Incremental clone detection for IDEs using dynamic suffix arrays

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Motivation

- Duplicated code is generally considered harmful to software quality
- Code clone detection, analysis and management is therefore important
- Incremental clone detection algorithms have not been thoroughly researched
- Incremental algorithms are useful in use-cases such as in IDEs

Our contribution

- CCDetect-LSP: An incremental clone detection tool for IDEs.
- Uses a novel application of dynamic extended suffix arrays for clone detection
- Language- and IDE agnostic via Tree-sitter and LSP

Code clones

Definition (Code snippet)

A code snippet is a piece of contiguous source code in a larger software system.

Definition (Code clone)

A code clone is a code snippet which is equal or similar to another code snippet. The two code snippets are both code clones, and together they form a clone pair. Similarity is determined by some metric such as number of equal lines of code.

Clone types

- Code clones are classified into four types
 - Type-1: Syntactically identical
 - Type-2: Structurally identical
 - Type-3: Structurally similar
 - Type-4: Functionally similar (generally)

Clone type examples: type-1 and type-2

```
print(i);
                    Figure: Type-1 clone pair
for (int i = 0; i < 10; i++) {
    print(i);
}</pre>
for (int j = 5; j < 20; j++) {
    print(j);
}
```

Clone type examples: type-3 and type-4

```
for (int i = 0; i < 10; i++) { | for (int i = 0; i < 10; i++) {
   print(i);
                                      print(i);
                                        print(i*2);
                         Figure: Type-3 clone pair
               print((n*(n-1))/2) \mid int sum = 0;
                                    for (int i = 0; i < n; i++) {
                                        for (int j = i+1; j < n; j++) {</pre>
                                            sum++;
                                     print(sum);
```

Code clone theory

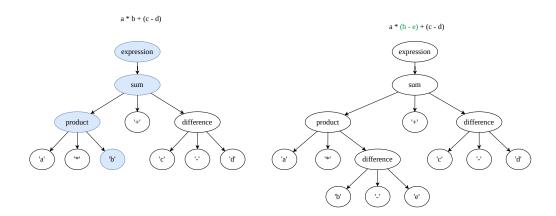
Clone detection



Clone matching techniques

- Text-based detection
 - Match based on raw source code
- Token-based detection
 - Match based on tokens
- Syntactic detection
 - Match based on AST
- Hybrid detection
 - Combine multiple approaches

Parsing and incremental parsing



Suffix tree

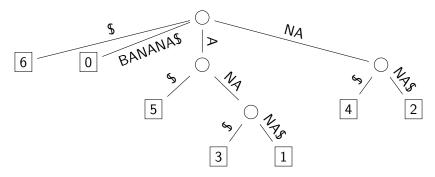


Figure: Suffix tree for S = BANANA\$

Suffix array

Index	Suffix			
0	BANANA\$			
1	ANANA\$			
2	NANA\$			
3	ANA\$			
4	NA\$			
5	A\$			
6	\$			
(a) Suffixes				

Index	Suffix
6	\$
5	A\$
3	ANA\$
1	ANANA\$
0	BANANA\$
4	NA\$
2	NANA\$

(b) Sorted suffixes

Index	SA	ISA	LCP
0	6	4	0
1	5	3	0
2	3	6	1
3	1	2	3
4	0	5	0
5	4	1	0
6	2	0	2

(c) SA, ISA and LCP

Burrows-Wheeler transform

Index	CS	Index	CS	L	F
0	BANANA\$	6	\$BANAN A	0	$Rank_A(0) + C[A] = 0 + 1 = 1$
1	ANANA\$B	5	A\$BANA N	1	$Rank_N(1) + C[N] = 0 + 5 = 5$
2	NANA\$BA	3	ANA\$BA N	2	$Rank_N(2) + C[N] = 1 + 5 = 6$
3	ANA\$BAN	1	ANANA\$ B	3	$Rank_B(3) + C[B] = 0 + 4 = 4$
4	NA\$BANA	0	BANANA\$	4	$Rank_{\$}(4) + C[\$] = 0 + 0 = 0$
5	A\$BANAN	4	NA\$BAN A	5	$Rank_A(5) + C[A] = 1 + 1 = 2$
6	\$BANANA	2	NANA\$B A	6	$Rank_{A}(6) + C[A] = 2 + 1 = 3$
		(e) Sorted	d cyclic shifts		(f) LF function

Table: S = BANANA, BWT = ANNBAA

CCDetect-LSP features

- CCDetect-LSP is implemented as an LSP server
 - List clones
 - Display clones inline with code
 - Jump between matching clones
 - Incremental updates on each edit



