

Digital Communications

SSY125, Lecture 7

Analysis of Linear Modulations (Chapter 6)

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Slides prepared by Alexandre Graell i Amat

November 13, 2022



CHALMERS

Analysis of Linear Modulations



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Analysis of Linear Modulations

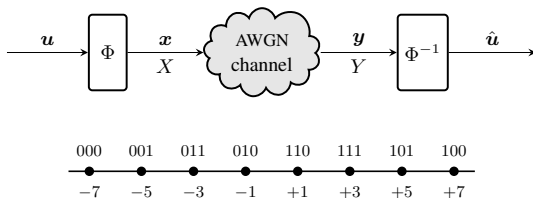


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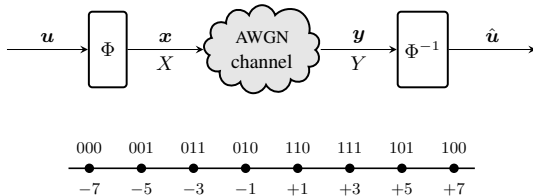


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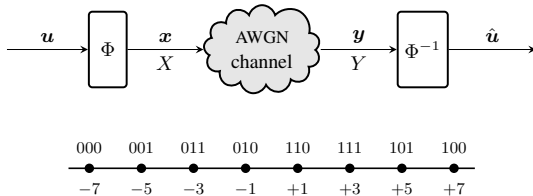


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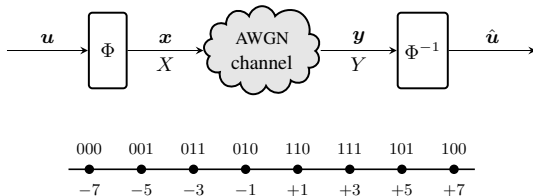
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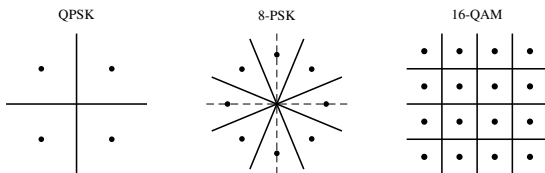
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The Union Bound

Given a number of events E_1, \dots, E_N

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The Q-function

For $Z \sim \mathcal{N}(\mu, \sigma^2)$,

$$\Pr(Z > \beta) =$$

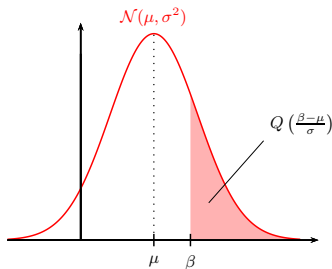
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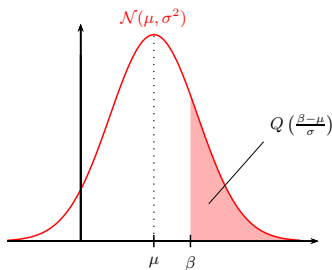


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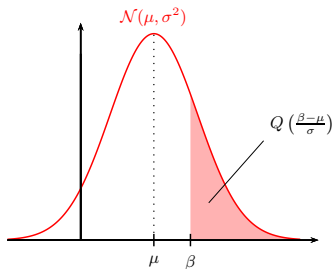
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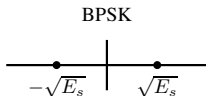
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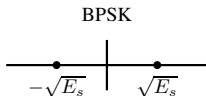
$$Q(a) + Q(b) + Q(c) \approx Q(a).$$

Symbol Error Probability of BPSK



- $\mathcal{X} = \{X_1, X_2\} \subset \mathbb{R}$, where $X_1 = -\sqrt{E_s}$ and $X_2 = \sqrt{E_s}$.

