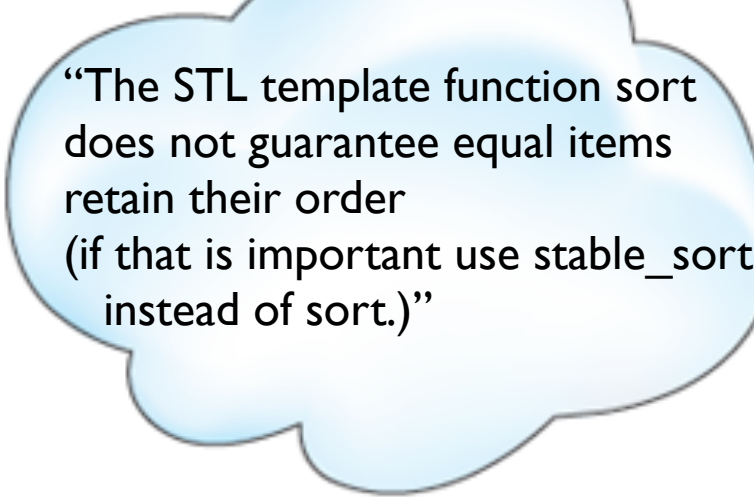# Sorting

Insertion Sort, Merge Sort, & Quick Sort

Comparison Based Sorting, only operations are

< operator

assignment operator

STL Sorting

"The STL template function sort does not guarantee equal items retain their order
(if that is important use stable_sort instead of sort.)"

```
void sort( Iterator begin, Iterator end );
void sort( Iterator begin, Iterator end, Comparator cmp );



sort( v.begin( ), v.end( ) );
sort( v.begin( ), v.end( ), greater<int>( ) );
sort( v.begin( ), v.begin( ) + ( v.end( ) - v.begin( ) ) / 2 );
```

function template

std::**sort**

<algorithm>

default (1)	`template <class RandomAccessIterator>` `  void sort (RandomAccessIterator first, RandomAccessIterator last);`
custom (2)	`template <class RandomAccessIterator, class Compare>` `  void sort (RandomAccessIterator first, RandomAccessIterator last, Compare comp);`

Sort elements in range

Sorts the elements in the range `[first,last)` into ascending order.

The elements are compared using `operator<` for the first version, and *comp* for the second.

Equivalent elements are not guaranteed to keep their original relative order (see `stable_sort`).

Parameters

`first, last`

> Random-access iterators to the initial and final positions of the sequence to be sorted. The range used is `[first,last)`, which contains all the elements between *first* and *last*, including the element pointed by *first* but not the element pointed by *last*.
> `RandomAccessIterator` shall point to a type for which `swap` is properly defined and which is both *move-constructible* and *move-assignable*.

`comp`

> Binary function that accepts two elements in the range as arguments, and returns a value convertible to `bool`. The value returned indicates whether the element passed as first argument is considered to go before the second in the specific *strict weak ordering* it defines.
> The function shall not modify any of its arguments.
> This can either be a function pointer or a function object.

Return value

http://www.cplusplus.com/reference/algorithm/sort/

Why Sort?

Sorted data is easier to work with. Historically 25% of computer time was spent on sorting!

grouping

Finding which pair of items are closest to each other is easy if the items are sorted! O(n) time

Testing if all the items are unique is easy if sorted!

Finding which item occurs most frequently is much faster if the items are sorted! O(n)

preprocessing:

Searching through 100,000,000 items (under the assumption of each item is equally likely to be searched for) takes ~50,000,000 comparisons if the items are unsorted vs ~ 27 comparisons using binary search. $O(n)$ vs $O(\log(n))$!

Finding the smallest, largest, K'th largest is O(1) if sorted.

subroutine in other algorithms

Optimal compression of text. Finding an optimal Huffman encoding starts by knowing the character frequency and sorting the characters by their frequencies.

Kruskal's algorithm for finding an optimal spanning tree starts by sorting the edge weights.

How can we sort the items?

A First Sort: Insertion Sort

Insertion Sort

17	20	43	25	4	72	15

a[0..p-1] sorted   rest of array

Insert element a[p] in a[0..p-1], sliding larger elements over to make room:

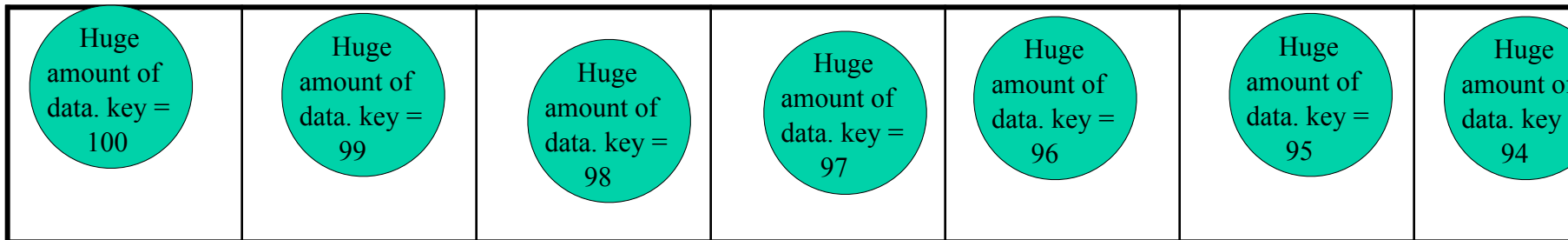| 17	20	25	43	4	72	15

a[0..p] sorted   rest of array

```cpp
// insertionSort: sort items in array a
// Comparable: must have copy constructor, operator=,
//   and operator<
template <class Comparable>
void insertionSort( vector<Comparable> & a )
{
    int j;

    for( int p = 1; p < a.size( ); p++ )
    {
        Comparable tmp = std::move(a[ p ]);
        for( j = p; j > 0 && tmp < a[ j - 1 ]; j-- )
            a[ j ] = std::move(a[ j - 1 ]);
        a[ j ] = std::move(tmp);
    }
}
```

Huge amount of data. key = 100	Huge amount of data. key = 99	Huge amount of data. key = 98	Huge amount of data. key = 97	Huge amount of data. key = 96	Huge amount of data. key = 95	Huge amount of data. key = 94

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   }
}
```
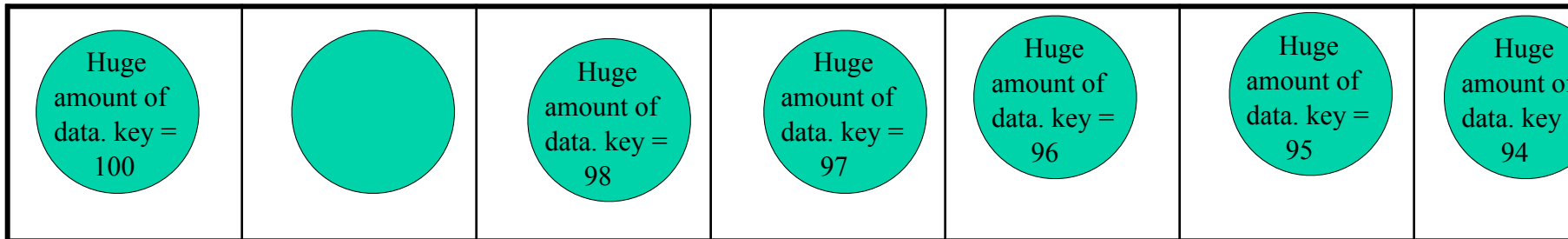
tmp

Huge amount of data. key = 99

| Huge amount of data. key = 100 | | Huge amount of data. key = 98 | Huge amount of data. key = 97 | Huge amount of data. key = 96 | Huge amount of data. key = 95 | Huge amount of data. key = 94 |

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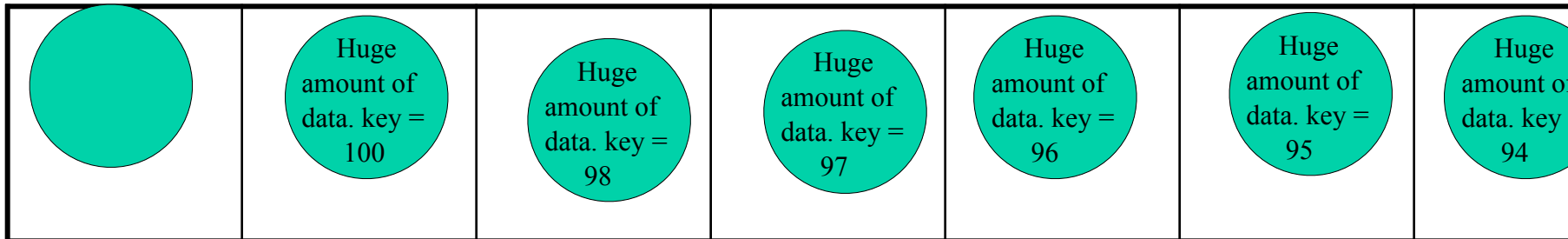
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| | Huge amount of data. key = 100 | Huge amount of data. key = 98 | Huge amount of data. key = 97 | Huge amount of data. key = 96 | Huge amount of data. key = 95 | Huge amount of data. key = 94 |

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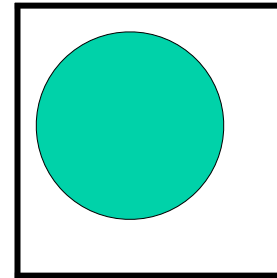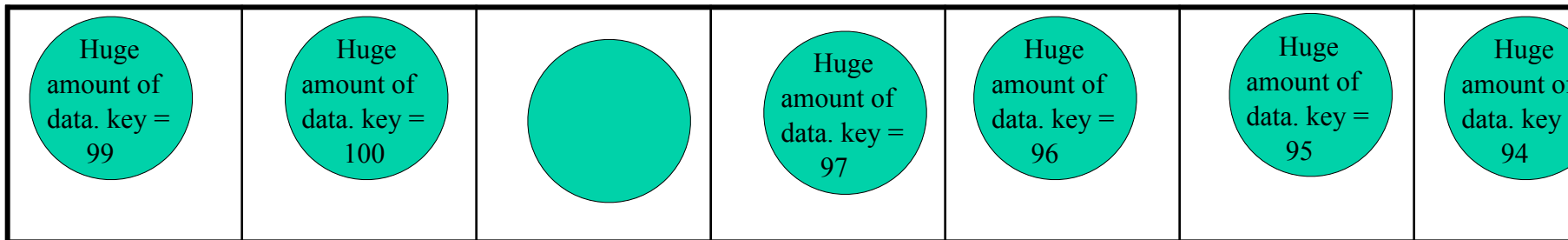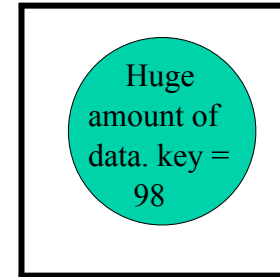
tmp

| Huge amount of data. key = 99 | Huge amount of data. key = 100 | Huge amount of data. key = 98 | Huge amount of data. key = 97 | Huge amount of data. key = 96 | Huge amount of data. key = 95 | Huge amount of data. key = 94 |

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        a[ j ] = std::move(tmp);
    }
}
```

tmp

Huge amount of data. key = 98

| Huge amount of data. key = 99 | Huge amount of data. key = 100 | | Huge amount of data. key = 97 | Huge amount of data. key = 96 | Huge amount of data. key = 95 | Huge amount of data. key = 94 |

