of the last 2 elements
- Stops once counter reaches 100

Question 5: If the current value of the PC is 0x01027B08 can you use a single RISC-V **jal** instruction to get to PC address 0x12345678? Explain your answer. (5 pts)

Jal x1, #0x 0x12345678 - 0x01027B08 = 0x11131D B70

Jal permits target to be within ± 1MByte only

So Jal cannot be performed here

Only 20 bits are allowed to specify displacement

Question 6: If the current value of the PC is 0x0FFFFFF0 can you use a single RISC-V **branch** instruction to get to PC address 0x10000080? Explain your answer. (5 pts)

Jalx1, #0x24

0x1 00000 80 - 0x 0FFFFFF0 = 0x 00000090

Should branch to 0x90 bytes ahead

So target address is 0x24 (0x90/4=0x24)

only