

# THE 18.821 MATHEMATICS PROJECT LAB REPORT [PROOFS]

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## 1. THEOREMS

### 1.1. Notation.

$t$ : Time  
 $l$ : Path length  
 $s$ : Speed  
 $a_t$ : Tangential acceleration  
 $a_c$ : Centripetal acceleration  
 $\vec{x}$ : Position  
 $\vec{v}$ : Velocity  
 $\vec{a}$ : Acceleration

## 2. COORDINATES

### 2.1. Scalar Calculus.

$$(1) \quad s = \frac{dl}{dt}$$

$$(2) \quad a_t = \frac{ds}{dt}$$

(3)

### 2.2. Vector Calculus.

$$(4) \quad \mathbf{x} = r\hat{\mathbf{r}}$$

$$(5) \quad \vec{v} = \dot{r}\hat{\mathbf{r}} + r\dot{\phi}\hat{\phi}$$

$$(6) \quad \vec{a} = \left(\ddot{r} - r\dot{\phi}^2\right)\hat{\mathbf{r}} + \frac{1}{r}\frac{d}{dt}\left(r^2\dot{\phi}\right)\hat{\phi}$$

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### 2.3. Relations.

$$(7) \quad l = \int_{t=0}^T \|\vec{v}\| \, dt$$

$$(8) \quad s = \vec{v}$$

$$(9) \quad a_t = \|\vec{a}\| \frac{\vec{v}}{\|\vec{v}\|}$$

### 3. TRAVELING BETWEEN POINTS

**Definition 3.1.** A description of a particle  $p$  is a set  $D(t)$  which contains a particle's position  $\vec{x} \in \mathbb{R}^2$  and its derivatives at some time  $t$ . In other words, we define  $D(t) = \{\vec{x}(t), \frac{d\vec{x}}{dt}(t), \frac{d^2\vec{x}}{dt^2}(t), \dots\}$ .

**Definition 3.2.** A condition  $c$  on a particle  $p$  is a boolean function  $c : D \rightarrow \{0, 1\}$  which takes as input a description  $D(t)$  of the particle  $p$  at some time  $t$  and outputs that either the condition is true or false.

**Definition 3.3.** A path  $\gamma : \mathbb{R} \rightarrow \mathbb{R}^2$  is a function which maps some time  $t$  to a position  $\gamma(t) \in \mathbb{R}^2$ . The path is defined from time  $t = 0$  until the end time of the path, denoted as  $T_{f,\gamma}$ .

**Definition 3.4.** A valid path  $\gamma : \mathbb{R} \rightarrow \mathbb{R}^2$  for some particle  $p$  and conditions  $C$  is some path which at all times  $t$  such that  $0 \leq t \leq T_{f,\gamma}$ , all conditions in  $C$  on the particle are satisfied.

**Definition 3.5.** A valid targetted path  $\gamma : \mathbb{R} \rightarrow \mathbb{R}^2$  for some point mass  $p$ , conditions  $C$ , starting point  $\vec{x}_1$ , and ending point  $\vec{x}_2$  is a valid path where  $\gamma(t) = \vec{x}_1$  and  $\gamma(T_f(\gamma)) = \vec{x}_2$ . In other words, it is a valid path which starts at  $\vec{x}_1$  and ends at  $\vec{x}_2$ .

**Definition 3.6.** A fastest path  $\hat{\gamma} : \mathbb{R} \rightarrow \mathbb{R}^2$  for a particular point mass  $p$ , a starting point  $\vec{x}_1$ , a destination point  $\vec{x}_2$ , and some set of conditions  $C$  is a valid targetted path  $\hat{\gamma}$  such that  $T_f(\hat{\gamma}) \leq T_f(\gamma)$  for all valid targetted paths  $\gamma$  with the same  $p$ ,  $\vec{x}_1$ ,  $\vec{x}_2$ , and  $C$ .

**Theorem 3.7.** Given points  $(x_1, y_1), (x_2, y_2) \in \mathbb{R}^2$  and a point mass  $p$  whose initial position is  $(x_1, y_1)$  which moves with acceleration bounded by  $\bar{a}$ , the fastest path  $\hat{\gamma}(t)$  which  $p$  can trace from  $(x_1, y_1)$  to  $(x_2, y_2)$  follows the straight line where all coordinates  $(x, y)$  on the straight line are given by:

$$(10) \quad y = \frac{y_2 - y_1}{x_2 - x_1}x + y_1$$

*Proof.* Let's transform the problem. We can reset our coordinate axes so that  $(x_1, y_1)$  is set to the origin and  $(x_2, y_2)$  is on the  $x$ -axis. In this new coordinate system, we have transformed the following:

$$(11) \quad (x_1, y_1) \rightarrow (0, 0)$$

$$(12) \quad (x_2, y_2) \rightarrow (x'_2, 0)$$

For convenience of notation, we will now refer to  $x'_2$  as  $x_2$ .

