

Longest Subarray zero sum

BRUTE FORCE

Generate all subarrays and calculate sum and check if it is equal to 0 if yes then store maxlen.

```
#include <bits/stdc++.h>

int LongestSubsetWithZeroSum(vector < int > arr) {

    // Write your code here
    int n=arr.size();
    int sum=0, len=0, maxlen=0;
    for(int i=0; i<n; i++)
    {

        for(int j=i; j<n; j++)
        {
            sum=0;
            for(int k=i; k<=j; k++)
            {
                sum+=arr[k];

            }
            if(sum==0)
                maxlen=max(j-i+1, maxlen);

        }
    }
    return maxlen;
}
```

- Time Complexity : $O(N^3)$
- Space Complexity : $O(1)$.

Improved Brute Force

Using two loops to calculate sum of subarrays and compare.

```
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int n=arr.size();
int sum=0, len=0, maxlen=0;
for(int i=0;i<n;i++)
{
    sum=0;
    for(int j=i;j<n;j++)
    {
        sum+=arr[j];
        if(sum==0)
            maxlen=max(j-i+1,maxlen);
    }
}
return maxlen;
}

```

- Time Complexity : $O(N^2)$
- Space Complexity : $O(1)$

Optimal Approach :

Will work only when sum is k not for zero sum.

We will use the concept of prefix sum here. We look for $\text{sum} == k$ if it is found then we compare maxlen and $i+1$ len whichever is maximum. Then moving forward if we have sum x then we find whether $x-k$ exist in already stored sum in the map , we will be storing each sum in each iteration in the map if the sum is found in map this means k is also found. We will store sum as well as index at which that sum was found subtracting $x-k$ sum index with current sum index will give us the len of array whose sum is k.

```

class Solution {
public:
    int subarraySum(vector<int>& arr, int k) {
        // Write your code here
        map<int,int> m;
        int cnt=0, sum=0;
        for(int i=0;i<arr.size();i++)
        {
            sum+=arr[i];
            if(sum==k)
            {
                cnt++;
            }
            int remain=sum-k;
            if(m.find(remain)!=m.end())

```

```

        {
            cnt+=m[remain];
        }
        m[sum]++;

    }
    return cnt;

}

};

```

- Time complexity : $O(N)$
- Space Complexity : $O(N)$

Optimal Approach :

Using two pointer if sum exceeds then we push left pointer and keep on moving right pointer and adding to sum. len is measured using $\text{right-left}+1$.

```

#include <bits/stdc++.h>

int LongestSubsetWithZeroSum(vector < int > arr,int k) {

    // Write your code here
    int left=0,right=0,n=arr.size(),maxlen=0;
    long long sum=arr[0];
    while(right<n)
    {

        while(left<=right && sum>k)
        {
            sum-=arr[left];
            left++;
        }

        if(sum==k)
        {
            maxlen=max(maxlen, right- left+1);
        }
        right++;
        if(right<n) sum+=arr[right];

    }
    return maxlen;

}

```

