



MAKE
SCHOOL

ALGORITHM ANALYSIS

Big O? What's the Big Deal?

WHAT IS AN ALGORITHM?

A sequence of steps to accomplish a task

...described precisely enough for a computer
to perform them

Algorithm != Code

Algorithms can be written in any language,
including *pseudocode*, English, Farsi, etc.

More info: [algorithm definition](#), [algorithm characterizations](#), [pseudocode](#)

WHO COINED ALGORITHM?

Al Gore... Rhythm?

After all, he *did** invent the Internet **not*

Muhammad ibn Mūsā al-Khwārizmī

9th century Persian mathematician

Also coined “*al-jabr*” – a.k.a. *algebra*

More info: [word origin](#), [etymology](#), [Muhammad ibn Mūsā al-Khwārizmī](#)

ALGORITHM ANALYSIS

Correctness – does it solve the problem?

Formal verification methods – e.g., NASA

Resource usage – how efficient is it?

Time – idealized work steps (CPU cycles)

Space – working memory (RAM, disk, etc.)

More info: [correctness](#), [formal verification](#), [algorithm analysis](#)

WHY IS THIS IMPORTANT?

Cost – computers and electricity aren't free

Comparison – choose best solution based on analysis *before* implementing and running

Tractability – not all problems can be solved in your lifetime, even with moderate size datasets

Non-deterministic Polynomial-time, NP-hard

More info: P versus NP problem, NP-hard, NP-complete

RUNTIME IN PRACTICE

Actual runtime depends on many factors:

CPU speed – pace of instruction execution

Language – abstractions on top of CPU instructions

Compiler – optimization of instruction ordering, etc.

Environment – resources available during runtime
(not used by other programs running concurrently)

RUNTIME COMPARISON

| | Computer A | Computer B |
|-----------------------------|----------------------------|----------------------------|
| Sorting Algorithm | Bubble Sort | Merge Sort |
| Time Complexity | n^2 | $n \cdot \log_2 n$ |
| Execution Speed | 10 GHz (billion instr/sec) | 10 MHz (million instr/sec) |
| Code Optimization | High (2 instr/step) | Low (50 instr/step) |
| Actual Complexity | $2n^2$ | $50n \cdot \log_2 n$ |
| Time to Sort 100K #s | 2 seconds | 8.3 seconds |
| Time to Sort 1M #s | 3.33 minutes | 1.67 minutes |
| Time to Sort 10M #s | 5.56 hours | 19.4 minutes |
| Time to Sort 100M #s | 23.1 days! | 3.69 hours |

Example from Algorithms Unlocked

RUNTIME ASSUMPTIONS

Need to make some assumptions to easily describe algorithm runtime in the abstract:

Time to execute basic operations is *constant*

Relative speed of operations is *not important*

Including arithmetic, logic, comparison, variable assignment, array indexing, function calls, etc.

RUNTIME ANALYSIS

What we really want to analyze is *complexity* – how algorithm runtime grows as the input gets larger

Describe runtime as a function of the size of the input

Input array/list length usually written as ***n***, ***m***, ***k***

Ignore constants and lower-order terms

e.g., **$3n \rightarrow n$** , **$6n^2 + 15n + 24 \rightarrow n^2$**

CLASSIC LINEAR SEARCH

Return index of **target** in **histogram**, or **None** if not found

```
1 def linear_search(target, histogram):
2     index = None # not found
3     i = 0
4     for item, count in histogram:
5         if item == target:
6             index = i # found
7             i += 1
8     return index
```

PYTHONIC LINEAR SEARCH

Return index of **target** in **histogram**, or **None** if not found

```
1 def linear_search(target, histogram):  
2     index = None # not found  
  
4     for i, (item, count) in enumerate(histogram):  
5         if item == target:  
6             index = i # found  
  
8     return index
```

EARLY EXIT LINEAR SEARCH

Return *first* index of **target** in **histogram**, or **None** if not found

```
1  def linear_search(target, histogram):  
  
4      for i, (item, count) in enumerate(histogram):  
5          if item == target:  
6              return i # found  
  
8      return None # not found
```

ASYMPTOTIC NOTATION

Worst case – upper bound

Algorithm is $O(f(n))$ – “big oh of $f(n)$ ”

Best case – lower bound

Algorithm is $\Omega(f(n))$ – “omega of $f(n)$ ”

If both bounds are the same, then

Algorithm is $\Theta(f(n))$ – “theta of $f(n)$ ”

RUNNING TIME SUMMARY

| | Best Case | Worst Case | Both Cases |
|--|-------------|------------|-------------|
| Exhaustive Linear Search (return last match) | $\Omega(n)$ | $O(n)$ | $\Theta(n)$ |
| Early Exit Linear Search (return first match) | $\Omega(1)$ | $O(n)$ | N/A |

RESOURCES

Introduction to Algorithms by Cormen, Leiserson, Rivest, and Stein – widely considered *the Bible of Algorithms*

Algorithms Unlocked by Thomas Cormen – introductory and more accessible, less technical detail than CLRS