Mazemake Sunday, November 3, 2019 Assume a maze of 2×2. Thus, where the internal walls are coloured blue. Quantity of internal walls: N = ((Ncols -1) × Nrows) + ((Nrows - 1) × Ncols)

vertical haizoutal Make poir of all cells w/ internal woll between them. "cell wall [N][z]" int k = 0; FOR EVERY CELL @ i, j... except in last column WALL[K][O] = cell(i ,j); // Crecout WALL[K][1] = cell(i+1,j); / RIGHT K++; FOR EVERY CELL @ i, j... except in last row WALL[K][O] = cell(i,j); / CHRECUT WALL[K][I] = cell(i,j+1); / BELOW K++> Thus WALL[K][N] is: 0 (0,0) (1,0) Pairs w/ wall between

Shuffle via randpern handout (you use regular "rand" w/ "srand (time (NULL))"

Call "dset ds" of size Nrow * Ncol (NOT N).

Thus, ds is as follows:

D		
2	3	

OBSERVE: If we merge 0 + 1,
we "break" the wall between
them.

0	>			4	sets	=>	3	5 e ³	ts
2 3	-/	2	3						

If we did 1 (in 0) + 3 next...

With this in mind, just keep merging until there is one set left.

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Which cells do we merge? Let's let wan help. Assume it's unshuffled.
0 (0,0) (1,0)
1 (0,1) (1,1)
2 (0,0) (0,1)
3(1,0)(1,1)
ALGORITHM
FOR I FROM O TO N ... {
  11 Convert 2D to 1D
   ii = WALL[i][o].x + (WALL[i][o].y x Ncol);
jj = SAME AS & BUT W/ WALL[i][1]
   / If not already merged ...
    IF (DS. FIND(ii) != DS. FIND(;))
      DS. MERLE (ii, jj);
    ELSE
      // Conit merge, so there's a wall
     PRINT WALL[i] 0 + 1
    IF (DS. SIZE() == 1)
      BREAK
3
i++5 // Just do it ... Seriously
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// Print out remaining walls from i to N FOR i FROM i to N PRINT Same as above (Ex i = O. WALL[i] = (0,0) + (1,0) NOT in same set. MERCE WALL[i] = (0,1) + (1,1)NOT in same set. MERCE i = 2. WALL[:] = (0,0) + (0,1)NOT in same set. MERCE i=3. WALL[:]= (1,0) + (1,1) ARE in same set. PRINT Oh, also if you do (2,4), flip when printing. That example was with when unshuffled. Just shuffle after you get this working.

Should write this: MAZE 2 2 O 1 1 1 Always 1 wall in a 2×2 maze.