

A Generalized 2D and 3D Hilbert Curve

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Abstract

The two and three dimensional Hilbert curves are fractal space filling curves that map the unit interval to the unit square or unit cube while preserving a notion of closeness, or locality, after the map. We present the Gilbert curve, a conceptually straight forward generalization of the Hilbert curve, that works on arbitrary rectangular regions, overcoming the limitation of the Hilbert curve that requires side lengths be exact powers of two. The construction provides reasonable worst case run-time guarantees for random access lookups for both two and three dimensions. We provide experiments showing comparable quality in locality of the Gilbert curve to the Hilbert curve for a range of cuboid regions and discuss some limitations of the construction algorithm.

Generalized Hilbert Curves

In a website application and scanned note [1], Tautenhahn provided the basis for a generalized 2D space filling curve. Tautenhahn also included policies for when to subdivide regions preferentially in only one dimension that help to create more *harmonious* curve realizations. Tautenhahn's exploration details the parity arguments necessary for when a subdivision scheme can be employed without creating diagonal moves (*notches*).

In this paper, we extend Tautenhahn's ideas to create a 2D and 3D generalized Hilbert curve. We further extend Tautenhahn's core ideas on when to use alternate subdivision schemes when one length is much larger than the rest, what we call *eccentric cases*, and apply them to a 3D generalized Hilbert curve. Tautenhahn's unpublished reasoning behind the constants used in the 2d eccentric split case are briefly discussed later ¹.

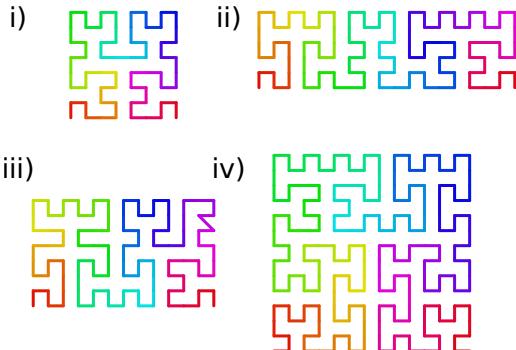


Figure 1: 2D Gilbert curves for i) 8×8 , ii) 18×6 ,
iv) 13×8 (with notch), iv) 14×14

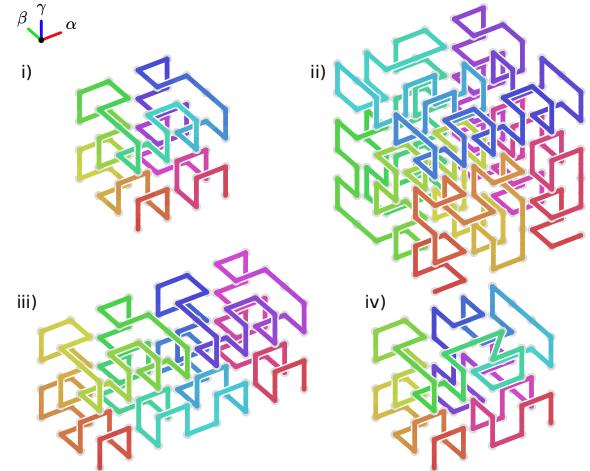


Figure 2: 3D Gilbert curves for i) $4 \times 4 \times 4$, ii)
 $6 \times 6 \times 6$, iii) $8 \times 4 \times 4$, iv) $5 \times 4 \times 4$
(with notch)