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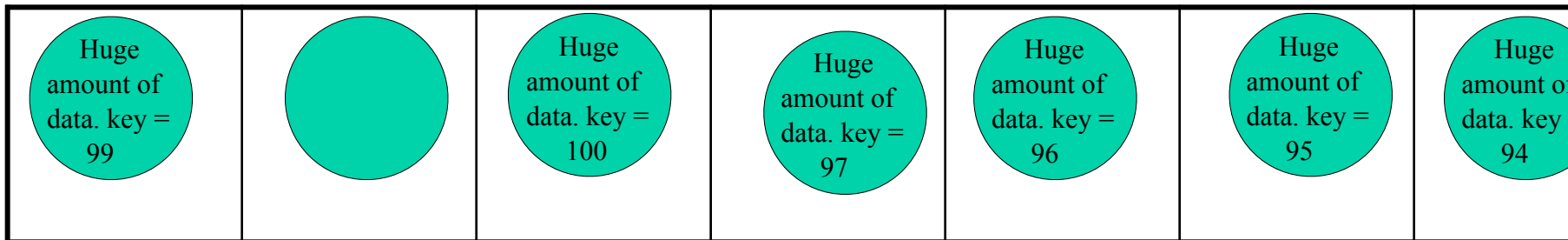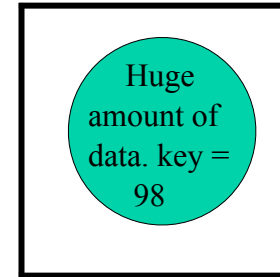
tmp

| Huge amount of data. key = 98 |

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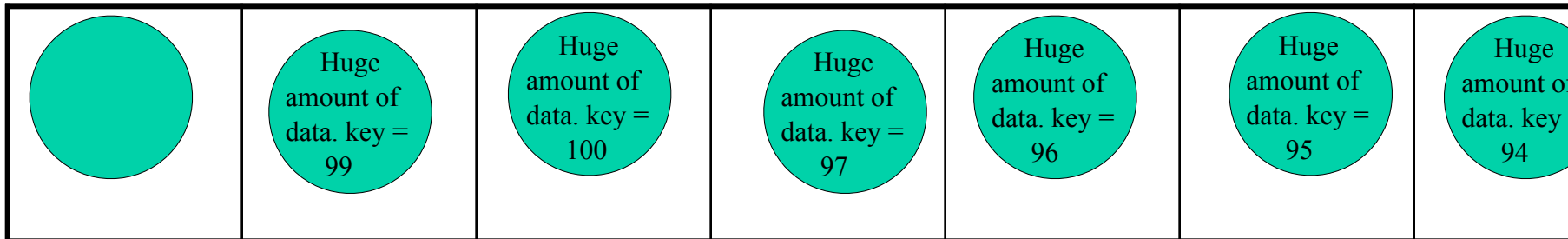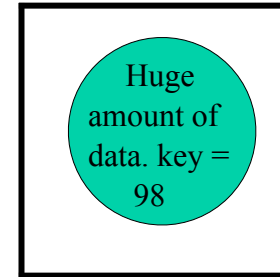
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    int j;

    for( int p = 1; p < a.size( ); p++ )
    {
        Comparable tmp = std::move(a[ p ]);
        for( j = p; j > 0 && tmp < a[ j - 1 ]; j-- )
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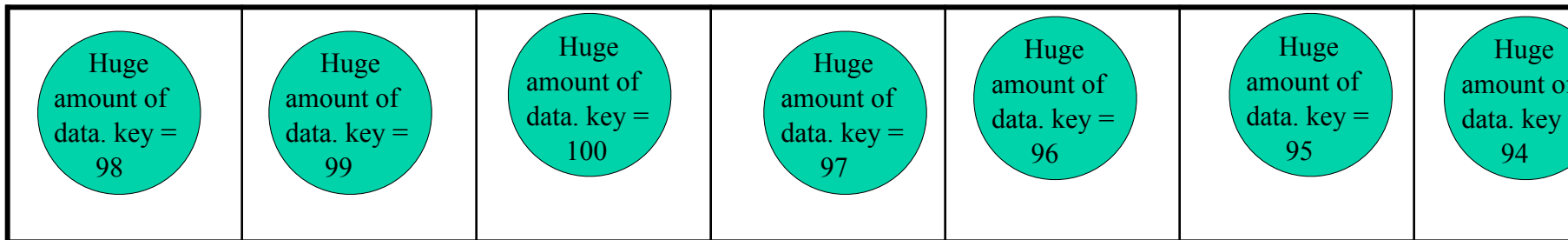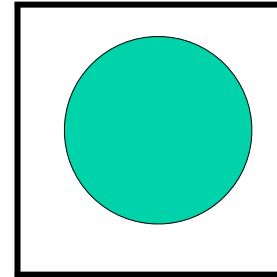
tmp

| Huge amount of data. key = 98 | Huge amount of data. key = 99 | Huge amount of data. key = 100 | Huge amount of data. key = 97 | Huge amount of data. key = 96 | Huge amount of data. key = 95 | Huge amount of data. key = 94 |

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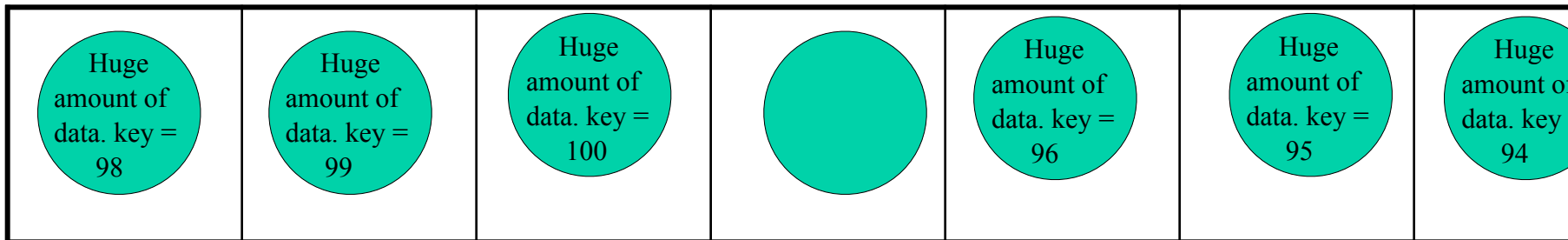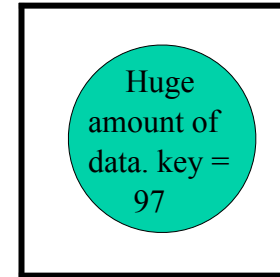
tmp

Huge amount of data. key = 97

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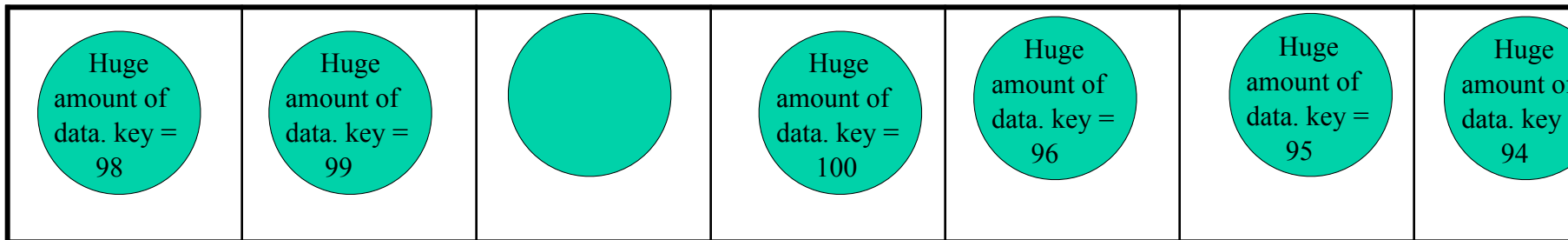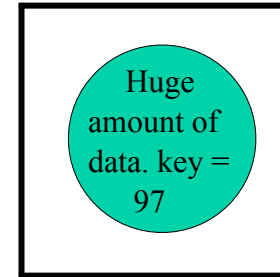
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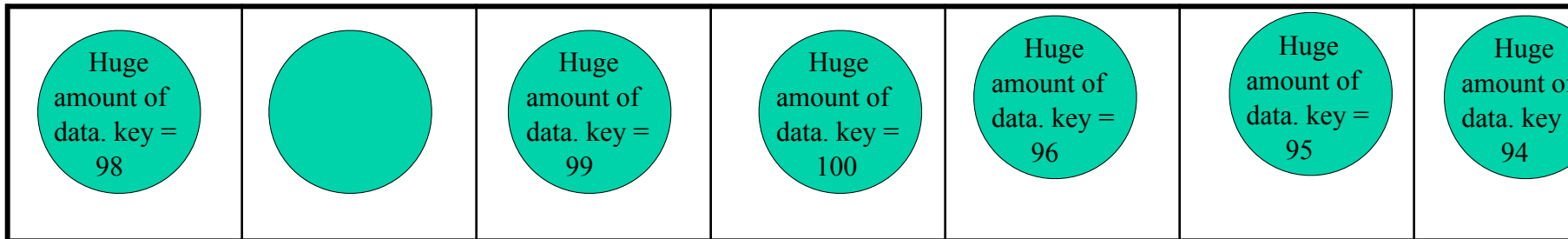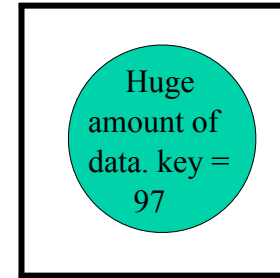
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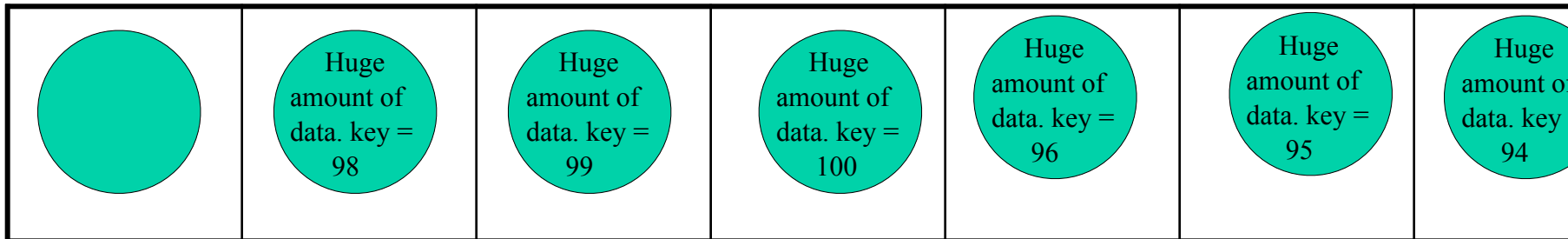
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   int j;