Now let us examine the particle's motion in the  $x$  direction. Let  $a_t(t)$  be the tangential acceleration at time  $t$  in the  $x$  direction. Then we can obtain the speed of the particle  $s(t)$  at time  $t$  in the  $x$  direction like so:

$$(13) \quad s(t) = \int_0^t a_t(t_1) dt_1$$

To find the distance  $d(t)$  travelled up to time  $t$  in the  $x$  direction, we can use the relation:

$$(14) \quad d(t) = \int_0^t s(t_2) dt_2$$

$$(15) \quad = \int_0^t \int_0^t a_t(t_1) dt_1 dt_2$$

Recall that the acceleration of the point mass  $p$  is bounded by  $\bar{a}$ . This means that  $a_t(t) \leq \bar{a}$  for all  $t$ . Therefore, we see:

$$(16) \quad d(t) \leq \int_0^t \int_0^t \bar{a} dt_1 dt_2$$

$$(17) \quad = \frac{\bar{a} t^2}{2}$$

Thus, in order to travel a distance of  $d(T_f) = x_2$ , it needs to be the case that  $T_f \geq \sqrt{\frac{2x_2}{\bar{a}}}$ . Moreover, equality holds if and only if  $a_t(t) = \bar{a}$  for all  $t \in [0, T_f(\gamma)]$ .

If the point mass travels for time  $t < \sqrt{\frac{2x_2}{\bar{a}}}$ , then it is impossible for the point mass to reach  $(x_2, 0)$  when starting at  $(0, 0)$ . This is because  $p$  cannot reach  $(x_2, 0)$  in the  $x$  direction when  $t < \sqrt{\frac{2x_2}{\bar{a}}}$  and any acceleration in the  $y$  direction would not enable this either.

This means that the fastest path is completed in time  $T_f(\hat{\gamma}) = \sqrt{\frac{2x_2}{\bar{a}}}$ . Let us examine the path taken by the point mass  $p$  on this fastest path. Recall that  $a_t(t) = \bar{a}$  for all  $t$  along the fastest path. This means that there was no centripetal acceleration  $|a_c| = 0$ . In other words, the point mass never turned on its way to reaching the destination point. The

only way this could have happened is if it travelled along the  $x$  axis in a straight line.

Now, we have seen that the fastest path in the transformed coordinates travels exactly on the  $x$  axis so that  $y = 0$  anywhere along the fastest path. Notice, however, that the  $x$  axis in the transformed coordinates is given exactly by the following line:

$$(18) \quad y = \frac{y_2 - y_1}{x_2 - x_1}x + y_1$$

Thus, we see that the fastest path in the original coordinates follows the above equation, which is what we wanted to show.  $\square$

**Corollary 3.8.** *The fastest path between two points  $(x_1, y_1), (x_2, y_2) \in \mathbb{R}^2$  is unique.*

#### 4. TURNING AROUND A CONE - CONSTANT VELOCITY

**4.1. Quadrilateral.** Before we attempt to tackle the problem of finding the optimal path around a cone at constant velocity, we may want to find some intuition for the problem with a simple, but related problem. In particular, if a particle is traveling around a cone, we would like to know what circle it traces as it goes around the cone.

To make a simple model of this problem, we will use a quadrilateral to model a path around a cone. We will construct a quadrilateral with particular constants that models the constraints of a path around a cone, and we will try to find the parameters that minimize the perimeter around the quadrilateral. Let us call the sides of the quadrilateral  $s_1, s_2, s_3$  and  $s_4$ . Imagine that  $s_1$  and  $s_2$  are on opposite sides of the quadrilateral with fixed lengths  $l_1$  and  $l_2$  respectively. Now we will constrain the problem to be related to the problem of a particle traveling around a cone: imagine  $s_1$  and  $s_2$  are connected by a bar  $s_b$  of length  $d$ . The bar will be perpendicular to  $s_1$ . The angle that  $s_b$  forms with  $s_2$  will be called  $\theta$ . We will try to find the optimal  $\theta$  that minimizes the perimeter around the quadrilateral.

**4.2. Symmetric Quadrilateral.** Let us begin by solving the simplest version of this problem. Let us imagine that  $s_b$  is connected to the midpoints of  $s_1$  and  $s_2$ . We know that the perimeter will be given by  $l_1 + l_2 + l_3 + l_4$  where  $l_3$  and  $l_4$  are the lengths of the sides of  $s_3$  and  $s_4$ , respectively. We are given  $l_1$  and  $l_2$ , but we will need to compute  $l_3$  and  $l_4$  as functions of  $l_1, l_2, d$ , and  $\theta$ .