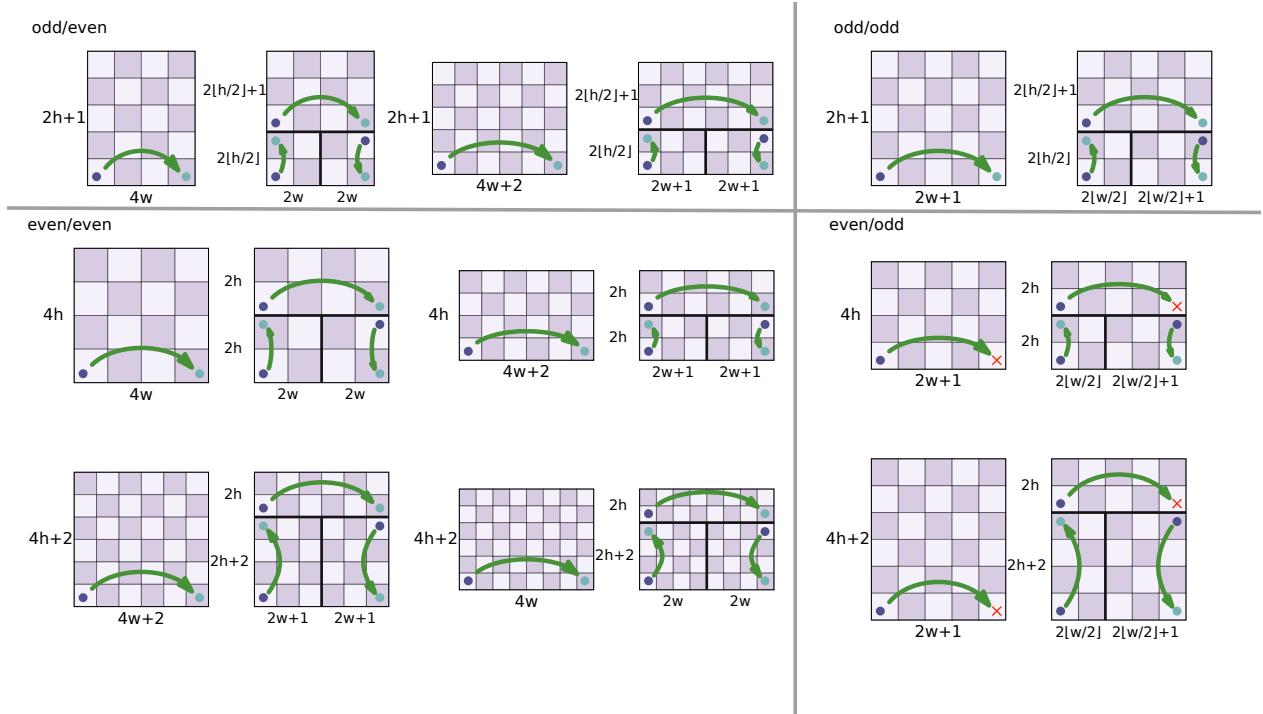


Figure 3: Enumeration of the subdivision template depending on different parities of α and β dimensions.

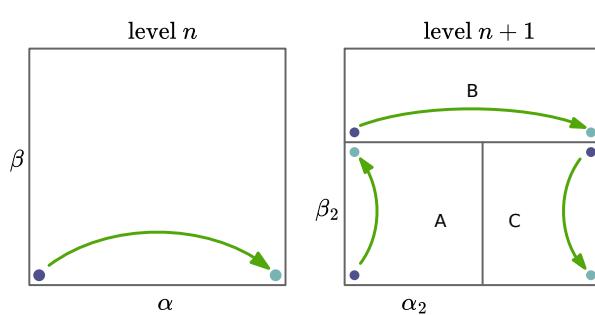


Figure 4: Subdivision strategy for the 2D Gilbert curve.

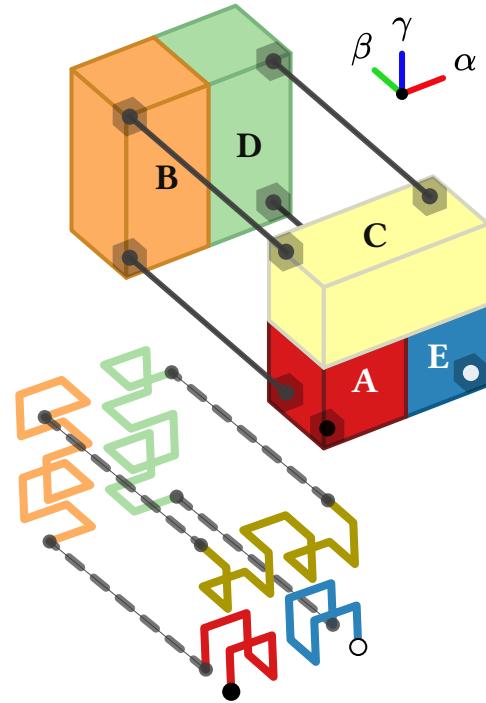


Figure 5: The main subdivision strategy for the 3D Gilbert curve.

Subdivision Strategies

For both the 2D and 3D Gilbert curve, a subdivision template is chosen to recursively partition the region. Figure 4 shows the 2D template and figure 5 shows the 3D template.

When dividing into sub-regions, side lengths are chosen to be integral and with even side lengths preferred for the first subdivided region. A path is recursively chosen for each of the subregions and, after resolution, endpoints are connected.

When regions have a side length that is much larger than the rest, another subdivision scheme is used. We call these cases *eccentric subdivisions* as the larger side length makes the region lopsided.

