Part 1: Set up recurrences, use the Master Theorem

Consider the following pseudo code sketches. For each, state its recurrence, using T(n) to denote the maximum number of steps the function makes on an input of size n. In particular, (1) next to the pseudo code above, highlight its non-recursive parts and estimate their running time in big-Oh notation, and (2) next to the pseudo code above, highlight the recursive parts and estimate their running times using the function T(). Then, evaluate it using the Master Theorem.

- (a) Function F (A an array of size n): if n=1: return A[1] use O(n) steps and compute arrays B, C, and D of size $\lfloor n/3 \rfloor$ each x=F(B) y=F(C) z=F(D) return x+y+z
 - (i) Set up a recurrence for this pseudo code sketch. Do not worry about the rounding. Do not forget the base case.

$$T(n) = 3T(n/3) + O(n) = 3(3T(n/n) + C \cdot \frac{1}{3}) + Cn = 9T(n/n) + 2cn$$

$$T(n) = 3kT(n/3) + kcn$$

$$T(n) = \begin{cases} O(1) & \text{if } n = 1 \\ 3T(n/3) + O(n) & \text{if } n > 1 \end{cases}$$

(ii) Analyze the recurrence using the Master Theorem.

- (b) Function G (A an array of size n): if n = 1: return A[1]use O(n) steps and compute arrays B and C of size $\lfloor n/3 \rfloor$ each x = F(B) y = F(C)
 - (i) Set up a recurrence for this pseudo code sketch. Do not worry about the rounding. Do not forget the base case.

$$T(n) = \begin{cases} O(n) & \text{if } n = 1 \\ 2T(n/3) + O(n) & \text{if } n > 1 \end{cases}$$

return x - y

(ii) Analyze the recurrence using the Master Theorem.

- (c) Function H (A an array of size n):

 if n = 1: return A[1]use O(n) steps and compute arrays B, C, D, and E of size $\lfloor n/3 \rfloor$ each x = H(B) y = H(C) z = H(D) w = H(E)return x + y + z + w
 - (i) Set up a recurrence for this pseudo code sketch. Do not worry about the rounding. Do not forget the base case.

$$T(n) = \begin{cases} O(1) & \text{if } n = 1 \\ 4T(N_3) + O(n) & \text{if } n \neq 1 \end{cases}$$

(ii) Analyze the recurrence using the Master Theorem.