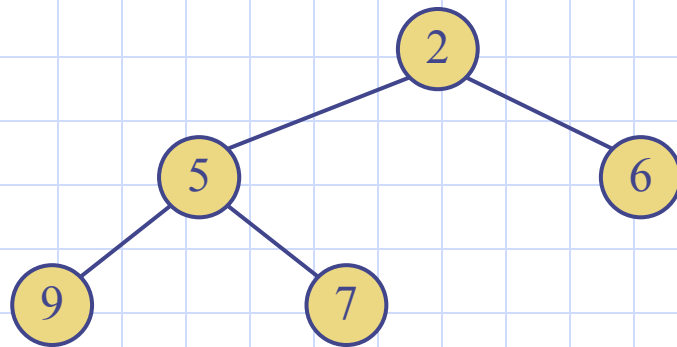


Heaps



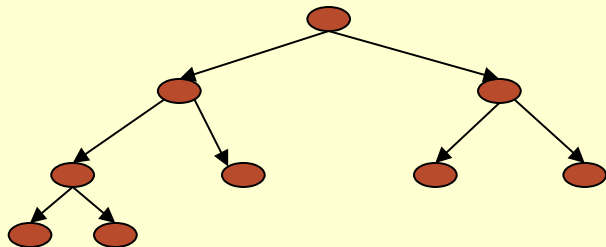
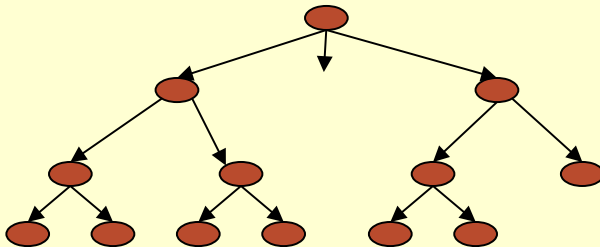
Definition of a heap

- A **heap** is a tree that satisfies the *heap property*.
 - **Minimum Heap Property**: The value stored at each node is less than or equal to the values stored at its children.
 - **OR Maximum Heap Property**: for greater

Heap Structure Property

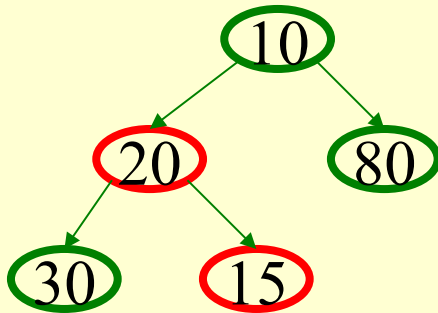
- All levels but the last (leaf) are full and the leaf nodes are left oriented, i.e. filled from left to right.

Examples:

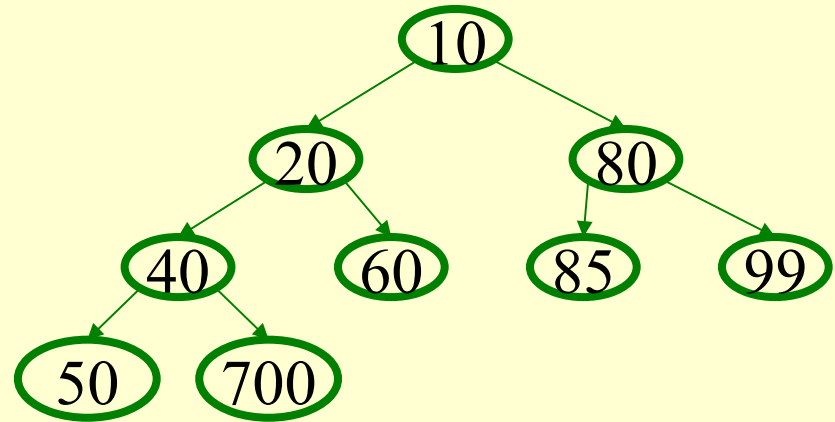


Heap Order Property

Heap order property: For every non-root node X , the value in the parent of X is less than (or equal to) the value in X .

