1 Monads and Comonads

1.1 Definition of Monads and Comonads

Definition 1.1 (Monad). A *Monad* (T, μ, η) in a Category $\mathcal{X}$ consists of

- an endofunctor $T\colon \mathcal{X} \to \mathcal{X}$
- a natural transformation $\eta\colon \mathrm{id}_{\mathcal{X}} \Rightarrow T$
- a natural transformation $\mu\colon T^2 \Rightarrow T$

such that the following diagrams commute:

(a)
$$\begin{tikzcd} T^3 \arrow[r] \arrow[d] & T^2 \arrow[d] \\ T^2 \arrow[r] & T \end{tikzcd}$$
and
$$\begin{tikzcd} T \arrow[r, "\eta T"] \arrow[dr, "\mathrm{id}_T"'] & T^2 \arrow[d, "\mu"] & T \arrow[l, "T\eta"'] \arrow[dl, "\mathrm{id}_T"] \\ & T & \end{tikzcd}$$

Example 1 (preorder). Recall: A *preorder* $(\mathcal{P}, \leq)$ is a category with $\mathcal{P}$ as objects and a morphism between X and Y iff $X \leq Y$. A functor $T\colon \mathcal{P} \to \mathcal{P}$ is thus a monotonic function $\mathcal{P} \to \mathcal{P}$ ($x \leq y \implies Tx \leq Ty$). The existence of the natural transformations η is equivalent to

$$x \leq Tx \ \forall x \in \mathcal{P}$$

and the existence of μ is equivalent to

$$T(Tx) \leq Tx \ \forall x \in \mathcal{P}$$

because there is at most one morphism $x \to y$, so the neccessary diagrams commute trivially.
Now suppose $\mathcal{P}$ is a *partial order*, i.e. $x \leq y \leq x \implies x = y \ \forall x, y \in \mathcal{P}$.
Then:

$$\begin{aligned} x \leq Tx &\implies Tx \leq T(Tx) \\ T(Tx) \leq Tx &\implies Tx = T(Tx) \end{aligned}$$

so a Monad T in a partial order $\mathcal{P}$ is a *closure operation* in $\mathcal{P}$, i.e. a monotonic function $T\colon \mathcal{P} \to \mathcal{P}$ with $x \leq Tx$ and $T(Tx) = Tx \ \forall x \in \mathcal{P}$.
Now every topological space X induces a partial order $\mathcal{P} = (\mathscr{P}(X), \subseteq)$. Here an example for a closure operation is taking the topological closure $A \mapsto \overline{A}$, since it holds for all $A \subseteq X$ that $A \subseteq \overline{A}$ and $\overline{\overline{A}} = \overline{A}$.

Definition 1.2 (Comonad). A *Comonad* (L, ε, ω) in a Category $\mathcal{A}$ consists of

- an endofunctor $L\colon \mathcal{A} \to \mathcal{A}$
- a natural transformation $\varepsilon\colon L \Rightarrow \mathrm{id}_{\mathcal{A}}$
- a natural transformation $\omega\colon L \Rightarrow L^2$

such that the following diagrams commute:

(a)
$$\begin{tikzcd} L \arrow[r, "L\omega"] \arrow[d, "\omega L"'] & L^2 \arrow[d, "L\omega"] \\ L^2 \arrow[r, "\omega L"] & L^3 \end{tikzcd} \quad \text{and} \quad \begin{tikzcd} L & L^2 \arrow[l, "\varepsilon L"'] \arrow[r, "L\varepsilon"] & L \\ & L \arrow[ul, "\mathrm{id}_L"] \arrow[u, "\omega"'] \arrow[ur, "\mathrm{id}_L"'] & \end{tikzcd}$$

diagram umdrehen

Definition 1.3 (Morphism of monads). Let $\mathcal{X}$ be a category, let (T, η, μ) and (T', η', μ') be monads in $\mathcal{X}$. We say that a natural transformation $\delta \colon T \Rightarrow T'$ is a *morphism of monads* if it preserves the unit and the multiplication, i.e. the following diagrams commute:

$$\begin{tikzcd} \mathrm{id}_x \arrow[r, "\eta_x"] \arrow[dr, "\eta'_x"'] & Tx \arrow[d, "\delta_x"] \\ & T'x \end{tikzcd}$$

$$\begin{tikzcd} T^2x \arrow[r, "\mu_x"] \arrow[d, "\delta T' \circ T\delta"'] & Tx \arrow[d, "\delta_x"] \\ T'^2 \arrow[r, "\mu'_x"] & T'x \end{tikzcd}$$

show that the other composition is the same(siehe iPad)

definition

Definition 1.4 (Morphism of comonads).