Symbol Error Probability of BPSK

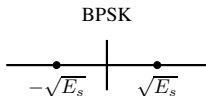


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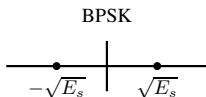
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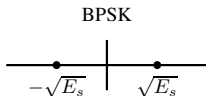
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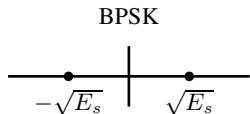
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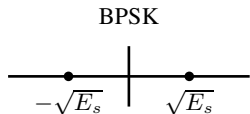
$$\begin{aligned} P_s^{\text{BPSK}} &= \sum_{x \in \mathcal{X}} \Pr(\hat{x} \neq x | x) P(x) \\ &= \Pr(X_2 | X_1) P(X_1) + \Pr(X_1 | X_2) P(X_2) \\ &= \frac{1}{2} \Pr(X_2 | X_1) + \frac{1}{2} \Pr(X_1 | X_2), \end{aligned}$$

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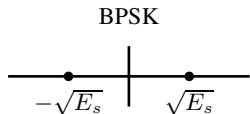
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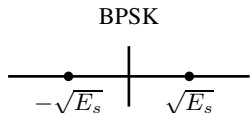
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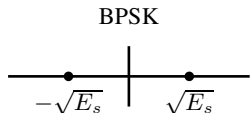
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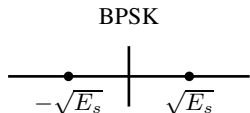
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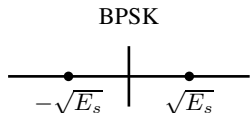
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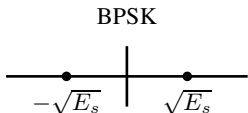
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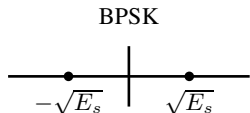
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Due to symmetry, $\Pr(\hat{x} = X_1 | x = X_2) = \Pr(\hat{x} = X_2 | x = X_1)$, and

$$P_s^{\text{BPSK}} = Q\left(\sqrt{\frac{2E_s}{N_0}}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right).$$

Symbol Error Probability and Euclidean Distance

For BPSK, $d_E(\mathbf{X}_1, \mathbf{X}_2) = |\mathbf{X}_1 - \mathbf{X}_2| = 2\sqrt{E_s}$, hence

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$$\begin{aligned}\Pr(\hat{x} = X_2 | x = X_1) &= \Pr\left(\tilde{Y} > \frac{d_E(X_1, X_2)}{2}\right) \Big|_{\tilde{Y} \sim \mathcal{N}(0, \sigma^2)} \\ &= Q\left(\frac{d_E(X_1, X_2)}{2\sigma}\right) = Q\left(\sqrt{\frac{d_E^2(X_1, X_2)}{2N_0}}\right).\end{aligned}$$

Symbol Error Probability and Euclidean Distance

For BPSK, $d_E(X_1, X_2) = |X_1 - X_2| = 2\sqrt{E_s}$, hence

$$P_s^{\text{BPSK}} = Q\left(\sqrt{\frac{d_E^2(X_1, X_2)}{2N_0}}\right).$$

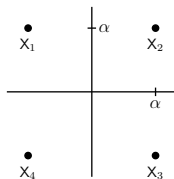
In the general case...

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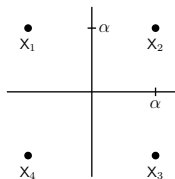
- $\Pr(\hat{x} = X_2 | x = X_1)$ depends on $d_E(X_1, X_2) \Rightarrow$ **Construct constellations with high distance between constellation points!**

Symbol Error Probability of 4-QAM



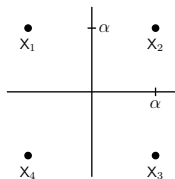
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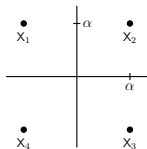


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$$\hat{x} = \begin{cases} X_1 & \text{if } y_I < 0 \text{ and } y_Q > 0 \\ X_2 & \text{if } y_I > 0 \text{ and } y_Q > 0 \\ X_3 & \text{if } y_I > 0 \text{ and } y_Q < 0 \\ X_4 & \text{if } y_I < 0 \text{ and } y_Q < 0 \end{cases},$$

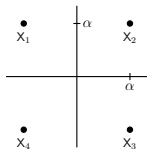
where y_I and y_Q are the in-phase and quadrature components of y .

