

Computer Algebra

Lecture 1

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Computer Algebra (Symbolic Computation)

“Getting computers to do algebra”

computers by themselves can't even do arithmetic: there are only finitely many `int` in C for example. And

$(1 + 10^{20}) - 10^{20} \xrightarrow{\text{double}} 0$ whereas

$1 + (10^{20} - 10^{20}) \xrightarrow{\text{double}} 1.$

to do Do we really know algorithms? Can you factor $x^5 - 2x^4 + 8x^3 + 3x^2 + 6x - 4$?

algebra and much geometry can be turned into algebra. So can much of calculus.

Some early history

[But I don't know anything about Chinese developments]

1948 Manchester Baby is first stored-program computer.

1951 (Cambridge, UK) $180(2^{127} - 1)^2 + 1$ is prime [MW51].

1952 Elliptic curve calculations on MANIAC [ST92, p. 119].

1953 two US theses [Kah53, Nol53] kicked off the 'calculus' side of computer algebra with programs to differentiate expressions.

1953 (Cambridge, UK) a group theory algorithm was implemented [Has53].

Thereafter, history rather splits into "polynomial/calculus" and "group theory".

“Polynomial/Calculus” systems

[Kah53, Nol53] could only do one thing, but soon people wrote systems, common by early 1970s.

SAC Written in Fortran [Col71], its descendant QEPCAD [Bro03] alive today

MACSYMA in LISP at MIT [MF71, PW85], available (as Maxima) in SAGE

Reduce in LISP at Utah [Hea71, Hea05] (our system)

SCRATCHPAD in LISP at IBM [GJY75], now dead, but inspired Axiom [JS92].

COBALG in COBOL at Cambridge [fHN76], now dead.

Maple Kernel in C, Waterloo (CA), 1983– [CGGG83].

Mathematica Kernel in C, Wolfram Research, 1988– [Wol88].

Are these some of the longest-lived programs?

I know less about these.

CAYLEY [BC90] was a major group theory system, replaced by

MAGMA [BCM94] a system built on Axiom-like ideas.

GAP [BL98], design inspired by Maple. Available in SAGE.

Why two different directions? Polynomial systems are concerned with a *specific* polynomial, whereas group theory systems are concerned with a *specific* group of permutations, rather than a specific permutation.

Algebraic geometry systems, such as SINGULAR [Sch03] (available in SAGE), Macauley [BS86] or COCOA [GN90], are concerned with *specific* sets of polynomials.

This is the first system I want to consider. It is relatively simple, but still current (last year won its division of the Satisfiability Modulo Theories competition [http:](http://smtcomp.sourceforge.net/2017/results-NRA.shtml)

[//smtcomp.sourceforge.net/2017/results-NRA.shtml](http://smtcomp.sourceforge.net/2017/results-NRA.shtml)).

It is free, and has a free manual [HS18].

Written in LISP, which means it's simple enough to understand.

This afternoon, you'll install and run it on your computers, but now I will demonstrate it.

Storage of Polynomials

How do we store polynomials

$$p_{\text{dense}} = \sum_{i=0}^n a_i x^i? \quad (1)$$

Your first thought might be a vector $[a_0, a_1, \dots, a_n]$ of all the coefficients. Such a method is called *dense*.

However, for $x^{1000000} + 1$ it would take megabytes of memory to store, which seems silly. After all, we didn't write down all those zeros, so why should a computer? Instead consider

$$p_{\text{sparse}} = \sum_{i=1}^m b_i x^{e_i} : b_i \neq 0 \quad (2)$$

and store the pairs (b_i, e_i) .

Such a storage method is sparse.

One of the oldest programming language, the fundamental data structure is the CONS cell, composed of two parts, known for historical reasons as CAR and CDR (though many people think of them as *first* and *rest*, which works if you have a list).



If we have such an object

A	B
---	---

, we write (A . B).

A list

A	
---	--

 →

B	
---	--

 →

C	/
---	---

 (where the last CDR points to a special object NIL), which should be (A. (B. (C. NIL))) is

written (A B C). If we had

A	
---	--

 →

B	
---	--

 →

C	D
---	---

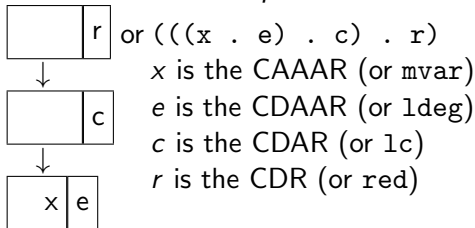
 we would write (A B C . D).

Also, instead of CAR(CDR(CDR(Z))) (the third element of the list Z), we write CADDR(Z).

I could say a lot more about LISP, but don't need to.

