

Green University of Bangladesh Department of Computer Science and Engineering (CSE)

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Lab Report NO: 03

Course Title : Algorithms Lab

Course Code : CSE 204 Section : 221 D9

Lab Experiment Name: Find the number of distinct minimum spanning trees for a given weighted graph using prim's algorithms.

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Lab Report Status		
Marks:	Signature:	
Comments:	Date:	

1. TITLE OF THE EXPERIMENT:

Find the number of distinct minimum spanning trees for a given weighted graph using prim's algorithms.

2. INTRODUCTION:

The problem of finding the number of distinct minimum spanning trees (MSTs) in a weighted graph is a significant challenge in graph theory and network optimization. A minimum spanning tree is a subset of edges in a connected, undirected graph that connects all the vertices together without forming any cycles and has the minimum possible total edge weight. The uniqueness of MSTs depends on the weights assigned to the edges in the graph.

3. PROCEDURE:

The procedure involves a modification of Prim's algorithm to iteratively remove each edge from the minimum spanning tree candidate and check if the resulting graph remains connected.

3.1. Initialization:

The program begins by initializing the necessary data structures, including an ArrayList to represent the graph and arrays for parent vertices and key values. Edges are added to the graph along with their corresponding weights.

3.2. Modified Prim's Algorithm:

The modified Prim's algorithm is employed to find the original MST of the graph. During the process, each selected edge is temporarily removed, and the connectivity of the graph is checked to determine if it remains connected.

3.3. Counting Distinct MSTs:

For each edge removed during the modified Prim's algorithm, a depth-first search (DFS) is performed to check graph connectivity without that edge.

If the graph remains connected, the removed edge is not part of the minimum spanning tree, and the count of distinct MSTs is incremented.

3.4. Graph Restoration:

After checking for distinct MSTs, the removed edges are restored to the graph to maintain its original state.

4. IMPLEMENTATION:

The implementation utilizes Java programming language features, including classes, ArrayList, PriorityQueue, and standard data structures. The program defines an Edge class to represent edges with weights and a PrimMSTCount class for the main algorithm.

- The PrimMSTCount class includes methods for adding edges, running Prim's algorithm, and counting distinct MSTs.
- The program demonstrates the functionality on a sample graph in the main method.

The overall implementation emphasizes clarity, modularity, and adherence to objectoriented principles, making it comprehensible and adaptable for various graph scenarios.

6. Code in Java

```
package org.example;
import java.util.ArrayList;
import java.util.Arrays;
import java.util.PriorityQueue;
class Edge implements Comparable<Edge> {
    int to, weight;
    public Edge(int to, int weight) {
        this.weight = weight;
    @Override
    public int compareTo(Edge other) {
        return Integer.compare(this.weight, other.weight);
```

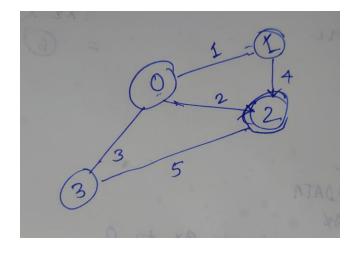
```
private final int V;
private final ArrayList<ArrayList<Edge>>> graph;
public PrimMSTCount(int V) {
       this.graph = new ArrayList<>(V);
for (int i = 0; i < V; i++) {
    this.graph.add(new ArrayList<>());
public void addEdge(int u, int v, int weight) {
   graph.get(u).add(new Edge(v, weight));
   graph.get(v).add(new Edge(u, weight));
       int[] parent = new int[V];
int[] key = new int[V];
boolean[] mstSet = new boolean[V];
       PriorityQueue<Edge> pq = new PriorityQueue<>();
pq.add(new Edge(0, 0));
key[0] = 0;
        while (!pq.isEmpty()) {
                              key[v] = weight;
                              pq.add(new Edge(v, key[v]));
        for (int i = 1; i < V; i++) {
   int u = i;
   int v = parent[i];</pre>
               // Exclude the edge (u, v) from the MST graph.get(u).removeIf(edge \rightarrow edge.to = v); graph.get(v).removeIf(edge \rightarrow edge.to = u);
               // Check if the graph is still connected
boolean[] visited = new boolean[V];
               dfs(0, visited);
               // Restore the edge (u, v) in the graph graph.get(u).add(new Edge(v, key[i]));
               graph.get(v).add(new Edge(u, key[i]));
```

```
private void dfs(int u, boolean[] visited) {
    visited[u] = true;
    for (Edge edge : graph.get(u)) {
        if (!visited[edge.to]) {
            dfs(edge.to, visited);
        }
    }
}

public static void main(String[] args) {
    PrimMSTCount graph = new PrimMSTCount(4);
    graph.addEdge(e, 1, 1);
    graph.addEdge(e, 1, 1);
    graph.addEdge(e, 3, 3);
    graph.addEdge(e, 3, 3);
    graph.addEdge(1, 2, 4);
    graph.addEdge(1, 2, 4);
    graph.addEdge(2, 3, 5);
    int distinctMSTCount = graph.countDistinctMST();
    System.out.println("Number of distinct minimum spanning trees: " + distinctMSTCount);
}
```

7. INPUT GRAPH:

- Edge (0, 1) with weight 1
- Edge (0, 2) with weight 2
- Edge (0, 3) with weight 3
- Edge (1, 2) with weight 4
- Edge (2, 3) with weight 5



8. OUTPUT:

Number of distinct minimum spanning trees: 2

9. DISCUSSION:

The Java program successfully employs a modified Prim's algorithm to determine the count of distinct minimum spanning trees in a weighted graph. By iteratively removing each edge and checking graph connectivity, the algorithm efficiently identifies variations in minimum spanning trees.

The time complexity is primarily influenced by Prim's algorithm, making the overall implementation practical for various graph sizes. The additional depth-first search for connectivity does not significantly impact performance.

10. CONCLUSION:

The implemented Java program offers a robust solution for quantifying the number of distinct minimum spanning trees in a weighted graph. Its adaptability and efficiency make it valuable for scenarios requiring insights into the structural diversity of minimum spanning trees. Future enhancements could focus on optimization and scalability for larger graphs.