   for( int p = 1; p < a.size( ); p++ )
   {
      Comparable tmp = std::move(a[ p ]);
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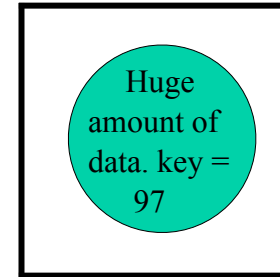
tmp

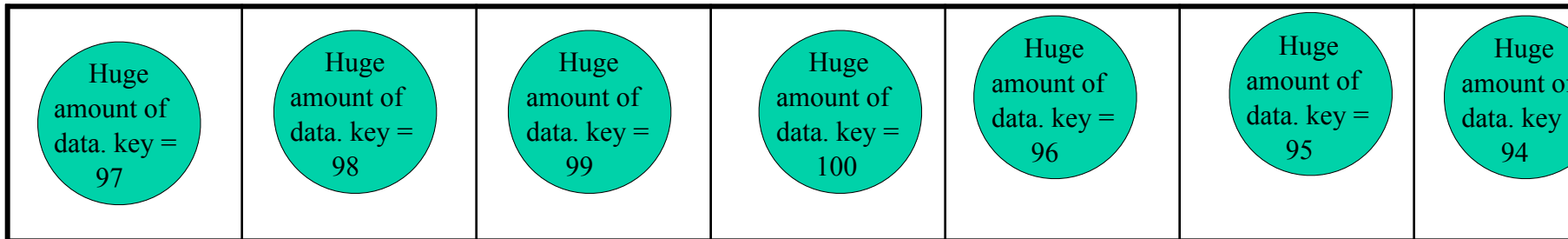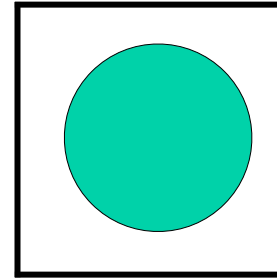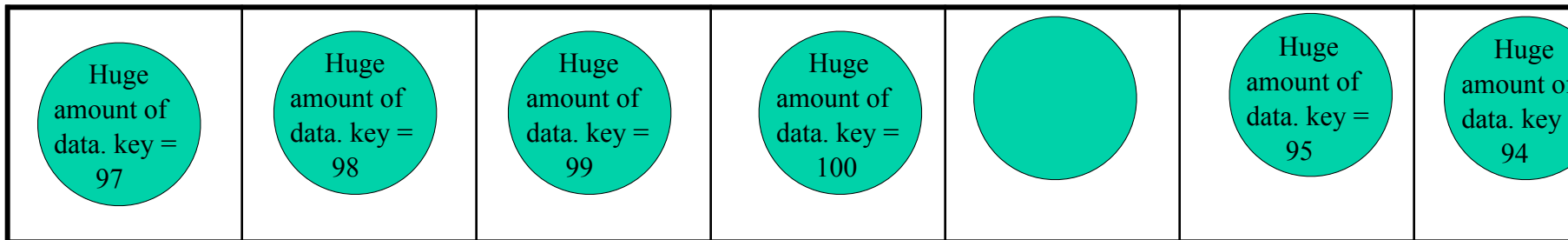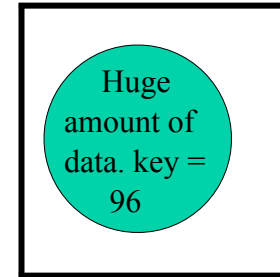| Huge amount of data. key = 97 | Huge amount of data. key = 98 | Huge amount of data. key = 99 | Huge amount of data. key = 100 | Huge amount of data. key = 96 | Huge amount of data. key = 95 | Huge amount of data. key = 94 |

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{
   int j;

   for( int p = 1; p < a.size( ); p++ )
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      a[ j ] = std::move(tmp);
   }
}
```

tmp

Huge amount of data. key = 96

Huge amount of data. key = 97

Huge amount of data. key = 98

Huge amount of data. key = 99

Huge amount of data. key = 100

Huge amount of data. key = 95

Huge amount of data. key = 94

Running time?

STL Style Insertion Sort

- Range of items to sort [begin, end)

- Array access using iterators

- Use auto to deduce type

- Use a functor to compare items

17	20	25	43	4	72	15

```cpp
template <class Comparable>
void insertionSort( vector<Comparable> & a )
{
    int j;

    for( int p = 1; p < a.size( ); p++ )
    {
        Comparable tmp = std::move(a[ p ]);
        for( j = p; j > 0 && tmp < a[ j - 1 ]; j-- )
            a[ j ] = std::move(a[ j - 1 ]);

        a[ j ] = std::move(tmp);
    }
}
```

```cpp
template <typename Iterator, typename Comparator>
void insertionSort( Iterator begin, Iterator end, Comparator lessThan )
{
    if( begin == end )
        return;

    Iterator j;

    for( Iterator p = begin+1; p != end; ++p )
    {
        auto tmp =  std::move(*p) ;
        for( j = p; j != begin && lessThan( tmp, *( j-1 ) ); --j )
            *j = std::move(*(j-1));

        *j = std::move(tmp);
    }
}
```

Lets do it again! This time with two parameters and <

void insertionSort(Iterator begin, Iterator end);

Using:
void insertionSort(Iterator begin, Iterator end, Comparator cmp);

```cpp
template <class Object>
class less                    An STL Functor
{ public:
    bool operator()(const Object& lhs, const Object& rhs) const
        {return lhs < rhs;}
};

/*
* The two-parameter version calls the three-parameter version,
* using C++11 decltype
*/
template <typename Iterator>
void insertionSort( const Iterator & begin, const Iterator & end )
{

    insertionSort( begin, end, less< decltype(*begin)>( ) );

}
```

"In the C++ programming language, **decltype** is an operator for querying the type of an expression. It was introduced in the current version of the C++ standard, C++11. Its primary intended use is in generic programming, where it is often difficult, or even impossible, to express types that depend on template parameters."
http://en.wikipedia.org/wiki/Decltype

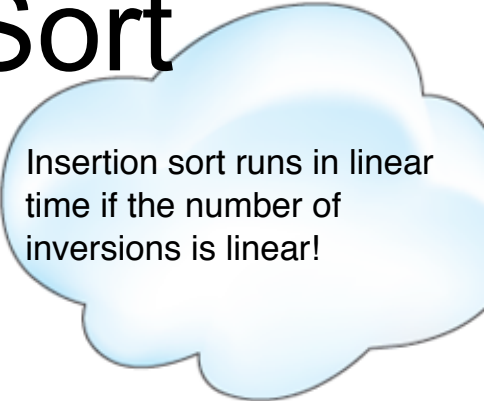Operator yields the declared type of the expression

Inversions

- an inversion is a pair of elements out of order

- e.g. [8, 5, 9, 2, 6, 3] has 10 inversions: (8,5), (8,2), (8,6), (8,3), (5,2), (5,3), (9,2), (9,6), (9,3), (6,3)

- # of inversions is the number of times we move an item

- in our analysis we will assume there are no duplicate items

Average Case

- Given an array {8, 5, 9, 2, 6, 3} and an array in reverse order {3, 6, 2, 9, 5, 8}

- For every pair of numbers in an array, (x,y) it is an inversion in exactly one of the arrays. E.g. (8, 5) in an inversion in the first array and not the second.

- If we count the number of inversions in both arrays it equals the number of pairs, n(n-1)/2.

- An average array there for has n(n-1)/4

Analysis of Insertion Sort
$O(I + n)$

Insertion sort runs in linear time if the number of inversions is linear!

Number of inversions:

- Worst case:

  $1+ 2 + ... + n -1 = O(n^2)$

- Average case: similar ... $O(n^2)$

- Best case (Array is already sorted)
  - inner loop iterates 0 times for each p (1 test of inner for loop)
  - 0 inversions so running time is $O(n)$ [!]

- Interesting special case: array is almost sorted ... similar to best case

Merge Sort
Divide and Conquer

Merge Sort

43	37	20	25	4	72	15	19

DIVIDE IN HALF

43	37	20	25

4	72	15	19

SORT EACH HALF

20	25	37	43

4	15	19	72

↑ left

↑ right

MERGE

4	15	19	20	25	37	43	72

Check out:

https://en.wikipedia.org/wiki/Merge_sort#/media/File:Merge-sort-example-300px.gif

http://www.sorting-algorithms.com/

https://www.youtube.com/user/AlgoRythmics

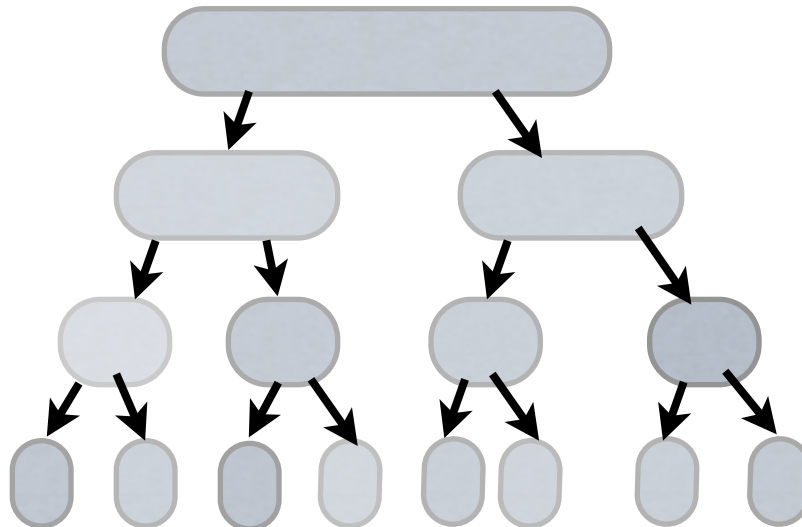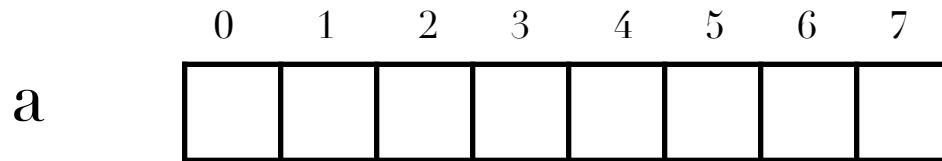https://www.youtube.com/watch?v=kPRA0W1kECg

http://sortvis.org/

Merge Sort

- Divide array into two (almost) equal pieces
- Conquer:
  - sort each half recursively
  - merge the sorted arrays
- code
- Analysis:
  - log n stages, each of which is linear
  - O(n log n)

- Divide array into two (almost) equal pieces
- Conquer:
  - sort each half recursively
  - merge the sorted arrays

$$0 \quad 1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 6 \quad 7$$

a

The Driver

