

2020/11/04 - Binary Search Tree (Part 2)

04 Tháng Mười Một 2020 3:19 SA

SYNOPSIS

- Go over Lab 6 (Parts 2 + 3)

LAB 6

- As always, on Canvas, go to the "Lab 6.2" assignment and read the "lab6.html" file attached. All lab details will be there.
- There are **three parts** split into **two submissions**. This will work in your favour in the long run.

SUBMISSION COMMAND

LAB 6.1

```
tar -cvf lab6-1.tar BST.h
```

LAB 6.2

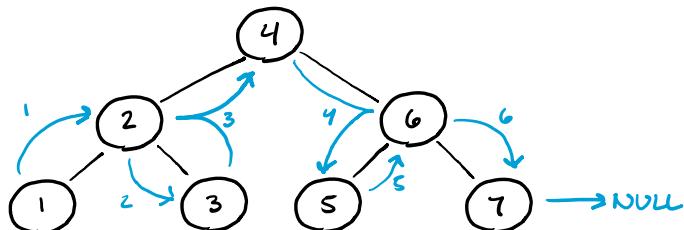
```
tar -cvf lab6-2.tar BST.h
```

SKETCH "KEY"

- Your mileage may vary. Different TAs may expect different information. A sketch + (correct) text should do.
- I write this in a way that will help you in next parts. Maybe you don't match. Don't panic or stress over it.

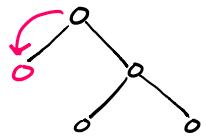
* OPERATOR ++

- Single step in in-order traversal. May be done via considering where node came from.



INITIAL

- Go as far left as possible. (bst::begin, not operator ++)

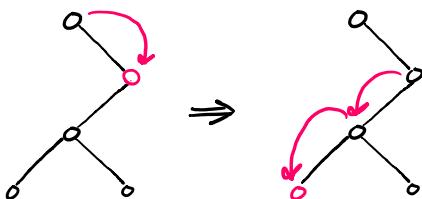


STEP

- If we can go right:

- Go right once

- Go left as much as possible



- Otherwise,

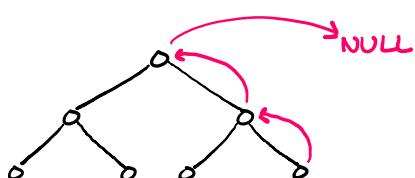
- Go up

- If the node we were just at is the right child of the node we are now at, go up again.

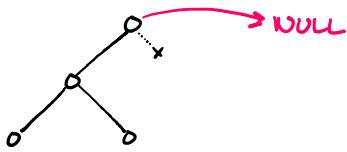
REPEAT UNTIL NO LONGER TRUE.



- If came from right and hit the root, set to NULL.

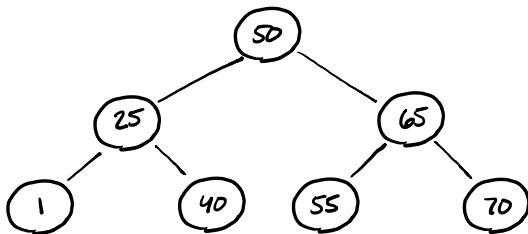


- If at root and right is NULL, set to NULL



* Lower-Bound

- lower-bound \leq key



$$\text{Lower-Bound}(1) = 1$$

$$\text{Lower-Bound}(25) = 25$$

$$\text{Lower-Bound}(26) = 40 \quad \leftarrow \text{It goes to the next one.}$$

$$\text{Lower-Bound}(52) = 55$$

$$\text{Lower-Bound}(0) = 1$$

$$\text{Lower-Bound}(99) = \text{NULL}$$

INITIAL

- Have two pointers :

- t - Does tree traversing

- n - Keeps track of last time t went left.

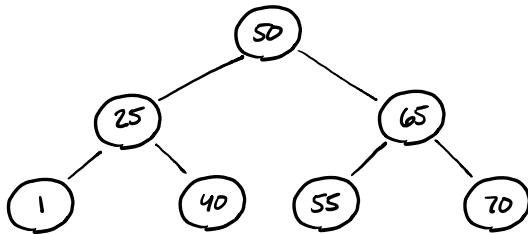
- n is NULL at the start. t is root.