Analysis of Linear Modulations



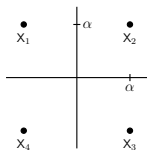
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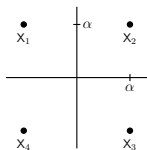
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Analysis of Linear Modulations



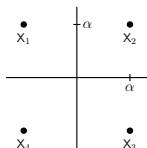
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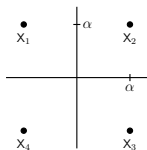
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If X_1 is transmitted, $Y_I \sim$ and $Y_Q \sim$.

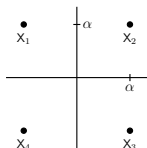
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Analysis of Linear Modulations

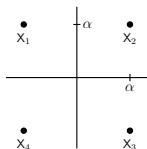


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Analysis of Linear Modulations

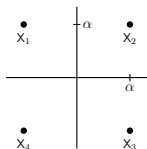


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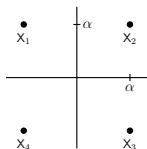


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Analysis of Linear Modulations

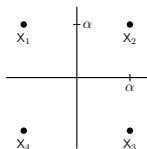


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By symmetry, $\Pr(Y_Q < 0 | X_1) = \Pr(Y_I > 0 | X_1)$, and

$$P_s^{4\text{QAM}} = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) - \left(Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\right)^2.$$

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For **general constellations**, the exact symbol error probability P_s is **hard to compute**

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Upper Bound on the Symbol Error Probability

Using

$$\Pr(\hat{x} = \mathbf{X}_j | x = \mathbf{X}_i) = Q\left(\sqrt{\frac{d_E^2(\mathbf{X}_i, \mathbf{X}_j)}{2N_0}}\right)$$

the symbol error probability of an M -ary constellation can be upperbounded as

$$P_s^{(M)} \leq \frac{1}{M} \sum_{i=1}^M \sum_{j \neq i} Q\left(\sqrt{\frac{d_E^2(\mathbf{X}_i, \mathbf{X}_j)}{2N_0}}\right).$$

Upper bound on P_s for 4-QAM

Example: 4-QAM

$$P_s^{4\text{QAM}} \leq \frac{1}{4} \sum_{i=1}^4 \sum_{j \neq i} Q \left(\sqrt{\frac{d_E^2(\mathbf{X}_i, \mathbf{X}_j)}{2N_0}} \right)$$

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Using $d_E^2(\mathbf{X}_1, \mathbf{X}_2) = d_E^2(\mathbf{X}_1, \mathbf{X}_4) = \|\alpha(-1+j) - \alpha(1+j)\|^2 = 4\alpha^2 = 2E_s$ and $d_E^2(\mathbf{X}_1, \mathbf{X}_3) = \|\alpha(-1+j) - \alpha(1-j)\|^2 = 8\alpha^2 = 4E_s$,

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$$\begin{aligned} P_s^{(M)} &\approx \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i) Q\left(\sqrt{\frac{d_{E,\min}^2(X_i)}{2N_0}}\right) \\ &\leq \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i) Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right) \\ &= \left(\frac{1}{M} \sum_{i=1}^M A_{\min}(X_i)\right) Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right) = \bar{A}_{\min} Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right), \end{aligned}$$

where $\bar{A}_{\min} = \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i)$ is the average number of nearest neighbors.