```cpp
// Mergesort algorithm (driver).
template<class Comparable>
void mergeSort( vector<Comparable> & a )
{
    vector<Comparable> tmpArray( a.size( ) );

    mergeSort( a, tmpArray, 0, a.size( ) - 1 );
}

void main()
{
    ...
    mergeSort(a);
    ...
}
```
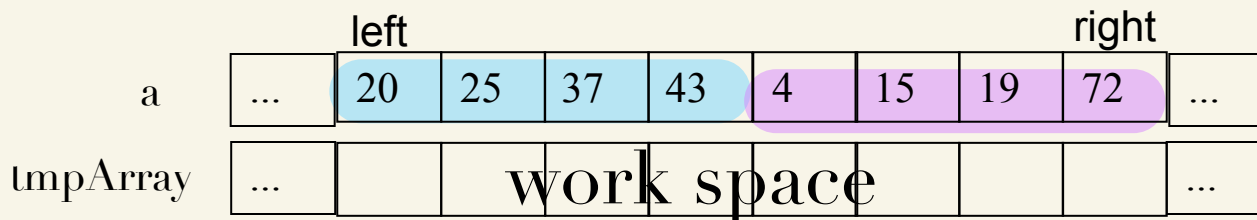
```cpp
template <class Comparable>
void mergeSort( vector<Comparable> & a,
        vector<Comparable> & tmpArray, int left, int right )
{
   if( left < right )
   {

     int center = ( left + right ) / 2;
      mergeSort( a, tmpArray, left, center );
      mergeSort( a, tmpArray, center + 1, right );
     myMerge( a, tmpArray, left, center + 1, right );

   }

}
```

		left						right	
a	...	20	25	37	43	4	15	19	72	...

tmpArray	...			work space					...

Merge

- Given two sorted sub arrays, combine them into a single sorted array
    - ~ Use auxiliary array
    - ~ successively compare smallest elements
    - ~ when one runs out, copy the rest of the other
- code for merge
- scans each element once: O(n) time
- also uses O(n) space

```cpp
template <class Comparable>
void myMerge( vector<Comparable> & a, vector<Comparable>& tmpArray,
int leftPos, int rightPos, int rightEnd )
{
    int leftEnd = rightPos - 1;
    int tmpPos = leftPos;
    int numElements = rightEnd - leftPos + 1;


// Main loop
  while( leftPos <= leftEnd && rightPos <= rightEnd )
    if( a[ leftPos ] <= a[ rightPos ] )
      tmpArray[ tmpPos++ ] = std::move(a[ leftPos++ ]);
    else
      tmpArray[ tmpPos++ ] = std::move(a[ rightPos++ ]);

  while( leftPos <= leftEnd )    // Copy rest of first half
    tmpArray[ tmpPos++ ] = std::move(a[ leftPos++ ]);

  while( rightPos <= rightEnd )  // Copy rest of right half
    tmpArray[ tmpPos++ ] =std::move(a[ rightPos++ ]);

    // Copy tmpArray back
  for( int i = 0; i < numElements; i++, - -rightEnd )
    a[ rightEnd ] = std::move(tmpArray[ rightEnd ];)
}
```

a

| | 15 | 25 | 37 | 43 | 4 | 12 | 18 | 19 | | | | | | | | | |

tmparray

| | | | | | | | | | | | | | | | | | |

```cpp
template <class Comparable>
void myMerge( vector<Comparable> & a, vector<Comparable>& tmpArray,
int leftPos, int rightPos, int rightEnd )
{
    int leftEnd = rightPos - 1;
    int tmpPos = leftPos;
    int numElements = rightEnd - leftPos + 1;


// Main loop
  while( leftPos <= leftEnd && rightPos <= rightEnd )
    if( a[ leftPos ] <= a[ rightPos ] )
      tmpArray[ tmpPos++ ] = std::move(a[ leftPos++ ]);
    else
      tmpArray[ tmpPos++ ] = std::move(a[ rightPos++ ]);

  while( leftPos <= leftEnd )    // Copy rest of first half
    tmpArray[ tmpPos++ ] = std::move(a[ leftPos++ ]);

  while( rightPos <= rightEnd )  // Copy rest of right half
    tmpArray[ tmpPos++ ] = std::::(move(a[ rightPos++ ]);

    // Copy tmpArray back
  for( int i = 0; i < numElements; i++, - -rightEnd )
    a[ rightEnd ] = std::move(tmpArray[ rightEnd ]);
}
```

a

tmparray

	4	12	15	18	19	25	37	43													

a

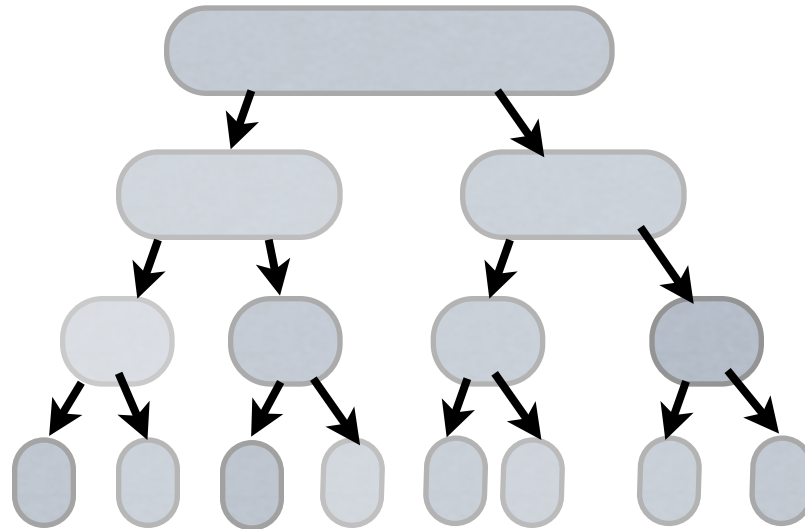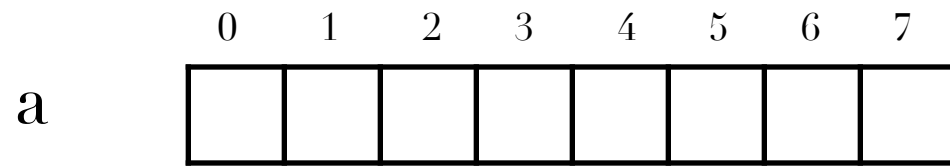| 43	37	20	25	4	72	15	19

tmpArray

mergeSort(a, tmpArray, 0, 0)

mergeSort(a, tmpArray, 0, 1)

mergeSort(a, tmpArray, 0, 3)

mergeSort(a, tmpArray, 0, 7)

a

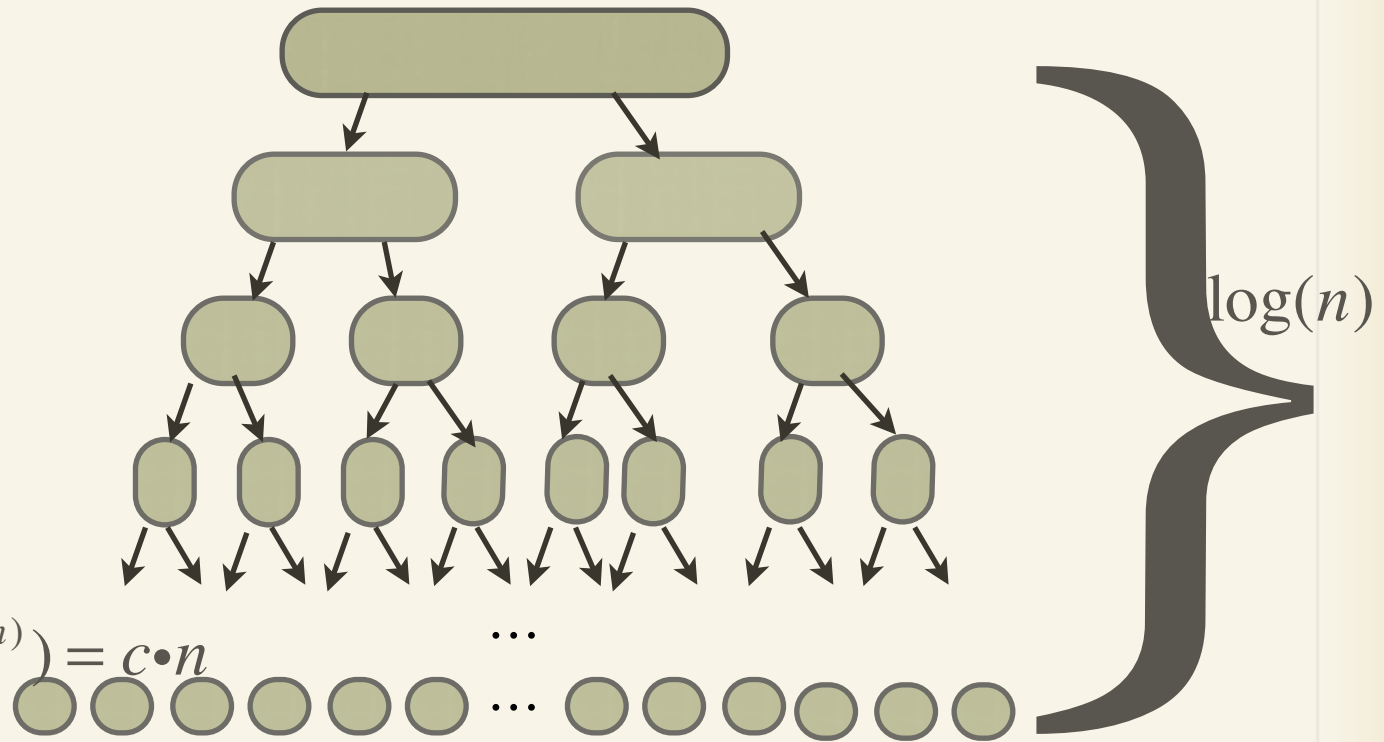|	0	1	2	3	4	5	6	7

$c \cdot n$

$2(c \cdot n / 2) = c \cdot n$

$4(c \cdot n / 4) = c \cdot n$

$8(c \cdot n / 8) = c \cdot n$

$2^{\log(n)}(c \cdot n / 2^{\log(n)}) = c \cdot n$

$\log(n)$

$O(n \cdot \log(n))$ time

Homework

