Assignment 2 - John Akujobi

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	Homework
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CSC 317 Spring 2024 | Assignment #2 | Due: 2-12-24

From the textbook (ARM edition) do the following problems that start on page 175.

Assignments are to be submitted to D2L by 2:00pm on the due date.

Exercise 2.1 (5 points)

For the following C statement, write the corresponding LEGv8 assembly code. Assume that the C variables f, g, and h, have already been placed in registers X0, X1, and X2 respectively. Use a minimal number of LEGv8 assembly instructions.

```
f = g + (h - 5);
```

```
; f = X0
; g = X1
; h = X2
SUBI X3, X2, #5 ; Subtract 5 from h(X2) and put result in ter
ADDI X0, X1, X3; Add temp to g(X1) and put result in f(X0)
```

2. Exercise 2.2 (5 points)

Write a single C statement that corresponds to the two LEGv8 assembly instructions below.

```
ADD f, g, h
```

```
// f = g + h

// f = f + i

// f = g + h + i

f = g + h + i;
```

3. Exercise 2.3 (5 points)

For the following C statement, write the corresponding LEGv8 assembly code. Assume that the variables f, g, h, i, and j are assigned to registers X0, X1, X2, X3, & X4, respectively.

Assume that the base address of the arrays A and B are in registers X6 and X7, respectively.

```
B[8] = A[i - j];
```

```
; f = X0
; g = X1
; h = X2
; i = X3
; j = X4
; A [0] = X6
; B [0] = X7
; B [8] = 8 * 8 = 64 Offset
;
```

4. Exercise 2.11 (all parts) (12 points)

Assume that registers X0 and X1 hold the values 0×80000000000000 and 0×D000000000000, respectively.

2.11.1

What is the value of X9 for the following assembly code?

```
ADD X9, X0, X1
```

X9 = 0x5000000000000000

2.11.2

Is the result in X9 the desired result, or has there been overflow?

There was overflow

2.11.3

For the contents of registers X0 and X1 as specified above, what is the value of X9 for the following assembly code?

```
SUB X9, X0, X1
```

$$X9 = 0x0B$$

2.11.4

Is the result in X9 the desired result, or has there been overflow?

```
There wasn't exactly an overflow.
But there was a wrap around
```

2.11.5

For the contents of registers X0 and X1 as specified above, what is the value of X9 for the following assembly code?

```
ADD X9, X0, X1
ADD X9, X9, X0
```

$$X9 = 0xD$$

```
ADD X9, X0, X1
0x800000000000000 = X0
```

2.11.6

Is the result in X9 the desired result, or has there been overflow?

```
The desired result is 0x1D but X9 does not contain it because of overflow
```

5. Exercise 2.14 (5 points)

Provide the instruction type and hexadecimal representation of the following instruction:

```
STUR X9, [X10, #32]
```

F8025200

```
STUR opcode is 1984ten
[X10, #32] means there is a 32 bit offset from X10
So the instruction wants to store the data from there into X9
x10, [x10, #16], [x10, #32]
That is, storing the third data from x10
Using D-format (opcode, op2, Rn, Rt, address)
opcode: 1984ten -> 0x7C0h -> 11111000000 in binary (11 bits)
op2: 00 in binary (2 bits)
Rn (X10): 1010 in binary (5 bits)
Rt (X9): 01001 in binary (5 bits)
address (#32): This needs to fit into 9 bits.
                        32 in decimal is 100000 in binary,
                        but we need to pad it to 9 bits: 0010000
= 11111000000 00 1010 01001 001000000
Splitting it into groups of 4
= 1111 1000 0000 0101 0010 0100 0000
= F8025200
```

6. Exercise 2.30 (18 points)

Implement the following C code in LEGv8 assembly. Hint: Remember that the stack pointer must remain aligned on a multiple of 16.

```
int fib (int n) {
if (n==0)
    return 0;
else if (n == 1)
```

```
return 1;
else
   return fib(n-1) + fib(n-2);
}
fib:
   //Lets assume that n is in x0
   //since that is where arguements go to
   sub sp, sp, #16 // Allocate stack space
   // Save frame pointer and return address
   stp x29, x30, [sp, #16]
   add x29, sp, #16 // Update frame pointer
   //if (n == 0) {return 0;}
   cmp x0, \#0 //check if n is 0
   beq end_if_zero //jump to return 0 if n is 0
   //else if (n == 1) { return 1;}
   cmp x0, #1 //check if n is 1
   beq end_if_one //jump to return 1 if n is 1
   //fib(n - 1)
   mov x1, x0 //copy n to x1 for keeping
   sub x0, x0, #1 //n = n - 1
                // call fib(n - 1)
   bl fib
   mov x19, x0 //store fib(n - 1) in x19
   //fib(n - 2)
   mov x0, x1 //copy n to x0
   sub x0, x0, #2 //subtract 2 from n
   bl fib
                   //call fib
   add x0, x19, x1 //add fib(n - 1) + fib(n - 2)
```

I used the below compiler alongside to translate and understand part of it

Compiler Explorer

Compiler Explorer is an interactive online compiler which shows the assembly output of compiled C++, Rust, Go (and many more) code.