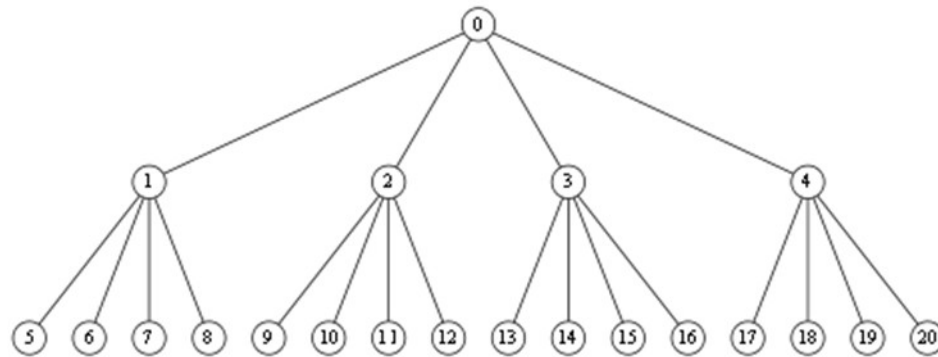


not a heap



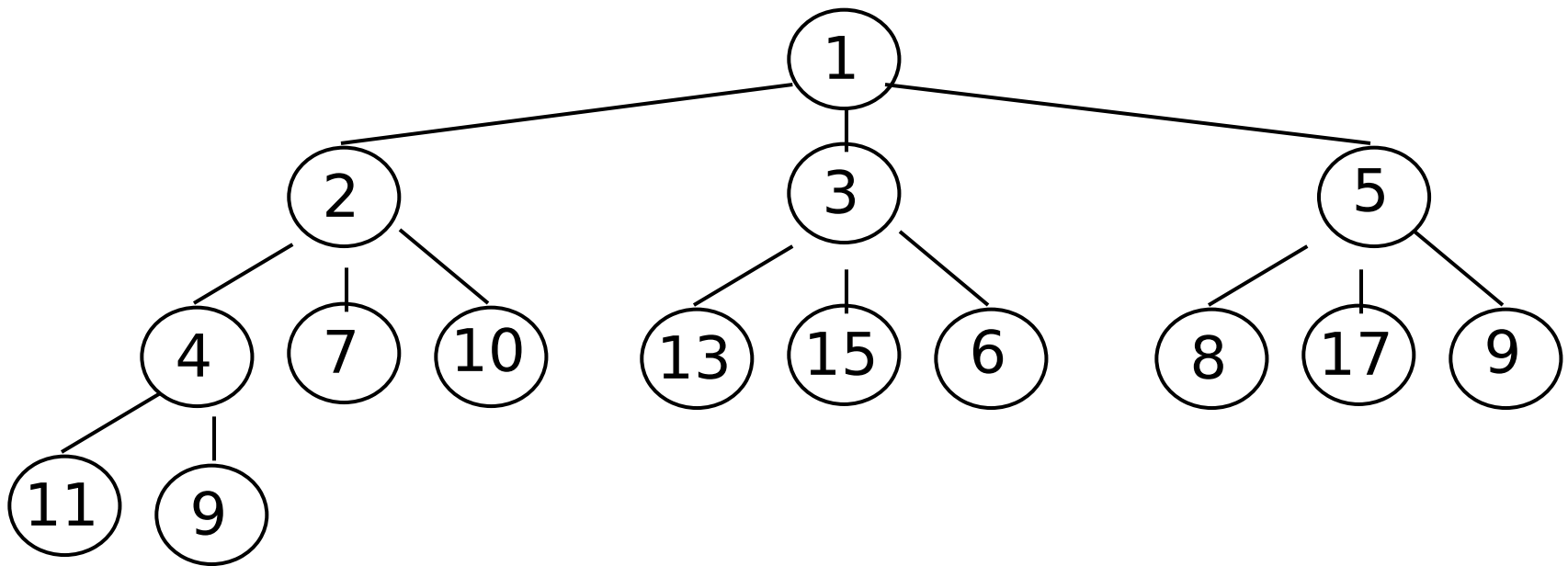
HEAPS

- **Example of a 4-heap**





d-Heap



An example of 3-
heap

Definition of a heap

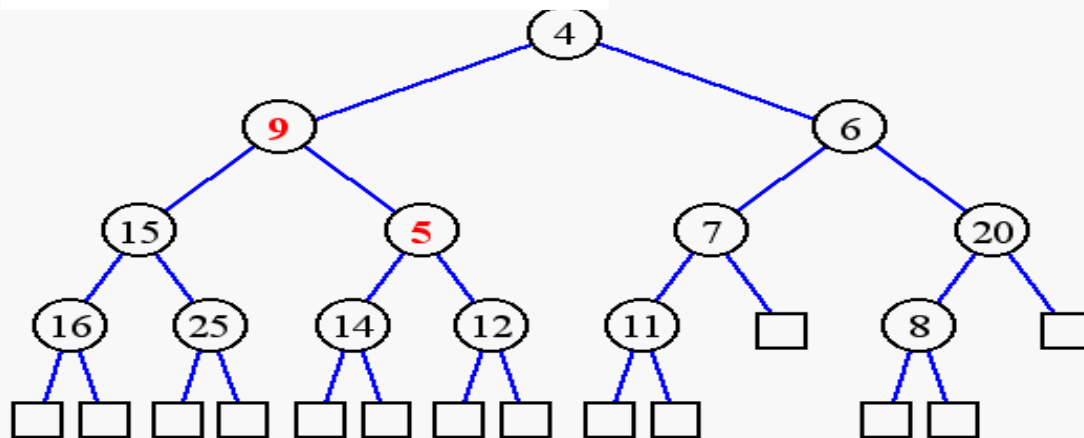
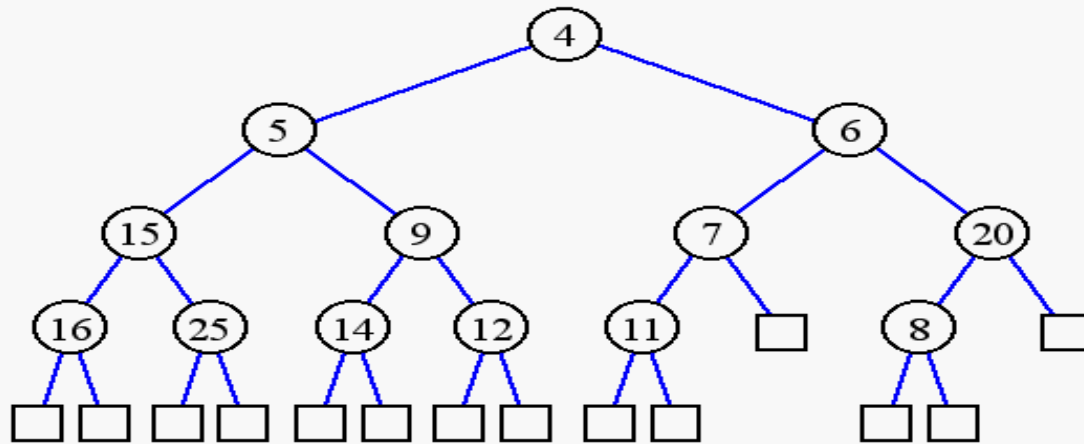
- A **heap** That has D children is called **D-Heap**.
- The most popular heap is **2-Heap**,
 - Also called **binary Heap Tree**

Definition of a heap

- **Difference between Binary Search Tree and Binary Heap Tree**



Heap or Not a Heap?



Height of a Heap

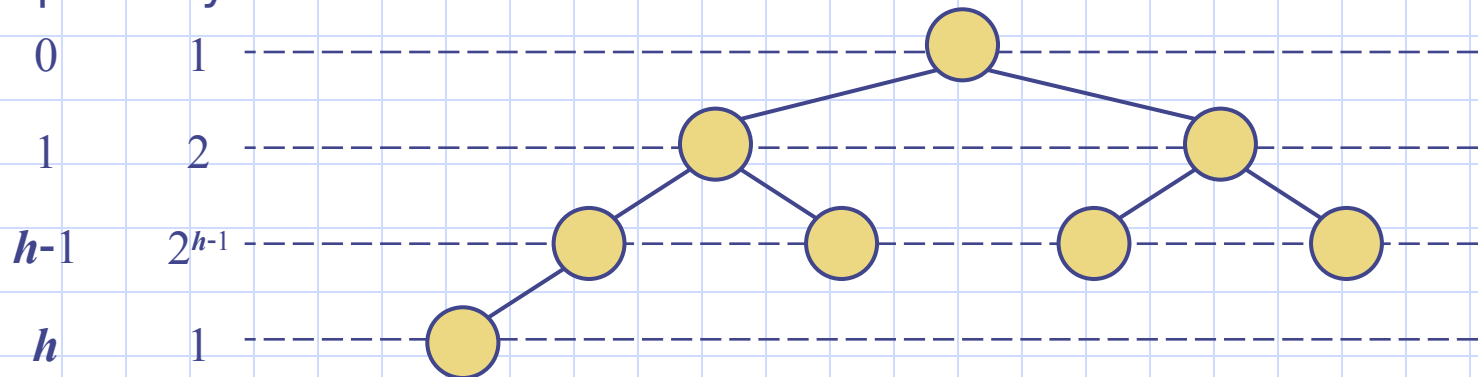


- **Theorem:** A heap storing n keys has height $O(\log n)$

Proof: (we apply the complete binary tree property)

- Let h be the height of a heap storing n keys
- Since there are 2^i keys at depth $i = 0, \dots, h-1$ and at least one key at depth h , we have $n \geq 1 + 2 + 4 + \dots + 2^{h-1} + 1$
- Thus, $n \geq 2^h$, i.e., $h \leq \log n$

depth keys

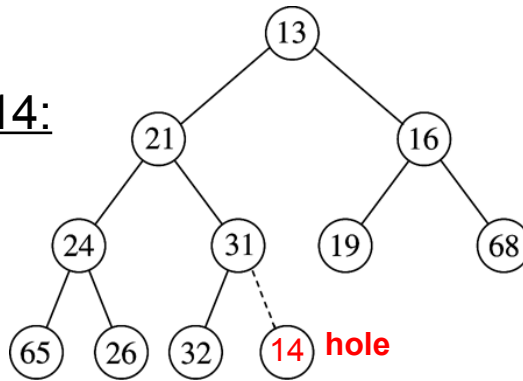


Operations on Heap

- **Insert**
- **Delete-Min**

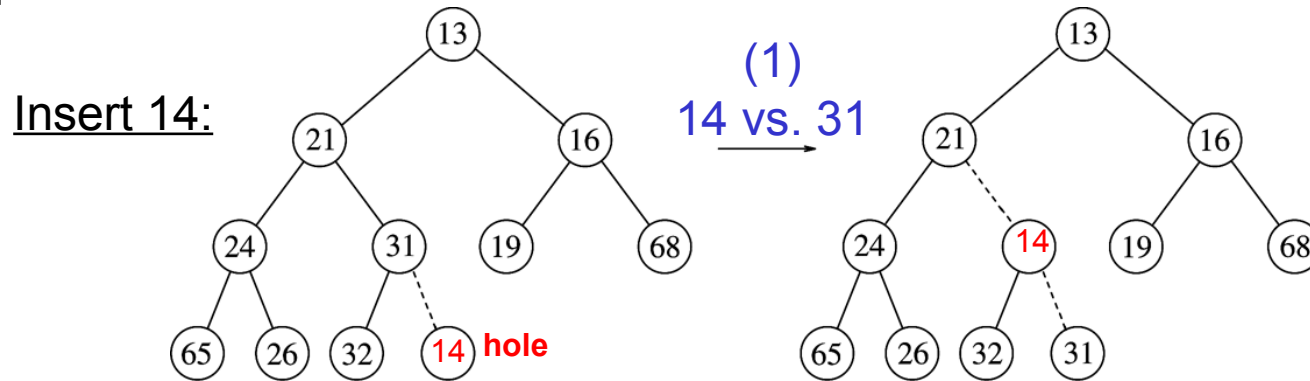
Heap Insert: Example

Insert 14:



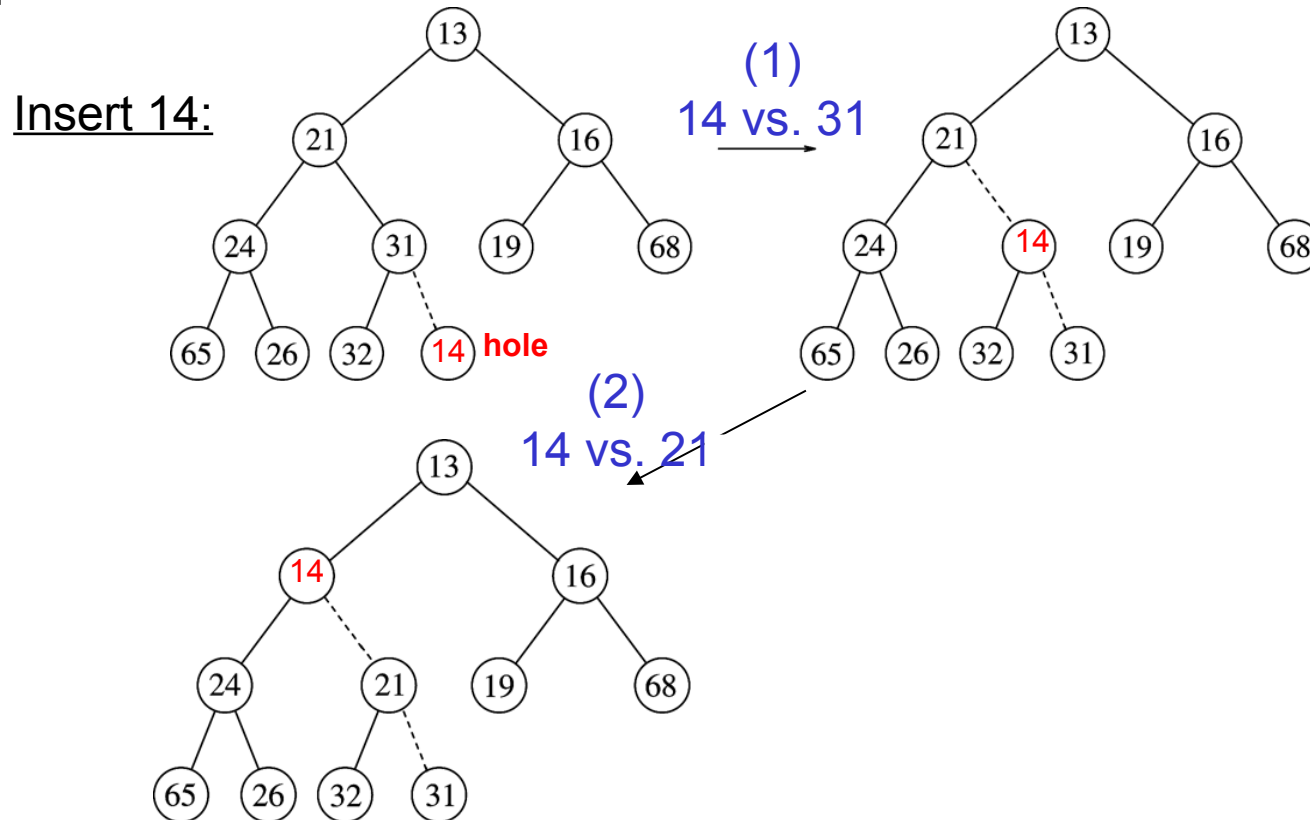
Percolating Up

Heap Insert: Example



Percolating Up

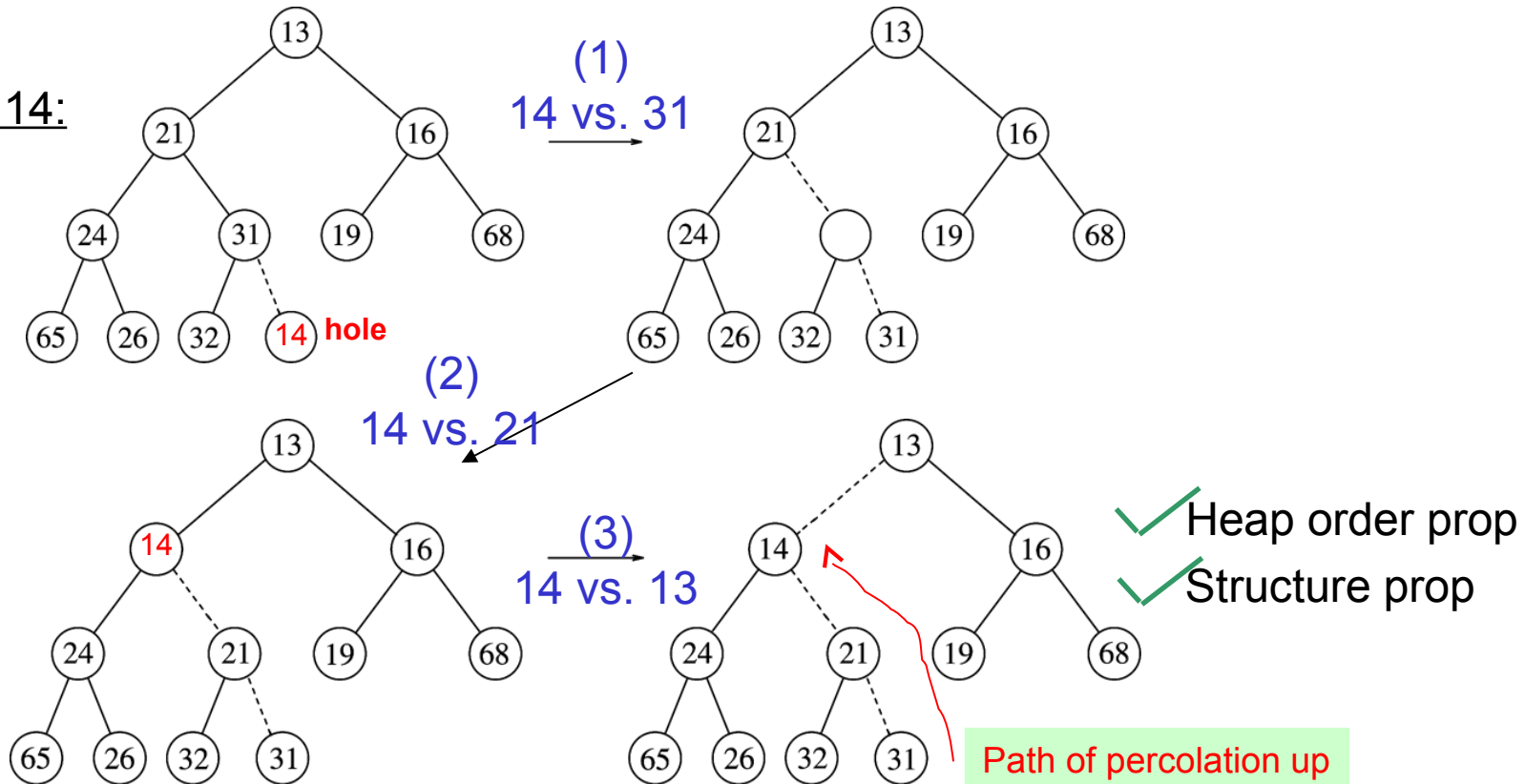
Heap Insert: Example



Percolating Up

Heap Insert: Example

Insert 14:



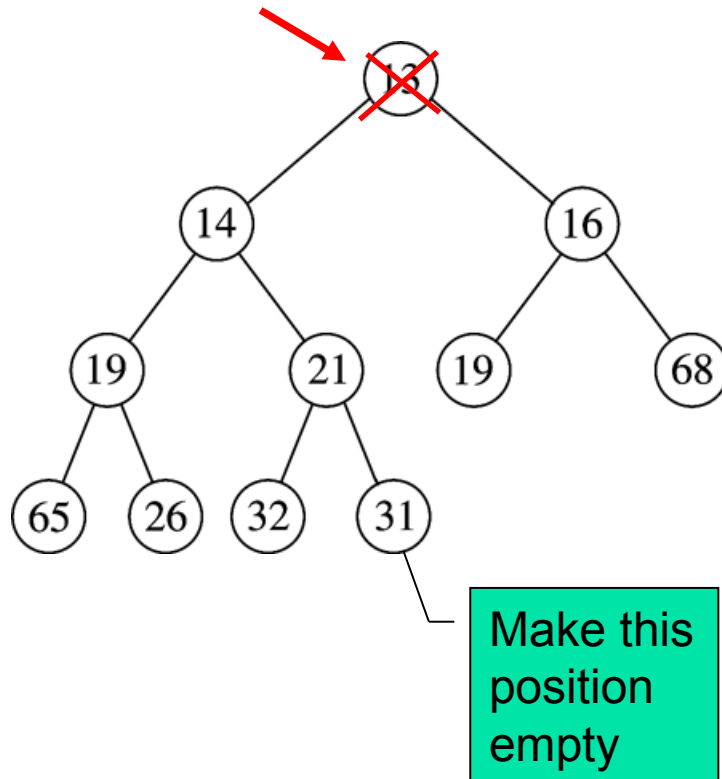


Heap DeleteMin

- Minimum element is always at the root
- Heap decreases by one in size
- Move last element into hole at root
- ***Percolate down*** while heap-order property not satisfied

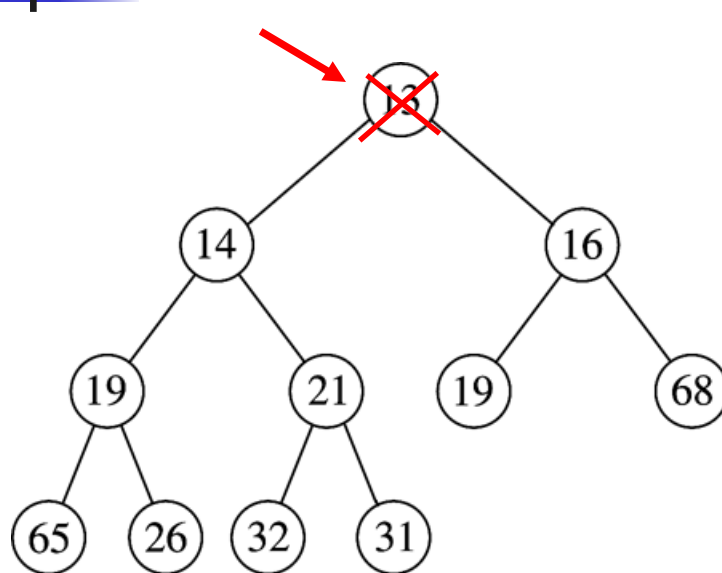
Percolating down...