If one side length is much larger than the rest, both the 2D and 3D Gilbert curve split the region in two. For the 3D case, if one side length is much smaller than the rest but with the remaining lengths roughly equal, the region is subdivided as if it were a 2D Gilbert curve, taking the smaller side length as the “flat” dimension.

The eccentric subdivision schemes provide a more visually pleasing solution as regimented subdivisions might paths that are long and skinny or wide and squat. We will briefly go over motivations for choosing how to choose these constants and why the particular value of constants was chosen in the next section.

For the 2D Gilbert curve, if $|\alpha|$ is odd, there will be a diagonal move, or *notch*, for the resulting path. The subdivision scheme will keep the notch in a single region, limiting the notch count to one.

For the 3D Gilbert curve, if any of the subdivided regions has endpoints that lay in a direction with odd length, a notch will appear. This means the number of notches for the 3D Gilbert curve isn’t limited to one.

Eccentric Subdivision

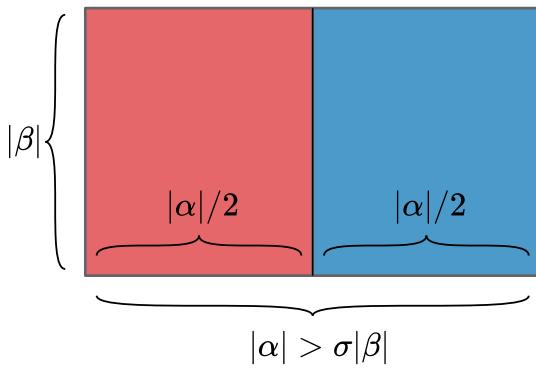


Figure 6: When the length of the width-like dimension ($|\alpha|$) gets larger than a threshold of the height-like dimension ($\sigma|\beta|$), then the width like dimension is split into two.

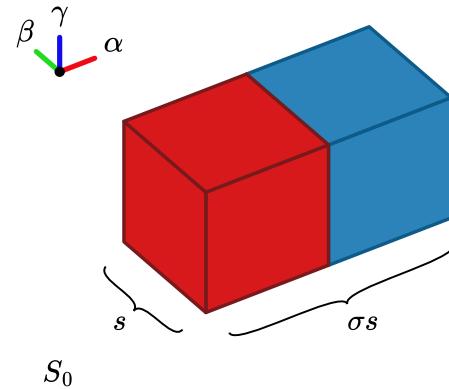


Figure 7: An S_0 eccentric subdivision showing the threshold, σ , used to determine when we should use this split.

When side lengths become too lopsided, an alternate *eccentric* subdivision scheme is chosen to split along the larger side length.

¹Through personal communication with the authors, Tautenhahn kindly provided the reasoning for the constants used in the eccentric split

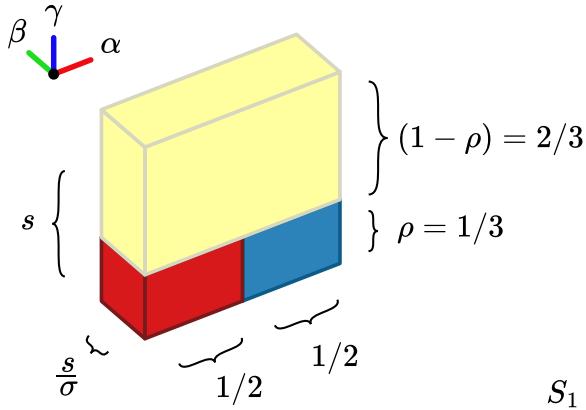


Figure 8: An S_1 eccentric subdivision showing the threshold, σ , used to determine when we should use this split and the split point ratio, ρ , to reduce the defect for a more harmonious subdivision.

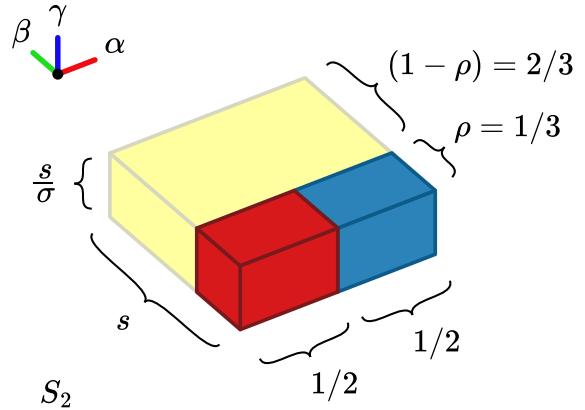


Figure 9: An S_1 eccentric subdivision showing the threshold, σ , used to determine when we should use this split and the split point ratio, ρ , to reduce the defect for a more harmonious subdivision.

