

Theory Assignment 01
On
**Cryptography & Network Security (CSE
3035)**

Submitted by

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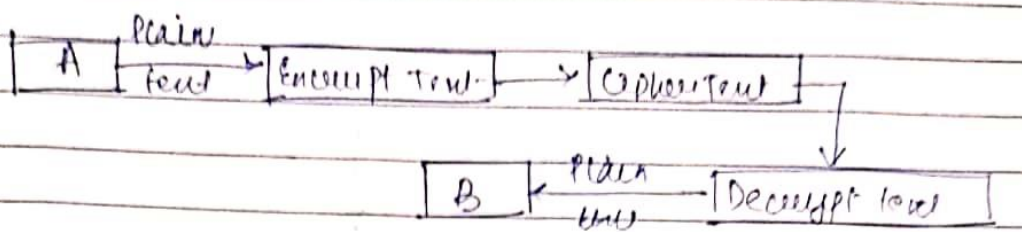


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(1) Define the following terms:

- (I) Cryptography (II) Cryptanalysis (III) Steganography (IV) network security
(V) computer security

i) Cryptography: The art of protecting info by transforming it into an unreadable format is called cryptography.



ii) Cryptanalysis: It is the process of studying cryptographic systems to look for weakness or ~~at~~ leak of information.

iii) Steganography: It means we are hiding the existence of data.

iv) Network security: It is any activity designed to protect the security and integrity of our network and data. It includes both hardware and software technologies.

v) Computer security: It is the protection of computer systems and information from harm, theft and unauthorized access.

(2) Formulate ceaser cipher for the cipher Text: PHHW PH DIWHU WKH WRJD SDUWB to identify the plain text with the default key $K=3$ and also give atleast three important characteristics of this problem that is enabled to bruteforce cryptanalysis.

2. From cipher to plaintext we will go for decryption.

$$P = D(K, C) = (C - K) \bmod 26.$$

$$D(P) = (15 \cdot 3) \bmod 26 = 12 = m.$$

$$D(H) = (7 \cdot 3) \bmod 26 = 4 = e$$

$$D(H) = e.$$

$$D(W) = (22 \cdot 3) \bmod 26 = t.$$

$$D(P) = (15 \cdot 3) \bmod 26 = m.$$

$$D(H) = e.$$

$$D(D) = (3 \cdot 3) \bmod 26 = a.$$

$$D(I) = (18 \cdot 3) \bmod 26 = f.$$

$$D(W) = t$$

$$D(H) = e$$

$$D(U) = (20 \cdot 3) \bmod 26 = u.$$

$$D(W) = t$$

$$D(K) = (10 \cdot 3) \bmod 26 = k.$$

$$D(H) = e$$

$$D(W) = t$$

$$D(R) = (17 \cdot 3) \bmod 26 = o.$$

$$D(J) = (9 \cdot 3) \bmod 26 = j.$$

$$D(D) = a$$

$$D(S) = (12 \cdot 3) \bmod 26 = p$$

$$D(D) = a$$

$$D(U) = u$$

$$D(W) = t$$

$$D(B) = (1 \cdot 3) \bmod 26 = (-2 + 26) = 24 \bmod 26 = y$$

to the plaintext is -

wee weafter the to 9 a party.

Three important characteristics of the plaintext
matrices enabled to build force cryptanalysis
1) Encryption and decryption algorithms
are known.

2. There are only 25 keys to try.
3. The language of the plaintext is known and easily recognisable.

(3) Tabulate the substitution and Transposition techniques in detail. Apply two stage transpositions Cipher on the "treat diagrams as single units" using the keyword "sequence".

3. (i) Substitution Technique: letters of the plaintext are replaced by other characters or numbers or symbol.

Eq: name - ISTO.

2. Transposition technique: Performing some sort of permutation rearranged the order of the symbol.

Eq: name $4! = 24 \rightarrow$ ANME
NMEA.
....

3. Plain text: Treat diagrams as single unit
Keyword: sequence.

s	e	q	u	e	n	c	e
7	2	6	8	3	5	1	4
t	r	e	a	t	d	i	a
g	r	a	m	s	a	s	
i	n	g	l	e	u	n	i
t	e	v	e	x	y	z	

Final stage:

Y	2	6	8	3	5	1	4
i	o	n	y	p	h	n	s
f	k	e	w	a	s	i	x
d	a	u	x	c	a	g	u
t	g	e	t	a	m	p	v.

