

# TCSS 343 - Week 3

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## Master Theorem Practice

### Master Theorem

The Master Theorem applies to recurrences of the following form:

$$T(n) = aT(n/b) + f(n)$$

where  $a \geq 1$  and  $b > 1$  are constants and  $f(n)$  is an asymptotically positive function.

There are 3 cases:

1. If  $f(n) = O(n^{\log_b a - \epsilon})$  for some constant  $\epsilon > 0$ , then  $T(n) = \Theta(n^{\log_b a})$ .
2. If  $f(n) = \Theta(n^{\log_b a} \log^k n)$  with  $k \geq 0$ , then  $T(n) = \Theta(n^{\log_b a} \log^{k+1} n)$ .
3. If  $f(n) = \Omega(n^{\log_b a + \epsilon})$  with  $\epsilon > 0$ , and  $f(n)$  satisfies the regularity condition, then  $T(n) = \Theta(f(n))$ .  
Regularity condition:  $af(n/b) \leq cf(n)$  for some constant  $c < 1$  and all sufficiently large  $n$ .

For the following problems show which case of the master theorem each problem goes to. The master theorem is applicable for each problem.

$$0. \ T(n) = 3T\left(\frac{n}{2}\right) + n^2$$

$$1. \ T(n) = 4T\left(\frac{n}{2}\right) + n^2 \log n$$

$$2. \ T(n) = 3T\left(\frac{n}{4}\right) + n \log n$$

$$3. \ T(n) = 2T\left(\frac{n}{4}\right) + 2$$

$$4. \ T(n) = T\left(\frac{n}{4}\right) + \log n$$

$$5. \ T(n) = 2T\left(\frac{n}{4}\right) + \sqrt{n}$$

$$6. \ T(n) = 2T\left(\frac{n}{4}\right) + n^{0.51}$$

$$7. T(n) = 3T\left(\frac{n}{2}\right) + n$$

$$8. T(n) = 4T\left(\frac{n}{2}\right) + n$$

$$9. T(n) = 3T\left(\frac{n}{3}\right) + \frac{n}{2}$$

$$A. T(n) = 4T\left(\frac{n}{2}\right) + \frac{n}{\log n}$$

$$B. T(n) = T\left(\frac{n}{3}\right) + n^2$$

$$C. T(n) = 8T\left(\frac{n}{3}\right) + 2^n$$

$$D. T(n) = 16T\left(\frac{n}{4}\right) + n$$

F.  $T(n) = 2T(\frac{n}{4}) + n!$

10.  $T(n) = 0.5T(\frac{n}{2}) + \frac{1}{n}$

11.  $T(n) = 16T(\frac{n}{4}) + n!$

12.  $T(n) = 9T(\frac{n}{3}) + n^2$

13.  $T(n) = 7T(\frac{n}{3}) + \cos n$

14.  $T(n) = 8T(\frac{n}{3}) + 1$

15.  $T(n) = T(\frac{n}{2}) + n^3$