

# DMFS: COMBINATORICS

C I

COMBINATORICS: Counting different objects/possibilities

Ex 1) Outcomes for gold, silver, bronze medals

for 10 competitors

Number ways  
2) ~~Sequence~~ of three dice can come up

3) Possible 5-card hands

- o Multiplication principle
- o Ordered / unordered; with/without repetition
- o When order matters
  - sequences (with repetitions)
  - permutations (without repetitions)
- o When order doesn't matter
  - subsets (without repetitions)
  - multisets (with repetitions)

## MULTIPLICATION PRINCIPLE

If  $n_1$  ways of performing task  $T_1$   
"  $n_2$  " "  $n_k$  "  $T_k$

then  $n_1 \cdot n_2 \cdots n_k$  ways of performing both tasks together

Ex CS researcher outfits

Blue or black jeans

white, black, or blue t-shirt

How many outfits?

## GENERAL MULTIPLICATION PRINCIPLE

$k$  tasks,  $n_i$  ways for task  $i$

Then  $\prod_{i=1}^k n_i = n_1 \cdots n_k$  ways of performing all tasks together

How many subsets of a set of  $n$  elements? CII

For each element  $s_i \in S$ , either include or not  
2 choices.

All in all  $2^n$  different choices = subsets  
(including empty set)

SEQUENCE ordered collection/list of elements

Ex  $(1, 2)$ ;  $(2, 2, 2)$ ;  $(2, 1)$ ;  $(1, 1, 2, 3, 5, 8)$

Note  $(1, 2) \neq (2, 1)$

Sequences without repetitions: PERMUTATION

SET unordered collection of elements without repetitions

$\{1, 2\}$ ;  $\{2, 1, 3\}$ ;  $\{2, 1\}$ ;  $\{1, 2, 3, 5, 8\}$

Note  $\{1, 2\} = \{2, 1\}$

MULTISET unordered collection of elements, possibly  
with repetitions

$[1, 2]$ ;  $[2, 1, 2]$ ;  $[2, 1]$ ;  $[1, 1, 2, 3, 5, 8]$

Note  $[1, 2] = [2, 1] \neq [2, 3, 2]$

	WITH REPETITIONS	WITHOUT REPETITIONS
ORDERED	SEQUENCE	PERMUTATION
UNORDERED	MULTISET	SET

How many different sequences of length  $r$   
choosing elements from set of size  $S$ ?

$r$  choices;  $n$  alternatives every time:  $n^r$  sequences  
by multiplication principle

How many length- $r$  sequences without repetitions?

$$n \cdot (n-1) \cdots (n-r+1) = {}^n P_r$$

"number of permutations of  $n$  objects taken  $r$  at a time"

$$\cdot {}^n P_r = \frac{n!}{(n-r)!}$$

10 competitions, # gold-silver-bronze outcomes?

$$10 \cdot 9 \cdot 8 = 720$$

How many permutations of whole set  $S$ ?

$${}^n P_n = n!$$

"permutation of  $S$ " tends to mean permutation of all of  $S$

What if we don't care about order?

# size- $r$  subsets of  $S$

$$\frac{n!}{(n-r)!} \text{ possibilities when order matters}$$

Each subset can be chosen in  $r!$  ways

$$\text{So } \frac{n!}{r!(n-r)!} = {}^n C_r \text{ subsets}$$

More common notation  $\binom{n}{r}$  "n-choose-r"

BINOMIAL COEFFICIENT

"number of combinations of  $n$  objects taken  $r$  at a time"

$$(x+y)^n = \sum_{i=0}^n \binom{n}{i} x^i y^{n-i}$$

Note  $\binom{n}{r} = \binom{n}{n-r}$

# ways to choose  $r$  elements =

# ways to leave  $n-r$  elements unchosen

How many different 5-card hands?

$$\binom{52}{5} = \binom{52}{5} = \frac{52!}{5!47!} = \frac{52 \cdot 51 \cdot 50 \cdot 49 \cdot 48}{5 \cdot 4 \cdot 3 \cdot 2 \cdot 1}$$

$$= 2,598,960$$

Combinations with repeats:

$S$  set with  $n$  elements

How many multisets of  $r$  elements?

If  $S$  is a set of  $n$  elements, then there are  $\binom{n+r-1}{r} = \binom{n+r-1}{n-1}$  multisets of size  $r$  with elements from  $S$ .

Proof

Need to show how to count so that

- each multiset counted once
- no multiset counted more than once
- nothing except multisets counted

Imagine we have  $n+r-1$  cells on a row  
Colour  $r$  cells red to encode multiset

CV



Assume  $|S|=n=7$  and  $r=5$ . Every multiset of 5 elements from  $S$  should result in some colouring of 5 out of  $7+5-1=11$  boxes as above

### FROM MULTISSET TO COLOURING

Let  $S = \{s_1, s_2, \dots, s_n\}$

MULTISSET 2 COLOUR (e will take values  $s_1, s_2, \dots, s_n$ )

pos := 1      in that order

for e in S      Note that we can have  $c=0$

    Let  $c := \#$  copies of e in multiset

    for i := 1 upto c

        colour box in position pos red

    pos := pos + 1

pos := pos + 1 // skip a cell

ACTUALLY NOT QUITE RIGHT. MORE CAREFUL ANALYSIS NEEDED

Algorithm will colour  $r$  cells red

Index pos will loop over  $1, 2, \dots, n+r-1$

and final increment  $pos := pos + 1$

will set  $pos = n+r$  (just to the right of boxes)

SEE  
NEXT  
PAGE

Different multisets will give different colourings  
For two multisets  $M \neq M'$ , look at first  
 $s_i$  that appears different # times in  $M$  and  $M'$   
These colourings will diverge.