Second stage:

Y	2	6	8	3	5	1	4
n	i	g	l	s	k	a	g
n	a	b	a	s	x	u	j
n	s	a	m	n	e	u	i
i	t	d	t	y	w	x	t.

Key =

NIGL	SSAG	RAEA	SZUV
RSAM	NEDT	ITDT	YWXT

(4) What is mono-alphabetic cipher? Examine how it differs from Caesar cipher?

Q. Monoalphabetic Cipher: A cipher where each symbol is replaced by another symbol, where the replacement does not vary in called monoalphabetic cipher.

How it differs from Caesar cipher?

- 1) If the replacement remains same throughout the message, both cipher is monoalphabetic as opposed to polyalphabetic cipher.
- 2) In Caesar cipher, the total encryption and decryption depends on the key.

(5) Solve using playfair cipher. Encrypt the word "Semester Result" with the keyword "Examination". List the rules used.

5. Plain text: Semester Result.
Keyword: Examination

E	X	A	M	I/J
N	T	O	B	C
D	F	G	H	K
L	P	Q	S	
U	V	W	Y	Z

SE = LI

ME → IX

ST → PC

RE → ML

HE → LM

SR → LX

LT → PN

Ciphertext = LIIXPCMLLM LXPN

Rules:

Plaintext is encrypted 2-letters at a time.

1) If a pair is replaced, a replaced letter, absent letter like 'X'.

2) If both letters fall in the same row, move each letter to the right (wrapping back from start to end).

3) If both letters fall in the same column, move each with the letter below it.

(again again wrapping from top to bottom).

4. Otherwise each letter is replaced by the letter in the same row and in the column of the other letter of the pair.

(6) Demonstrate the encryption of the message "PAY" using hill cipher with the following key matrix and show the decryption.

$$K = \begin{bmatrix} 17 & 17 & 5 \\ 21 & 18 & 21 \\ 2 & 2 & 19 \end{bmatrix}$$

6. Plaintext : PAY.
Key Matrix : $K = \begin{bmatrix} 17 & 17 & 5 \\ 21 & 18 & 21 \\ 2 & 2 & 19 \end{bmatrix}$
 $C = K * P \text{ mod } 26.$

$$\begin{bmatrix} 17 & 17 & 5 \\ 21 & 18 & 21 \\ 2 & 2 & 19 \end{bmatrix} \begin{bmatrix} 5 \\ 80 \\ 24 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 375 \\ 819 \\ 486 \end{bmatrix} \pmod{26} = \begin{bmatrix} 11 \\ 13 \\ 18 \end{bmatrix} = \begin{bmatrix} L \\ N \\ S \end{bmatrix}$$

Ciphertext = LNS.

Now decryption:

$$\text{Ciphertext} = \text{LNS.} \quad \text{Key } K = \begin{bmatrix} 17 & 17 & 5 \\ 21 & 18 & 21 \\ 2 & 2 & 19 \end{bmatrix}$$

$$\text{Plaintext} = P = K^{-1} C \pmod{26}$$

$$\text{adj}(K) = \begin{bmatrix} 306 & 25 & 267 \\ 7 & 313 & 8 \\ 6 & 6 & 1 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 4 & 9 & 15 \\ 15 & 19 & 6 \\ 24 & 0 & 17 \end{bmatrix}$$

$$\Rightarrow \begin{bmatrix} 4 & 9 & 15 \\ 15 & 19 & 6 \\ 24 & 0 & 17 \end{bmatrix} \begin{bmatrix} 11 \\ 13 \\ 18 \end{bmatrix} \pmod{26} = \begin{bmatrix} 15 \\ 0 \\ 24 \end{bmatrix}$$

Plaintext = PAY.