Implementation: Initial clone detection

- Algorithm which initially detects type-1 and optionally type-2 clones
- Pipeline of 5 phases, returns a list of clones
- Uses an extended suffix array for match detection
- Starting point: Assume documents are indexed

Detection algorithm overview

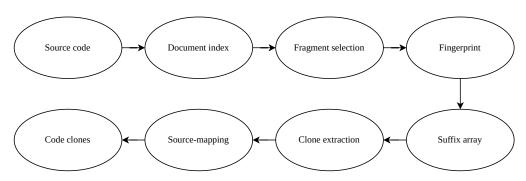


Figure: Overview of detection algorithm phases

Phase 1: Fragment selection

- Parse files using Tree-sitter
- Use a configurable Tree-sitter query to "capture" nodes
- Extract and store the tokens of captured nodes

```
(method_declaration) @method (constructor_declaration) @constructor
```

Phase 2: Fingerprinting

- Consistently hash each token value with an increasing integer counter
- Store the fingerprint of each fragment in the document index
- For type-2 detection, hash the token type instead

Phase 2: Fingerprinting

```
public class Math() {
    public int multiplyByTwo(int param) {
        return param * 2;
    }
    public int addTwo(int param) {
        return param + 2;
    }
}
```

Token	Fingerprint
public	2
int	3
multiplyByTwo	4
(5
param	6
)	7
{	8
return	9
*	10
2	11
;	12
}	13
addTwo	14
+	15

[2, 3, 4, 5, 3, 6, 7, 8, 9, 6, 10, 11, 1, 2, 3, 14, 5, 3, 6, 7, 8, 9, 6, 15, 11, 1, 0]

Figure: Example fingerprint of Java source-code

Phase 3: Suffix array construction

- Concatenate the fingerprints of each document in the index
- Construct SA, ISA and LCP array of the full fingerprint
- Uses "Induced sorting variable-length LMS-substrings" (SA-IS) algorithm
- LCP algorithm slightly modified

```
F:[2, 3, 4, 5, 3, 6, 7, 8, 9, 6, 10, 11, 1, 2, 3, 14, 5, 3, 6, 7, 8, 9, 6, 15, 11, 1, 0] SA:[26,25,12,0,13,1,4,17,14,2,3,16,5,18,9,22,6,19,7,20,8,21,10,24,11,15,23] ISA:[3,5,9,10,6,12,16,18,20,14,22,24,2,4,8,25,11,7,13,17,19,21,15,26,23,1,0] LCP: [0, 0, 0, 0, 2, 0, 1, 6, 1, 0, 0, 7, 0, 5, 1, 1, 0, 4, 0, 3, 0, 2, 0, 0, 1, 0, 0]
```

Phase 4: Clone extraction

- Use SA, ISA and LCP to find clone positions
- Linear scan through the fingerprint
- Skip contained clones

Phase 5: Source-mapping

- Map the positions of clones back to the original source-code
- Find the correct file and position of an index in the fingerprint

Implementation: Incremental clone detection

- Convert the algorithm to an incremental one
- Input is now the file which has changed and potentially the range
- Dynamic suffix array with edit operations as input

Phase 1: Update document index and fragment selection

- Store the AST of the opened files
- If range available, incrementally parse the changed file
- Mark changed files
- Fragment selection unchanged

Phase 2: Fingerprinting

- Each document stores its fingerprint
- Only need to fingerprint (and fragment select) changed files

Phase 2.5: Edit operations

- Input to phase 3: Edit operations
- How to determine edit operations?
- Edit distance algorithm!
- "Batched" operations preferred

		D	E	М	0	С	R	Α	T
	0	1	2	3	4	5	6	7	8
R	1	1	2	3	4	5	5	6	7
Е	2	2	1	2	3	4	5	6	7
Р	3	3	2	2	3	4	5	6	7
U	4	4	3	3	3	4	5	6	7
В	5	5	4	4	4	4	5	6	7
L	6	6	5	5	5	5	5	6	7
I	7	7	6	6	6	6	6	6	7
С	8	8	7	7	7	6	7	7	7
Α	9	9	8	8	8	7	7	7	8
N	10	10	9	9	9	8	8	8	8

Table: REPUBLICAN → DEMOCRAT

Optimize edit distance

- Standard algorithm memory usage is too high
- Need to optimize
 - Compare new/old fingerprint of changed document only
 - Remove trivial part at each end of matrix
 - Hirschberg's algorithm

