

# CSE-217: Theory of Computation

## REGULAR LANGUAGES

Lec Md Jakaria

Department of Computer Science and Engineering  
Military Institute of Science and Technology

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# Computational Model

What is a computer?



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**Computational Model:** An idealized computer



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**Computational Model:** An idealized computer

finite state machine *or* finite automaton.



# Automata



# Finite Automata

**Finite automata are good models for computers with an extremely limited amount of memory.**



# Finite Automata

**Finite automata are good models for computers with an extremely limited amount of memory.**

**What can a computer do with such a small memory?**



# Example - 1

Hopcroft, Motowani and Ullman: Figure 1.1

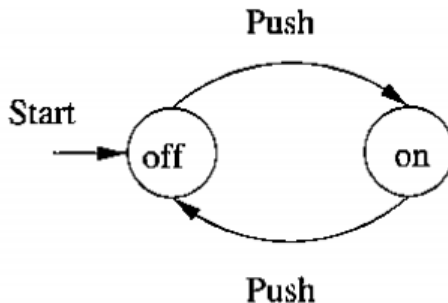


Figure: A finite automaton modeling an on/off switch





## Example - 2

Michael Sipser: Figure 1.1

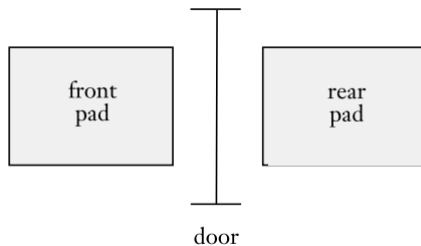


Figure: Top view of an automatic door



## Example - 2

Michael Sipser: Figure 1.2

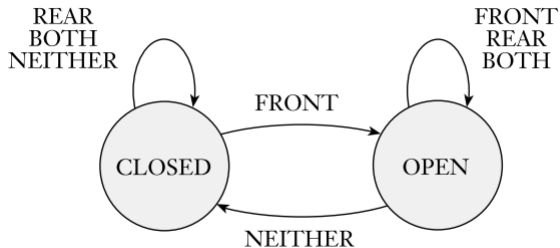


Figure: State diagram for an automatic door controller



## Example - 3

Michael Sipser: Figure 1.4

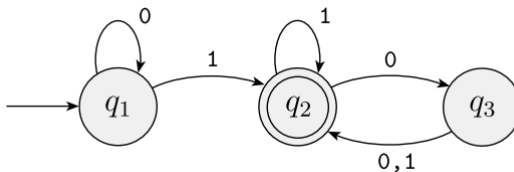


Figure: A finite automaton that has three states



# Finite Automata

## State diagram

- States
- Start State
- Accept State
- Transitions



# Automata

## Automata

- **Finite Automata**
- **Infinite Automata**



# Automata

## Automata

- **Finite Automata**
- **Infinite Automata**

## Finite Automata

- Deterministic
- Non-deterministic



# Finite Automata



# Formal Definition

## DEFINITION 1.5

A *finite automaton* is a 5-tuple  $(Q, \Sigma, \delta, q_0, F)$ , where

1.  $Q$  is a finite set called the *states*,
2.  $\Sigma$  is a finite set called the *alphabet*,
3.  $\delta: Q \times \Sigma \rightarrow Q$  is the *transition function*,<sup>1</sup>
4.  $q_0 \in Q$  is the *start state*, and
5.  $F \subseteq Q$  is the *set of accept states*.<sup>2</sup>





# Formal Definition

## Language

- $A$  is the set of all strings that machine  $M$  accepts.
- We say that  $A$  is the language of machine  $M$ .
- Write  $L(M) = A$ .
- We say that  $M$  recognizes  $A$  or that  $M$  accepts  $A$ .



## Example - 3 continued

Michael Sipser: Figure 1.4

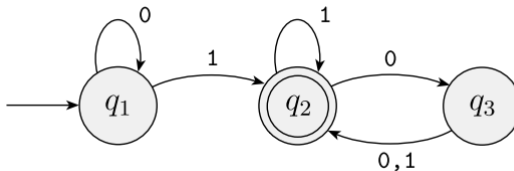


Figure: A finite automaton called  $M_1$  that has three states



## Example - 3 continued

We can describe  $M_1$  formally by writing  $M_1 = (Q, \Sigma, \delta, q_1, F)$ , where

1.  $Q = \{q_1, q_2, q_3\}$ ,
2.  $\Sigma = \{0, 1\}$ ,
3.  $\delta$  is described as

	0	1
$q_1$	$q_1$	$q_2$
$q_2$	$q_3$	$q_2$
$q_3$	$q_2$	$q_2$ ,

4.  $q_1$  is the start state, and
5.  $F = \{q_2\}$ .



## Example - 3 continued

$A = \{w \mid w \text{ contains at least one 1 and}$   
an even number of 0s follow the last 1}.

Then  $L(M_1) = A$ , or equivalently,  $M_1$  recognizes  $A$ .



## Example - 4

Michael Sipser: Figure 1.9

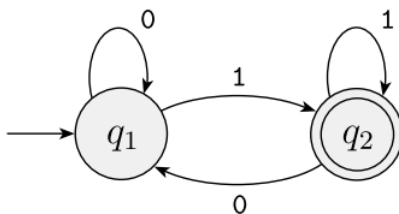


Figure: State diagram of the two-state finite automaton  $M_2$



# Example - 5

Michael Sipser: Figure 1.10

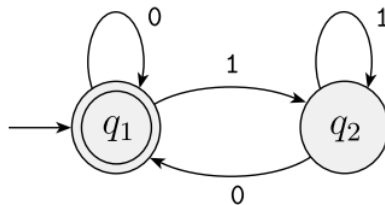


Figure: State diagram of the two-state finite automaton  $M_3$



# Example - 5

Michael Sipser: Figure 1.11

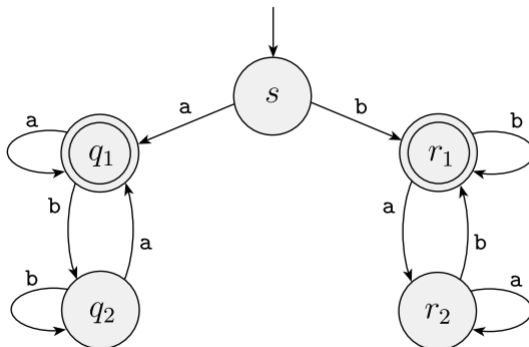


Figure: State diagram of the two-state finite automaton  $M_4$



# Example - 6

Michael Sipser: Figure 1.12

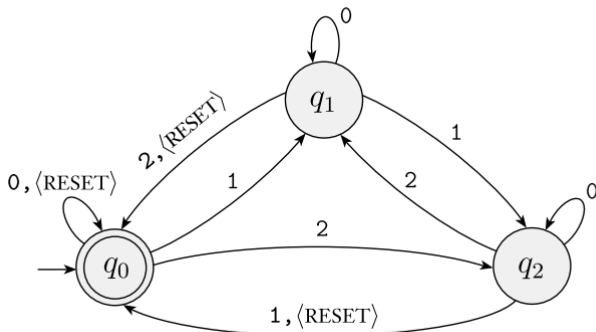


Figure: State diagram of the two-state finite automaton  $M_5$