For each eccentric subdivision, a proportion threshold needs to be chosen to determine when a side lengths become too lopsided to warrant an alternate subdivision. Here we motivate the eccentric subdivision constants for the 2D case.

Consider a *defect* function, $\lambda_2 : \mathbb{N}^2 \mapsto \mathbb{N}$, that measures the area relative to what the area would be if it just the minimum side length were taken:

$$\lambda_2(|\alpha|, |\beta|) = \frac{|\alpha| \cdot |\beta|}{\min(|\alpha|, |\beta|)^2}$$

If there is a disjoint subdivision of a volume V_0 to $V_1 = (V_{0,0}, V_{0,1}, \dots, V_{0,m-1})$, $V_0 = \cup_k V_{0,k}$, define the *average defect* of the subdivided volume to be the sum of defects weighted by their proportional volumee:

$$\lambda_s(V_1) = \sum_k \frac{\text{Vol}(V_{0,k})}{\text{Vol}(V_0)} \cdot \lambda_2(V_{0,k})$$

The defect gives a coarse idea of how lopsided or *eccentric* a cuboid region is. If the defect is too high, we might want to split the larger sides while keeping the smaller sides the same size. Reducing the average defect attempts to make each subdivided region more square-like. When subdivided areas are more square-like, we say that the subdivided regions are more *harmonious*.

If the α side length is significantly longer than the β side length, we want to subdivide the rectangle into two nearly equal regions and recursively find a Gilbert curve in each region. We will justify the constant $(3/2)$ as the ratio threshold to split on using an argument originally developed by L. Tautenhahn ².

The defect of a rectangle of side length $|\alpha|$ and $|\beta|$ is, with $|\alpha| > |\beta|$:

$$\begin{aligned} \lambda_2(|\alpha|, |\beta|) &= |\alpha| \cdot |\beta| / |\beta|^2 \\ &= |\alpha| / |\beta| \end{aligned}$$

After a subdivision, if we assume $|\alpha| < 2|\beta|$, the defect is:

²Given to us through personal communication with L. Tautenhahn

$$\begin{aligned}\lambda_2(|\alpha|/2, |\beta|) &= \frac{|\alpha|}{2} \cdot |\beta| / (\frac{|\alpha|}{2})^2 \\ &= 2|\beta|/|\alpha|\end{aligned}$$

We're looking for the condition when there's a defect reduction, so

$$\begin{aligned}\lambda_2(|\alpha|/2, |\beta|) &< \lambda_2(|\alpha|, |\beta|) \\ \rightarrow 2|\beta|/|\alpha| &< |\alpha|/|\beta| \\ \rightarrow \sqrt{2} &< |\alpha|/|\beta|\end{aligned}$$

Since $\sqrt{2} \approx 1.4142 < (3/2)$, if we choose $|\alpha|/|\beta| > (3/2)$ we can be assured a defect reduction.

In the case when $|\alpha| > 2|\beta|$, it is easy to verify that the defect is always reduced ³.

Defect reductions for the 3D case can be done but are more complicated. The constants for the 3D subdivision in this paper were observed to yield visually pleasing results.

Algorithm

Algorithm 1

```
#  $p, \alpha, \beta \in \mathbb{Z}^3$ 
function GILBERT2D( $p, \alpha, \beta$ )
     $\alpha_2, \beta_2 = \text{div}(\alpha, 2), \text{div}(\beta, 2)$ 
    if ( $|\beta| \equiv 1$ ) then
        yield  $p + i \cdot \delta(\alpha)$  forall  $i \in |\alpha|$ 
    else if ( $|\alpha| \equiv 1$ ) then
        yield  $p + i \cdot \delta(\alpha)$  forall  $i \in |\beta|$ 
    else if ( $2|\alpha| > 3|\beta|$ ) then
        if ( $|\alpha_2| > 2$ ) and
            ( $|\alpha_2| \bmod 2 \equiv 1$ ) then
                 $\alpha_2 \leftarrow \alpha_2 + \delta(\alpha)$ 
            end if
        yield Gilbert2D( $p, \alpha_2, \beta$ )
        yield Gilbert2D( $p + \alpha_2, \alpha - \alpha_2, \beta$ )
    else
        if ( $|\beta_2| > 2$ ) and
            ( $|\beta_2| \bmod 2 \equiv 1$ ) then
                 $\beta_2 \leftarrow \beta_2 + \delta(\beta)$ 
            end if
        yield Gilbert2D( $p, \beta_2, \alpha_2$ )
        yield Gilbert2D( $p + \beta_2, \alpha, (\beta - \beta_2)$ )
        yield Gilbert2D( $p + \alpha - \delta(\alpha) + \beta_2 - \delta(\beta), \beta_2, -(\alpha - \alpha_2)$ )
    end if
end function
```

