

Board Work for Lecture 2: The Toolkit

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1 Bertrand

We will solve this problem several ways to get a feel for what is going on. It is a powerful example from game theory, and the logic underpins a lot of the economic forces we will discuss especially later in this class.

Before proceeding, first notice that although there are technically 3 players in this game, there are really only 2 interesting players. Does anyone want to guess which of the players is not interesting (and volunteer a reason for their guess)? The worker is not really interesting. This is because the worker is mechanical: they take the job that pays the most. Once we realize this, we can just pretend like there are two players: the two firms. The two players get the worker if they post a higher wage than the other, get a 50-50 shot at the worker if they tie, and get 0 if they post too low a wage.

With the worker out of the way, this boils down to a static game, because the firms both post a wage at the same time (they do not react to the wage posted by the other). Thus to solve this game we just specify two numbers (one wage for each firm). If one firm moved first (a sequential version of this game), then we would need to specify a function which states the wage one firm will post given every wage the other firm could post. We will look at this later.

Now, we want to look for a Nash equilibrium. An informal definition of an NE is a strategy for each firm such that neither firm can profitably change their strategy given knowledge of the other firm's strategy. To build some intuition, let's first suppose productivity (p) is 20. Suppose further both firms offer a wage of 10. If both firms do this, they each get the worker with probability $1/2$, and they get a profit of $1/2 * (20 - 10) = 5$. Is this a Nash equilibrium?

Well, no: either firm could offer some wage slightly more than 10 (say, 11). Then they get the worker for sure, and they get a profit of $1 * (20 - 11) = 9 > 5$, clearly much better than before. This is a profitable deviation. What if we repeat this? Suppose firm 1 now offers 11, and firm 2 offers 10. If they both do this, firm 1 gets profit of 9 and firm 2 gets 0. Who has a profitable deviation? Well, firm 2 can now out-bid 1 and offer slightly more than 11 (say, 12). Now firm 2 gets the worker for sure and gets profit $1 * (20 - 12) = 8 > 0$. We can continue repeat this process for any wages that are less than $p = 20$.

This loose discussion suggests a guess: what if both firms offer a wage of 20? Perhaps this is a Nash equilibrium (and a solution for the purposes of this class). To be general, let's guess that both firms play $w_1 = w_2 = p$.

- To prove this is an NE, we just need to check that given the other firm is playing $w_{-i} = p$, firm i does not gain by playing something other than $w_i = p$.
- First note that when $w_1 = w_2 = p$ profit for both is 0 because whichever firm gets the worker pays out exactly their productivity so: $0.5 * (p - w_i) = 0$.
- Suppose one firm attempts to outbid by offering some higher wage $w'_i > p$ while the other stays. Does the deviating firm gain from doing this?
- No: the deviator wins the worker for sure, but now the worker is being paid more than their productivity, so the firm makes negative profit: $1 * (p - w'_i) < 0$.
- What if one firm tries to undercut with a lower wage: $w'_i < p$? Do they gain?
- No: the deviating firm now misses out on the worker for sure, and continues to get a profit of 0: $0 * (w'_i - p) = 0$

- Thus $w_1 = w_2 = p$ is an NE.

For the purposes of this class, this is enough for a complete answer, as you have just proven that $w_1 = w_2 = p$ is a Nash Equilibrium. Now, you may wonder, is this the only Nash Equilibrium? It turns out it is. Showing this requires only a bit more work. We will prove this by contradiction.

- Suppose, for sake of contradiction, there is another NE where at least one wage is not equal to productivity p .
- For clarity, just assume that the high wage firm is 1: $w_1 \geq w_2$.
- **Case 1:** $p < w_1$. In this case, firm 1 gets the worker for sure but makes negative profit because wages are more than productivity. This is clearly worse than not getting the worker at all (0 profit) so it cannot be an NE.
- **Case 2:** $p > w_1$. Now firm 1 is making positive profit (because $w_1 \geq w_2$). Firm 2 is either making 0 profit, or is splitting the positive profit with firm 1. But firm 2 can just slightly outbid firm 1 and offer $p > w'_2 > w_1$ and make more profit!
- **Case 3:** $p = w_1 > w_2$: Now firm 1 wins the worker and makes 0 profit. Firm 2 loses and also makes 0 profit. This is therefore equivalent in terms of profit for the firms as our solution. But this is still not an NE! To see why, notice that firm 1 can slightly reduce the wage to $p > w'_1 > w_2$. Firm 1 still wins, but now makes positive profit!
- We have considered every other possible strategy of the two firms, and every other possible strategy has a profit deviation and is therefore not an NE. Therefore, $w_1 = w_2 = p$ is the unique NE!

2 Company Call List - Static

- Write down the utility-maximization problem for worker i . We can focus on some generic worker i because both workers are the same. This is called solving by symmetry.

Solution)

$$u_i = \max_{q_i} (120 - q_i - q_{-i})q_i$$

- Solution)**

To find the optimal choices, we take the FOC of the utility expression. This is the partial derivative with respect to q_i :

$$\frac{\partial u_i}{\partial q_i} = 120 - 2q_i - q_{-i} = 0 \quad \Longleftrightarrow \quad q_i = \frac{120 - q_{-i}}{2} \quad \dots (*)$$

Notice that this depends on what each worker expects the other to do (this is why q_{-i} is in the expression).

