

Time Complexity

• Multiple $O(N)$

$$= O(N) + O(N) \approx O(N)$$

• $O(K * N)$

$$K \ll N = O(N)$$

• $O(3), O(4), O(5) \approx O(1)$ }

* Number of iterations in 1sec - $10^7 - 10^8$

$$10^7 \rightarrow 1\text{sec}$$

$$10^8 \rightarrow 10\text{sec}$$

$$10^9 \rightarrow 100\text{sec}$$

⋮

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बुधवार Example: $\#n \leq 10^{15}$

```
1) int n; cin >> n; int ct = 0;
   while (n > 0)
   { n = n / 2; ct++; }
```

$$2^a = N$$

$$\log_2(2)^a = \log_2 N$$

$$a = \log_2(N)$$

$$\text{Complexity} = O(\log(N))$$

```
2) for (int i = 0; i <= n; i++) {
    for (int j = 0; j < i; j++) {
        ct++;
    }
}
```

$$\begin{array}{cc} 0 & 0 \\ 1 & 1 \\ \vdots & \vdots \\ n & n \end{array} \quad \downarrow \quad \frac{n(n+1)}{2}$$

$$TO = O\left(\frac{n^2}{2} + \frac{n}{2}\right)$$

$$= O\left(\frac{n^2}{2}\right) + O\left(\frac{n}{2}\right)$$

$n^2 \gg n$

$$\approx O(n^2) \text{ Ans}$$

③ Sum of N over all test cases $n < 10^{17}$

$1 \leq T \leq 10^5$
 $1 \leq N \leq 10^5$

while(t--)

for (int i=0; i<n; i++)

{ ct++; }

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$O(n_0) \dots O(n_n) \leq 10^7$

↳ $ct < 10^7$

↳ runs accord. to value of n .

Print Answer Modulo

$$(a+b) \% M = ((a \% M) + (b \% M)) \% M.$$

$$(a * b) \% M = ((a \% M) * (b \% M)) \% M$$

$$(a-b) \% M = ((a \% M) - (b \% M) + M) \% M.$$

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$$(a/b) \% M = ((a \% M) * (b^{-1}) \% M) \% M.$$

↳ Multiplicative inverse

Eg: int main()

{

int M = 47;

for (int i=0; i<n; i++)

{ fact = (fact * i) % M; }

cout << fact;

}

$1 \cdot 10^9 + 7 \rightarrow$ very close to integer maximum.