Heap DeleteMin: Example

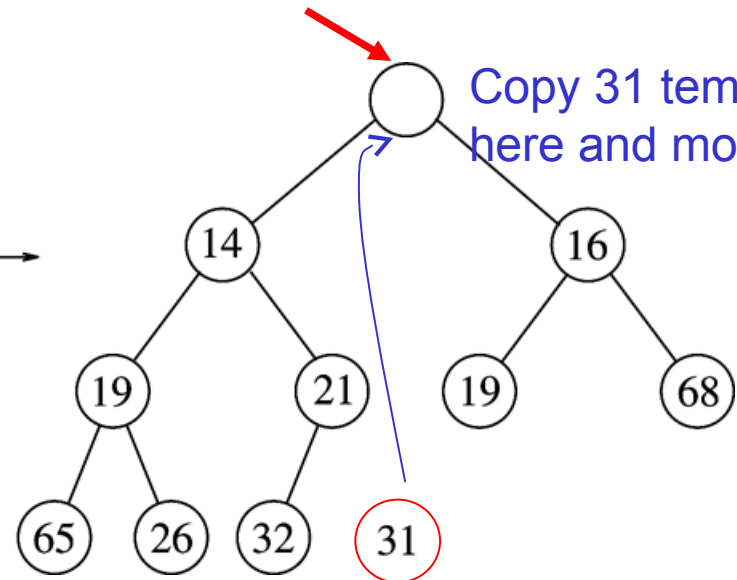


Percolating down...

Heap DeleteMin: Example



Make this
position
empty



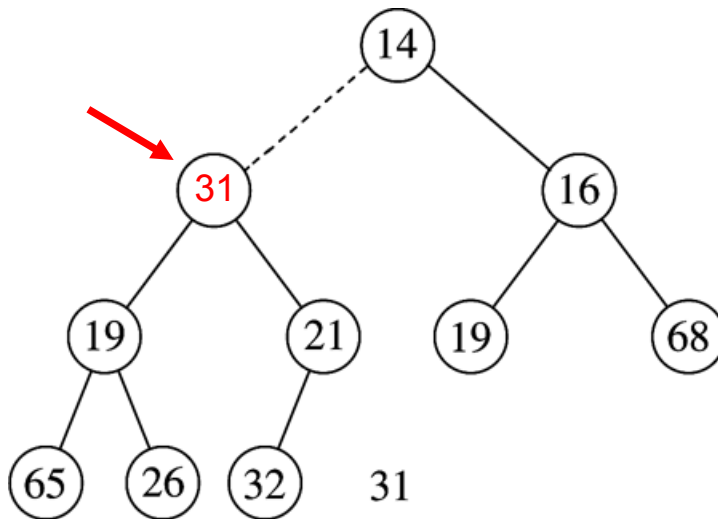
Copy 31 temporarily
here and move it down

Is $31 > \min(14, 16)$?

• Yes - swap 31 with $\min(14, 16)$

Percolating down...

Heap DeleteMin: Example

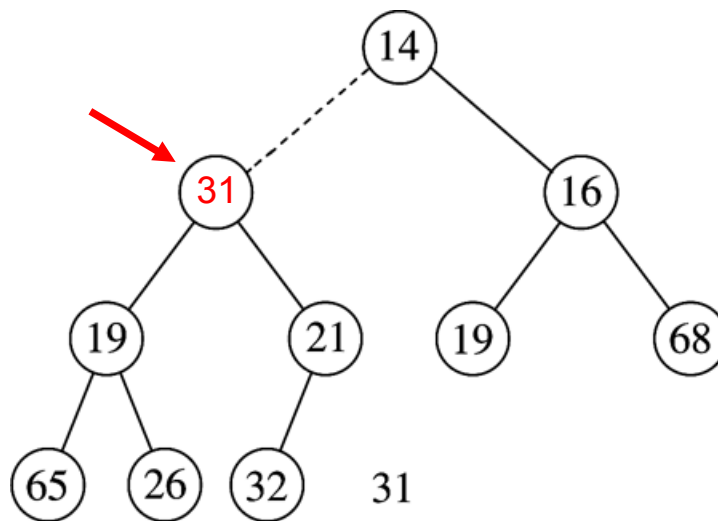


Is $31 > \min(19, 21)$?

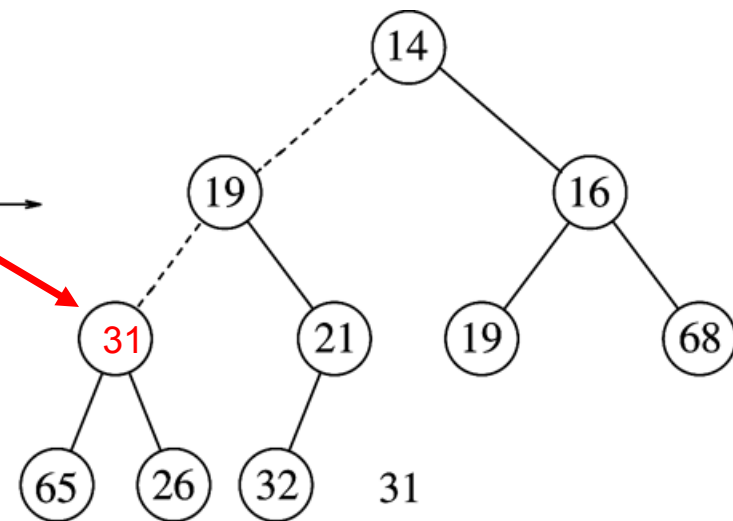
• Yes - swap 31 with $\min(19, 21)$

Percolating down...

Heap DeleteMin: Example



Is $31 > \min(19, 21)$?
• Yes - swap 31 with $\min(19, 21)$

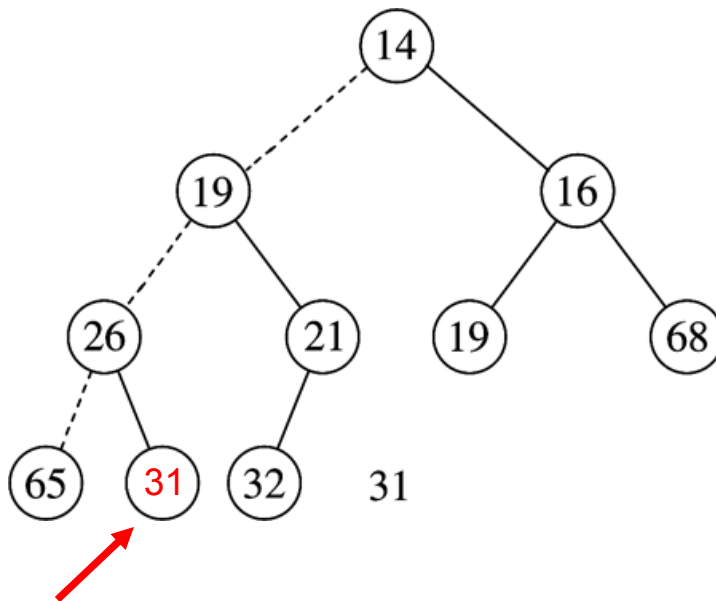


Is $31 > \min(65, 26)$?
• Yes - swap 31 with $\min(65, 26)$

Percolating down...

Percolating down...

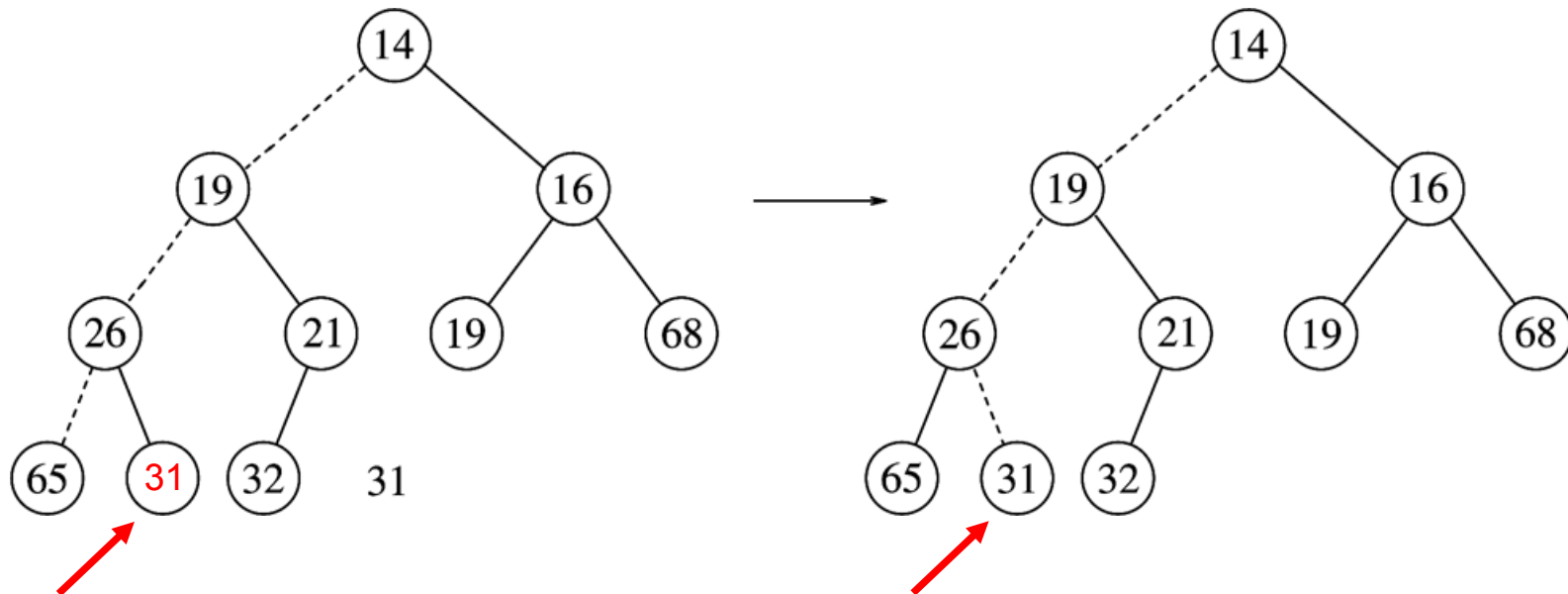
Heap DeleteMin: Example



Percolating down...