STEP

- As long as t isn't **NULL** ...
 - If key (search key) $\leq t \rightarrow \text{key}$, go **left**. But before you do, log t in n .
 - If not \leq , have t go to the **right**.
 - No, n will not change here.
- Return iterator pointing to n .

* UPPER-BOUND

- $\text{key} < \text{upper-bound}$



$\text{UPPER-BOUND}(1) = 25$ ← It goes to the next one.
 $\text{UPPER-BOUND}(25) = 40$
 $\text{UPPER-BOUND}(40) = 55$
 $\text{UPPER-BOUND}(55) = 65$
 $\text{UPPER-BOUND}(65) = 70$
 $\text{UPPER-BOUND}(70) = \text{NULL}$

PROCEDURE

- Same as **lower_bound**, but $<$ instead of \leq .

BST.h - PART 2

- Implement the `iterator` public subclass in `bst`.
- Implement the `begin` and `end` functions in `bst`.
- All but 2 functions are one-liners.
- Iterator prototype:

```
class iterator {
public:
    //Constructor (Public)
    iterator();

    //Operator Overloads
    iterator & operator ++();
    They & operator *();
    bool operator == (const iterator &) const;
    bool operator != (const iterator &) const;

private:
    friend class bst<They>;

    //Constructor (Private)
    iterator (node *);

    //Interior Pointer
    node * p;
}
```

- Other prototypes (public functions of `bst` introduced after `iterator` subclass):

```
iterator begin();
iterator end();
```

- So... all except `bst<They>::iterator::operator++` and `bst<They>::begin` are one-liners.

- You might figure out that there are **two** possible ways to **operator++** via Google or something. This is because of things like **++i** (pre-increment) and **i++** (post-increment). This lab wants the **former**.

- Function Definition Outlines:

//Constructors (Public & Private)

```
template <class TKey>
bst<TKey>::iterator::iterator() {
| |
| |
}
```

template <class TKey>

```
bst<TKey>::iterator::iterator(bst<TKey>::node *ptr) {
| |
| |
}
```

//Operator Overloads

template <class TKey>

```
typename bst<TKey>::iterator & bst<TKey>::iterator::operator++() {
| |
| |
}
```

template <class TKey>

```
TKey & bst<TKey>::iterator::operator*() {
| |
| |
}
```

template <class TKey>

```
bool bst<TKey>::operator==(const bst<TKey>::iterator &rhs) const {
| |
| |
}
```

```
template <class TKey>
bool bst<TKey>::iterator::operator !=(const bst<TKey>::iterator &rhs) const {
|
|}
```

//BST Functions

```
template <class TKey>
typename bst<TKey>::iterator bst<TKey>::begin() {
|
|}
```

```
template <class TKey>
typename bst<TKey>::iterator bst<TKey>::end() {
|
|}
```

- Function Breakdown (One-Liners)

- Both iterator constructors just set p. Either to **NULL** or to a specified value. Should be obvious what to do there.

- **operator *** returns p's key.
- **operator ==** returns TRUE if p is the same as rhs.p.

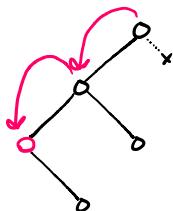
- **operator !=** is the opposite of operator ==...
- **end()** returns an iterator pointing to **NULL**. See sketch key to see why. (It's essentially one element past the last node. Hence **NULL**).

- Function Breakdown (Others)

- **operator ++**. Refer to sketch key. Algorithm is discussed there with diagrams.

- `begin()`. If `Troot` is `NULL`, return an iterator pointing `NULL`. Otherwise, start at `Troot` and go as far left as possible.
Return an iterator pointing to left-most node.
 - Since it was asked in lab: **No**. Do not take a right.

(Ex



BST.h - PART 3

- Implement the `lower-bound` and `upper-bound` functions in `bst`.
- If done correctly, this is a **two-for-one**.
A.k.a. write one and you have both (with a minor adjustment).
- Function Definition Outlines:

```

template <class TKey>
typename bst<TKey>::iterator bst<TKey>::lower_bound(const TKey &key) {
    :
}
  
```

```

template <class TKey>
typename bst<TKey>::iterator bst<TKey>::upper_bound(const TKey &key) {
    :
}
  
```

- Refer to sketch key. Procedure and examples are shown. **#TeachBySledgehammer**.