Nearest Neighbor Approximation

- Further approximate P_s considering only the **minimum Euclidean distance**,

$$d_{E,\min} = \min_{X_i} d_{E,\min}(X_i).$$

Then,

$$\begin{aligned} P_s^{(M)} &\approx \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i) Q\left(\sqrt{\frac{d_{E,\min}^2(X_i)}{2N_0}}\right) \\ &\leq \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i) Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right) \\ &= \left(\frac{1}{M} \sum_{i=1}^M A_{\min}(X_i)\right) Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right) = \bar{A}_{\min} Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right), \end{aligned}$$

where $\bar{A}_{\min} = \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i)$ is the **average number of nearest neighbors**.

- Requires only **knowledge of** \bar{A}_{\min} and $d_{E,\min}$!

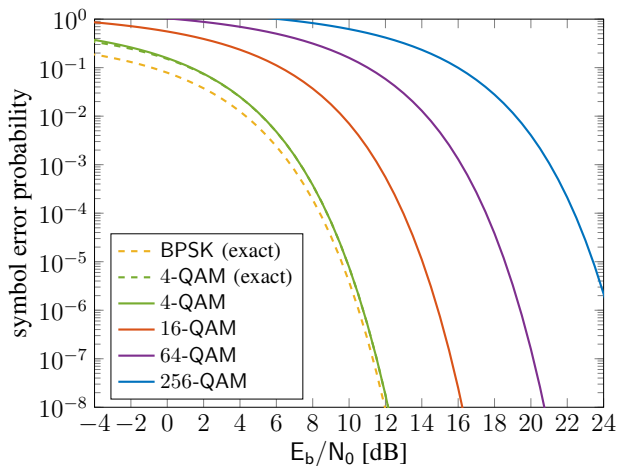
Nearest Neighbor Approximation

Nearest Neighbor Approximation for squared M -QAM

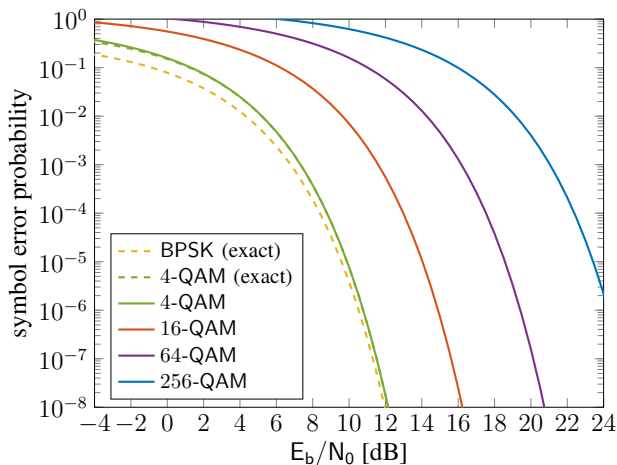
For squared M -QAM constellations, there are 4 constellation points with 2 neighbors, $4(\sqrt{M} - 2)$ points with 3 neighbors and the remaining points have 4 neighbors, hence $\bar{A}_{\min} = 4 - 4/\sqrt{M}$. Furthermore, $d_{E,\min} = \sqrt{\frac{6E_s}{M-1}}$. Hence,

$$\begin{aligned} P_s^{MQAM} &\approx \left(4 - \frac{4}{\sqrt{M}}\right) Q\left(\sqrt{\frac{3E_s}{(M-1)N_0}}\right) \\ &= \left(4 - \frac{4}{\sqrt{M}}\right) Q\left(\sqrt{\frac{3E_b \log M}{(M-1)N_0}}\right). \end{aligned}$$