Algorithm 2

```
#  $p, \alpha, \beta, \gamma \in \mathbb{Z}^3$ 
function GILBERT3D( $p, \alpha, \beta, \gamma$ )
    # Parity of dimensions
     $\alpha_0 \leftarrow (|\alpha| \bmod 2)$ 
     $\beta_0 \leftarrow (|\beta| \bmod 2)$ 
     $\gamma_0 \leftarrow (|\gamma| \bmod 2)$ 
    # Base cases
    if (( $|\alpha| \equiv 2$ ) and ( $|\beta| \equiv 2$ ) and ( $|\gamma| \equiv 2$ ))
        return Hilbert3D( $p, \alpha, \beta, \gamma$ )
    return Gilbert2D( $p, \beta, \gamma$ ) if ( $|\alpha| \equiv 1$ )
    return Gilbert2D( $p, \alpha, \gamma$ ) if ( $|\beta| \equiv 1$ )
    return Gilbert2D( $p, \alpha, \beta$ ) if ( $|\gamma| \equiv 1$ )
    # Eccentric cases
    if ( $3|\alpha| > 5|\beta|$ ) and ( $3|\alpha| > 5|\gamma|$ )
        return  $S_0(p, \alpha, \beta, \gamma)$ 
    if ( $2|\beta| > 3|\gamma|$ ) or ( $2|\beta| > 3|\alpha|$ )
        return  $S_2(p, \alpha, \beta, \gamma)$ 
    if ( $2|\gamma| > 3|\beta|$ )
        return  $S_1(p, \alpha, \beta, \gamma)$ 
    # Bulk recursion
    return  $J_0(p, \alpha, \beta, \gamma)$  if ( $\gamma_0 \equiv 0$ )
    return  $J_1(p, \alpha, \beta, \gamma)$  if ( $\alpha_0 \equiv 0$  or ( $\beta_0 \equiv 0$ ))
    return  $J_2(p, \alpha, \beta, \gamma)$ 
end function
```

³The ratio of the defects before and after the split is (1/2)

Algorithm 1 shows the pseudo-code for computing the 2D Gilbert curve. Note that α and β are taken to be vectors in 3D, where the third dimension can be ignored if a purely 2D curve is desired. The generalization to 3D allows the Gilbert2D function to be used unaltered when the 3D Gilbert curve needs to trace out in-plane sub-curves.

The $\delta(\cdot)$ function returns one of the six directional vectors indicating which of the major signed axis aligned directions the input vector points to $((\pm 1, 0, 0), (0, \pm 1, 0), (0, 0, \pm 1))$.

Algorithm 1 assumes standard Euclidean two norm ($|v| = \sqrt{v_0^2 + v_1^2 + v_2^2}$) and abuses notation by allowing scalar to vector multiplication ($i \in \mathbb{Z}, v \in \mathbb{Z}^3, i \cdot v \rightarrow (i \cdot v_0, i \cdot v_1, i \cdot v_2)$). where the context is clear.

Conclusion and Future Directions

Generalizing a Hilbert space filling curve to arbitrary side lengths is not a well posed problem and is open to interpretation. The Gilbert curve presented in this paper provides one interpretation for generalizing space filling curves to arbitrary side lengths.

When designing an algorithm to generalize the Hilbert curve, four features that might be desirable are:

- *Hilbert-esque*
The resulting approximation yields the Hilbert curve on side lengths that are exact powers of two
- *Stability*
Curve realizations converge as the cuboid side lengths increase proportionally
- *Harmony*
Realizations use eccentric splits to reduce the *defect*
- *Notch-Limited*
Curve realizations are notch free unless a single notch is necessary

The 2D Gilbert curve presented here has all of these properties while the 3D Gilbert curve lacks harmony and the notch-limited feature.

Future work could focus on creating 3D Gilbert curve that has all of these features, perhaps by relaxing strict endpoint requirements at the cuboid corners. Some other natural ideas are to extend the concepts above to other space filling curves or circuits, or, more generally, to try and create an arbitrary space filling curve in a cuboid region with arbitrary endpoints.

A libre/free implementation for the Gilbert curve in 2D and 3D has been developed and can be downloaded from its repository ⁴.

References

- [1] L. Tautenhahn. “Draw a space-filling curve of arbitrary size.” 2003.
https://lutanho.net/pic2html/draw_sfc.html.

⁴ <https://github.com/jakubcerveny/gilbert>

