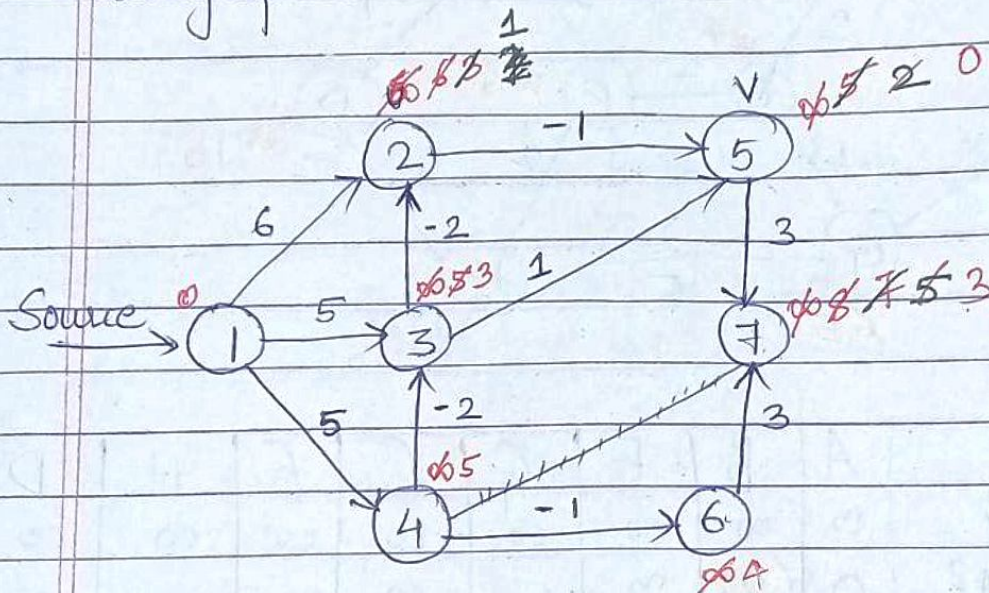


11/2/11

KNOWLEDGE

BELLMAN FORD :-

Bellman-Ford algorithm is an algorithm that computes shortest path from a single source vertex to all of the other vertices in a weighted diagraph.



Relax all edges, $(n-1)$ ie. (Vertices - 1)

if $(d[u] + c(u,v) < d[v])$
 $d[v] = d[u] + c(u,v)$

i) Prepare list of all edges.

$(1,2), (1,3), (1,4), (2,5), (3,2), (3,5), (4,3), (4,6)$
 $(5,7), (6,7)$

for $n=1$.

Initially Source = 1 = 0, all other vertices are ∞

$(1,2)=6$, $(1,3)=5$, $(1,4)=5$, $(2,5)=5$,
 $(3,2)=3$, $(3,5)=$ No modification⁵, $(4,3)=3$, $(4,6)=4$,
 $(5,7)=8$, $(6,7)=7$

Edges	(1,2)	(1,3)	(1,4)	(2,5)	(3,2)	(3,5)	(4,3)	(4,6)	(5,7)	(6,7)
Cost	6	5	5	5	3	5	3	4	8	7

(Red)

for $n=2$.

Edges	(1,2)	(1,3)	(1,4)	(2,5)	(3,2)	(3,5)	(4,3)	(4,6)	(5,7)	(6,7)
Cost	3	3	5	2	1	2	3	4	8 5	7 5

(Black)

for $n=3$.

Edges	(1,2)	(1,3)	(1,4)	(2,5)	(3,2)	(3,5)	(4,3)	(4,6)	(5,7)	(6,7)
Cost	1	3	5	0	1	0	3	4	3	3

(Red)

for $n=4$

Edges	(1,2)	(1,3)	(1,4)	(2,5)	(3,2)	(3,5)	(4,3)	(4,6)	(5,7)	(6,7)
Cost	1	3	5	0	1	0	3	4	3	3

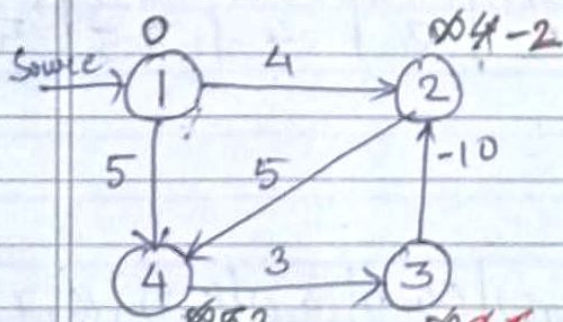
Now, there is no change for $n=3$ and $n=4$, so stop here.

1/8/11

Path cost :-

	Vertex : Cost
$(1,2) = 1$	$1 = 0$
$(1,3) = 3$	$2 = 1$
$(1,4) = 5$	$3 = 3$
$(2,5) = 0$	$4 = 5$
$(3,2) = 1$	$5 = 0$
$(3,5) = 0$	$6 = 4$
$(4,3) = 3$	$7 = 3$
$(4,6) = 4$	
$(5,7) = 3$	
$(6,7) = 3$	

DRAWBACK OF BELLMAN FORD ALGORITHM



for $n=1$

$(3,2)$	$(4,3)$	$(1,4)$	$(1,2)$	$(2,4)$
∞	∞	5	4	5

for $n=2$

$(3,2)$	$(4,3)$	$(1,4)$	$(1,2)$	$(2,4)$
∞	8	5	4	5

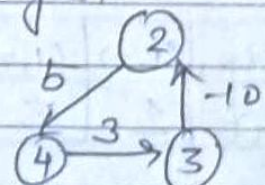
for $n=3$

$(3,2)$	$(4,3)$	$(1,4)$	$(1,2)$	$(2,4)$
-2	8	5	-2	3

for $n=4$

$(3,2)$	$(4,3)$	$(1,4)$	$(1,2)$	$(2,4)$
-2	6	3	-2	3

if we try $(n=4)$ it will also change, as there is cycle of edges,



$$5 + 3 - 10 = -2$$

then cycle is -ve, It is disadvantage. Total weight is -ve so, not work.