```cpp
template<class InputIterator, class OutputIterator>
  OutputIterator copy (InputIterator first, InputIterator last, OutputIterator result)
{
  while (first!=last) {
    *result = *first;
    ++result; ++first;
  }
  return result;
}
```

```cpp
template <class InputIterator1, class InputIterator2, class OutputIterator>
  OutputIterator merge (InputIterator1 first1, InputIterator1 last1,
                        InputIterator2 first2, InputIterator2 last2,
                        OutputIterator result)
{
  while (true) {
    if (first1==last1) return std::copy(first2,last2,result);
    if (first2==last2) return std::copy(first1,last1,result);
    *result++ = (*first2<*first1)? *first2++ : *first1++;
  }
}
```
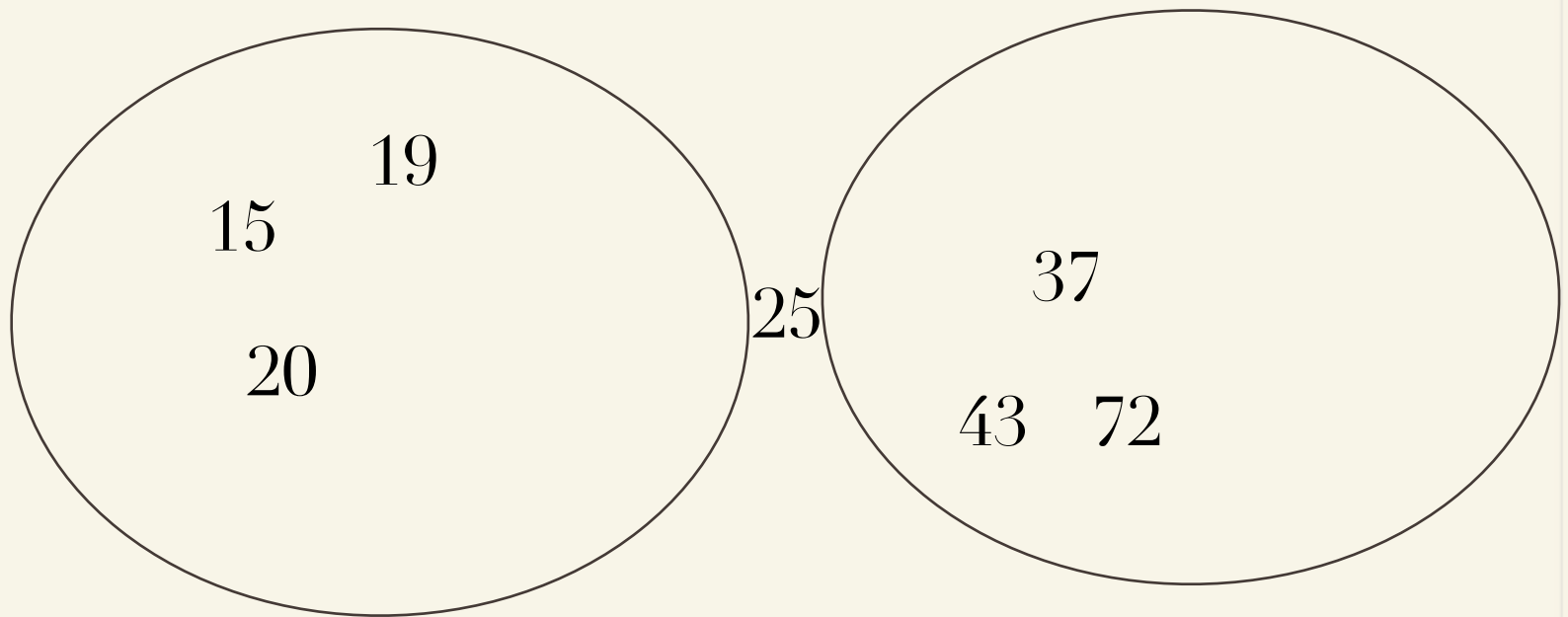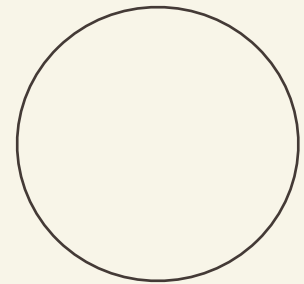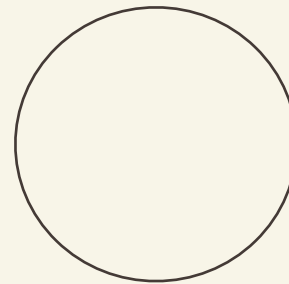
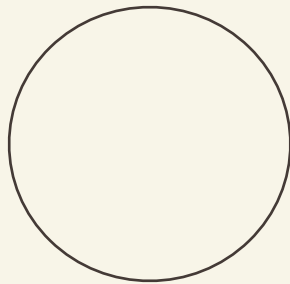From: http://www.cplusplus.com/reference/algorithm/
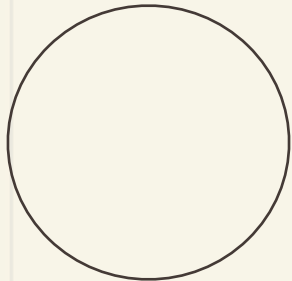
Quick Sort
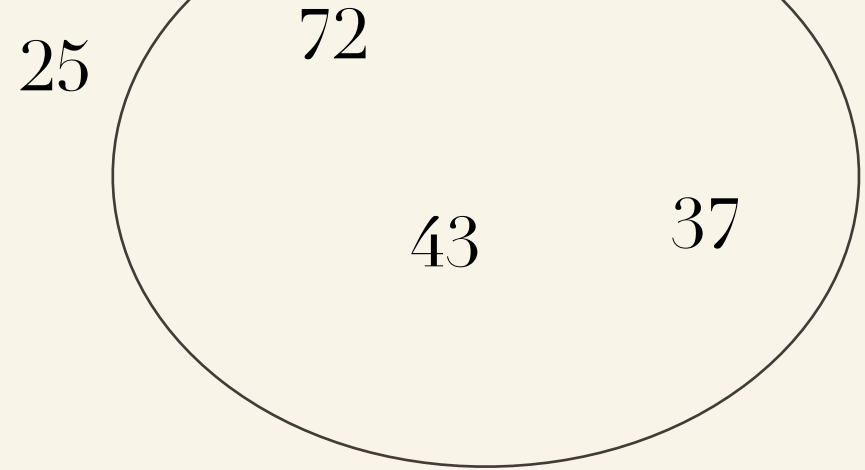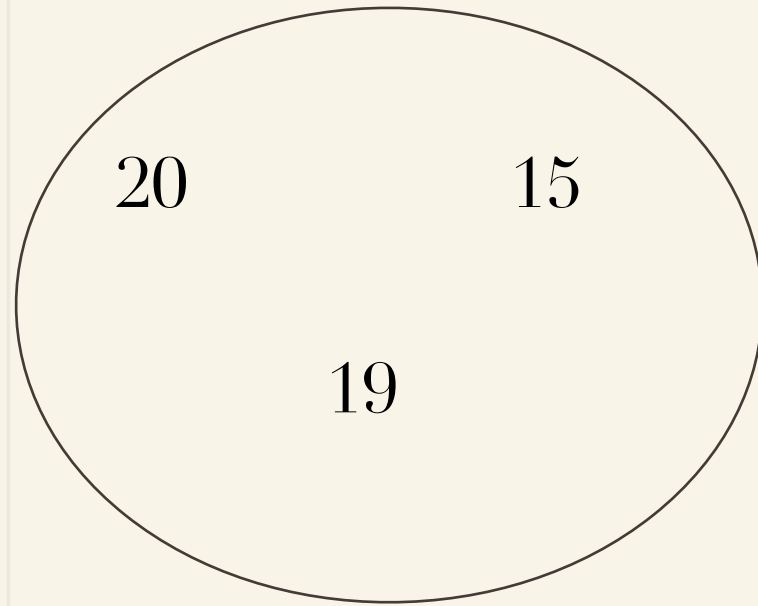# Divide and Conquer

Reduce the original problem into subproblems that are easier to solve.

15   37   25   19   20   72   43
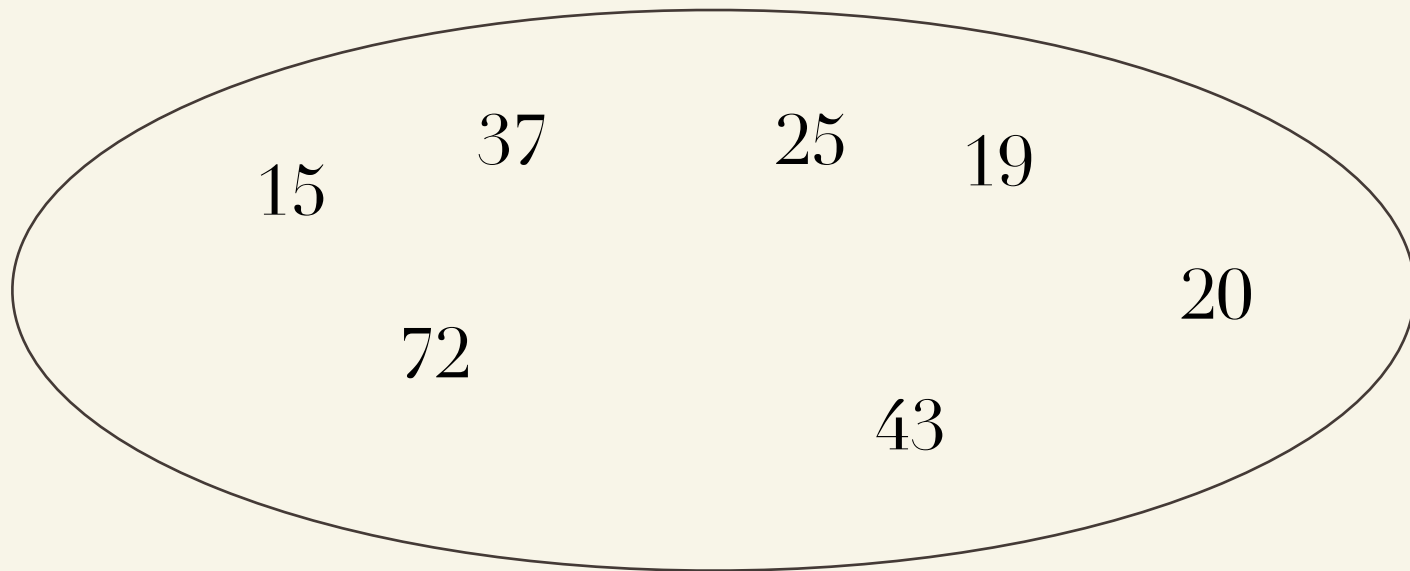
Solve the subproblems, and the combine to solve the original problem

The subproblems are solved recursively

20　15

19

25

72

43　37

QuickSort

43	37	25	20	4	72	15	19

SELECT PIVOT; PARTITION

19	15	4		20		25	72	37	43

<= pivot   pivot   >= pivot

SORT EACH PIECE (Recursively)

4	15	19		20		25	37	43	72

done!

Quick Sort

- Partition array into two pieces:
  - (elements <= pivot) , (elements >= pivot)
- Conquer:
  - sort each half recursively
- Analysis:
  - partition step can be done in linear time
  - Running time depends on choice of pivots at each stage
  - Equal sized piece → situation similar to merge sort
  - Very unequal sizes → poor performance

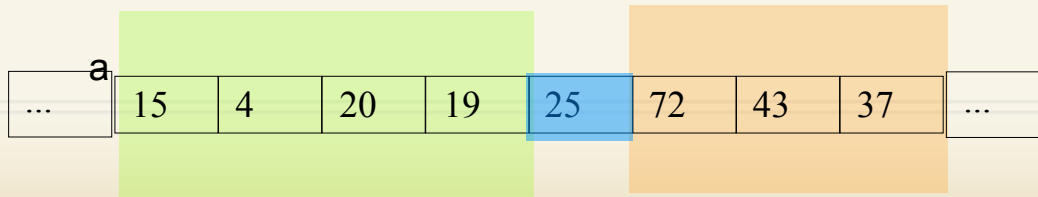# Implementation details on how we partition the items

First we select a pivot

...	43	37	20	25	4	72	15	19	......

In our implementation we move the pivot to the the last position

......	43	37	20	19	4	72	15	25	...

Then we partition the items using the pivot