- Because the problem is the same for both players, we can now solve by finding a solution to the system of equations:

Solution)

$$\begin{aligned} q_1 &= \frac{120 - q_2}{2} \\ q_2 &= \frac{120 - q_1}{2} \end{aligned}$$

Plugging in the 2nd equation into the 1st one, we get:

$$q_1 = \frac{120 - \frac{120 - q_1}{2}}{2} = \frac{60}{2} + \frac{q_1}{4}$$

$$\iff \frac{3}{4}q_1 = 30$$

So $q_1^* = 40$ and $q_2^* = 40$ (by symmetry).

- d. I will not require this on exams but we can also prove that this NE is the unique (or only) pure strategy Nash Equilibrium using a bit of calculus.

Solution)

To prove uniqueness, it is sufficient for us to show that the objective function $u_i(q_i)$ is strictly concave so that the solution of the FOC is not only indeed a maximum but also the global maximum of the function. This will be true if $u_i(q_i)$ has a negative 2nd derivative:

$$u'(q_i) = 120 - 2q_i - q_{-i}$$

$$u''(q_i) = -2 < 0$$

Thus, there is a unique best response for every action of the other player, i.e. there is a unique best response to the Nash equilibrium strategy of the other player. This is also called checking the second order condition (SOC).

- e. Now suppose there is a social planner that maximizes total surplus (the sum of both worker's utility). The planner controls both q_1, q_2 . Find the values the social planner would choose¹. Show that one of the values is a Pareto improvement over the equilibrium actions (both players are better off than in NE).

Solution)

The social surplus (the objective function of the social planner) is:

$$u_1 + u_2 = (120 - q_1 - q_2)q_1 + (120 - q_1 - q_2)q_2 = 120q_1 + 120q_2 - q_1^2 - q_2^2 - 2q_1q_2$$

So the maximization problem is:

$$u_T = \max_{q_1, q_2} 120q_1 + 120q_2 - q_1^2 - q_2^2 - 2q_1q_2$$

By FOC, we have:

$$\frac{\partial u_T}{\partial q_1} = 120 - 2q_1 - 2q_2 = 0 \quad \dots (1)$$

$$\frac{\partial u_T}{\partial q_2} = 120 - 2q_2 - 2q_1 = 0 \quad \dots (2)$$

Note that the equation (1) includes $2q_2$ rather than just q_2 (same for equation (2)). This can be interpreted as the social planner taking into account externalities when choosing q_1 and q_2 . This does not occur in the Nash Equilibrium we solved for earlier. (Why?) Solving the above system, we get:

$$120 - 2q_2 = 120 - 2q_1 \quad \Rightarrow \quad q_1 = q_2.$$

This means that there are multiple (in fact, infinitely many) solutions. We can see this also by realizing that equations (1) and (2) are the same, so we have two variables and 1 equation (meaning many solutions). So any q_1, q_2 that solves equation (1) will maximize the social surplus. Thus, q_1^{**}, q_2^{**} are any values such that:

$$q_1^{**} + q_2^{**} = 60.$$

¹There may be more than one (or infinitely many) values.

Among these, one natural solution would be equal division of the calls: $q_1^{**} = q_2^{**} = 30$. Then accordingly, the utility for each worker under the social planner's equal division solution is:

$$\begin{aligned} u_1(30) &= (120 - 30 - 30)30 = 1800 \\ u_2(30) &= 1800. \end{aligned}$$

The utility from the Nash equilibrium is lower:

$$u_1(q_1^*) = u_2(q_2^*) = u_1(40) = (120 - 40 - 40)40 = 1600.$$

Thus, we see that both workers would be strictly better off if they could commit to only making 30 calls each. A firm could make this happen by having a supervisor who assigns calls.

3 Company Call List - Sequential

Suppose everything about that problem is the same, except that player 1 moves first and then player 2 moves.

- Draw the extensive form of the game (game tree). Solution.
- Let's solve using backwards induction.

Start with the second stage and worker 2, taking q_1 as fixed. Then worker 2 solves:

$$\max_{q_2} q_2(120 - q_1 - q_2)$$

FOC:

$$120 - q_1 - q_2 - q_2 = 0 \implies q_2(q_1) = \frac{120 - q_1}{2}$$

Now roll back. Worker 1 anticipates the above strategy by worker 2 in the next stage and solves:

$$\max_{q_1} q_1(120 - q_1 - q_2(q_1)) = \max_{q_1} q_1(120 - q_1 - \frac{120 - q_1}{2}) = \max_{q_1} q_1(60 - q_1/2)$$

FOC:

$$\frac{du_1}{dq_1} = 60 - q_1/2 - q_1/2 = 0 \implies q_1 = 60$$

Then the SPNE strategy is:

$$q^* = 60 \quad q_2^*(q_1) = \frac{120 - q_1}{2}$$

Notice that worker 2's strategy is a function which states the number of calls to make for every possible number of calls made by worker 1. For example, if worker 1 made 120 calls, worker 2 makes 0. But if worker 1 makes none, worker 2 makes 60.

- Compute profit for worker 1 and worker 2. Some games are said to exhibit a "first-mover advantage." Does this game? Solution.

$$\begin{aligned} u_1 &= 60(120 - 60 - 30) = 1800 \\ u_2 &= 30(120 - 60 - 30) = 900 \end{aligned}$$

Since $u_1 > u_2$ and worker 1 moved first this game has a first-mover advantage. This is because the first-mover (worker 1) can call more people knowing that worker 2 will compensate by calling less people. We see that relative to the static problem, worker 1 is calling more while worker 2 is calling less.