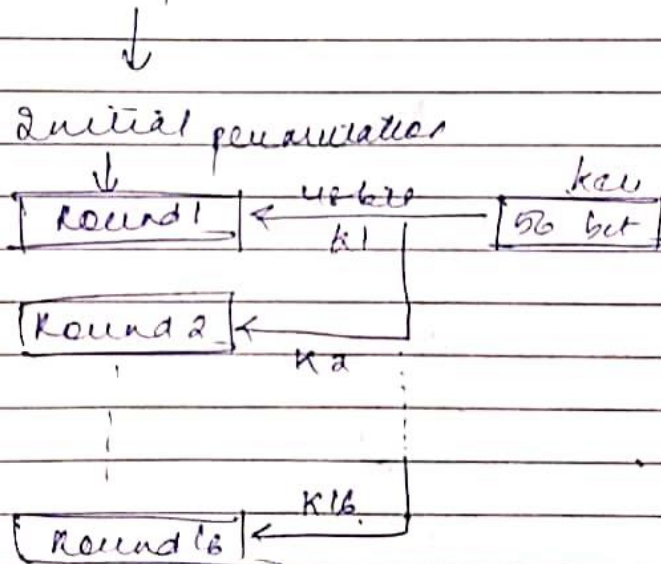
(7) Formulate the single round of DES algorithm and Design the key generation process of DES.

7. DES Algorithm

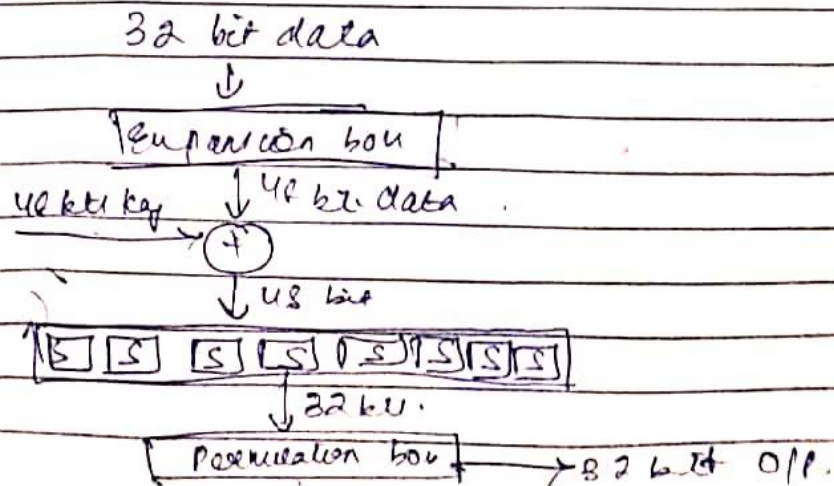
- Step :
- 1) Initial Permutation
 - 2) 16 Feistel Rounds.
 - 3) Swapping / left-right swap
 - 4) Final permutation.

Basic Structure

We have 64 bit plaintext.



Function definition:



Expansion box

Shon

32 bit data \rightarrow 48 bit data

48 bit data \rightarrow 32 bit data

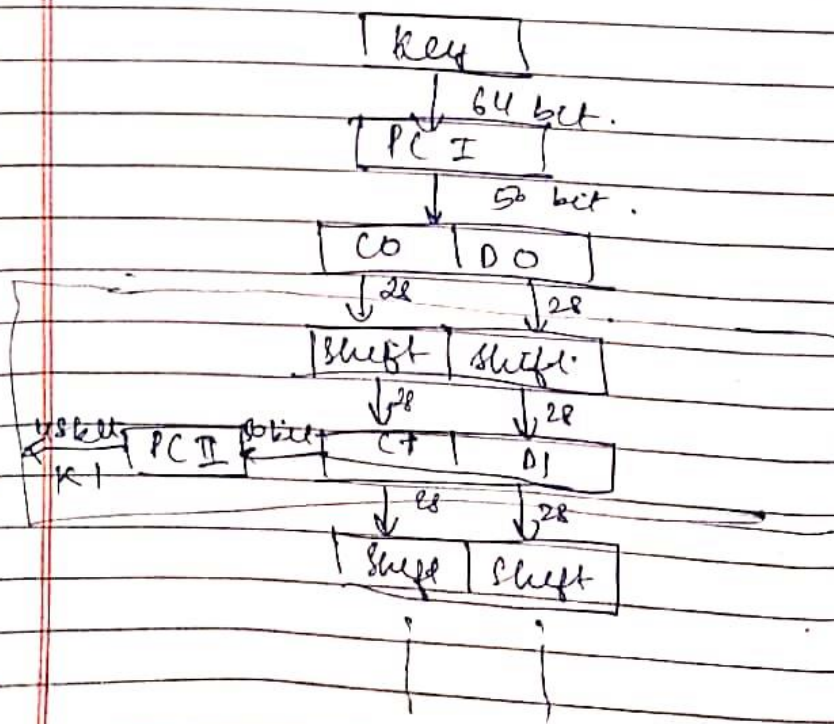
Definition of DES S-Boxes.

