

Multiplicative order $o_r(n)$ [slide 7]

$o_r(n)$ is the *order* of $n \bmod r$,

i.e. the smallest number $o_r(n)$ such that $n^{o_r(n)} \equiv 1 \pmod{r}$.

- e.g. $o_7(3) = 6$:
 $3^1 = 3, \quad 3^2 = 2, \quad 3^3 = 6, \quad 3^4 = 4, \quad 3^5 = 5, \quad \underline{3^6 = 1} \pmod{7}$

The inefficient primality test [slide 12]

$X + a^n \equiv (X + a)^n$ below is an equality of *polynomials*, not numbers.

For any $a \in \mathbb{Z}$ such that $\gcd(a, n) = 1$:

$$X + a^n \equiv (X + a)^n \pmod{n} \iff n \text{ is prime}$$

(For the \Leftarrow direction you don't need the $\gcd(a, n) = 1$ assumption.)

AKS pseudocode [slide 18]

- if $(n = a^b \text{ for } a, b \in \mathbb{N} \text{ and } b \geq 2)$, return COMPOSITE.
- Find the smallest r such that $o_r(n) > \log^2 n$.
- If some $a \leq r$ is not coprime with n , return COMPOSITE.
- If $r \geq n$, return PRIME.
- For $a = 1$ to $\lfloor \sqrt{\varphi(r)} \log n \rfloor$ do:
 - if $X + a^n \not\equiv (X + a)^n \pmod{X^r - 1, n}$, return COMPOSITE.
- Return PRIME.

Main challenge: show that if Line 6. is reached, then n must be prime.

Plan of the final attack [slide 25]

- Define a group \mathcal{G} of polynomials.
- Prove that $|\mathcal{G}| \geq \text{LowerBound}$.
- Prove that if $n \neq p^k$, then $|\mathcal{G}| < \text{LowerBound}$.
- Deduce that $n = p^k$ for some k .

But recall Line 1. of AKS: if $(n = a^b \text{ for } a, b \in \mathbb{N} \text{ and } b \geq 2)$, return COMPOSITE.

- Therefore $k = 1 \dots$ and we get that $n = p^1$ is prime. \square

Introspectiveness [slide 26]

Definition. $m \in \mathbb{N}$ is *introspective* for $f(X) \iff$

$$f(X^m) \equiv [f(X)]^M \pmod{X^r - 1, p}$$

Corollary. n and p are introspective for $X + a$ (for all $a \in \{0, 1, \dots, l\}$).

G_1 and \mathcal{G} [slide 29]

$$G_1 := \{n^i \cdot p^j \mid i, j \geq 0\} \pmod{r}$$

$$\mathcal{G} := \left\{ \prod_{a=0}^l (X + a)^{e_a} \mid e_a \geq 0 \right\} \pmod{h(X), p}$$