		F	1	N	1	S	Н	1	N	G
	0	1	2	3	4	5	6	7	8	9
F	1	0	1	2	3	4	5	6	7	8
Α	2	1	1	2	3	4	5	6	7	8
S	3	2	2	2	3	3	4	5	6	7
С	4	3	3	3	3	4	4	5	6	7
	5	4	3	4	3	4	5	4	5	6
N	6	5	4	3	4	4	5	5	4	5
Α	7	6	5	4	4	5	5	6	5	5
Т	8	7	6	5	5	5	6	6	6	6
I	9	8	7	6	5	6	6	6	7	7
N	10	9	8	7	6	6	7	7	6	7
G	11	10	9	8	7	7	7	8	7	6

Table: FASCINATING \rightarrow FINISHING ASCINAT \rightarrow INISH

Phase 3: Dynamic suffix array

- Update suffix array based on edit operations
- "Four-stage algorithm for updating a Burrows-Wheeler transform"
- Updates to the BWT correlates with updates to the SA and ISA
- Inserting a single character leads to:
 - A new character in the BWT
 - A changed character in the BWT
 - 0 or more reordering of characters

Phase 3: Dynamic suffix array

Order	F	L
6	\$	Α
5	Α	N
3	Α	N
1	Α	В
0	В	\$
4	N	Α
2	N	Α
(a) Origi	nal R	\/\T

al	BWT

Order	F	L	
7	\$	Α	
6	Α	N	
4	Α	N	
1	Α	В	
0	В	\$	
2	В	Α	Inserted
5	В	Α	
3	N	В	$A \rightarrow B$
	'	'	

(b) After change and insert

F Order \$ Α 6 Α Ν В Α Α Ν В В Α 5 В Α 3 Ν В

(c) After reordering

Table: Updating BWT dynamically for the string BANANA $\$ \rightarrow BABNANA\$$

Dynamic extended suffix array

Updating suffix array is slow

■ Inserting: O(n)

■ Incrementing: O(n)

■ We therefore change the data structure which stores SA, ISA and LCP

Dynamic extended suffix array

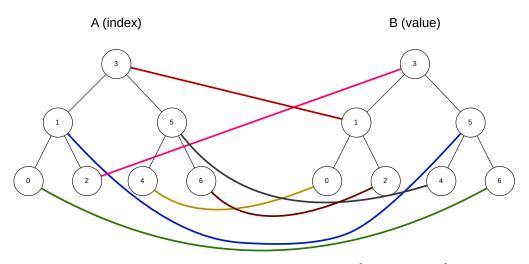


Figure: Dynamic permutation for the permutation [6, 5, 3, 1, 0, 4, 2].

Updating LCP values

■ SA updates correlate with LCP values which need to be updated

			+	\rightarrow		INS		
BWT	Α	Ν	В	Ν	\$	Α	Α	
SA	7	6	4	1	0	2	5	3
Old LCP	0	0	1	3	0	<u>N</u>	<u>0</u>	2
New LCP	0	0	1	1	0		0	2

Phase 4 and 5: Clone extraction and source-mapping

- Very similar to the initial detection
- Accessing SA, ISA and LCP is now a bit slower, but this is optimized

Evaluation

- CCDetect-LSP evaluation:
 - Verify correctness with BigCloneEval
 - Informal complexity analysis
 - Benchmark performance
 - Benchmark memory usage
 - Test multiple languages and IDEs

BigCloneBench

- A large database of clones in a Java dataset
- BigCloneEval can evaluate detection tools on BigCloneBench

```
-- Recall Per Clone Type (type: numDetected / numClones = recall) --
Type-1: 23209 / 23210 = 0.999956915122792
Type-2: 3542 / 3547 = 0.9985903580490555
Type-2 (blind): 242 / 245 = 0.9877551020408163
Type-2 (consistent): 3300 / 3302 = 0.9993943064809206
```

Figure: BigCloneEval evaluation report for CCDetect-LSP

Performance evaluation

- CCDetect-LSP was evaluated on multiple codebases
- Performance compared against SACA detection and iClones
- Incremental updates were randomly generated
- 10×10 or 10×100 tokens inserted/deleted

Code base	LOC	Clones detected	LCP _{avg}	LCP _{≥100}
WorldWind	550KLOC	1517	18	63967
neo4j	1MLOC	1313	9	27557
graal	2.2MLOC	2012	28	154452
flink	2.3MLOC	4729	13	155754
elasticsearch	3.2MLOC	9986	14	289511
intellij-community	5.8MLOC	3585	19	336190

Table: Properties of code bases

WorldWind (550KLOC)

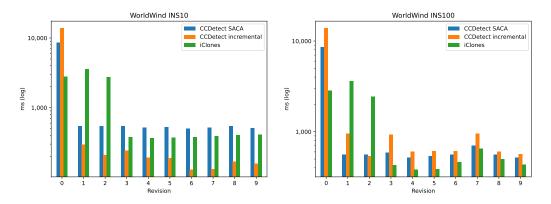


Figure: WorldWind performance benchmark

neo4j (1MLOC)