~~10 4 13 8 2 15 11 3 10 6 12 5 9 0 7~~

14	4	13	1	2	15	11	8	3	10	6	12	5	9	0	7
0	15	7	4	14	2	13	1	10	6	12	11	9	5	3	8
4	1	14	8	13	6	2	11	15	12	9	7	3	10	5	0
15	12	8	2	4	9	1	7	5	11	3	14	10	0	6	13

How to generate key.

The 64 bit key which goes as input to PC-I and we get 16 as 56 bit keys.



(8) The following message was encrypted with a columnar transposition cipher using a full rectangular array and keyword *mathematics*. Decrypt it.

RIUGS IPNCT MSPAL AUNCY SOOCH UEYSA RTE

8. Key = MATHEMATICS

Rearranging key in alphabetic Order -
 - A A C E H I M N S T.

Plaintext size = 33 letters.

Key size = 11.

Matrix dimension: $(33/11) \times 11$.

7.

A	A	C	E	H	I	M	N	S	T	
1	2	3	4	5	6	7	8	9	10	11
R	I	U	G	S	I	P	N	C	T	
M	S	P	A	L	A	U	N	C	Y	
S	O	O	C	H						
U	E	Y	S	A	R	T	E			

Rearranging the matrix key as per MATHEMATICS

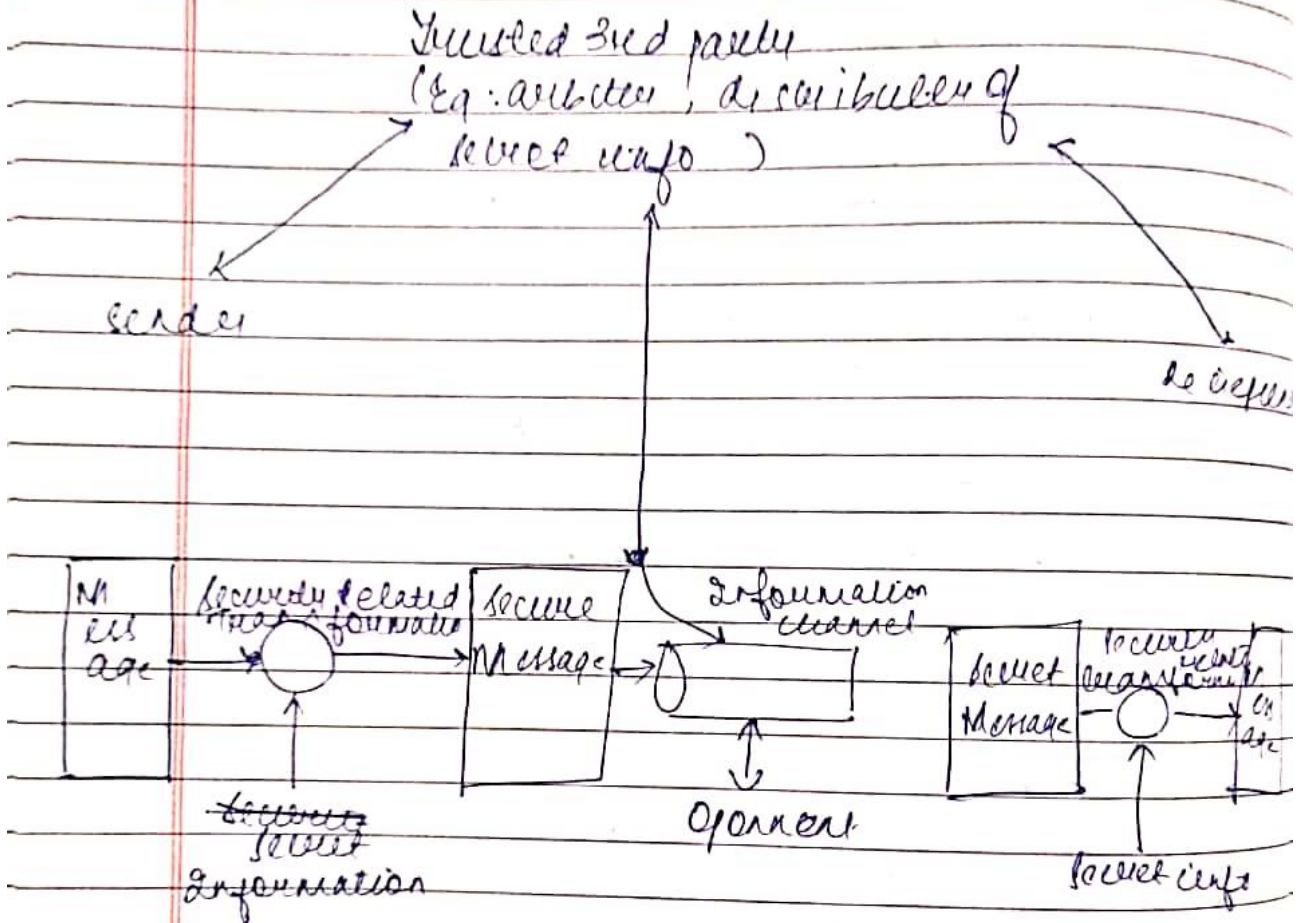
M	A	T	H	E	M	A	T	I	C	S
1	C	U	R	P	T	O	G	R	A	P
4	I	S	A	M	O	S	T	U	N	U
S	U	A	L	S	C	I	E	N	C	E