Nearest Neighbor Approximation of P_s of QAM

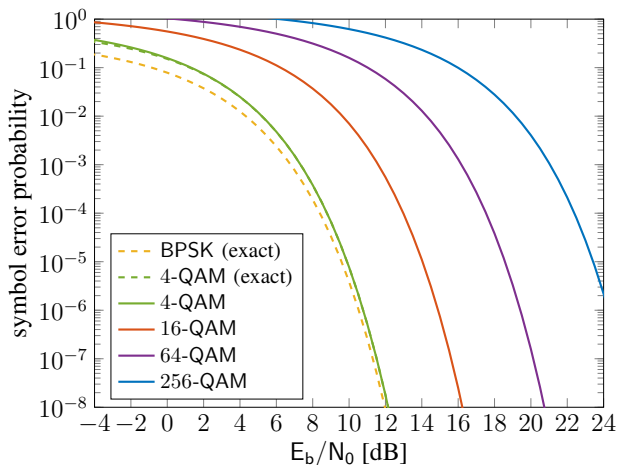


Nearest Neighbor Approximation of P_s of QAM



- The **power efficiency** decreases with M .

Nearest Neighbor Approximation of P_s of QAM



- The **power efficiency** decreases with M .
- The **spectral efficiency** increases with M .

Bit Error Probability of Linear Modulations

- The bit error probability depends on the binary labeling, i.e., how tuples of $m = \log M$ bits are mapped to the constellation symbols.

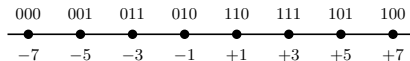
Bit Error Probability of Linear Modulations

- The bit error probability depends on the **binary labeling**, i.e., **how tuples of $m = \log M$ bits are mapped to the constellation symbols**.
- Let

$$\mathbf{L}(x) = (b_1(x), \dots, b_m(x))$$

be the m -bit labeling associated to constellation symbol $x \in \mathcal{X}$.

Bit Error Probability of Linear Modulations

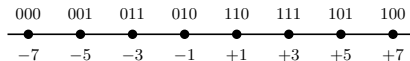


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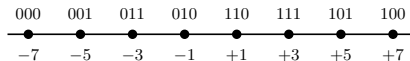
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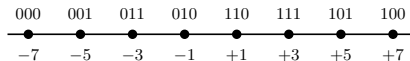
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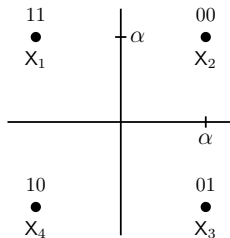
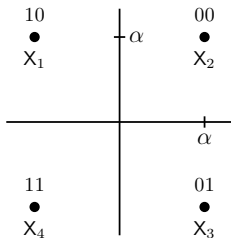
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Bit Error Probability of 4-QAM with Gray and Lexicographic Labeling



Gray Labeling

- e_{b_i} : event that bit i is decoded in error.

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$$p_1 = \Pr(e_{b_1})$$

Gray Labeling

- e_{b_i} : event that bit i is decoded in error.

$$p_1 = \Pr(e_{b_1}) = \sum_{x \in \mathcal{X}} \Pr(e_{b_1} | x) P(x)$$

Gray Labeling

- e_{b_i} : event that bit i is decoded in error.

$$p_1 = \Pr(e_{b_1}) = \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x)P(x) = \frac{1}{4} \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x)$$

Gray Labeling

- e_{b_i} : event that bit i is decoded in error.

$$\begin{aligned} p_1 &= \Pr(e_{b_1}) = \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x)P(x) = \frac{1}{4} \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x) \\ &= \frac{1}{4}(\Pr(e_{b_1}|X_1) + \Pr(e_{b_1}|X_2) + \Pr(e_{b_1}|X_3) + \Pr(e_{b_1}|X_4)) \end{aligned}$$

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- e_{b_i} : event that bit i is decoded in error.