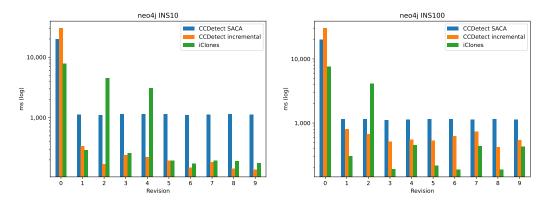


Figure: neo4j performance benchmark

graal (2.2MLOC)

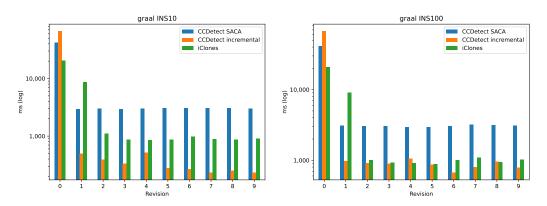


Figure: graal performance benchmark

flink (2.3MLOC)

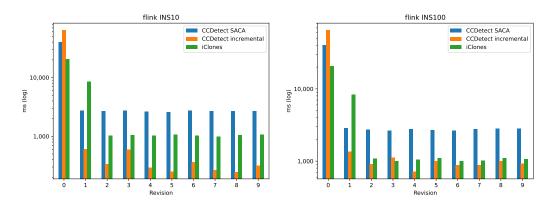


Figure: flink performance benchmark

elasticsearch (3.2MLOC)

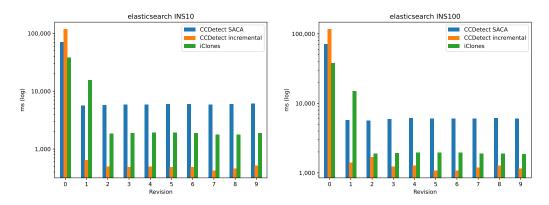


Figure: elasticsearch performance benchmark

intellij-community (5.8MLOC)

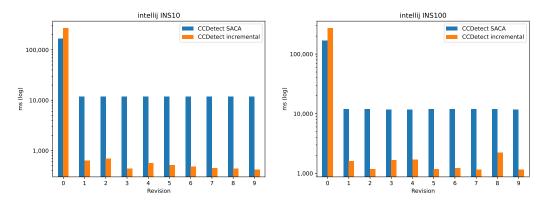


Figure: intellij-community performance benchmark

elasticsearch (3.2MLOC)

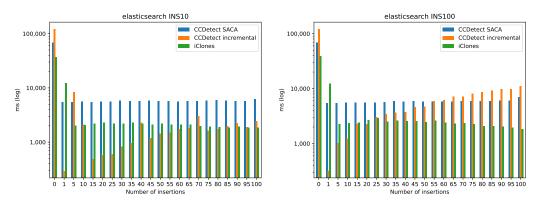


Figure: Elasticsearch performance benchmark with increasing number of edits

Memory usage

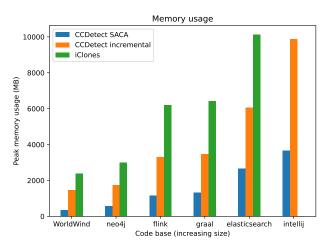


Figure: Peak memory usage of each tool when running the DEL10 test for each code base.

L Evaluation

Language and IDE agnostic

- CCDetect tested on 6 languages
 - Java
 - Python

 - Rust
 - Javascript
 - Go
- And 2 IDEs
 - Neovim
 - VSCode
- Works well in our experience

Discussion

- Performance
 - Incremental detection has best performance if edits are small
- Memory usage
 - SACA detection best, incremental detection better than iClones
 - Memory usage is probably a bottle-neck for usage
- Language agnostic
 - Works well if grammar is correct
- IDE agnostic
 - Functionality should work in most IDEs
 - Setup and configuration depends on IDE

Conclusion

■ Why do people use clone detection tools?

L Demo?

Demo

■ Demo!

Implementation: LSP architecture and functionality

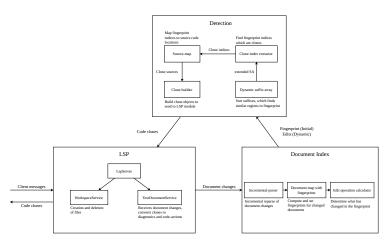


Figure: Architecture of CCDetect-LSP

LSP

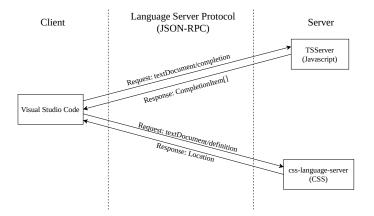


Figure: Example LSP server communication