## Detour

CI + 8/2

More careful analysis of  
MULTISET 2 COLOUR:

By construction, if multiset has  $r$  elements,  
then algorithm will colour  $r$  cells red.  
Great!

But we need to check that it never writes  
out of bounds, i.e., in position  $\text{pos} \geq n+r$

Note that at end of algorithm we have  
 $\text{pos} = n+r+1$ , so there is reason to worry

Consider when  $e = s_n$  at top of  
for loop. Then

$$\boxed{\text{pos} = 1 + (n-1) + (\text{#elements in multiset except } s_n)}$$

Also note that pos was incremented after colouring, so all coloured cells have indices

$$\leq (n-1) + (\text{# elements in multiset except } s_n)$$

$$\leq n-1+r \quad \text{OK so far!}$$

If we have seen all of multiset, then

$$\boxed{\text{pos} = n+r}$$

but no further colouring will be made.

Otherwise,  $c$  copies of  $s_n$  will colour positions  $n+r-c, n+r-c+1, \dots, n+r-2, n+r-1$   
so we are OK. GOOD, LET'S GO BACK TO THE PROOF...

Since all multisets yield different colourings, this proves that

CV+E

$$\boxed{\# \text{multisets} \leq \binom{n+r-1}{r}}$$

But it could be that there are more colourings than multisets! To show this is not so, describe how to translate any set of coloured boxes (like on the previous page) to a multiset!

### FROM COLOURING TO MULTISET

$$S = \{s_1, s_2, \dots, s_n\}$$

#### COLOUR 2 MULTISET

pos := 1

i := 1 // index of element in S

while pos < n + r

if cell in position pos is red

output copy of si to multiset

else

i := i + 1

pos := pos + 1

There are r red cells (by assumption) so

i will be incremented  $(n+r-1) - r = n-1$  times,

i.e., loops over 1, 2, ..., n so S; ALWAYS VALID ELEMENT

For same reason, OUTPUTS EXACTLY r ELEMENTS

So we get valid multiset. Can also see that different colourings yield different multisets, so

$$\boxed{\# \text{multisets} \geq \binom{n+r-1}{r}}$$



WHY? →

ILLUSTRATION OF PROOF | CV 1/2  
 set of  $n$  elements,  $|S|=n$   
 Choose multiset of  $r$  elements  
 (order doesn't matter, but count copies)  
 Claim:  $\binom{n+r-1}{r}$  such sets

= Given row with  $n+r-1$  cells, how many ways to colour  $r$  cells red

Present "translations"

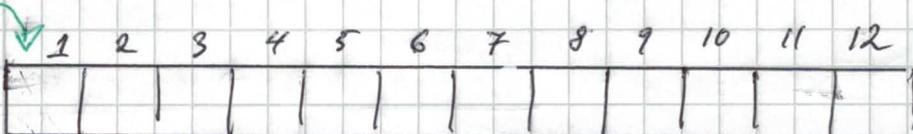
- (a) from multiset to colouring
- (b) from colouring to multiset

"COMBINATORIAL PROOF"

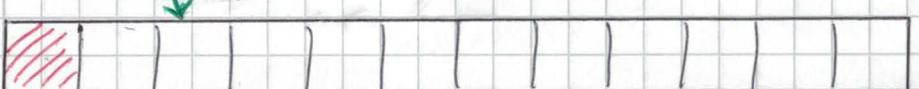
Example  $n=7, r=6$

Consider multiset  $[1, 3, 3, 3, 4, 4]$

START HERE  $n+r-1 = 12$  cells



Take care of 1



No 2, take care of 3,3,3



Take care of 4,4



No 5, no 6, so done!

no 7

For 5: skip from position 11 to 12

For 6: skip from position 12 right

For 7: Do nothing — last element in universe

How to translate back?

C V 2/3

We have  $n = 7$ ,  $r = 6$

Consider

1	2	3	4	5	6	7	8	9	10	11	12
█			█	█	█	█	█	█			

Translate back to multiset

① Current position 1

1 red, then blank

$\Rightarrow$  One copy of 1 in multiset

② Current position 3

Blank

$\Rightarrow$  No 2 in multiset

③ Current position 4

3 red, then blank

$\Rightarrow$  Three copies of 3 in multiset

④ Current position 8

2 red, then blank

$\Rightarrow$  Two copies of 4 in multiset

⑤ Current position 11

Blank

$\Rightarrow$  No 5 in multiset

⑥ Current position 12

Blank

$\Rightarrow$  No 6 in multiset

⑦ Current position 13  
Now out of bounds

$\Rightarrow$  No 7 in multiset

Makes sense, since we have already chosen our 6 elements

Constructed multiset:  $[1, 3, 3, 3, 4, 4]$