$$\begin{aligned} p_1 &= \Pr(e_{b_1}) = \sum_{x \in \mathcal{X}} \Pr(e_{b_1} | x) P(x) = \frac{1}{4} \sum_{x \in \mathcal{X}} \Pr(e_{b_1} | x) \\ &= \frac{1}{4} (\Pr(e_{b_1} | X_1) + \Pr(e_{b_1} | X_2) + \Pr(e_{b_1} | X_3) + \Pr(e_{b_1} | X_4)) \\ &= \frac{1}{4} (\Pr(b_1(\hat{x}) = 0 | X_1) + \Pr(b_1(\hat{x}) = 1 | X_2) \\ &\quad + \Pr(b_1(\hat{x}) = 1 | X_3) + \Pr(b_1(\hat{x}) = 0 | X_4)), \end{aligned}$$

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- All terms are $Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$, thus $p_1 = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$.

Gray Labeling

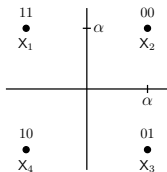
- e_{b_i} : event that bit i is decoded in error.

$$\begin{aligned} p_1 &= \Pr(e_{b_1}) = \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x)P(x) = \frac{1}{4} \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x) \\ &= \frac{1}{4}(\Pr(e_{b_1}|X_1) + \Pr(e_{b_1}|X_2) + \Pr(e_{b_1}|X_3) + \Pr(e_{b_1}|X_4)) \\ &= \frac{1}{4}(\Pr(b_1(\hat{x}) = 0|X_1) + \Pr(b_1(\hat{x}) = 1|X_2) \\ &\quad + \Pr(b_1(\hat{x}) = 1|X_3) + \Pr(b_1(\hat{x}) = 0|X_4)), \end{aligned}$$

- All terms are $Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$, thus $p_1 = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$.
- By symmetry, $p_2 = p_1$, hence,

$$P_b^{4\text{QAM-Gray}} = \frac{1}{2}(p_1 + p_2) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right).$$

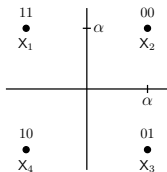
Lexicographic Labeling



$$p_1 = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

$$p_2 =$$

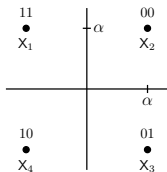
Lexicographic Labeling



$$p_1 = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

$$p_2 = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\left(1 - Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\right).$$

Lexicographic Labeling



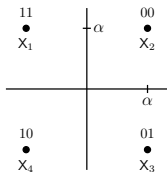
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Thus,

$$P_b^{4\text{QAM-Lex}} = \frac{1}{2}(p_1 + p_2) = \frac{3}{2}Q\left(\sqrt{\frac{2E_b}{N_0}}\right) - \left(Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\right)^2.$$

Lexicographic Labeling



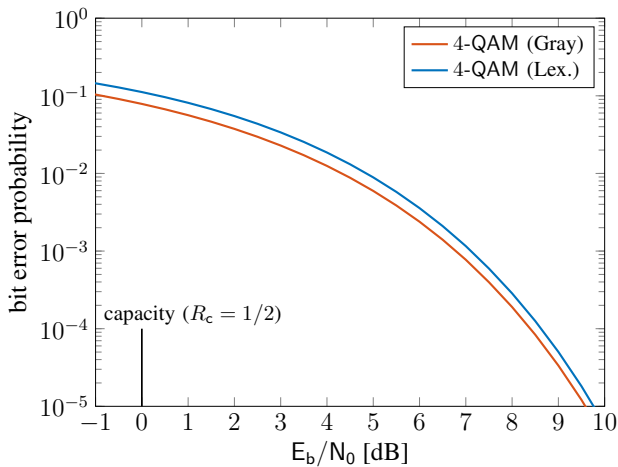
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Thus,

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- Due to the **lack of symmetry**, $p_1 \neq p_2$.

Bit Error Probability for 4-QAM



Bit Error Probability of M -ary Constellations

- For general M -ary constellations and arbitrary labelings, the computation of P_b is cumbersome \rightarrow **Nearest neighbor approximation**.