We can find  $l_3$  by using the fact that it forms a right triangle. We know that  $l_3^2 = m^2 + n^2$ . Finding  $l_4$  is similar. Therefore, we just need to

find  $m$  and  $n$ . This can be done by using the fact that  $m = \frac{l_1}{2} - \frac{l_2}{2} \sin \theta$ . This is just the upper half of  $s_1$  minus the projection of  $s_2$  onto  $s_1$ . We can find  $n$  similarly:  $n = d - \frac{l_2}{2} \cos \theta$ . This is just the bar  $s_b$  minus the projection of  $s_2$  onto the bar. Thus by substituting, we have:

$$(19) \quad l_3 = \sqrt{\left(\frac{l_1}{2} - \frac{l_2}{2} \sin \theta\right)^2 + \left(d - \frac{l_2}{2} \cos \theta\right)^2}$$

We can do a similar analysis on  $l_4$ , only remembering that  $n$  for  $l_4$  gets extended by the projection onto  $s_4$  instead of shrunk. We therefore have:

$$(20) \quad l_4 = \sqrt{\left(\frac{l_1}{2} - \frac{l_2}{2} \sin \theta\right)^2 + \left(d + \frac{l_2}{2} \cos \theta\right)^2}$$

Now, to minimize the perimeter with respect to  $\theta$ , we want to minimize  $l_1 + l_2 + l_3 + l_4$ . Since we know that  $l_1$  and  $l_2$  are fixed, we really want to minimize  $l_3 + l_4$  with respect to  $\theta$ . The other thing to note is that we're minimizing positive distances. We will invoke the following lemma so that we can simplify our expression for  $\min l_3 + l_4$ :

**Lemma 4.1.** *If  $f(t), g(t) > 0$  and  $k(t) > 0$  is a strictly monotonically increasing function for all  $t$ , then  $\arg \min_t k(f(t)) + k(g(t)) = \arg \min_t f(t) + g(t)$ .*

*Proof.* Let  $t_1, t_2$  be such that  $f(t_1) + g(t_1) < f(t_2) + g(t_2)$ . In this proof, we will show that  $k(f(t_1)) + k(g(t_1)) < \sqrt{f(t_2)} + \sqrt{g(t_2)}$ . Since  $k(t)$  is a strictly monotonically increasing function when  $t > 0$ , we know that  $k(x) < k(y)$  if and only if  $x < y$  (assuming we can confine  $x, y$  to be non-negative).

Because this is the case, we see that  $k(f(x)) < k(f(y))$  if and only if  $f(x) < f(y)$  (the same goes for  $g$ ). Thus, we see that if we have found the minimum  $t_m$  to  $\arg \min_t k(f(t)) + k(g(t))$ , then it is the case that  $k(f(t_m)) + k(g(t_m)) < k(f(t)) + k(g(t))$  for all  $t \neq t_m$  (again where  $t > 0$ ). Following our train of logic, we see that  $f(t_m) + g(t_m) < f(t) + g(t)$  for all  $t \neq t_m$ , which means that  $t_m$  is a minimum of  $f(t) + g(t)$ . Thus by finding a minimum  $t_m$  to  $k(f(t)) + k(g(t))$ , we also found a minimum to  $f(t) + g(t)$ .  $\square$

Since we've proven this lemma, we can invoke it upon  $\min l_3 + l_4$ . Since  $l_3 = \sqrt{z_3}$  and  $l_4 = \sqrt{z_4}$ , we can use  $k(t) = \sqrt{t}$  and we can write  $\min l_3 + l_4 = \min z_3 + z_4$  by using our lemma. Thus, we now want to solve the problem:

$$(21) \quad \arg \min_{\theta} 2 \left( \frac{l_1}{2} - \frac{l_2}{2} \sin \theta \right)^2 + \left( d - \frac{l_2}{2} \cos \theta \right)^2 + \left( d + \frac{l_2}{2} \cos \theta \right)^2$$

Now, we can expand out our expression and use the fact that  $\sin^2 \theta + \cos^2 \theta = 1$  to obtain a much simpler (but equivalent) minimization problem:

$$(22) \quad \arg \min_{\theta} \frac{l_1^2}{2} + \frac{l_2^2}{2} + 2d^2 - l_1 l_2 \sin \theta$$

We note that  $l_1, l_2$ , and  $d$  are all constants which are given to us in the problem. Therefore, the minimization problem really boils down to

$$(23) \quad \arg \min_{\theta} -l_1 l_2 \sin \theta = \arg \max_{\theta} \sin \theta$$

Thus, we see that  $\theta = \frac{\pi}{2}$  minimizes the perimeter of the quadrilateral in the symmetric case.

#### REFERENCES

- [1] [http://en.wikipedia.org/wiki/Polar\\_coordinate\\_system](http://en.wikipedia.org/wiki/Polar_coordinate_system)