

Hashing

example

- Suppose you use open hashing (chaining) and the following keys are inserted:
- 5, 28, 19, 15, 20, 33, 12, 17, 10
- and $m = 9$.
- For simplicity, here we do not distinguish a key from its hashcode, so we assume $h(\text{key}) = \text{key}$.

example

- Suppose you use separate chaining and the following keys are inserted: 5, 28, 19, 15, 20, 33, 12, 17, 10 and $m = 9$.
- For simplicity, here we do not distinguish a key from its hashcode, so we assume $h(\text{key}) = \text{key}$.
- $H(5) = 5 \bmod 9 = 5$
- $H(28) = 28 \bmod 9 = 1$
- $H(19) = 19 \bmod 9 = 1$
- $H(15) = 15 \bmod 9 = 6$
- $H(20) = 20 \bmod 9 = 2$
- $H(33) = 33 \bmod 9 = 6$
- $H(12) = 12 \bmod 9 = 3$
- $H(17) = 17 \bmod 9 = 8$
- $H(10) = 10 \bmod 9 = 1$

example

- Suppose you use linear open addressing and the following keys are inserted: 5, 28, 19, 15, 20, 33, 12, 17, 10 and $m = 9$.
- For simplicity, here we do not distinguish a key from its hashcode, so we assume $h(\text{key}) = \text{key}$.
- $H(5) = 5 \bmod 9 = 5$
- $H(28) = 28 \bmod 9 = 1$
- $H(19) = 19 \bmod 9 = 1$
- $H(15) = 15 \bmod 9 = 6$
- $H(20) = 20 \bmod 9 = 2$
- $H(33) = 33 \bmod 9 = 6$
- $H(12) = 12 \bmod 9 = 3$
- $H(17) = 17 \bmod 9 = 8$
- $H(10) = 10 \bmod 9 = 1$

example

- Suppose you use linear open addressing and the following keys are inserted: 5, 28, 19, 15, 20, 33, 12, 17, 10 and $m = 9$.
- For simplicity, here we do not distinguish a key from its hashcode, so we assume $h(\text{key}) = \text{key}$.
- $H(5) = 5 \bmod 9 = 5$ at 5
- $H(28) = 28 \bmod 9 = 1$ at 1
- $H(19) = 19 \bmod 9 = 1$ at 2
- $H(15) = 15 \bmod 9 = 6$ at 6
- $H(20) = 20 \bmod 9 = 2$ at 3
- $H(33) = 33 \bmod 9 = 6$ at 7
- $H(12) = 12 \bmod 9 = 3$ at 4
- $H(17) = 17 \bmod 9 = 8$ at 8
- $H(10) = 10 \bmod 9 = 1$ at 0

example

- Suppose you use linear open addressing and the following keys are inserted: 5, 28, 19, 15, 20, 33, 12, 17, 10 and $m = 9$.
- For simplicity, here we do not distinguish a key from its hashcode, so we assume $h(\text{key}) = \text{key}$.
- $H(5) = 5 \bmod 9 = 5$ at 5
- $H(28) = 28 \bmod 9 = 1$ at 1
- $H(19) = 19 \bmod 9 = 1$ at 2
- $H(15) = 15 \bmod 9 = 6$ at 6
- $H(20) = 20 \bmod 9 = 2$ at 3 Now rehash to table size 19

example

- Suppose you use linear open addressing and the following keys are inserted: 5, 28, 19, 15, 20, 33, 12, 17, 10 and $m = 9$.
- For simplicity, here we do not distinguish a key from its hashcode, so we assume $h(\text{key}) = \text{key}$.
- $H(5) = 5 \bmod 19 = 5$ at 5
- $H(28) = 28 \bmod 19 = 9$ at 9
- $H(19) = 19 \bmod 19 = 0$ at 0
- $H(15) = 15 \bmod 19 = 15$ at 15
- $H(20) = 20 \bmod 19 = 1$ at 1
- $H(33) = 33 \bmod 19 = 14$ at 14
- $H(12) = 12 \bmod 19 = 12$ at 12
- $H(17) = 17 \bmod 19 = 17$ at 17
- $H(10) = 10 \bmod 19 = 10$ at 10

Exercises

- Table Size 15
 - Hash(k) = k
 - Insert 16, 7, 28, 31, 67,
38, 29, 73, 99, 43, 218
- 1) Linear Probe
 - 2) Quadratic Probe
 - 3) Chaining
 - 4) Double hash – $\%13 + 1$

General Idea



Example