Percolating down...

Heap DeleteMin: Example



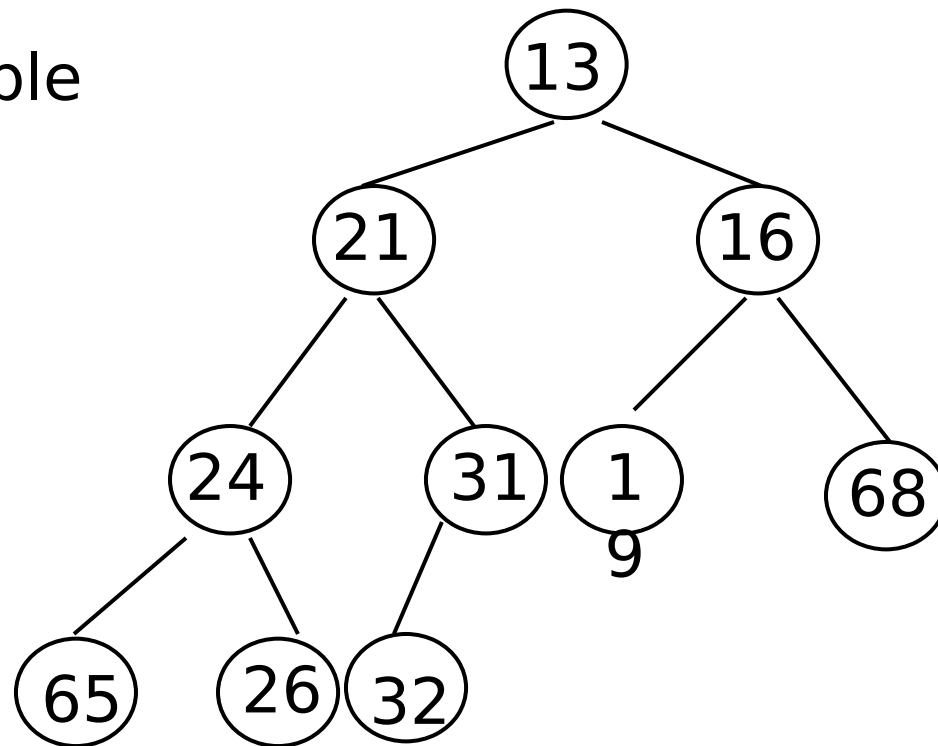
- ✓ Heap order prop
- ✓ Structure prop



Basic Heap Operations

Example

:

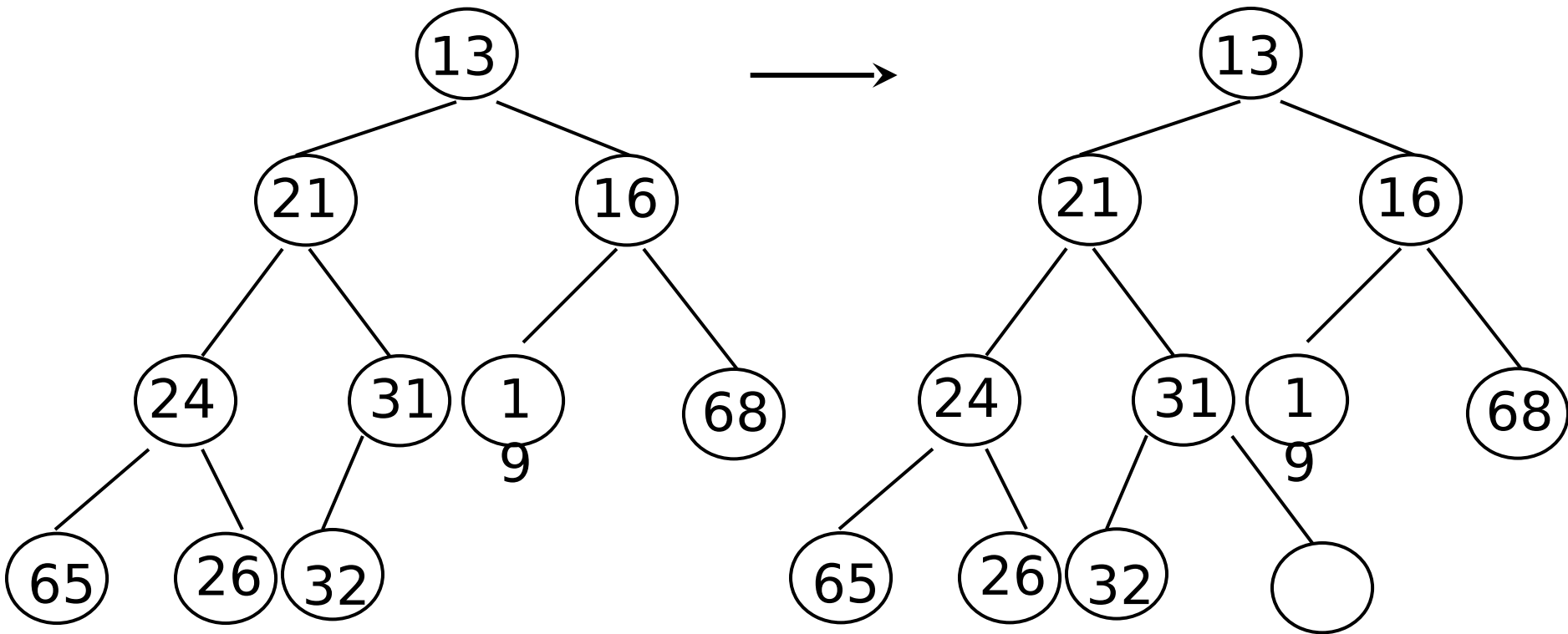


Original Tree

Insert 14

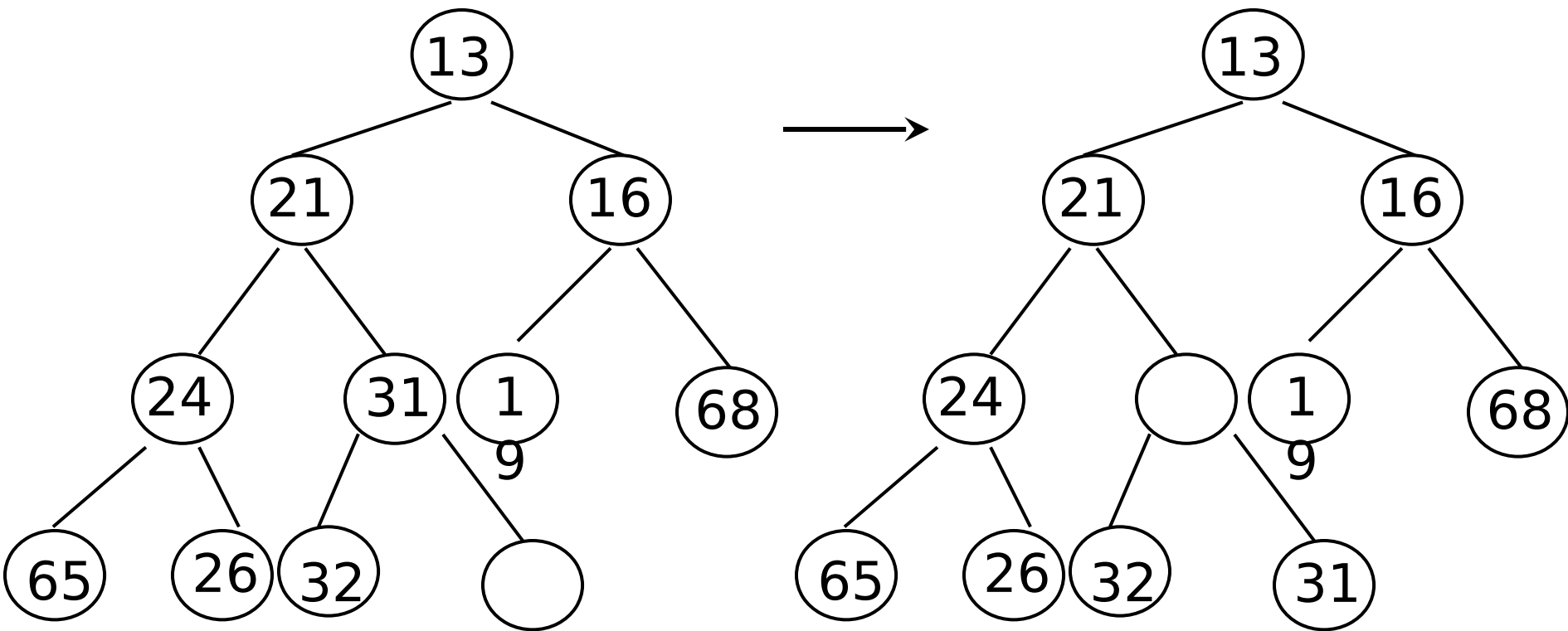


Basic Heap Operations



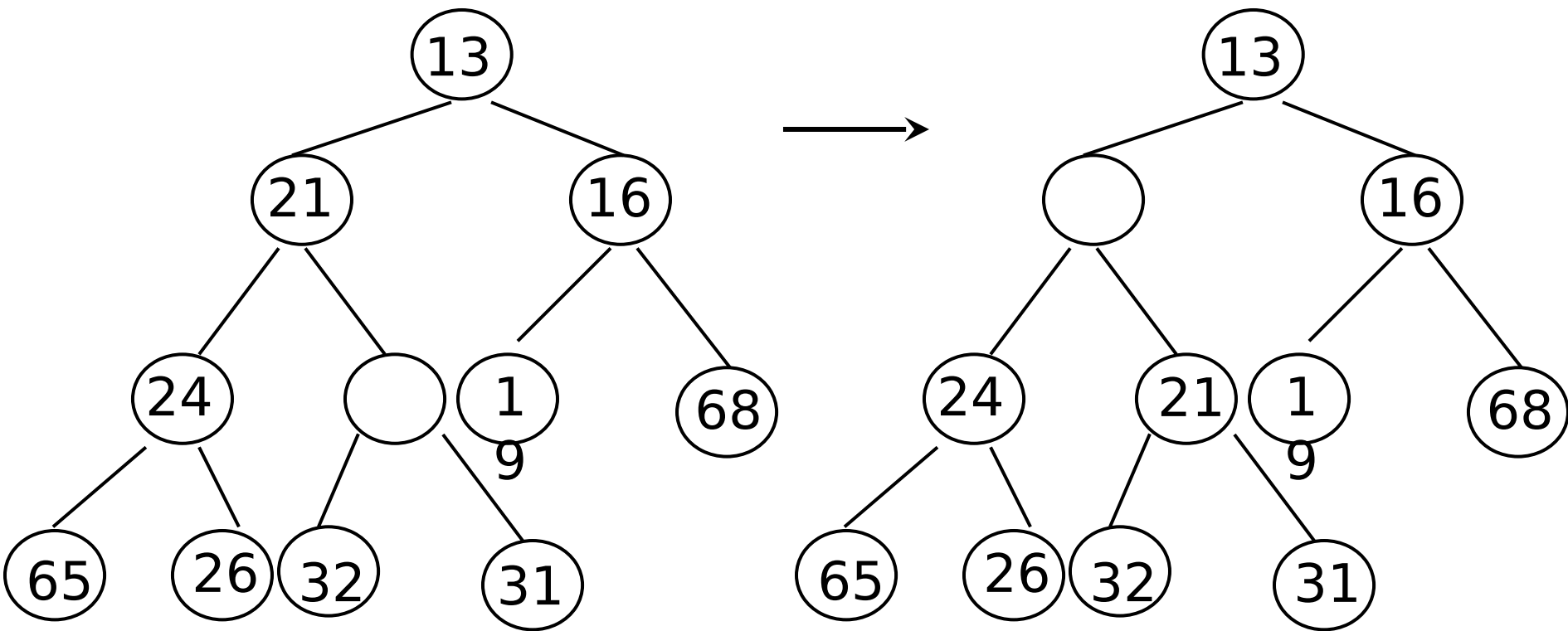


Basic Heap Operations



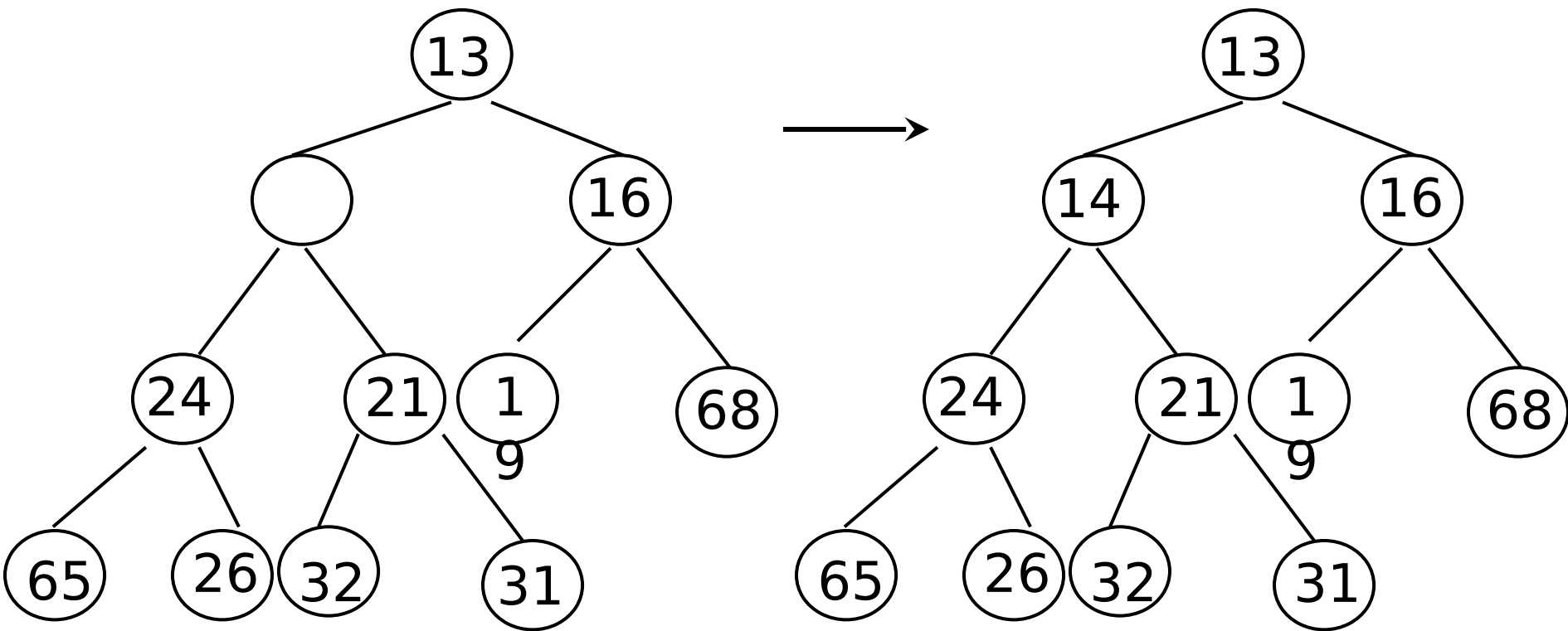


Basic Heap Operations



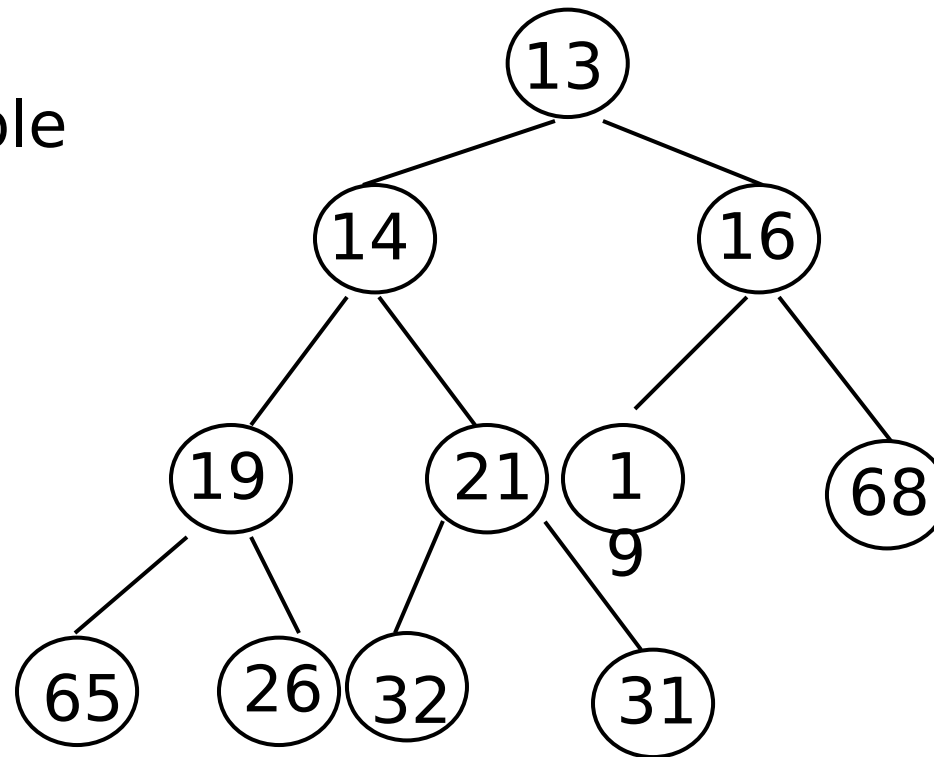


Basic Heap Operations



Basic Heap Operations

Example
:

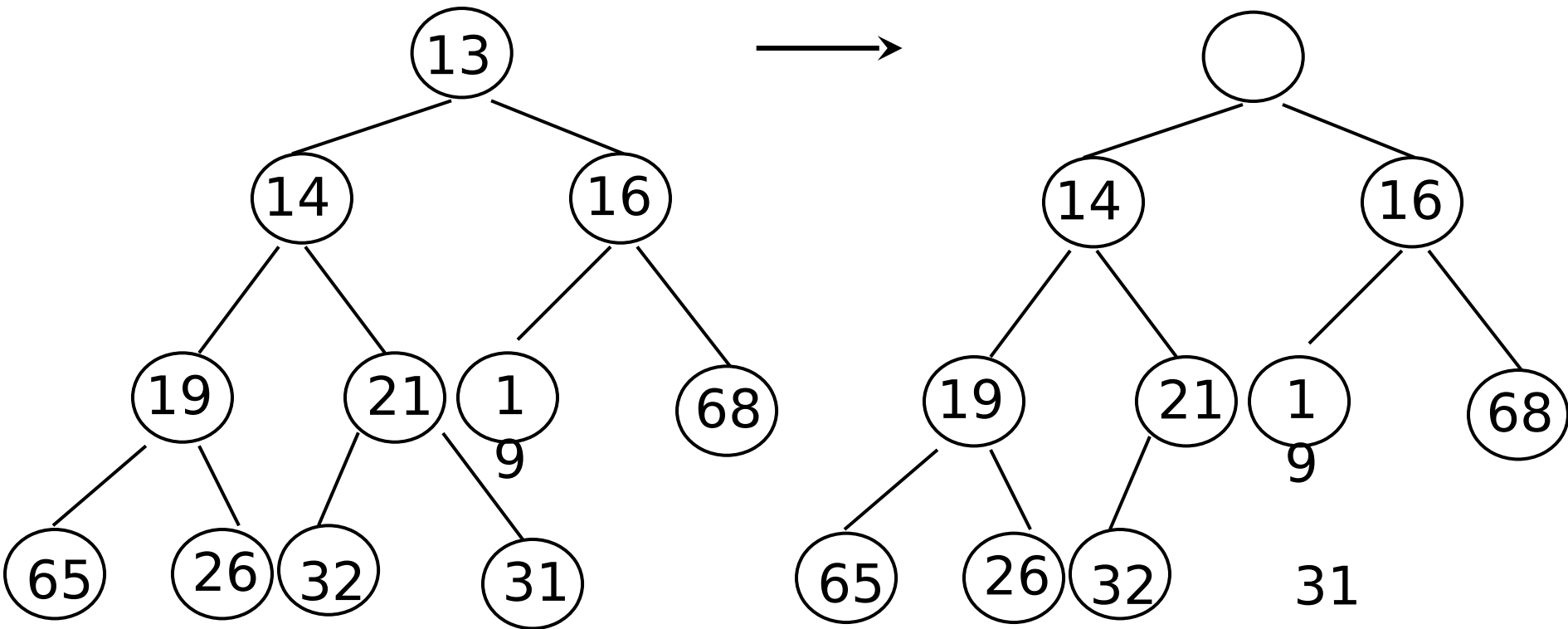


Delete the
minimum
key

Original Tree

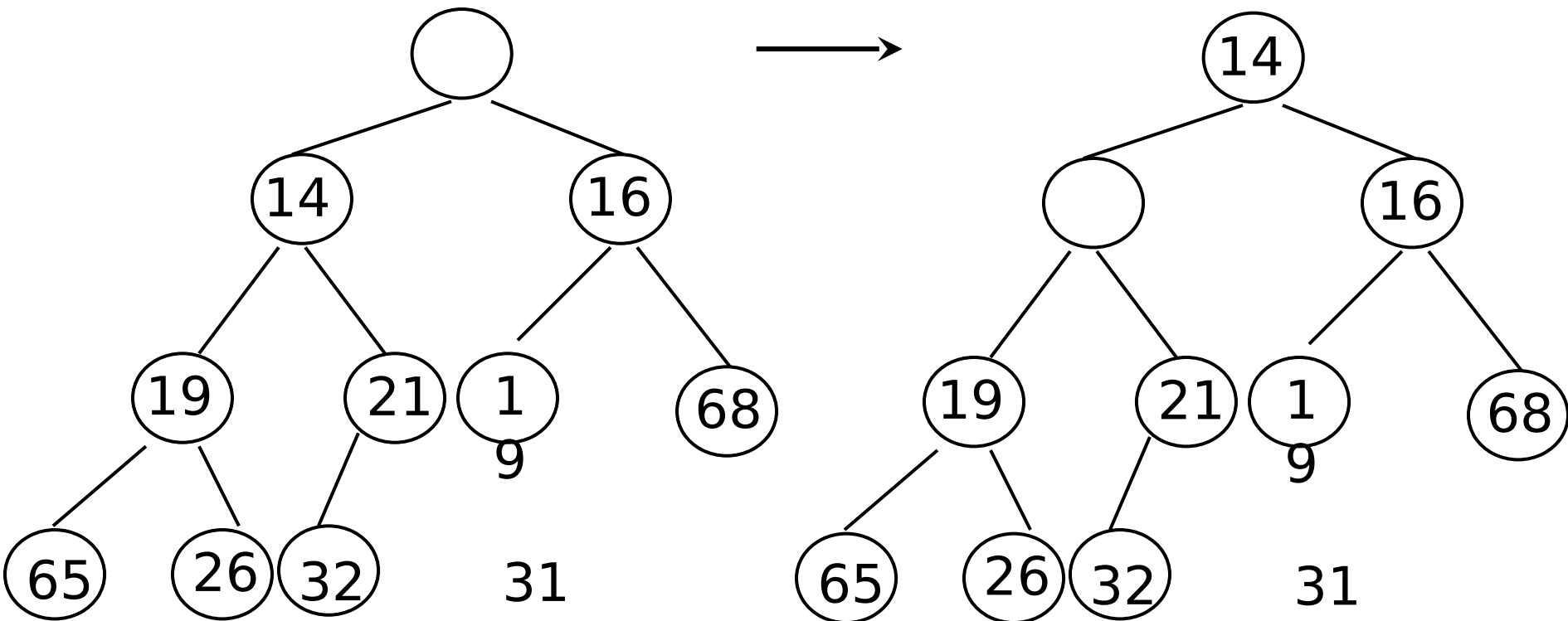


Basic Heap Operations



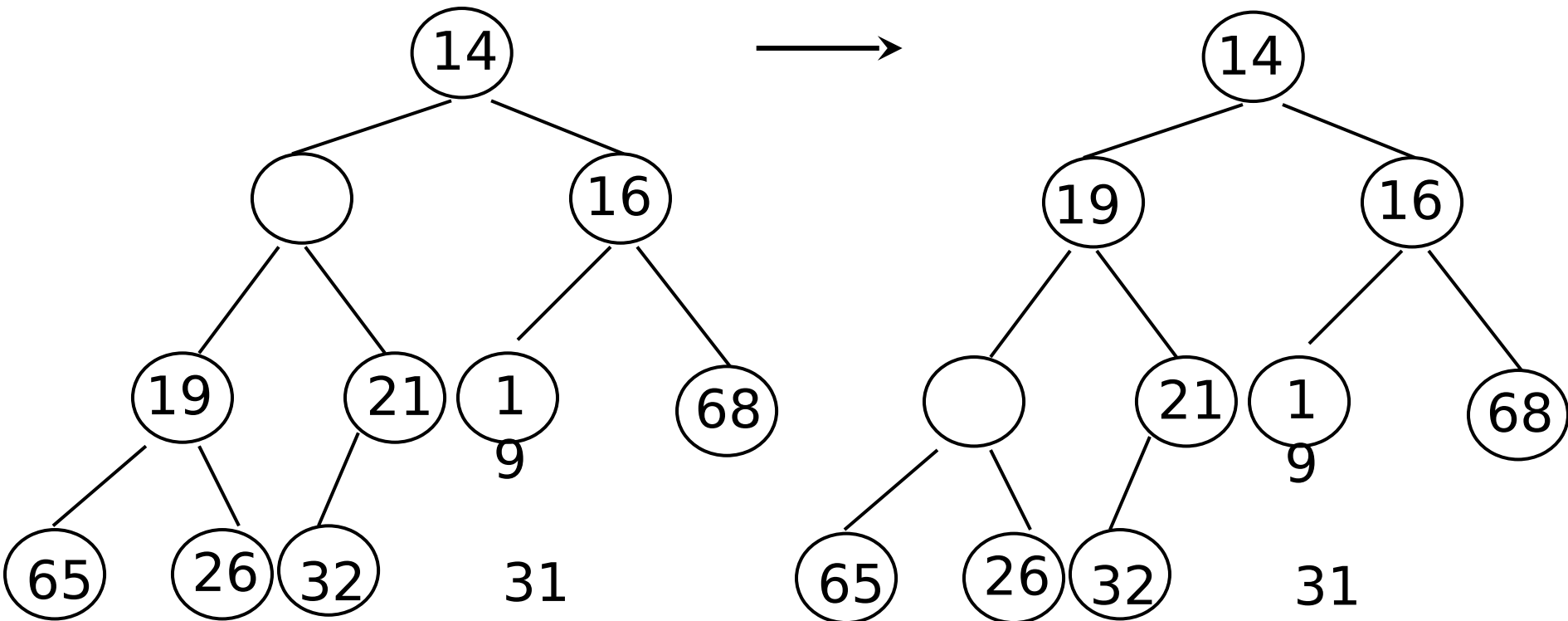


Basic Heap Operations



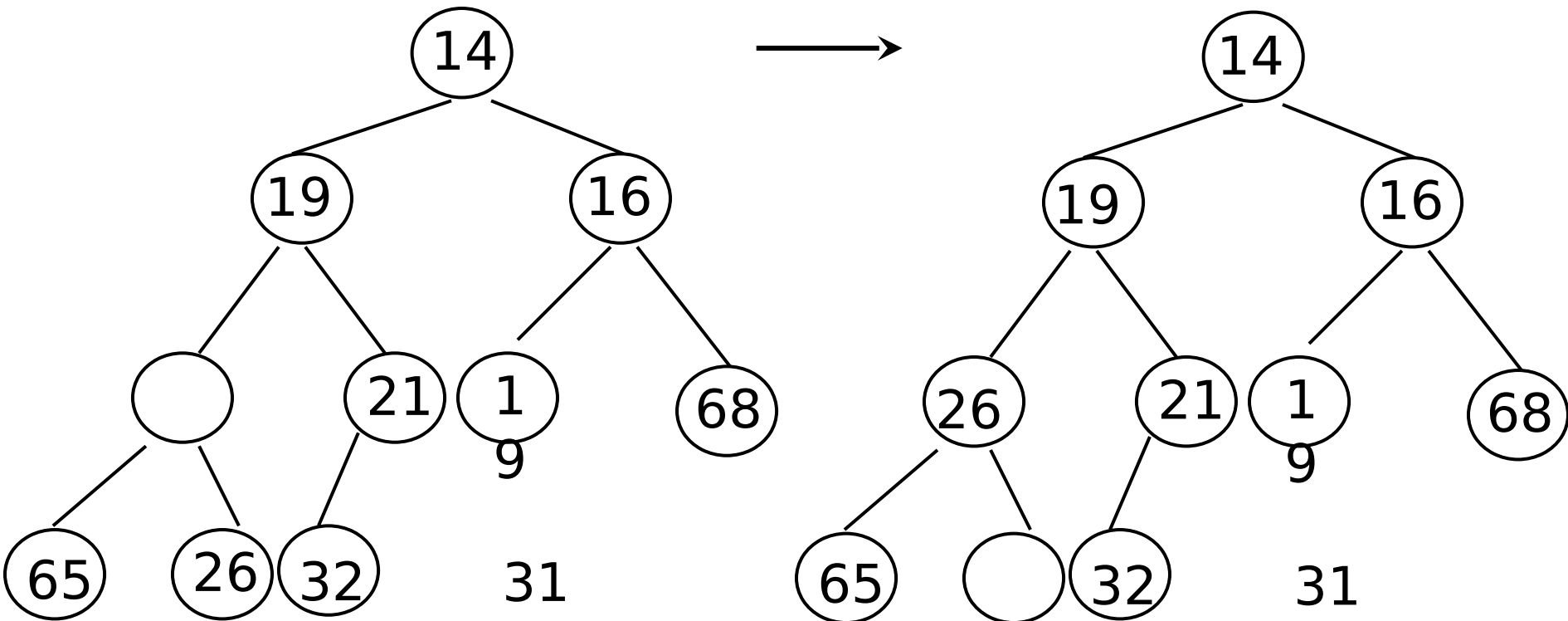


Basic Heap Operations



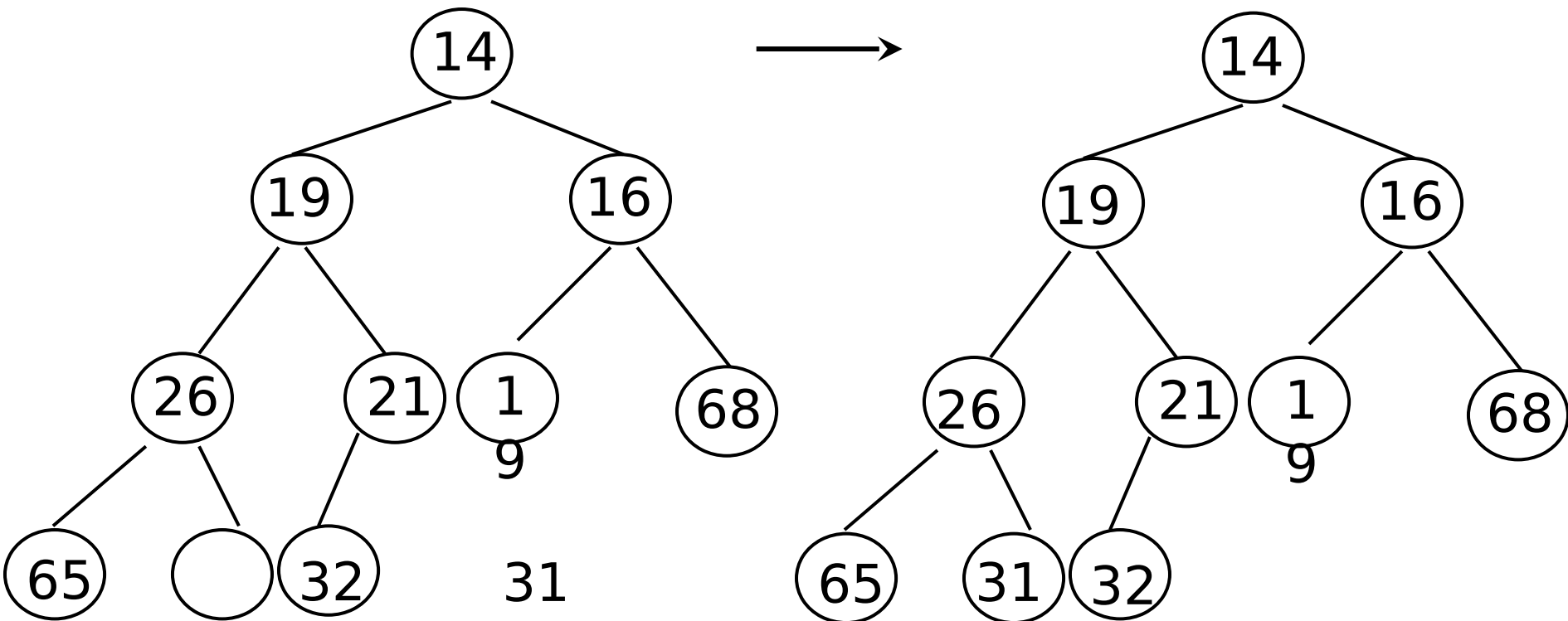


Basic Heap Operations





Basic Heap Operations



What are Heaps Useful For ?

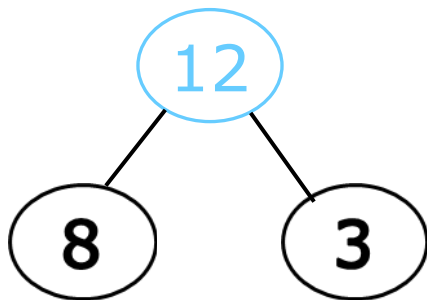
- **To implement priority queues**
- **Priority queue = a queue where all elements have a “priority” associated with them**
- **Remove in a priority queue removes the element with the smallest priority**
- **insert**
- **removeMin**

What are Heaps Useful For ?

- A stack is first in, last out
- A queue is first in, first out
- A priority queue is least-first-out
- The “smallest” element is the first one removed
- The definition of “smallest” is up to the programmer (for example, you might define it by implementing Comparator or Comparable)
- If there are several “smallest” elements, the implementer must decide which to remove first
- Remove any “smallest” element (don’t care which)
- Remove the first one added

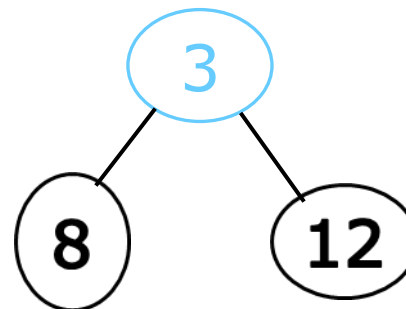
Heap Implementation of PQ

- A priority queue can be implemented as a heap
- In order to do this, we have to define the heap property
- In Heapsort, a node has the heap property if it is at least as large as its children
- For a priority queue, we will define a node to have the heap property if it is at least as small as its children (since we are using smaller numbers to represent higher priorities)



Heapsort: Blue node
has the heap property

09/25/22



Priority queue: Blue node
has the heap property

Heap 36

Heaps and Priority Queues

- We can use a heap to implement a priority queue
- We store a (key, element) item at each internal node
- We keep track of the position of the last node