Hence, plaintext on decrypt sent =
 CRYPTOGRAPHY IS A MOST UNUSUAL SCIENCE.

(9) (i) With a neat block diagram, explain the network security model and the important parameters associated with it.

(ii) Differentiate active and passive security attacks. Categorize these attacks and explain one examples of each.

9ki) Model for network security



1) Transformation: Transformation of the info which has to be sent to the receiver, so that any opponent present at the information channel is unable to read the message. This initiates the encryption of the message.

2) Secret Info: Sharing of the secret info between sender and receiver of which the opponent must not have any clue. ~~as the~~

Notarization: There should be a trusted third party which should take the responsibilities of distributing the secret info to both the communicating parties and also present it from any opponent.

(ii) Security Attack

Security Attacks.

Passive Attack

Attempt to make use of info from the system but does not affect the system resources.

→ Release of message content - The attackers will easily be able to understand the details of the message.

→ Traffic Analysis: The attackers will understand the pattern.

Active Attack

~~At the~~ Attempt to alter the system resources or affect their operation

→ Masquerade: When one impersonates to be another entirely

→ Replay: Involves passive capture of a message and its subsequent re-transmission

→ Modification of message.

→ Denial of Service - It prevents normal use of communication facilities

(10) Explain why Modular arithmetic has been used in cryptography.

- 1) Modular arithmetic allows us to easily create group, rings and fields which are fundamental building blocks of most modern public key - crypto system.
- 2) Eg: Diffie - Hellman uses the multiplicative group of integers modulo a prime p .

(11) What is the difference between :

- (I) A block cipher and a stream cipher
- (II) Diffusion and Confusion
- (III) Differential and Linear cryptanalysis
- (IV) Active and Passive attack

11) Block Cipher

(i) Plaintext \rightarrow Ciphertext.
By taking plaintext block at a time.

- 2) Uses 64 bits or more.
- 3) Complexity is simple.
- 4) Uses confusion and diffusion.
- 5) Reverse encryption is difficult.
- 6) Eg: ECB, CBC

Stream Cipher

1 Byte Plaintext \rightarrow Ciphertext

- (i) 8 bits.
- (ii) Complexity is more.
- 4) Uses only confusion.

5) Reverse encryption is easy.

6) Eg: LFB, OPB.

(ii) Diffusion

- 1) It hides the relationship between ciphertext and plaintext.
- 2) If a single symbol is changed in the plaintext then all symbols will change in the ciphertext.

Confusion

- 1) Hides the relationship between ciphertext and the key.
- 2) If a single ~~key~~ with key is changed, all bits in the ciphertext will change.

(iii) Differential

1) Differential analysis focuses on statistical analysis of two 1/b and two 0/b of cryptographic algorithm.

Linear

1) Linear cryptanalysis focuses on statistical analysis against one round of decrypted ciphertext.

