

$$i/p \rightarrow \frac{29}{7} \rightarrow 4 \dots$$

↓  
ans

$$\frac{28}{2} \rightarrow (14)$$

↓  
ans

$$\text{divisor} \overline{) \text{dividend}} \left[ \text{Quotient} \right.$$

$$\underline{\text{Remainder}}$$

$$\boxed{\text{Quotient} * \text{Divisor} + \text{Remainder} = \text{Dividend}}$$

⇒

$$\text{Quotient} * \text{Divisor} \leq \text{Dividend}$$

$$\frac{29}{7} \rightarrow [0 \longleftrightarrow 29]$$

$$\text{search space} \rightarrow [0 \longrightarrow \text{dividend}]$$

$$0 \longleftarrow \text{-----} \longrightarrow 29$$

$$s=0, e=29, \text{mid} = \frac{0+29}{2} = \boxed{14}$$

is  $\boxed{14}$  a possible answer?

$$Q * \text{Divisor} \Rightarrow 14 * 7 = (98) \neq 29 \times$$

$$98 < 29 \times$$

$$(98 > 29)$$

→ left me  
jao

i/p  $\rightarrow$  2 no.  $\rightarrow$  Divide  $\rightarrow$  Binary search

(a, b)

$\boxed{\frac{a}{b}}$   $\rightarrow$  Quotient

$$\frac{28}{7} \geq 4$$

divisor  $\sqrt{\text{Dividend}}$  Quotient

Remainder

Quotient  $\times$  divisor + remainder = Dividend

$$\boxed{\text{Quotient} \times \text{divisor} \leq \text{Dividend}}$$

①

0 ————— 29

$$s = 0, e = 29$$

$$\text{mid} = \frac{0 + 29}{2} \Rightarrow \boxed{14}$$

is this a possible answer

dividend  $\rightarrow$

$$\left( \frac{29}{7} \right) \rightarrow 4$$

divisor  $\downarrow$

$$Q \times \text{divisor} \leq \text{Dividend}$$

$$14 \times 7 \Rightarrow \boxed{98} \neq \boxed{29} \rightarrow \text{false}$$

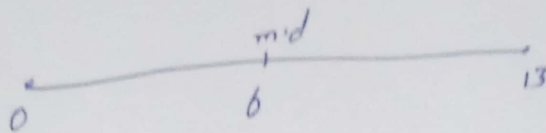
$$\boxed{98} > 29$$

$\rightarrow$  Left  $\rightarrow e = \text{mid} - 1$

$$e = 14 - 1$$

$$e = \boxed{13}$$

2



$$s = 0, e = 13$$

$$mid = 6$$

Q x divis.

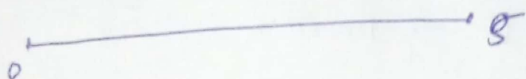
$$6 \times 7$$

$$29$$

$$42 > 29$$

→ left →  $e = mid - 1$   
 $e = 6 - 1 \Rightarrow 5$

3



$$s = 0, e = 5$$

$$mid = 2$$

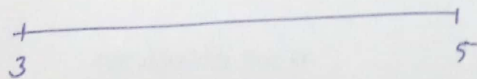
Q x divisor  $\leq$  dividend

$$2 \times 7 \leq 29$$

$$14 \leq 29 \rightarrow \text{valid ans}$$

→ ans store  
→ right move  
 $s = mid + 1$   
 $s = 3$

4



$$s = 3$$

$$mid = 4$$

$$e = 5$$

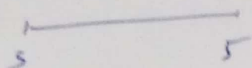
Q x divisor  $\leq$  dividend

$$4 \times 7 \leq 29$$

$$28 \leq 29 \rightarrow \text{valid ans}$$

→ ans store  
→ right move

5



$$s = 5$$

$$e = 5$$

$$mid = 5$$

$$Q \times \text{divisor} \leq \text{dividend}$$

$$5 \times 7 \leq 29$$

$$(35 \leq 29) \rightarrow \text{false} \rightarrow \text{left}$$

### Conditions

$$Q \times \text{divisor} == \text{dividend}$$

$\rightarrow Q \rightarrow \text{finalAns}$

$$Q \times \text{divisor} < \text{dividend}$$

$\rightarrow \text{ans store}$   
 $\rightarrow \text{right}$

$$Q \times \text{divisor} > \text{dividend}$$

$\rightarrow \text{left}$

$$T.C \Rightarrow O(\log n)$$

$n \rightarrow \text{dividend}$



$$\left( \frac{+ve}{+ve} \right) \rightarrow +ve$$

$$\left( \frac{-ve}{+ve} \right) \rightarrow -ve$$

$$\left( \frac{-ve}{-ve} \right) \rightarrow +ve$$

$$\left( \frac{+ve}{-ve} \right) \rightarrow -ve$$

# Binary search on nearly sorted array

nearly  
sorted  
array

20	10	30	50	40	70	60
0	1	2	3	4	5	6

sorted  
array



$i^{\text{th}}$  index

choice

$(i-1)$  index

$(i)$  index

$(i+1)$  index

nearly sorted  
array

S.A  
colony

(10)