```
int i = low;
int j = high -1;
for( ; ; )
{
  while ( a[i ] < pivot )  ++i;
  while(  j > i && pivot < a[j ] ) --j;

    if( i < j )
        std::swap( a[ i++ ], a[ j--] );
    else
        break;

}
swap( a[ i ], a[ high  ] );
```

Implementation detail:
The pivot is moved to the last position
to get it out of the way during
the partitioning step.

a

| ... | 15 | 4 | 20 | 19 | 25 | 72 | 43 | 37 | ... |

The Driver

```
/* driver for quicksort */
void quicksort(vector<int> &a)
{
    quicksort(a, 0, a.size()-1);
}



void main()
{
    ...
    quicksort(a);
    ...
}
```
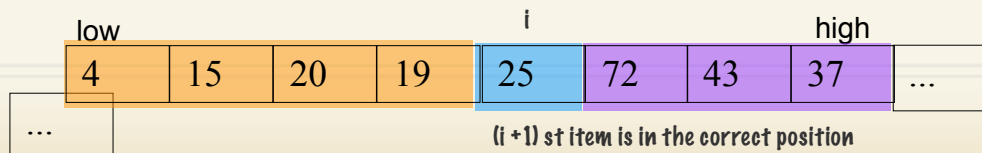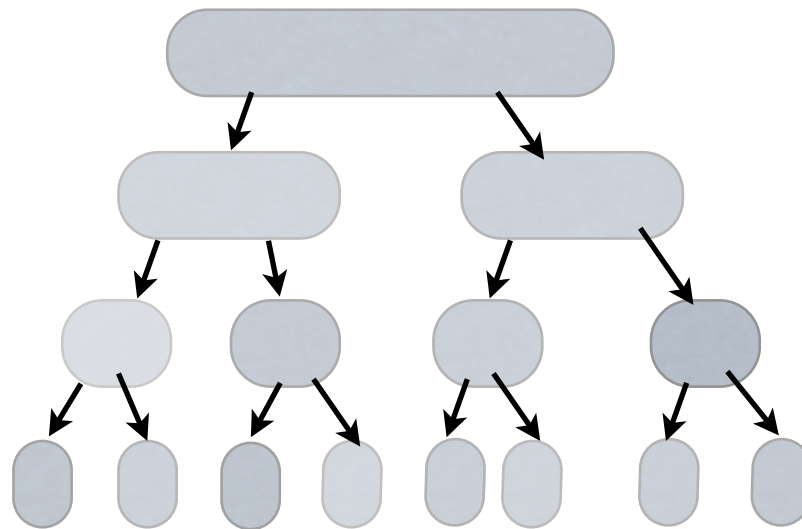
```cpp
void quicksort( vector<int> & a, int low, int high)
{
    if (low < high)
    {
        // select pivot to be element in middle position
        int mid = (low + high)/2;
        int pivot = a[ mid ];
        // put pivot in a[high]
        swap(a[high], a[mid]);
        // Begin partitioning
        int i = low;
        int j = high - 1;
        for( ; ; )
        {
            while ( a[i ] < pivot )  ++i;
            while( j > i && pivot < a[j ] ) --j;
            if( i < j )
                swap( a[ i++], a[ j-- ] );
            else
                break;
        }
        swap( a[ i ], a[ high  ] );  // Restore pivot

        quicksort( a, low, i - 1);
        quicksort( a, i + 1, high );