<u>Differential</u>	<u>Linear</u>
2. In differential crypt analysis, the changes to the immediate ciphertext are obtained between multiple rounds of encryption. The attack can be combined and this is referred as differential cryptanalysis.	2. The crypt analyst decrypts each cipher using all possible subkeys for one round of encryption and studies the resulting unencrypted cipher text to analyse the random result.
<u>Active attack</u>	<u>Passive attack.</u>
1) The attacker can see the data but can't modify.	1) Attempts to make change to data on target.
2) It is in a danger of confidentiality.	2) Danger for session as messages are intercepted.
3) System has no impact.	3) Systems damaged.
4) System resources are not changed.	4) System resources can be changed.

(12) This problem provides a numerical example of encryption using a one-round version of DES. We start with the same bit pattern for the key K and the plaintext, namely:

in hexadecimal notation: 0 1 2 3 4 5 6 7 8 9 A B C D E F

in binary notation: 0000 0001 0010 0011 0100 0101 0110 0111

1000 1001 1010 1011 1100 1101 1110 1111

- Derive K_1 , the first-round subkey.
- Derive L_0 , R_0 .
- Expand R_0 to get $\text{EXP}(R_0)$.
- Calculate $A = \text{EXP}(R_0) \text{ XOR } K_1$.
- Group the 48-bit result of (d) into sets of 6 bits and evaluate the corresponding S-box substitutions.
- Concatenate the results of (e) to get a 32-bit result, B .
- Apply the permutation to get $P(B)$.

12) a) In Binary notation: 0000 1011 0000 0010 0110
0111 1001 1011 0100 1001 1010 0101.

in hexadecimal notation: 0 B 0 2 6 - 1 9 B 4 9 A 5.

b) LO, RO are derived by passing the 64-bit operands through initial permutation.

LO = 1100 1100 0000 0000 1100 1100 1111 1111
RO = 1111 0000 10 10 10 10 1111 0000 1010 1010.

c) $EXP(K_0) = 011110 \ 100001 \ 010101 \ 010101$
 $011110 \ 100001 \ 010101 \ 010101.$

d) $A = 011100 \ 010001 \ 011100 \ 110010 \ 111000$
 $010101 \ 110011 \ 110000.$

e) 0 (base 10) = 0000 (base 2)
12 (base 10) = 1100 (base 2)
2 (base 10) = 0010 (base 2)
1 (base 10) = 0001 (base 2)
6 (base 10) = 0110 (base 2)
13 (base 10) = 1101 (base 2)
5 (base 10) = 0101 (base 2)
0 (base 10) = 0000 (base 2)

f) $B = 0000 \ 1100 \ 0010 \ 0001 \ 0110 \ 1101 \ 0101 \ 0000$

g) $P(B) = 1001 \ 0010 \ 0001 \ 1100 \ 0010 \ 0000 \ 1001 \ 1100.$

(13) Show how to convert the given text "VALLIAMMAI" into cipher text using Rail fence Technique.

13) Plaintext: VALLIAMMAI.

Using depth = 2.

V L I M A
 A L A M I.

Ciphertext: VLIMAAALAMI.

(14) Which parameters and design choices determines the actual algorithm of a Feistel cipher?

14. Parameters and Algorithms.

1. Block size: Larger block size means greater security but reduced encryption/decryption speed for a given algorithm.

2. Key size: Larger key size means greater security but may decrease encryption/decryption speed.

3) Number of Rounds: A single round offers inadequate security but that multiple rounds offer increasing security. A typical size is 16 rounds.

4) Key Generation Algorithm: Generate complexity in the algorithm should lead to greater difficulty of cryptanalysis.

5) Round Function (F): Greater complexity generally means greater resistance to cryptanalysis.

(15) Define the type of security attack in each of the following cases:

- (i) A student breaks into university's office to get a copy of exam paper to be held on the next day.
- (ii) A person sends hundreds of e-mails every day to another person using a phony return e-mail address.
- (iii) John gives a cheque of \$14 to the shopkeeper to buy a book. Later he finds that the cheque was cashed for \$140.

15) (i) This is a type of confidentiality attack because an unauthorised student introduced into university office to get a copy of exam paper.
(ii) This is a type of phishing attack.
(iii) This is a type of integrity attack because the check has been already given by John and does not have the permission to do that. So, the question is shopkeeper's honesty.