NSA  
9/10/11

sorted  
array

-101 ↑↑↑	012 ↑↑↑	123 ↑↑↑	234 ↑↑↑	345 ↑↑↑	456 ↑↑↑	567 ↑↑↑
10	20	30	40	50	60	70
0	1	2	3	4	5	6

normal sorted  
array

→ if ( $arr[mid] == target$ )  
    return mid

→ if ( $target > arr[mid]$ )  
    ↳ right

— else  
    ↳ left

nearly sorted array

→ if ( $arr[mid-1] == target$ )  
    return ( $mid-1$ )

if ( $arr[mid] == target$ )  
    return mid

if ( $arr[mid+1] == target$ )  
    return ( $mid+1$ )

→ if ( $target > arr[mid]$ )  
    ↳ right

Day 800

①

				mid		
20	10	30	50	40	70	60
0	1	2	3	4	5	6

$s=0, e=6$

$mid=3$

$arr[mid-1] \rightarrow 30 == 70 \times$

$arr[mid] \rightarrow 50 == 70 \times$

$arr[mid+1] \rightarrow 40 == 70 \times$

if (target > arr[mid])

$70 > 50$

→ right →  $s = mid + 2$  <sup>catch</sup>

②

40	70	60
4	5	6

$s=4, e=6$

$mid=5$

$arr[mid-1] \Rightarrow$

70	60
5	6

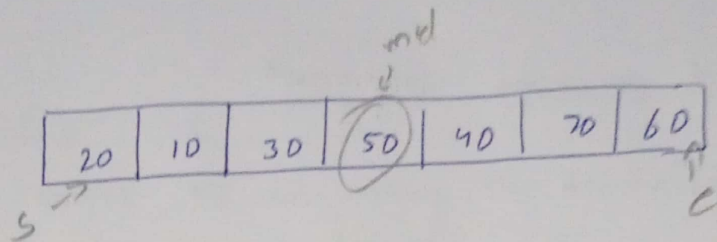
$s=5, e=6$

$arr[mid-1]$

$70 == 70$

→ True





$s=0, e=6$

$mid=3$

target = 20

$arr[mid-1] \rightarrow 30 == 20 \quad \times$

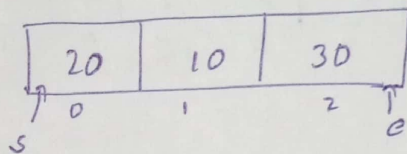
$arr[mid] \rightarrow 50 == 20 \quad \times$

$arr[mid+1] \rightarrow 40 == 20 \quad \times$

if (target > arr[mid])

$20 > 50 \rightarrow \text{false}$

catch  
 $\rightarrow \text{Left} \Rightarrow e = mid - 1$   
 $e = mid - 2$



$s=0, e=2$

$mid=1$

$arr[mid-1] \rightarrow 20 ==$

Imp

Find the odd occurring element

- find element that occurs odd times
- all elements → even no. of times occur except one → odd
  - all repeating no. → pairs repeat & pairs are not repeated
  - ek baar me koi bhi no. 2 se jada baar nhi aa skta

1	1	5	5	2	2	3	3	2	4	4
0	1	2	3	4	5	6	7	8	9	10

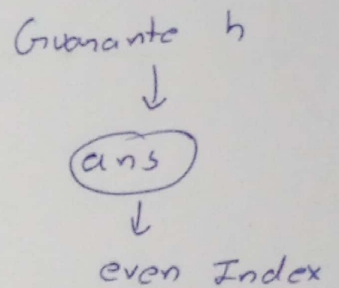
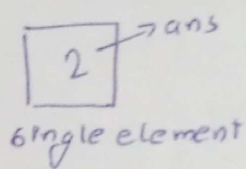
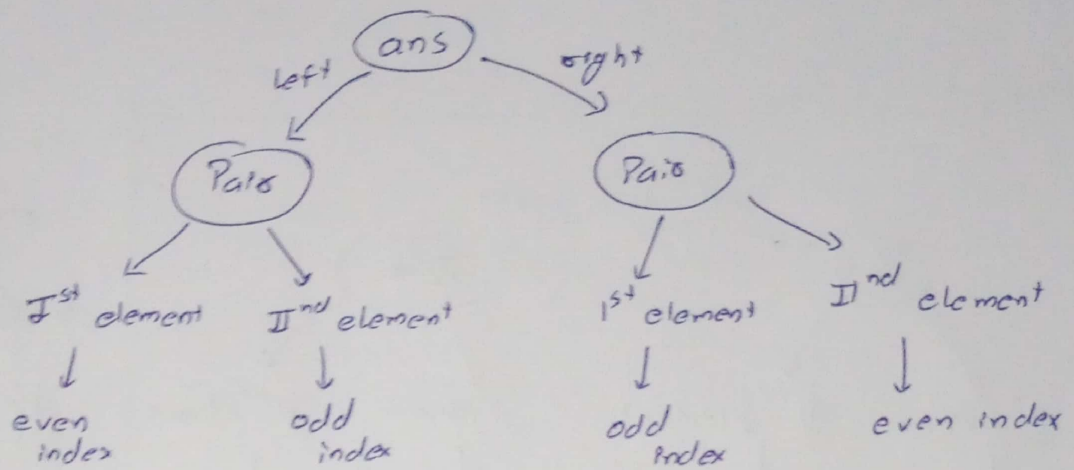
Approach