    }
}
```

low				i			high	
4	15	20	19	25	72	43	37	...

...

(i +1) st item is in the correct position

Best Case: 50/50 Split!

Best Case: 50/50 Split!

$\log(n)$

...

...
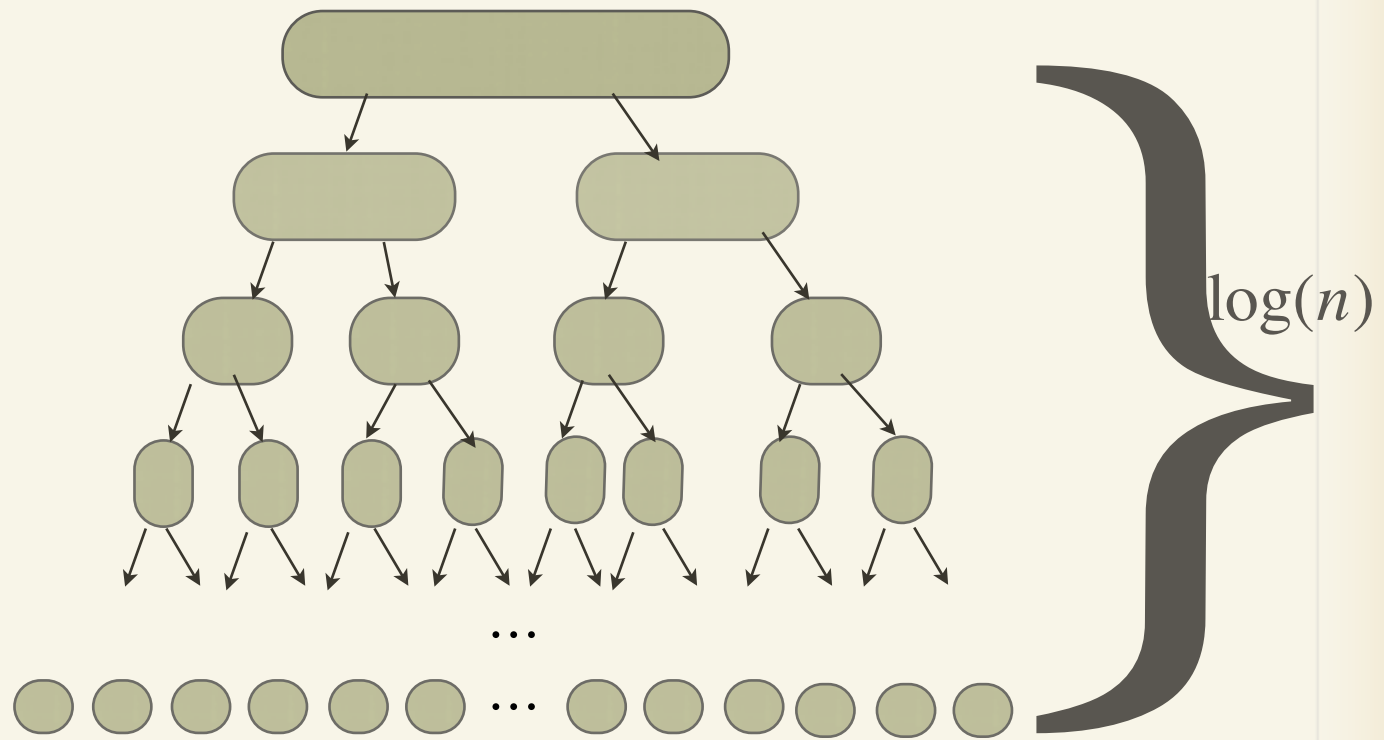
$O(n \cdot \log(n))$ time

Suppose: 90/10 Split...

$n(0.9)$ $n(0.1)$

$n(0.81)$ $n(0.081)$ $n(0.09)$ $n(0.01)$

$n(0.73)$ $n(0.081)$ ... $n(0.073)$ $n(0.081)$ ...

... ...

1

$c \log(n)$
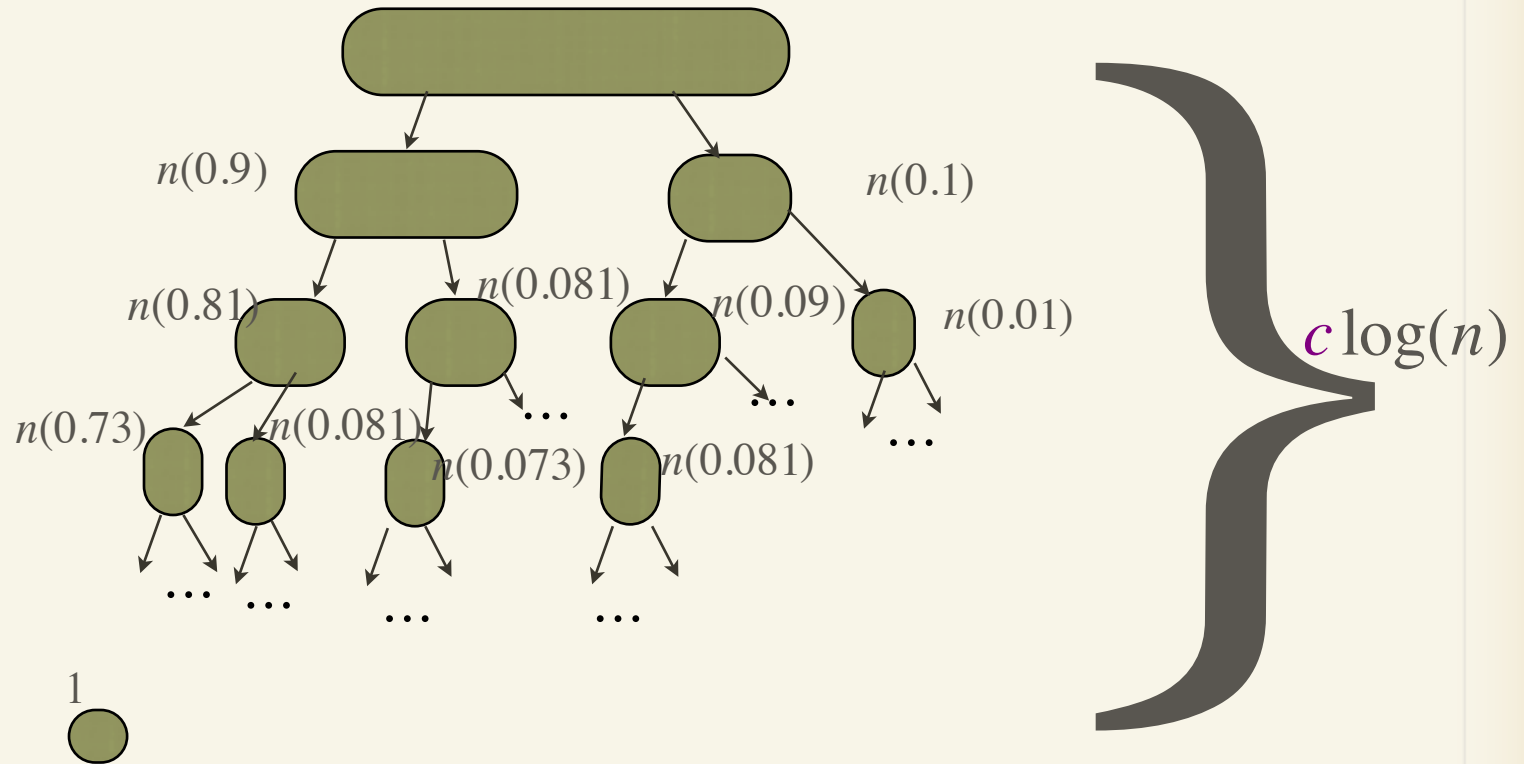
$O(n \cdot \log(n))$ time

Worst Case: n-1/0 Split...

$n-1$

$n-2$

$n-3$

1

n
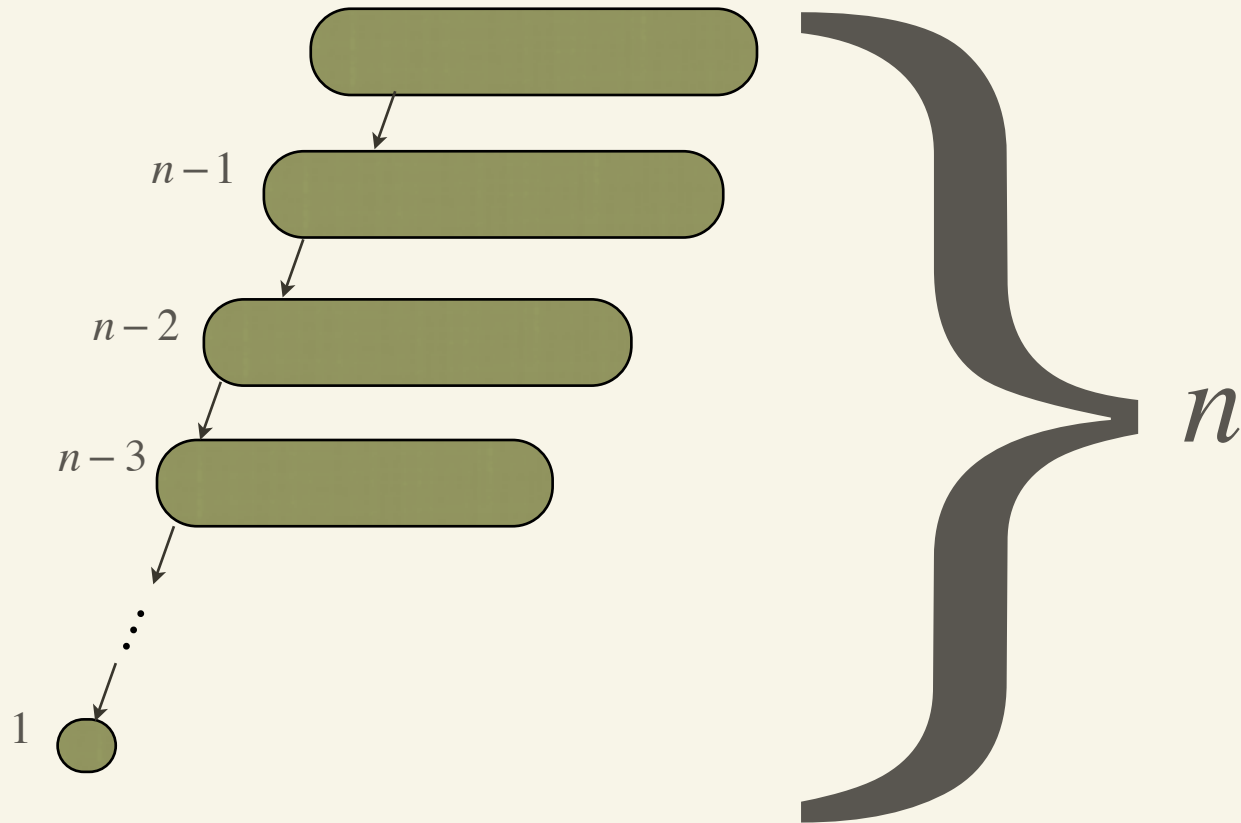
$O(n^2)$ time

Running Time

- Worst case:

  - suppose largest element is chosen as pivot on each call to Quicksort:

  - Quicksort(a, 0, n-1) calls

    Quicksort(a, 0, n-2) calls

    Quicksort(a, 0, n-3) calls

    ..

    Quicksort(a, 0, 0) calls

- $1 + 2 + ... + n = O(n^2)$

Running Time

- Best case:

  - suppose median element is chosen as pivot on each call to Quicksort:

  - Each instance calls Quicksort on two pieces of size roughly n/2

  - Picture similar to MergeSort analysis

- $T(n) = 2T(n/2) + n$

- $O(n \log n)$

- Luckily AVERAGE CASE is close to BEST CASE.

  - With good pivot selection strategy, you have to be very unlucky to get worst case partition on many calls.

Pivot Selection Strategies

- Pivot = a[low]
  - bad if array is already sorted
- Pivot = a[(high +low)/2]
  - OK, unless you're unlucky
- Pivot = median element of a[low .. high]
  - hard to compute
- Pivot = median of { a[high], a[middle],a[low]}
  - good compromise in practice

Comparison of Sorts

	worst	average	best
insertion sort	$O(n^2)$	$O(n^2)$	$O(n)$
MergeSort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
QuickSort	$O(n^2)$	$O(n \log n)$	$O(n \log n)$

Can we do better?

No! Any comparison based sorting algorithm has worst case at least n log n.

Running Times

n	insertion O(n^2)	merge O(nlog(n))	quick O(nlog(n))ave
512	0.0018	0.000166	7.4e-05
1024	0.00747	0.000356	0.000157
2048	0.030801	0.000765	0.000325
4096	0.120905	0.001644	0.000697
8192	0.474183	0.003504	0.001444
16384	1.87247	0.007584	0.002947
32768	7.56883	0.015684	0.005933
65536	29.9251	0.033017	0.012211

Lower Bound on ANY comparison based sorting algorithm
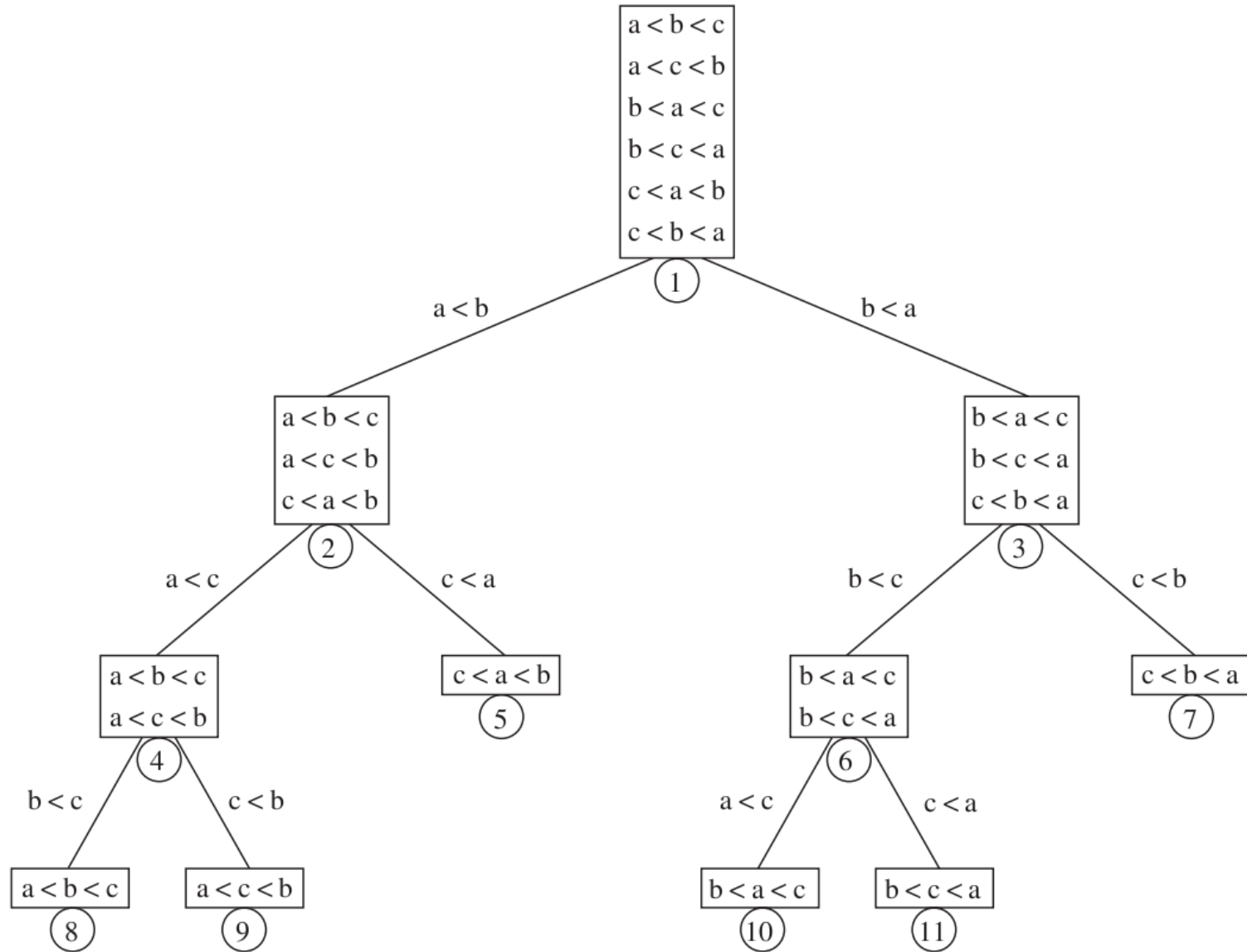