Reduce implementation in LISP: Polynomials

If we have a non-trivial (i.e. not just a number) polynomial p , we can write it as $p = cx^e + r$, with $e \neq 0$, e.g. $3x^2 + 5$ or $(3y + 2)x^7 + (2y^2 - 4)x^6 + 6y^4 + 7$. c is a *non-zero* polynomial *not* involving x , and r might involve x , but only to a power $< e$. r is called the *reductum*, c the *leading coefficient*, x the *main variable* and e the *exponent*.



Reduce implementation in LISP continued

Of course, not everything is a polynomial.

But we can view everything as a fraction of two polynomials $\frac{\text{numerator}}{\text{denominator}}$, known as a *standard quotient* (SQ).

This is easy:

numerator	denominator
-----------	-------------

.

Once Reduce has simplified something to being an SQ, it likes to keep it like that. Hence a standard quotient z is stored as $(!*SQ\ z\ NIL)$, or $(!*SQ\ z\ T)$ if the system knows it has been simplified.

```
69: z;
```

$$x + y$$

```
70: lisp get('z,'avalue);
```

```
(scalar (*sq (((x . 1) . 1) ((y . 1) . 1)) . 1) nil))
```

Reduce

```
73: sub(n=1023,f*g);
```

$$a^b x^{1023} + a^b y^{1023} + b^{1023} x^{1023} + b^{1023} y^{1023} + c^{1023} x^{1023} + c^{1023} y^{1023}$$

```
74: f:=a^n+b^n+c^n;
```

$$f = a^n + b^n + c^n$$

```
75: g:=x^n+y^n;
```

$$g = x^n + y^n$$

```
76: sub(n=1023,f*g);
```

$$a^{1023} x^{1023} + a^{1023} y^{1023} + b^{1023} x^{1023} + b^{1023} y^{1023} + c^{1023} x^{1023} + c^{1023} y^{1023}$$

```
77: lisp ws;
```

```
(!*sq (((a . 1023) ((x . 1023) . 1) ((y . 1023) . 1)) ((b . 1023) ((x . 1023) . 1) ((y . 1023) . 1)) ((c . 1023) ((x . 1023) . 1) ((y . 1023) . 1))) . 1)
```

Maple's poly format [MP14]

$$f := d^n + b^n + c^n;$$

$$f := d^n + b^n + c^n$$

$$g := x^n + y^n;$$

$$g := x^n + y^n$$

dismantle(expand(eval(fg, n = 1023)));

```
POLY(14)|
EXPSEQ(6)
  NAME(4): a
  NAME(4): b
  NAME(4): c
  NAME(4): x
  NAME(4): y
DEGREES(HW): ^2046 ^1023 ^0 ^0 ^1023 ^0
INTPOS(2): 1
DEGREES(HW): ^2046 ^1023 ^0 ^0 ^0 ^1023
INTPOS(2): 1
DEGREES(HW): ^2046 ^0 ^1023 ^0 ^1023 ^0
INTPOS(2): 1
DEGREES(HW): ^2046 ^0 ^1023 ^0 ^0 ^1023
INTPOS(2): 1
```

Maple's sum format

```
dismantle( expand( eval( f.g, n = 1024 ) ) );
```

```
SUM(13)
```

```
  PROD(5)
```

```
    NAME(4) : c
```

```
    INTPOS(2) : 1024
```

```
    NAME(4) : y
```

```
    INTPOS(2) : 1024
```

```
  INTPOS(2) : 1
```

```
  PROD(5)
```

```
    NAME(4) : c
```

```
    INTPOS(2) : 1024
```

```
    NAME(4) : x
```

```
    INTPOS(2) : 1024
```

```
  INTPOS(2) : 1
```

```
  PROD(5)
```

```
    NAME(4) : b
```

```
    INTPOS(2) : 1024
```

```
    NAME(4) : y
```

```
    INTPOS(2) : 1024
```

```
  INTPOS(2) : 1
```

Taxonomy of representations

Reduce's representation is

- **sparse** (only non-zero c are stored)
- **recursive** (polynomials in x whose coefficients are polynomials in the rest)



Reduce *prints* distributed, but stores recursive

- **sparse in variables** (we don't store any intermediate y^0)

Maple's representation is

- **sparse** (only non-zero c are stored)
- **distributed** (multiplied out)

sum sparse in variables (we don't store any intermediate y^0)

poly dense in variables (y^0 stored)

But all of them require some order on the variables to know which is “main”.

How big are they?

$$\underbrace{(a^n + b^n + c^n)}_f * \underbrace{(x^n + y^n)}_g$$

Category	Maple sum	Maple poly	Reduce
Variables	12	5	< 9
Exponents	12	(6)	< 9
Coefficients	6	6	6
Data Structure	$13 + 6 * 5 = 43$	$14 + 6 = 20$	$< 2 * 27 = 54$
Generalised	$7t_f t_g + 1$	$2t_f t_g + t_f + t_g + 3$	$< 6t_f(t_g + 1)$

We can see why [MP14] invented poly!

The $<$ for Reduce is because of sharing of CONS cells.

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