- (A) XOR →  $O(n)$
- (B) Count → map
- (C) sorting → 3 times

Binary search

- classical
- Search space
- Predicate function
- index  
↓  
Logic

## Observation

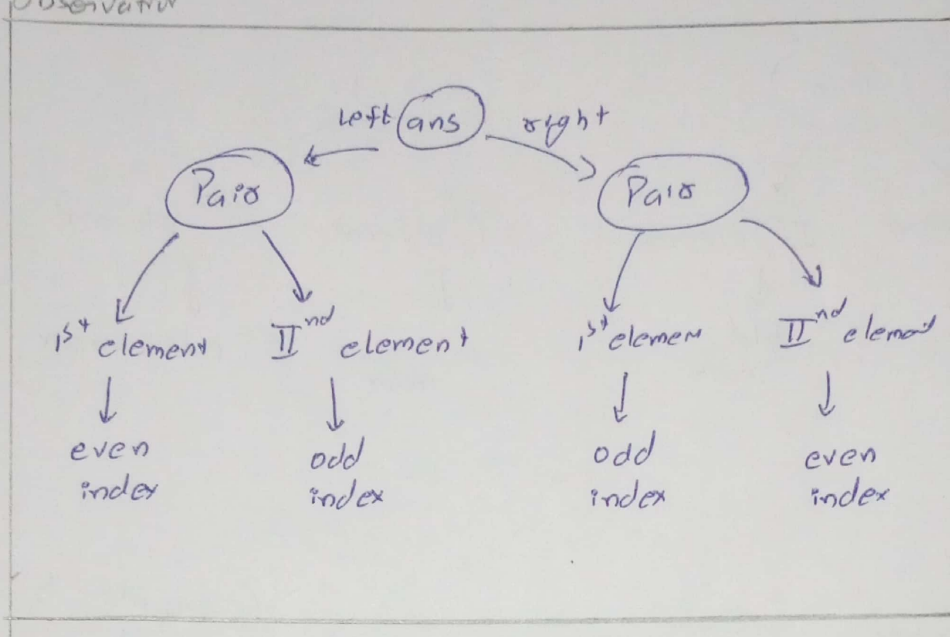


if( $\text{mid} \& 1$ )  $\rightarrow$  true  $\rightarrow$  odd number



10	10	2	2	5	5	2	5	5	20	20	11	11	10	10
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

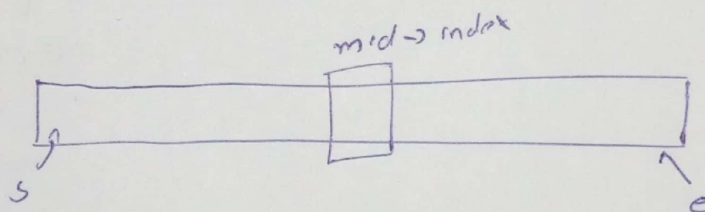
observation



single element

$s \rightarrow [2] \leftarrow e$

if  $(s == e)$   
return  $s$



if  $(mid \% 2 == 0) \Rightarrow$  even

$\hookrightarrow$  if  $(arr[mid] == arr[mid+1])$

$\hookrightarrow$  right  $\rightarrow s = mid + 2$

else

$\hookrightarrow$  right / ans

$e = mid$

catch

why

Peak in array



if (mid % 2 == 1)

↳ if (arr[mid] == arr[mid-1])

↳ right →  $s = mid + 1$

else

↳ left →  $e = mid - 1$

ans right me hai

### Pseudocode

int  $s = 0$ ,  $e = n - 1$ ,  $mid = s + (e - s) / 2$

while ( $s \leq e$ ) {

if ( $s == e$ )

return  $s$ ;

if ( $mid \& 1$ )

↳ if (arr[mid] == arr[mid-1]) →  $s = mid + 1$

else  $e = mid - 1$

else

↳ if (arr[mid] == arr[mid+1]) →  $s = mid + 2$

else →  $e = mid$