Decision Tree for sorting 3 items

Figure 7.20 A decision tree for three-element sort

Any Algorithm that sorts by using comparisons between elements uses at least cn log(n) comparisons for some input.

(Worst case lower bound for sorting using comparisons)

There are n! ways n items can be arranged.

After 0 comparisons all n! possible ways to order the items are possible.

After 1 comparison, the set of possible orders is divided into two groups: the passed the test group and the didn't pass the test group. One of these groups has at least half the permutations. The size of the larger group is at least n!/2

After j comparisons, one group will have at least $n!/2^j$ still-possible permutations.

The number of comparisons till we have only one possible ordering is thus log(n!)   (this is about n log n - 1.44n)

10240

x^2

(80, 6,400)

5120

(80, 506)

$x\log(x)$

x

(80, 80)

40

80
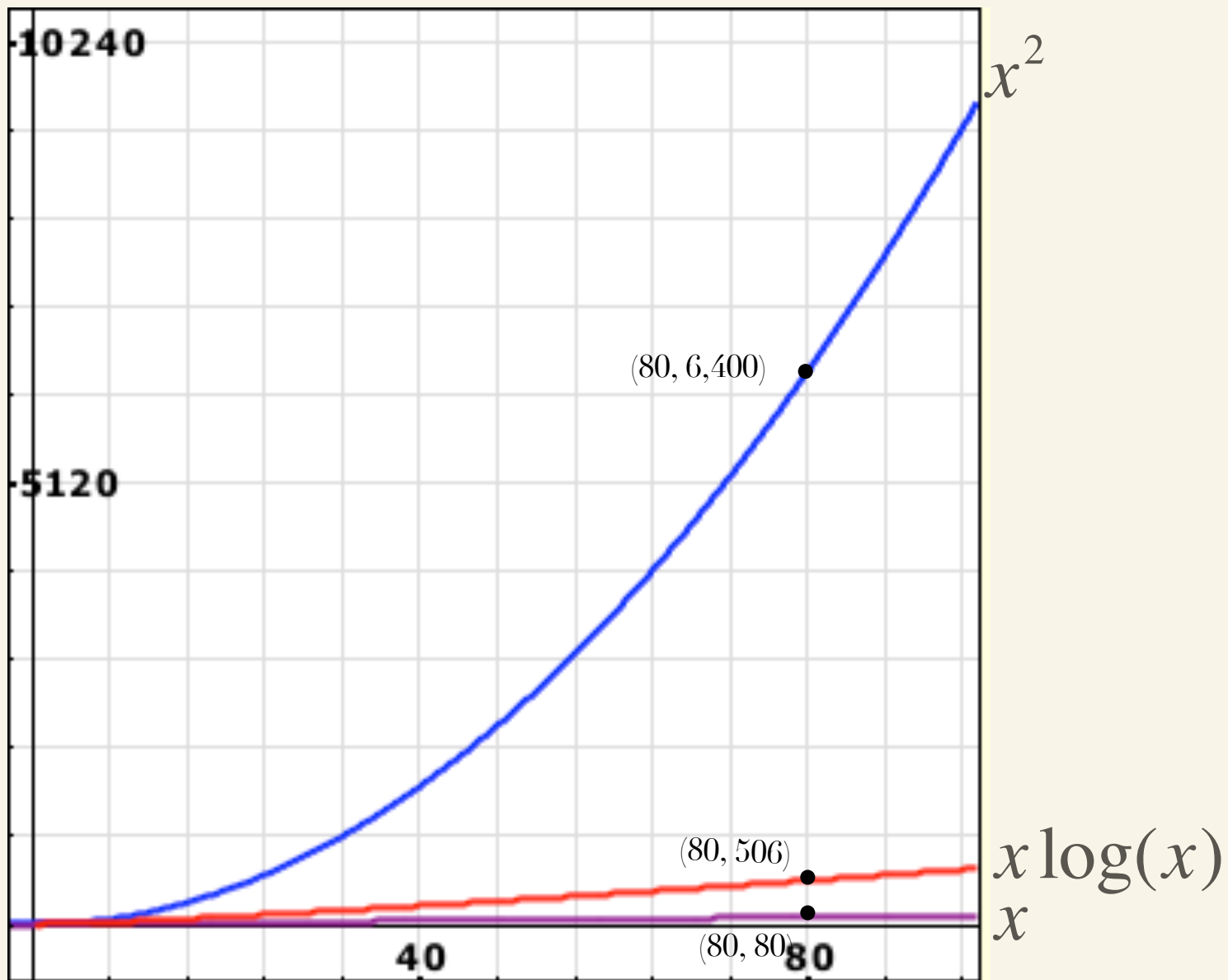
Finding Order Stats

- kth order statistic is the kth smallest element in a collection

- special case: median is the (n/2)th order statistic

- Could find it by sorting, but this takes O(n log n) time.

- modification of QuickSort can be used to find kth order stat more efficiently (in average case).

The Driver

```
template <class Comparable>
void quickSelect( vector<Comparable> & a, int k )
{
    quickSelect( a, 0, a.size( ) - 1, k );
}



void main()
{
    ...
    quickSelect(a, k);
    ...
}
```

```cpp
void quickselect( vector<int> & a, int low, int high, int k )
{
  if (low < high)
  {
     // select pivot to be element in middle position
     int mid = (low + high)/2;
     int pivot = a[ mid ];
     // put pivot in a[high]
     swap(a[high], a[mid]);
      // Begin partitioning
     int i = low;
     int j = high -1;
     for( ; ; )
     {
       while ( a[i ] < pivot )  ++i;
       while( j > i && pivot < a[j ] ) --j;
       if( i < j )
           swap( a[ i++], a[ j-- ] );
       else
           break;
     }
     swap( a[ i ], a[ high  ] );  // Restore pivot

     if (k <= i)
        quickselect( a, low, i - 1, k );
     else if(k > i + 1)
        quickselect( a, i + 1, high, k );
  }
}
```


low   i   high

| ... | 45 | 43 | 20 | 19 | 25 | 72 | 43 | 32 | ... |

(i +1) st item is in the correct position